

# COA Assignment U1

1. a) Explain the logical Micro operations.
- b) Discuss about Arithmetic Logic Shift Unit.

## 1. a) Logic Micro-operations

- Logic micro-operations specify binary operations for bit strings in registers.
- Each bit in the register is treated as a binary variable.
- Example: Exclusive-OR micro-operation between registers R1 and R2 when the control variable P = 1.

TABLE 4-5 Truth Tables for 16 Functions of Two Variables

$x$	$y$	$F_0$	$F_1$	$F_2$	$F_3$	$F_4$	$F_5$	$F_6$	$F_7$	$F_8$	$F_9$	$F_{10}$	$F_{11}$	$F_{12}$	$F_{13}$	$F_{14}$	$F_{15}$
0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1	
0	1	0	0	0	0	1	1	1	1	0	0	0	0	1	1	1	
1	0	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	
1	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1	1	

## List of Logic Microoperations

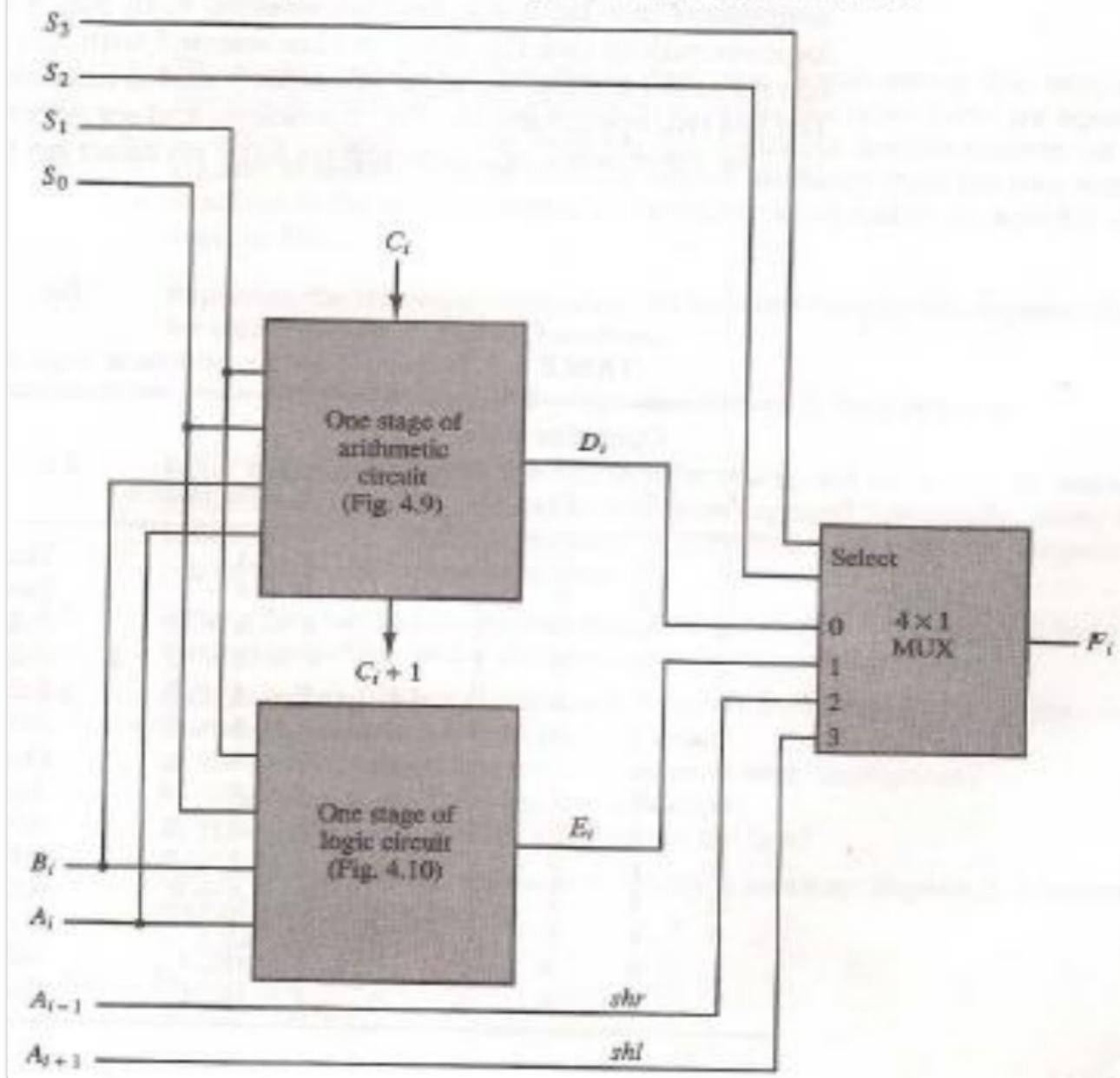
- 16 different logic operations for two binary variables.
- Derived from truth tables with two binary variables (see Table 4-5).
- Expressions in Table 4-6 show Boolean functions for variables x and y.
- Logic micro-operations derived by replacing x with content of register A and y with content of register B.

