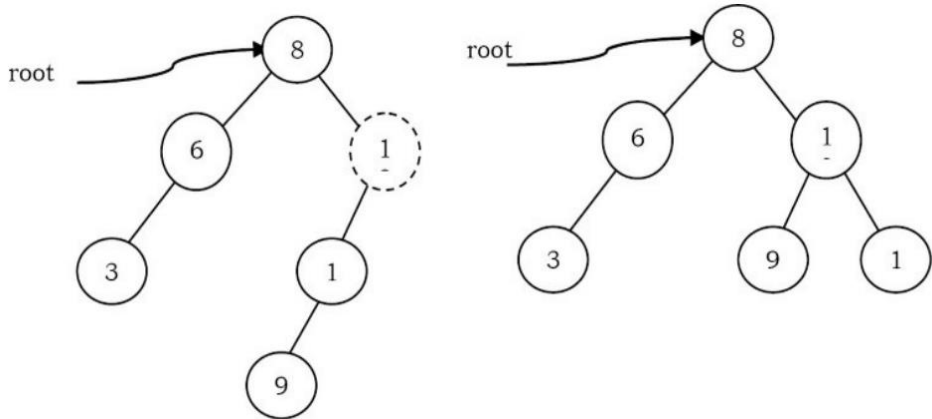


| | |
|----------------|-------------------|
| Name | Hatim Yusuf Sawai |
| UID no. | 2021300108 |
| Experiment No. | 7 |

| | |
|------|---------------------|
| AIM: | Program on AVL Tree |
|------|---------------------|

Program 1

| | |
|--------------------|--|
| PROBLEM STATEMENT: | Write a program to demonstrate insertion into an avl tree using single & double rotations |
|--------------------|--|

| | |
|---------|---|
| THEORY: | <p>AVL (Adelson-Velskii & Landis) Tree:</p> <p>AVL tree is a self-balancing Binary Search Tree (BST) where the difference between heights of left and right subtrees cannot be more than 1 for all nodes. That means, the balancing factor should be < 2 or > -2 for every node of the tree.</p> <p>Example:</p> <p>From below binary search trees, the left one is not an AVL tree, whereas the right binary search tree is an AVL tree:</p>  |
|---------|---|

Advantages of AVL Tree:

Most **BST** operations (insertion, deletion or traversal) take $O(h)$ time where h is the height of the BST. If we make sure that the height of the tree remains $O(\log(n))$ after every insertion and deletion, then we can guarantee an upper bound of $O(\log(n))$ for all these operations. The height of an **AVL** tree is always $O(\log(n))$ where n is the number of nodes in the tree.

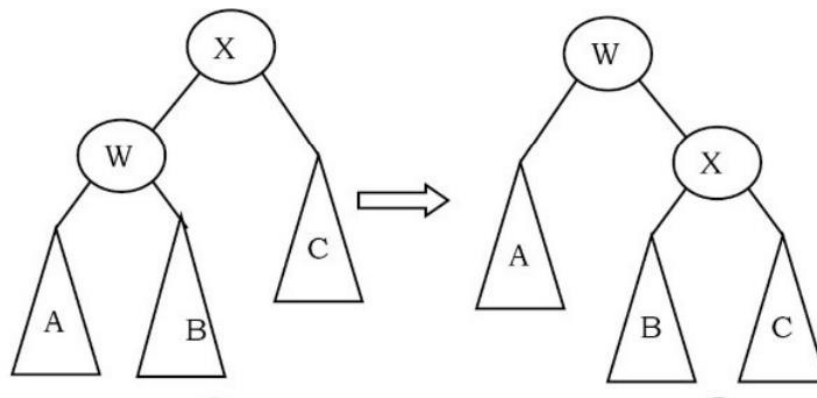
Rotations in AVL Tree:

If the AVL tree property is violated at a node X , it means that the heights of $\text{left}(X)$ and $\text{right}(X)$ differ by exactly 2. This is because, if we balance the AVL tree every time, then at any point, the difference in heights of $\text{left}(X)$ and $\text{right}(X)$ differ by exactly 2. Rotations is the technique used for restoring the AVL tree property. This means, we need to apply the rotations for the node X .

