Course Notes

CSCA67 - Discrete Mathematics

Instructors:

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1 Propositions, Implications

Definitions:

A **proposition** is a statement that evaluates to True or False. In computer science, its often referred to as a **Boolean expression**.

A **compound roposition** is a proposition statementt that involves multiple propositions joined by connectives. It takes multiple truth values as input and returns a single truth value as output.

A **connective** corresponds to English conjunctions such as "and", "or", "not" etc.

Basic connectives and truth tables:

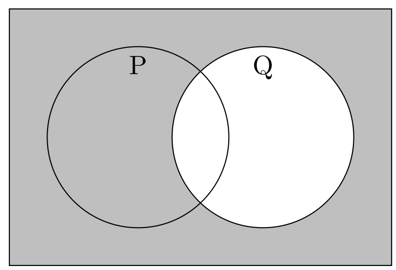
Symbol	Meaning
$\wedge$	"AND"
$\vee$	"OR"
$\rightarrow$	"IF...THEN"
$\leftrightarrow$	"IF AND ONLY IF"
$\neg$	"NOT"

P	Q	$P \wedge Q$	$P \vee Q$	$P \rightarrow Q$	$P \leftrightarrow Q$
T	T	T	T	T	T
T	F	F	T	F	F
F	T	F	T	T	F
F	F	F	F	T	T

Implication:

Different ways of writing $P \rightarrow Q$:

1. If P then Q
2. If P, Q
3. Q, if P
4. P only if Q
5. P is sufficient for Q
6. Q is neccesary for P
7. If not Q, then not P
8. Not P or Q



Logical Equivalences:

Commutative	$p \wedge q \iff q \wedge p$	$p \vee q \iff q \vee p$
Associative	$(p \wedge q) \wedge r \iff p \wedge (q \wedge r)$	$(p \vee q) \vee r \iff p \vee (q \vee r)$
Distributive	$p \wedge (q \vee r) \iff (p \wedge q) \vee (p \wedge r)$	$p \vee (q \wedge r) \iff (p \vee q) \wedge (p \vee r)$
Identity	$p \wedge T \iff p$	$p \vee F \iff p$
Negation	$p \vee \neg p \iff T$	$p \wedge \neg p \iff F$
Double Negative	$\neg(\neg p) \iff p$	
Idempotent	$p \wedge p \iff p$	$p \vee p \iff p$
Universal Bound	$p \vee T \iff T$	$p \wedge F \iff F$
De Morgan's	$\neg(p \wedge q) \iff (\neg p) \vee (\neg q)$	$\neg(p \vee q) \iff (\neg p) \wedge (\neg q)$
Absorption	$p \vee (p \wedge q) \iff p$	$p \wedge (p \vee q) \iff p$
Conditional or ($\rightarrow$) Law	$(p \rightarrow q) \iff (\neg p \vee q)$	$\neg(p \rightarrow q) \iff (p \wedge \neg q)$
Biconditional	$(p \leftrightarrow q) \iff (p \rightarrow q) \wedge (q \rightarrow p)$	

Order of Operations:

1. NOT($\neg$) 2. AND($\wedge$) 3. OR($\vee$) 4. Quantifiers($\forall$/$\exists$) 5. ($\rightarrow$ / $\leftrightarrow$)

2 Predicates and Quantifiers

Forall:	$\forall$
There exists:	$\exists$
Negations:	
$\neg\forall = \exists$	$\neg\exists = \forall$

Prove statement in the form of $\exists x \in S, \ni: P(x)$
We simply need to find **one** value of x in the set S, that makes $P(x)$ true.
One value is enough.

Example:
There exists an integer n, such that n^2 is even.
$\exists n \in \mathbb{Z}, \ni: n^2 \in 2\mathbb{Z}$
Let $n = 2$, then $(2)^2 = 4$ which is an even number

Prove statemnet in the form of $\forall x \in S, \ni: P(x)$
This means we must use techniques such as algebraic manipulation to show that:
$P(x)$ holds for every arbitrary $x \in S$

Example:
Forall integers n, if n is odd, then n^2 is odd.
$\forall n \in \mathbb{Z}, n \in 2\mathbb{Z} \rightarrow n^2 \in 2\mathbb{Z}$
Let $n = 2k, k \in \mathbb{Z}$
then $n^2 = (2k)^2 = 4k^2$ which is an even number
Therefore: Forall integers n, if n is odd, then n^2 is odd. QED

2.1 Modulus