TABLE 4-6 Sixteen Logic Microoperations

Boolean function	Microoperation	Name
$F_0 = 0$	$F \leftarrow 0$	Clear
$F_1 = xy$	$F \leftarrow A \wedge B$	AND
$F_2 = xy'$	$F \leftarrow A \wedge \bar{B}$	
$F_3 = x$	$F \leftarrow A$	Transfer A
$F_4 = x'y$	$F \leftarrow \bar{A} \wedge B$	
$F_5 = y$	$F \leftarrow B$	Transfer B
$F_6 = x \oplus y$	$F \leftarrow A \oplus B$	Exclusive-OR
$F_7 = x + y$	$F \leftarrow A \vee B$	OR
$F_8 = (x + y)'$	$F \leftarrow \overline{A \vee B}$	NOR
$F_9 = (x \oplus y)'$	$F \leftarrow \overline{A \oplus B}$	Exclusive-NOR
$F_{10} = y'$	$F \leftarrow \bar{B}$	Complement B
$F_{11} = x + y'$	$F \leftarrow A \vee \bar{B}$	
$F_{12} = x'$	$F \leftarrow \bar{A}$	Complement A
$F_{13} = x' + y$	$F \leftarrow \bar{A} \vee B$	
$F_{14} = (xy)'$	$F \leftarrow \overline{A \wedge B}$	NAND
$F_{15} = 1$	$F \leftarrow \text{all } 1\text{'s}$	Set to all 1's

## 1. b) Arithmetic Logic Shift Unit

Figure 4-13 One stage of arithmetic logic shift unit.



- Computer systems use an Arithmetic Logic Unit (ALU) for microoperations.
- ALU is a combinational circuit connected to storage registers.
- Shift microoperations can be part of the ALU.
- One stage of an arithmetic logic shift unit shown in Fig. 4-13.
- Selection variables  $S_1, S_0, S_3, S_2$ , and  $Cin$  determine the operation.
- Table 4-8 lists 14 operations of the ALU (8 arithmetic, 4 logic, 2 shift).

**TABLE 4-8** Function Table for Arithmetic Logic Shift Unit

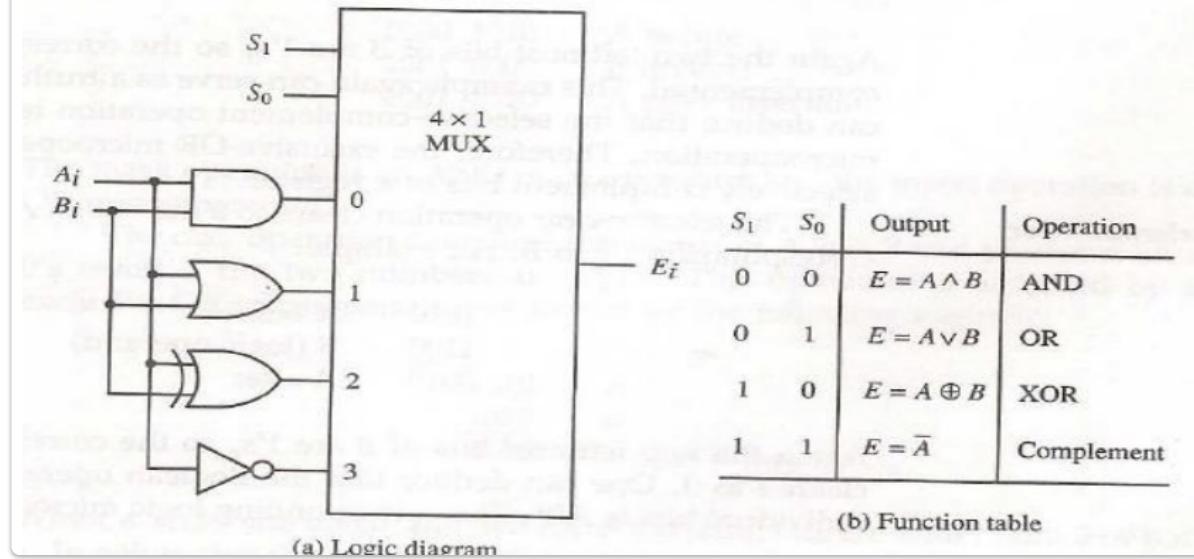
Operation select					Operation	Function
$S_3$	$S_2$	$S_1$	$S_0$	$C_{in}$		
0	0	0	0	0	$F = A$	Transfer $A$
0	0	0	0	1	$F = A + 1$	Increment $A$
0	0	0	1	0	$F = A + B$	Addition
0	0	0	1	1	$F = A + B + 1$	Add with carry
0	0	1	0	0	$F = A + \bar{B}$	Subtract with borrow
0	0	1	0	1	$F = A + \bar{B} + 1$	Subtraction
0	0	1	1	0	$F = A - 1$	Decrement $A$
0	0	1	1	1	$F = A$	Transfer $A$
0	1	0	0	$\times$	$F = A \wedge B$	AND
0	1	0	1	$\times$	$F = A \vee B$	OR
0	1	1	0	$\times$	$F = A \oplus B$	XOR
0	1	1	1	$\times$	$F = \bar{A}$	Complement $A$
1	0	$\times$	$\times$	$\times$	$F = \text{shr } A$	Shift right $A$ into $F$
1	1	$\times$	$\times$	$\times$	$F = \text{shl } A$	Shift left $A$ into $F$