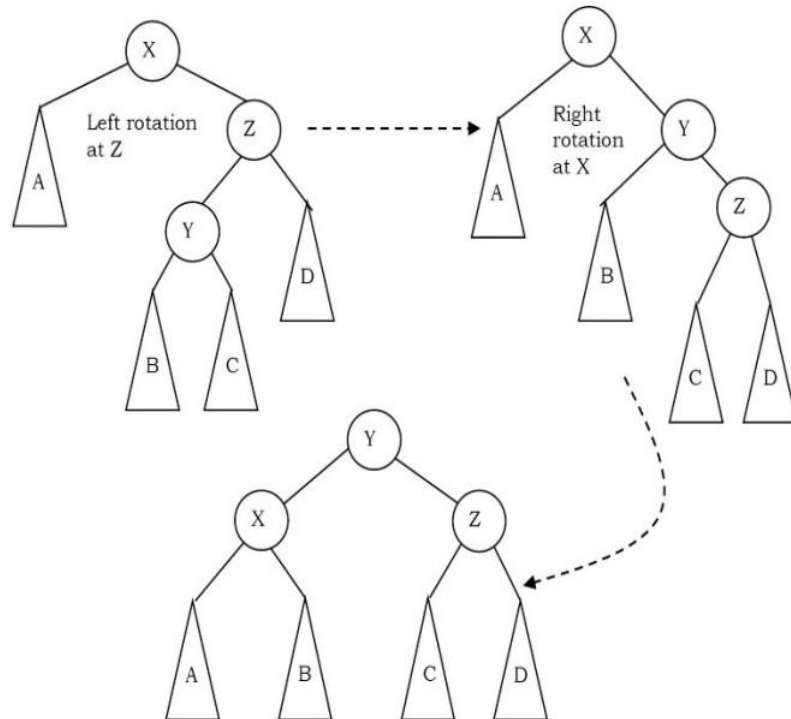
Types of Violations & Corresponding Rotations:

1. An insertion into the **left** subtree of the **left** child of X . (**LL**)

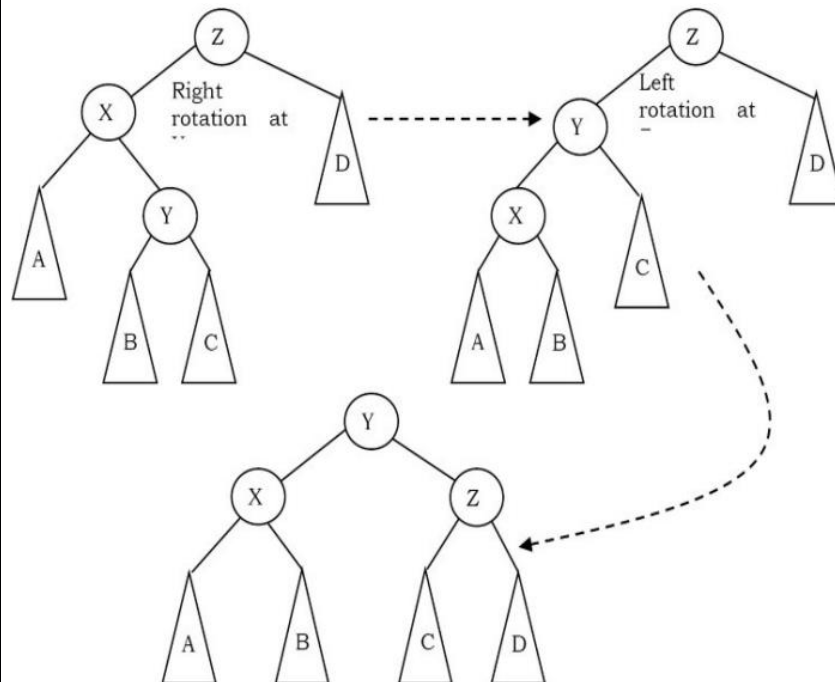
After performing LL rotation, we get tree on right:



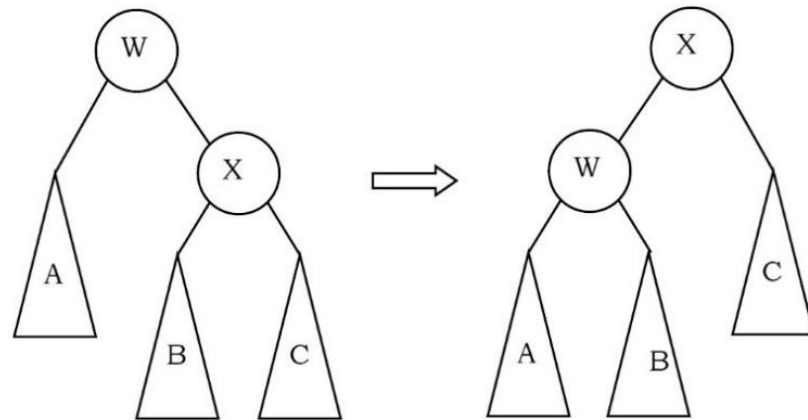
2. An insertion into the **right** subtree of the **left** child of X. (RL)
 After performing RL rotation, we get tree on right:



3. An insertion into the **left** subtree of the **right** child of X. (LR)
 After performing LR rotation, we get tree on right:



4. An insertion into the **right** subtree of the **right** child of X. (**RR**)
After performing RR rotation, we get tree on right:



ALGORITHM:

1. Create `AvlTree` Class with inner `Node` class
2. Create 2 ref vars: `right`, `left`, int var `data` & `height`
3. Initialize `right`, `left` to null & `height` = 1
4. Initialize root node to null

height Method:

1. if `root=null`:
2. return 0
3. else return `root.height`

SingleLL Method:

1. Initialize node `temp` to `root.left`
2. `root.left = temp.right`
3. `temp.right = root`
4. Update **height** for both root & temp nodes
5. return `temp`

SingleRR Method:

1. Initialize node `temp` to `root.right`
2. `root.right = temp.left`
3. `temp.left = root`
4. Update **height** for both root & temp nodes
5. return `temp`

DoubleLR Method:

1. root.left = **SingleRR**(root.left);
2. return **SingleLL**(root);

DoubleRL Method:

1. root.right = **SingleLL**(root.right);
2. return **SingleRR**(root);

getBF Method:

1. if root=null:
2. return 0
3. else return **height**(root.left) - **height**(root.right)

insert Method:

1. This method helps to make the recursive call to **inserter** Method

inserter Method:

1. if root is null:
2. root = new node(data)
3. return root
4. if data is less than root.data
5. root.left = **inserter**(root.left, data)
6. else if data is more than root.data
7. root.right = **inserter**(root.right, data)
8. Else return the root
9. Update **height** for root
10. initialize bal = **getBF**(root)
11. if bal>1 and data<root.left.data:
12. return **SingleLL**(root)
13. if bal<-1 and data>root.right.data:
14. return **SingleRR**(root)
15. if bal>1 and data>root.left.data:
16. return **DoubleLR**(root)
17. if bal<-1 and data<root.right.data:
18. return **DoubleRL**(root)
19. return root

minValue Method:

1. initialize minv equal to root.data
2. loop till root.left is not null:
 - a. minv = root.left.data
 - b. root = root.left
3. return minv

PreOrder Method:

1. if root is not null
2. print root.data
3. recurr: PreOrder(root.left)
4. recurr: PreOrder(root.right)

InOrder Method:

1. if root is not null
2. recurr: PreOrder(root.left)
3. print root.data
4. recurr: PreOrder(root.right)

PostOrder Method:

1. if root is not null
2. recurr: PreOrder(root.left)
3. recurr: PreOrder(root.right)
4. print root.data

