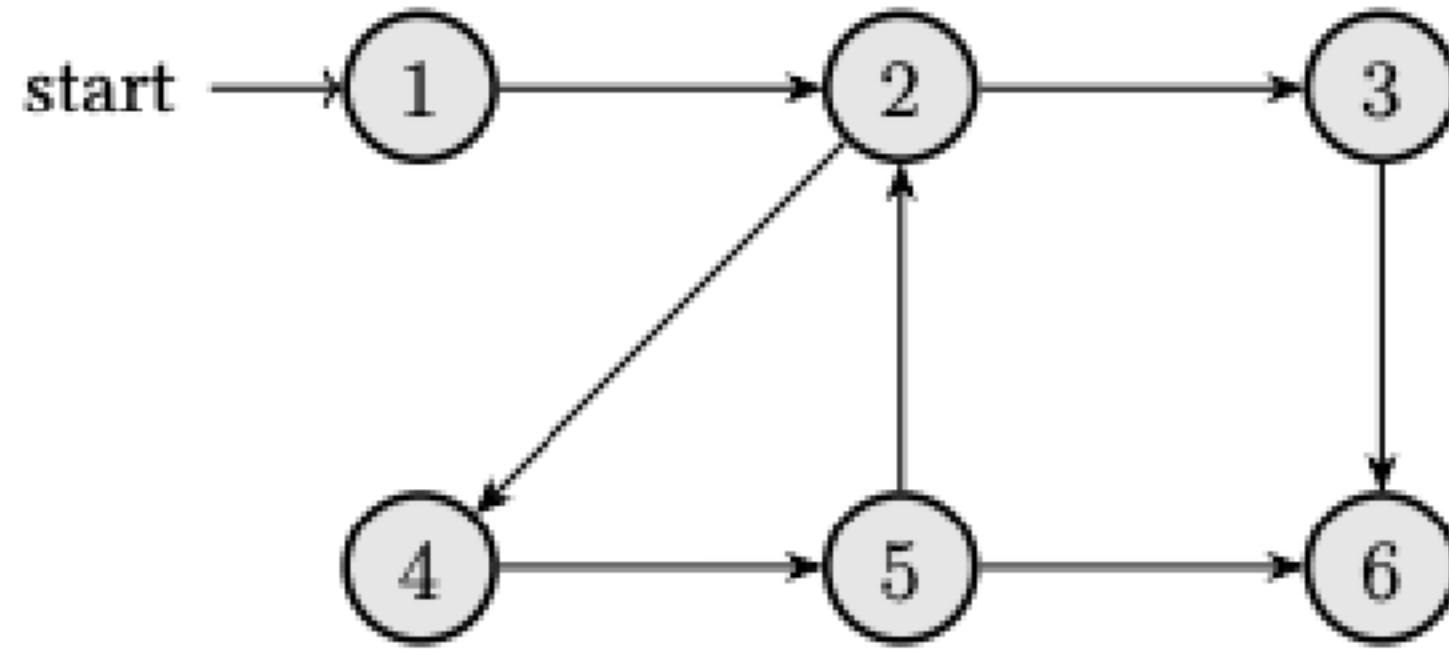


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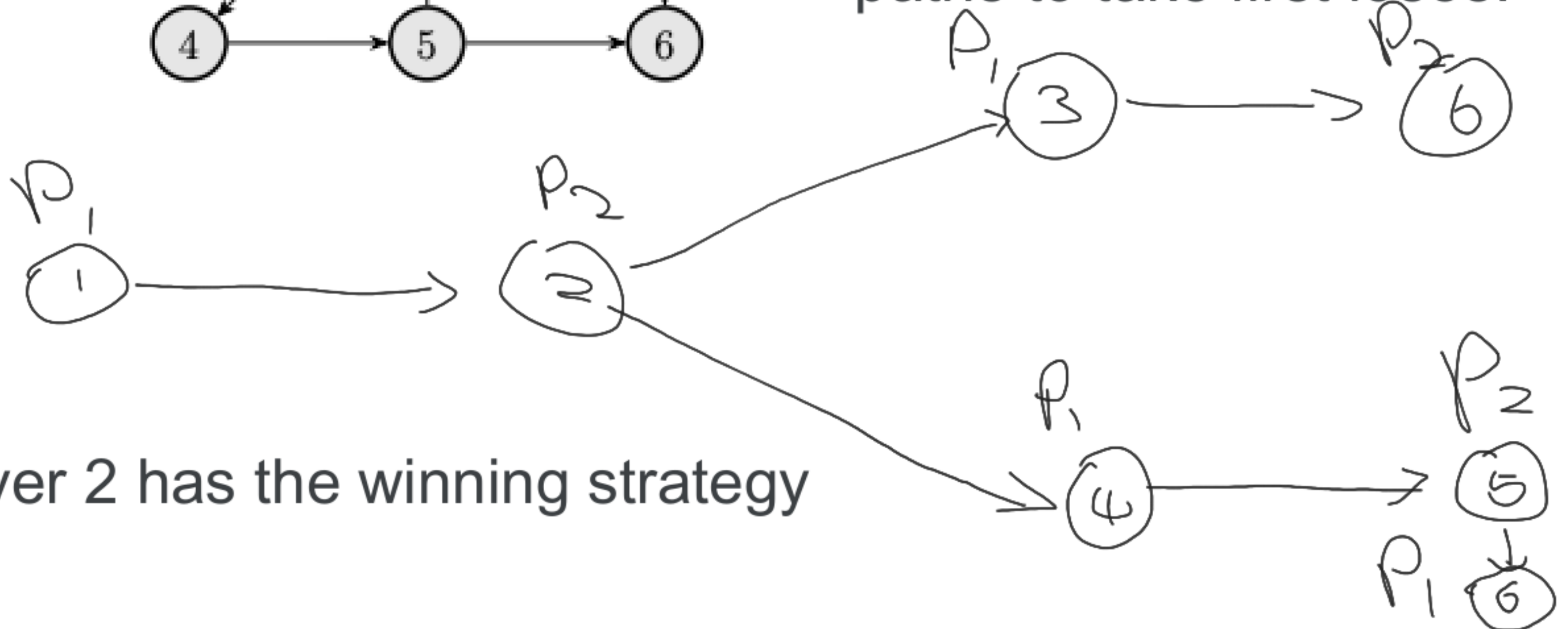
Tutorial 9