2. a) Explain the hardware implementation of Logic Micro operations.  
 b) Draw the instruction code Format.

## 2. a) Hardware Implementation

- Logic microoperations hardware implementation uses logic gates for each bit or pair of bits in registers.
- Most computers use four basic logic microoperations: AND, OR, XOR (exclusive-OR), and complement.
- Circuit example for generating basic logic microoperations using gates and a multiplexer.

Figure 4-10 One stage of logic circuit.



## Some Applications

- Logic micro-operations useful for manipulating individual bits or portions of a word in a register.
- Examples of applications:
  - Selective set
  - Selective complement
  - Selective clear
  - Mask
  - Insert
  - Clear

### Selective Set

- Sets to 1 the bits in register A where there are corresponding 1's in register B.
- OR microoperation used for selective setting.

### Selective Complement

- Complements bits in A where there are corresponding 1's in B.
- Exclusive-OR microoperation used for selective complementing.

## Selective Clear

- Clears to 0 the bits in A where there are corresponding 1's in B.
- Corresponding logic microoperation is AND.

## Mask

- Similar to selective clear, but clears bits in A only where there are corresponding 0's in B.
- Mask operation is an AND microoperation.

## Insert

- Inserts a new value into a group of bits by masking unwanted bits and ORing with the required value.
- Example provided for inserting a value into a register.

## Clear

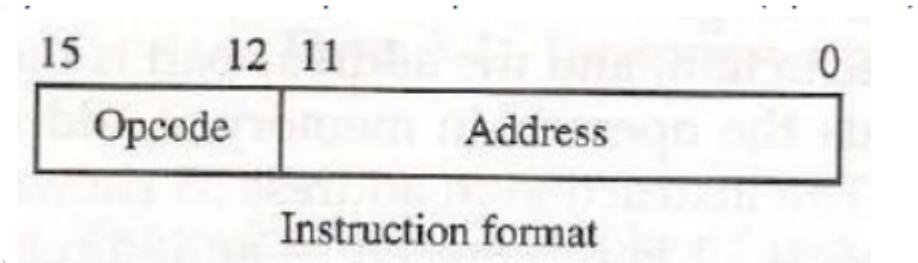
- Compares words in A and B, produces all 0's result if the two numbers are equal.
- Achieved by an exclusive-OR microoperation.

## 2. b) Instruction Codes

### Instruction Codes

- The organization of the computer is defined by its internal registers, timing and control structure, and the set of instructions.
- Instructions are stored in memory and executed based on their opcode and operand.