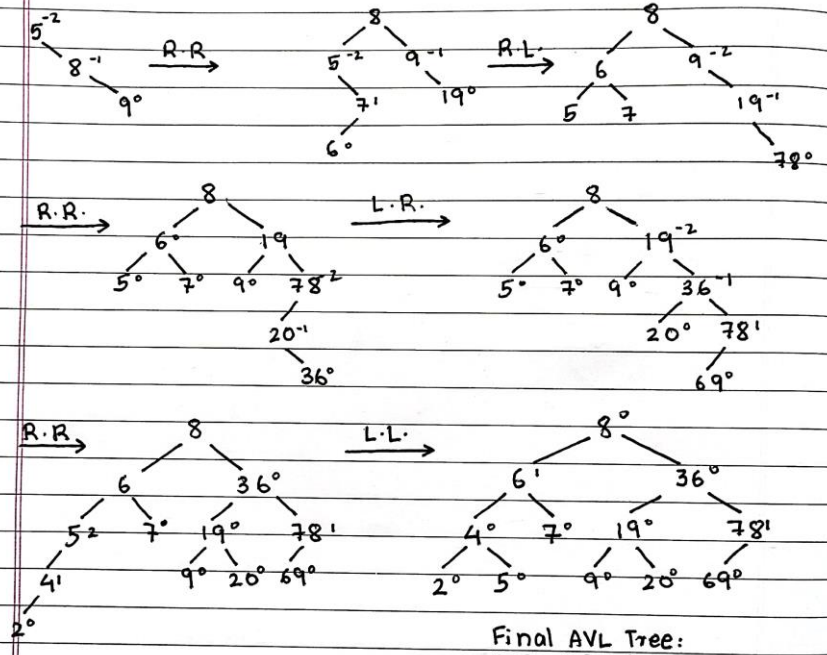
PROBLEM SOLVING:

AVL Tree:

→ Hatim Sawai

→ Insertion using all 4 rotations:

Input: 5, 8, 9, 19, 7, 6, 78, 20, 36, 69, 4, 2



PreOrder(Final) = 8, 6, 4, 2, 5, 7, 36, 19, 9, 20, 78, 69

PROGRAM:

AVLCheck.java:

```
import java.util.Scanner;
import avltreeds.AvlTree;
public class AVLCheck {
    public static void main(String[] args) {
        Scanner sc = new Scanner(System.in);
        AvlTree avl = new AvlTree();
        int choice, flag;
        int n, d;
        while(true) {
            System.out.println("Select 1 operation:\n1. Insert\t2.
Delete\n3. PreOrder\t4. InOrder\t5. PostOrder");
            choice = sc.nextInt();
            switch(choice) {
                case 1:
```

```

        System.out.print("Enter no. of elements to insert: ");
        n = sc.nextInt();
        System.out.print("Enter the elements: ");
        for (int i = 0; i < n; i++) {
            avl.insert(sc.nextInt());
            System.out.print("Tree (pre): ");
            avl.PreOrder(avl.root);
            System.out.println();
        }
        break;
    case 2:
        System.out.print("Enter the element to delete: ");
        d = sc.nextInt();
        avl.delete(avl.root,d);
        System.out.print("InOrder: ");
        avl.InOrder(avl.root);
        System.out.println();
        break;
    case 3:
        System.out.print("PreOrder: ");
        avl.PreOrder(avl.root);
        System.out.println();
        break;
    case 4:
        System.out.print("InOrder: ");
        avl.InOrder(avl.root);
        System.out.println();
        break;
    case 5:
        System.out.print("PostOrder: ");
        avl.PostOrder(avl.root);
        System.out.println();
        break;
    case 6:
        System.out.println("Size of the tree is " +
avl.size(avl.root));

```



```

        break;
    default:
        System.out.println("Invalid choice!");
    }
    System.out.println("Do you want to continue?\n1. Yes\t2. No");
    flag = sc.nextInt();
    if (flag == 2) {
        break;
    }
}
sc.close();
}
}

