CSCC63 TUT 0002

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Proving decision problems are in NP/co-NP

Show that L is in NP

- 1. Show that L is a decision problem
- 2. Provide a certificate for the verifier
- 3. Justify why the certificate is polynomial in terms of the size of the input
- 4. Provide a Verifier that runs on the input/certificate
- 5. Justify the Verifier halts and verifies solutions in polynomial time

Prove that NP is closed under the concatenation operation

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Prove that if L1, L2 are in NP, then L1 + L2 is in NP L1 + L2 = \{x \mid x = yz, y \text{ in L1}, z \text{ in L2}\}
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Assume that L1, and L2 are in NP

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L1 = { x | exists c (certificate), |c| polynomial in terms of |x|, V1(x, c) accepts in O(|x|^k) } L2 = { x | exists c (certificate), |c| polynomial in terms of |x|, V2(x, c) accepts in O(|x|^j) }
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- 1. L1 + L2 is a Decision problem, either elements belong to the set or it doesn't
- 2. Let x be the input, C = (i, c1, c2) such that c1 verifies x[0:i], c2 verifies x[i:]
- 3. |C| = |c1| + |c2|, by assumptions, C is polynomial in terms of |x|.
- 4. define V on input <x, C = (i, c1, c2)>: assert i in [0, |x|] # O(|x|) assert V1(x[0:i], c1) and V2(x[i:], c2) both accept # O(|x|^max(k, j)) if both all assertions pass accept, otherwise reject
- 5. All assertions are done in polynomial time, since V1 and V2 run in polynomial time.

Prove that co-NP is closed under the concatenation operation

Prove that if L1, L2 are in co-NP, then L1 + L2 is in co-NP $L1 + L2 = \{ x \mid exists y in L1, z in L2, such that x = y + z \}$ $co-(L1 + L2) = \{ x \mid forall y in L1, z in L2, such that x != y + z \}$ $co-(L1 + L2) = \{ x \mid for i in [0, |x|], x[0:i] in co-L1 or x[i:] in co-L2 \}$ Assume that L1, and L2 are in co-NP co-L1 = $\{x \mid exists c (certificate), |c| polynomial in terms of |x|, V1(x, c) accepts in O(|x|^k) \}$ co-L2 = $\{x \mid exists c (certificate), |c| polynomial in terms of |x|, V2(x, c) accepts in O(|x|^i) \}$ 1. This is a decision problem, either an element is in L1 + L2 or its not. 2. Let x be the input, let $C = [C_0, C_1, ... C_x]$, each $C_i = (i, c_i, c_i)$, i in [0, x]3. |C| = |x|, we know each $|C_i|$ is polynomial from the previou question say $O(|x|^m)$, then we have $O(|x| * |x|^m)$ 4. define V on input <x, C = [C_0, C_1, ... C_|x|]>: for each $C_i = (i, c1, c2)$ in C: # O(|C|) = O(|x|)assert that i in [0, |x|] # O(|x|)assert that atleast one of V1(x[0:i], c1) or V2(x[i:], c2) accepts # O(n^c) if any assertions fail, reject

accept # we accept within $O(|x| * (|x| + n^c))$