### Instruction Code Format

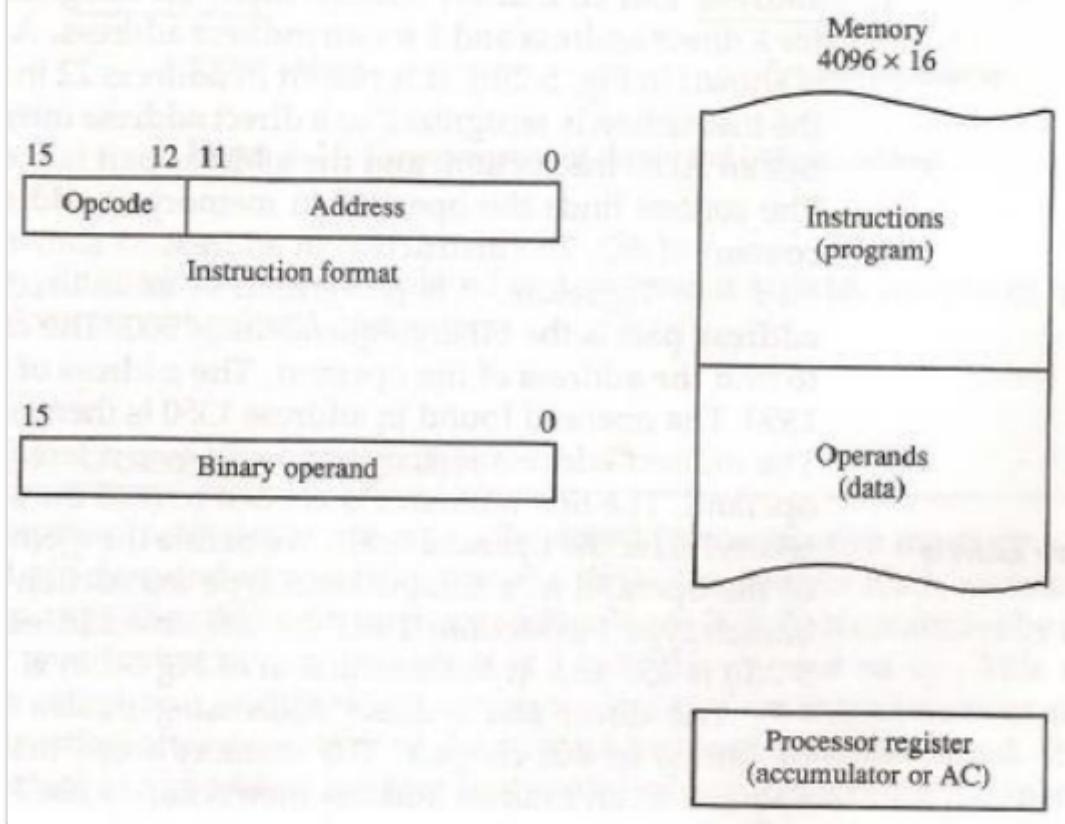


- Instruction codes are divided into two parts: Opcode and Address (Operand).
- Operation Code (Opcode):**
  - Specifies the operation (add, subtract, etc.).
  - Number of bits required depends on the number of operations.
- Address (Operand):**
  - Specifies the location of operands (registers or memory words).

## Stored Program Organization

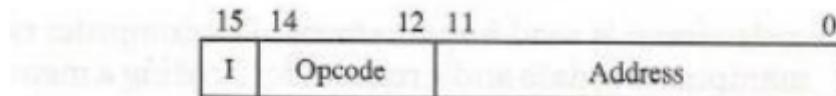
- Stored program organization allows storing instructions in memory and executing them sequentially.
- Example instruction format: 16-bit word with 3-bit opcode and 12-bit address.

Figure 5-1 Stored program organization.