Generalized Geography Example
NP Oracles

1. Generalized Geography Game



Player 1 goes first,
whoever runs out of
paths to take first loses.



Player 2 has the winning strategy

Clique = $\{ \langle G, k \rangle \mid G \text{ is an undirected graph with a } k \text{ clique} \}$

Max Clique:

Input: An undirected graph G

Output: Largest clique of G (set of nodes that make up the max-clique)

WTS given oracle for CLIQUE, we can solve MAX-CLIQUE in polytime and vice versa.

Suppose we have an oracle solver SOLVE-MAX-CLIQUE (takes input G , returns set of nodes) that solves Max-Clique in one step, write a function that solves Clique in polytime.

P on input $\langle G = (V, E), k \rangle$:

$\text{max} = |\text{SOLVE-MAX-CLIQUE}(G)| \# O(|V|)$

accept if $\text{max} \geq k$, otherwise reject $\# O(1)$

Suppose we have an oracle solver SOLVE-CLIQUE (takes input G , number k , returns whether G has a k sized CLIQUE) that solves Clique in one step, write a function that solves Max-Clique in polytime.

```
P on input  $\langle G=(V, E) \rangle$ : # We want to return some set of nodes that form the max clique
max = 0
for i = |V| ... 0: #  $O(|V|)$ 
    if SOLVE-CLIQUE( $\langle G, i \rangle$ ):
        max = i
        break
let  $V'$  be a copy of  $V$  #  $O(|V|)$  duplicate the set of nodes
let  $G' = (V', E')$  be a copy of  $G$ 
for v in  $V'$ : #  $O(|V|)$ 
    let  $V'' = V' - \{ v \}$  #  $O(|V|)$ 
    let  $E'' = E' - \{ (u', v') \mid u' = v \mid v' = v \}$  #  $O(|E|)$ 
    if SOLVE-CLIQUE( $\langle V'', E'' \rangle, \text{max} \rangle$ ): #  $O(1)$ 
         $V' = V''$ 
         $E' = E''$ 
return  $V'$ 
```

Q2 Help:

Multiset: set that contains duplicates, your $\text{sum}(S')$ will probably take in the form of $a + a^2$

Consider 1-2 Subset

Input: A multiset S

Question: Whether or not S can be divided into subsets A and B where $2 \text{sum}(A) = \text{sum}(B)$

Consider the following reduction from subset sum

P on input $\langle S, t \rangle$:

return $S' = S \cup \{ \text{sum}(S), 3t + \text{sum}(S) \}$ # $\text{sum}(S')$ takes the $a + 2a$

Suppose that $\langle S, t \rangle$ in Subset-Sum

there must exist S' subset S such that $\text{sum}(S') = t$

consider $A = S' \cup \{ \text{sum}(S) \}$, $B = (S - S') \cup \{ 3t + \text{sum}(S) \}$

$\text{sum}(A) = \text{sum}(S') + \text{sum}(S) = t + \text{sum}(S)$

$\text{sum}(B) = \text{sum}(S) - \text{sum}(S') + (3t + \text{sum}(S)) = \text{sum}(S) - t + 3t + \text{sum}(S) = 2t + 2\text{sum}(S)$

Suppose that S' in 1-2 Subset (How come $\langle S, t \rangle$ has to be in subset sum?)

Q1 Help:

Colorability

We say that a graph is k -Colorable, if every node can be colored with one of the k colors with no two neighbouring nodes sharing the same color.

Example: $k = 3$