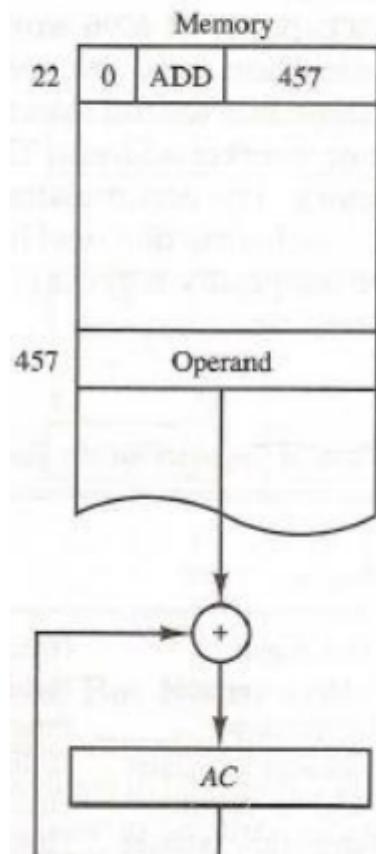


## Addressing of Operand

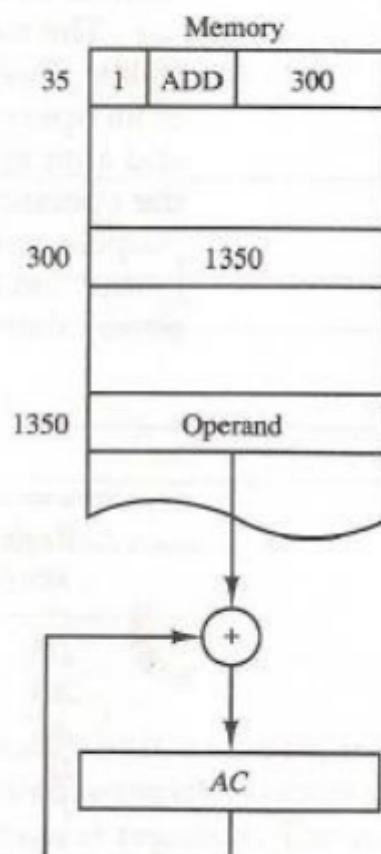
- **Immediate Operand:** Operand is part of the instruction.
- **Direct Address:** Operand is specified by the address part of the instruction.
- **Indirect Address:** Address of the operand is found in another memory location.
- **Addressing modes** are determined by a mode bit (0 for direct, 1 for indirect).



(a) Instruction format



(b) Direct address



(c) Indirect address

Figure 5-2 Demonstration of direct and indirect address.

3. a) Explain the computer registers.
- b) Explain the timing and control signal with help of timing diagram

3. a)

## Computer Registers

- Registers play a crucial role in instruction sequencing and data manipulation.

- Common registers include Accumulator (AC), Instruction Register (IR), Program Counter (PC), etc.

TABLE 5-1 List of Registers for the Basic Computer

Register symbol	Number of bits	Register name	Function
<i>DR</i>	16	Data register	Holds memory operand
<i>AR</i>	12	Address register	Holds address for memory
<i>AC</i>	16	Accumulator	Processor register
<i>IR</i>	16	Instruction register	Holds instruction code
<i>PC</i>	12	Program counter	Holds address of instruction
<i>TR</i>	16	Temporary register	Holds temporary data
<i>INPR</i>	8	Input register	Holds input character
<i>OUTR</i>	8	Output register	Holds output character

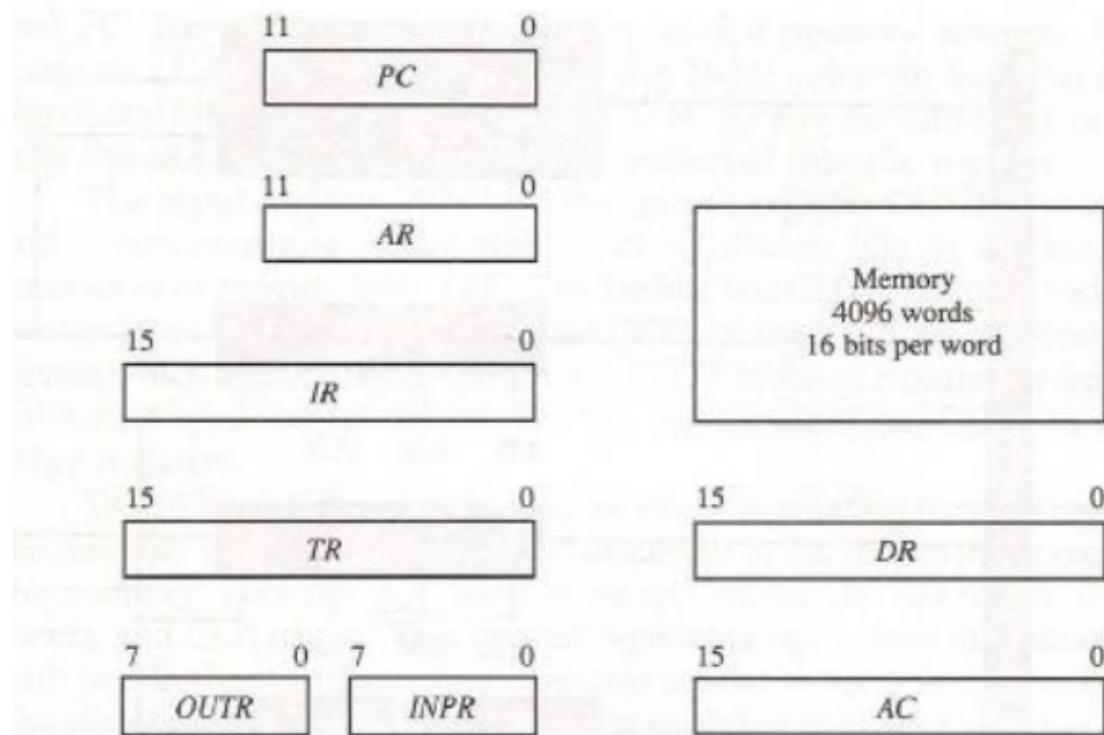


Figure 5-3 Basic computer registers and memory.

## Common Bus System

- A common bus system efficiently transfers information between registers and memory.
- Selection variables determine which register's output is connected to the bus.

- Load (LD), Increment (INR), and Clear (CLR) control inputs are used for register operations.

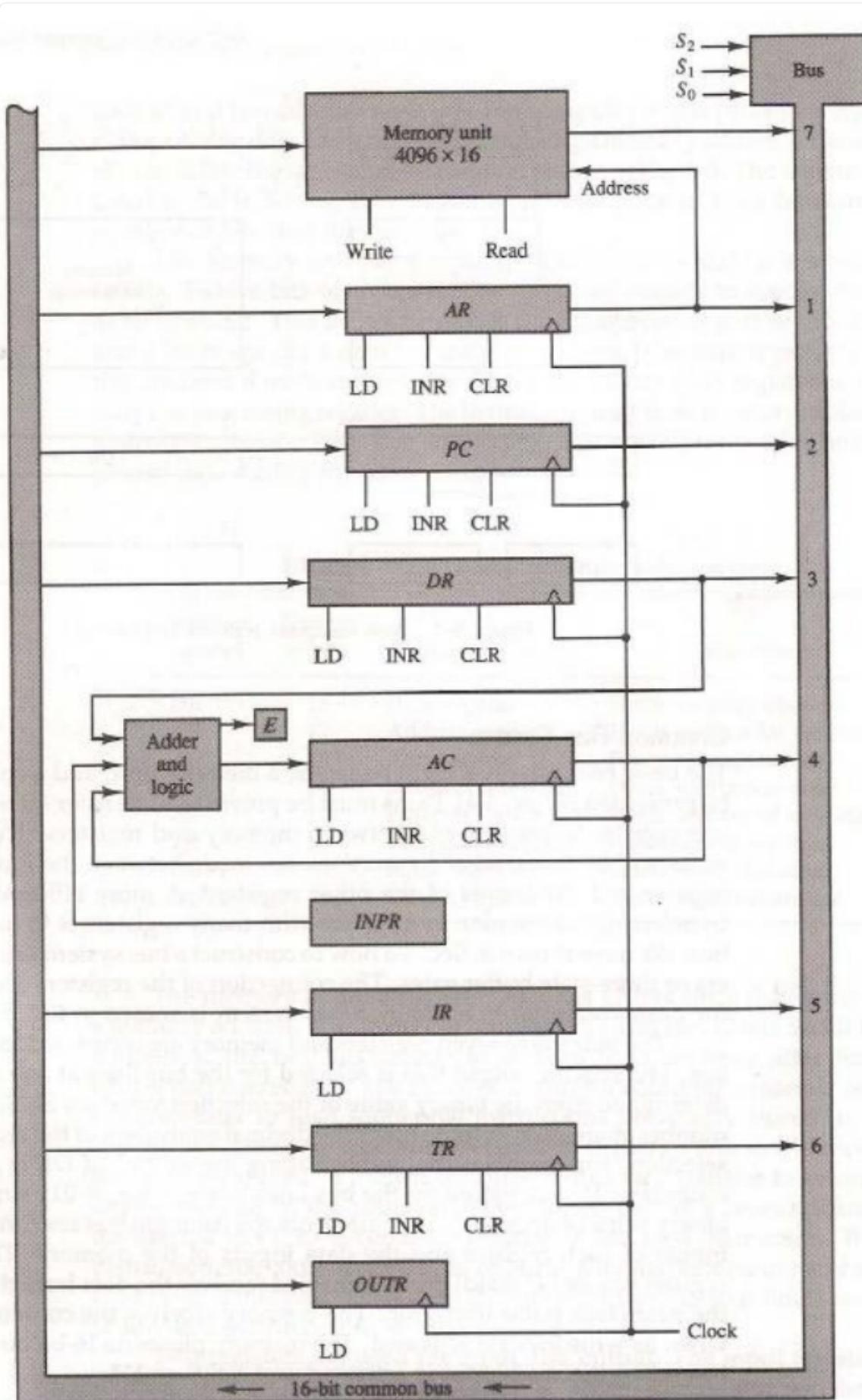


Figure 5-4 Basic computer registers connected to a common bus.

## Example Registers

- Example registers include Data Register (DR), Temporary Register (TR), Memory Address Register (AR), etc.
- Input and output registers (INPR and OUTR) communicate with external devices.

## Adder and Logic Circuit

- The Adder and Logic Circuit perform arithmetic operations and logical functions.
- Multiple inputs, including content from registers, enable versatile operations.
- Operations can be executed in parallel during the same clock cycle.

### 3. b)

## Timing and Control

- The timing for all registers in the basic computer is controlled by a master clock generator.
- Clock pulses are applied to all flip-flops and registers in the system when enabled by a control signal.
- Control signals, generated in the control unit, control multiplexers, processor registers, and microoperations.

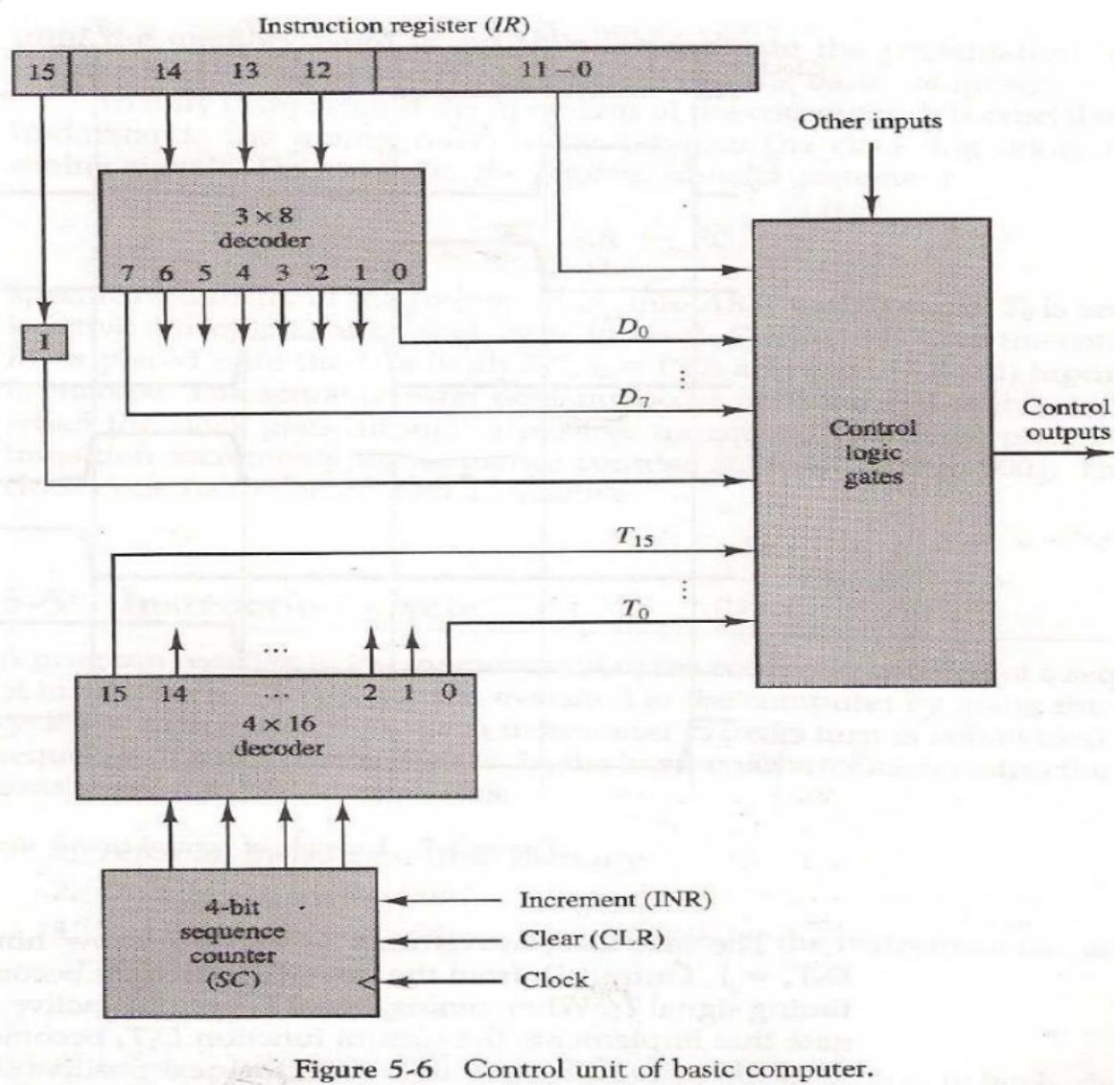


Figure 5-6 Control unit of basic computer.

## Control Organizations

### Hardwired Control

- Implemented with gates, flip-flops, decoders, and other digital circuits.
- Control information is stored in the control memory, programmed for the required sequence of microoperations.
- Advantages include optimization for faster operation.
- Changes in design require modifications in wiring among components.

### Microprogrammed Control

- Control logic is stored in control memory.

- The control memory is programmed to initiate the required sequence of microoperations.
- Offers flexibility for modifications by updating the microprogram.

## Hardwired Control Unit

- Block diagram includes two decoders, a sequence counter, and control logic gates.
- Instruction from memory is placed in the instruction register (IR), divided into I bit, operation code, and bits 0-11.
- Operation code (bits 12-14) is decoded with a 3x8 decoder, producing outputs D0 through D7.
- Bit 15 of the instruction is transferred to a flip-flop (I).
- Bits 0-11 are applied to control logic gates.
- Sequence counter (SC) can count in binary from 0 through 15.
- Outputs of SC are decoded into 16 timing signals (T0-T15).
- SC can be incremented or cleared synchronously.
- Timing diagram shows the time relationship of control signals.

### Timing Diagram

- SC responds to the positive transition of the clock.
- Initially, the CLR input of SC is active.
- SC is incremented with every positive clock transition unless CLR input is active.
- The timing diagram illustrates the sequence of timing signals and the clearing of SC based on control signals.

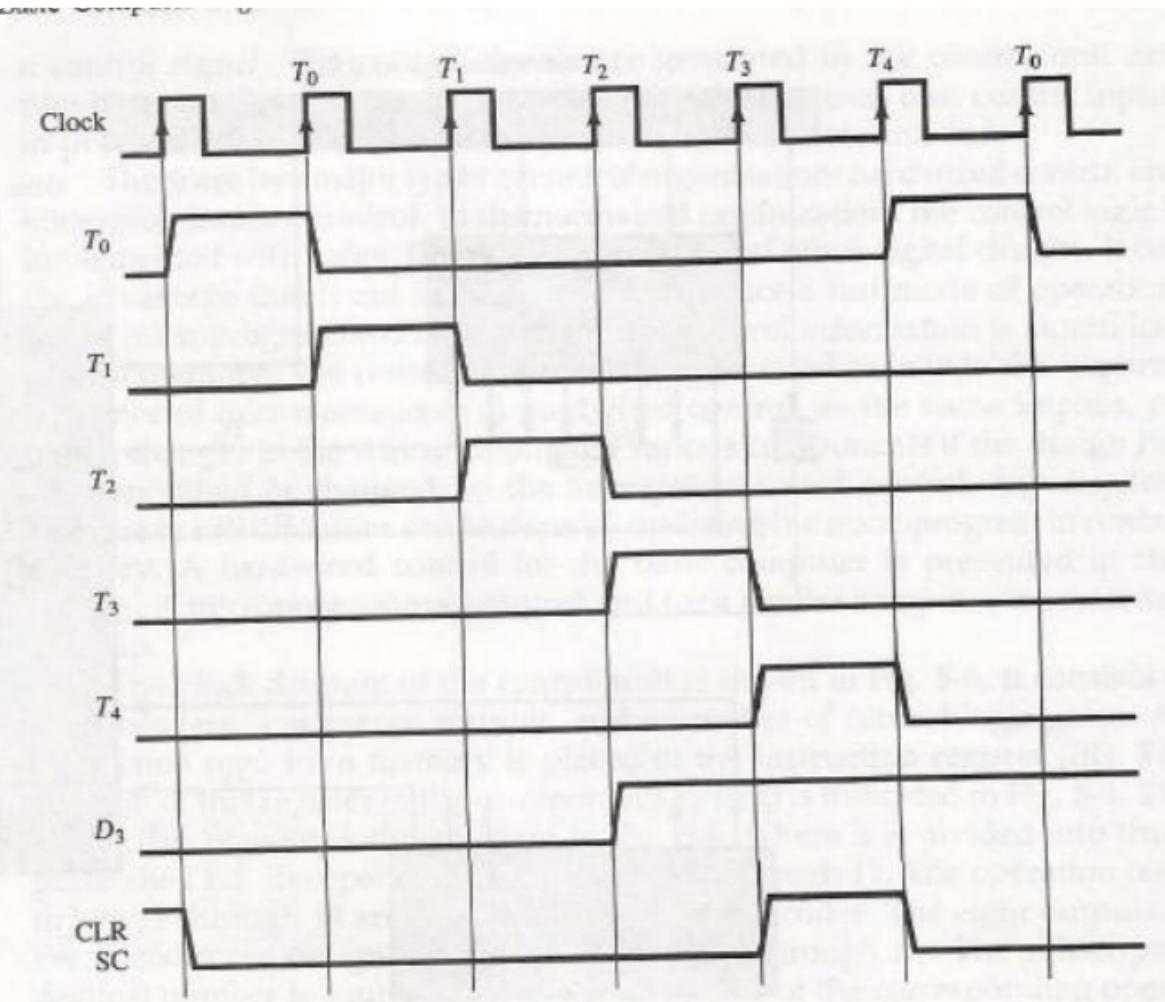


Figure 5-7 Example of control timing signals.