```

AvlTree.java:

```

package avltreeds;
public class AvlTree {
    class Node {
        int data;
        Node left, right;
        int height;
        Node(int data) {
            this.data = data;
            left = right = null;
            height = 1;
        }
    }
    public Node root=null;
    public int height(Node root) {
        if (root == null)
            return 0;
        return root.height;
    }
    public Node SingleLL(Node root) {
        Node temp = root.left;
        root.left = temp.right;
    }
}

```

```

        temp.right = root;
        root.height = Math.max(height(root.left), height(root.right)) + 1;
        temp.height = Math.max(height(temp.left), height(temp.right)) +
1;
        return temp;
    }
    public Node SingleRR(Node root) {
        Node temp = root.right;
        root.right = temp.left;
        temp.left = root;
        root.height = Math.max(height(root.left), height(root.right)) + 1;
        temp.height = Math.max(height(temp.left), height(temp.right)) +
1;
        return temp;
    }
    public Node DoubleLR(Node root) {
        root.left = SingleRR(root.left);
        return SingleLL(root);
    }
    public Node DoubleRL(Node root) {
        root.right = SingleLL(root.right);
        return SingleRR(root);
    }
    public int getBF(Node root) {
        if (root == null)
            return 0;
        return height(root.left) - height(root.right);
    }
    public void insert(int data) {
        root = inserter(root, data);
    }
    public Node inserter(Node root, int data) {
        if (root == null) {
            root = new Node(data);
            return root;
        }

```

```

    if (data < root.data) {
        root.left = inserter(root.left, data);
    } else if (data > root.data) {
        root.right = inserter(root.right, data);
    } else {
        return root;
    }
    root.height = 1 + Math.max(height(root.left), height(root.right));
    int bal = getBF(root);
    if (bal > 1 && data < root.left.data) {
        System.out.println("LL Rotation performed");
        return SingleLL(root);
    }
    if (bal < -1 && data > root.right.data) {
        System.out.println("RR Rotation performed");
        return SingleRR(root);
    }
    if (bal > 1 && data > root.left.data) {
        System.out.println("LR Rotation performed");
        return DoubleLR(root);
    }
    if (bal < -1 && data < root.right.data) {
        System.out.println("RL Rotation performed");
        return DoubleRL(root);
    }
    return root;
}

public void delete(Node root, int data) {
    root = deleter(root, data);
}

public Node deleter(Node root, int data) {
    if (root == null) {
        return root;
    }
    if (data < root.data) {
        root.left = deleter(root.left, data);

```

```

    } else if (data > root.data) {
        root.right = deleter(root.right, data);
    } else {
        if (root.left == null) {
            return root.right;
        } else if (root.right == null) {
            return root.left;
        }
        root.data = minValue(root.right);
        root.right = deleter(root.right, root.data);
    }
    root.height = Math.max(height(root.left), height(root.right)) + 1;
    int bal = getBF(root);
    if (bal > 1 && getBF(root.left) >= 0) {
        return SingleLL(root);
    }
    if (bal > 1 && getBF(root.left) < 0) {
        return DoubleLR(root);
    }
    if (bal < -1 && getBF(root.right) <= 0) {
        return SingleRR(root);
    }
    if (bal < -1 && getBF(root.right) > 0) {
        return DoubleRL(root);
    }
    return root;
}

public int minValue(Node root) {
    int minv = root.data;
    while (root.left != null) {
        minv = root.left.data;
        root = root.left;
    }
    return minv;
}

public int size(Node root) {

```

```

        if (root == null)
            return 0;
        else
            return (size(root.left) + 1 + size(root.right));
    }

    public void PreOrder(Node root) {
        if (root!=null) {
            System.out.print(root.data + " ");
            PreOrder(root.left);
            PreOrder(root.right);
        }
    }

    public void InOrder(Node root) {
        if (root!=null) {
            InOrder(root.left);
            System.out.print(root.data + " ");
            InOrder(root.right);
        }
    }

    public void PostOrder(Node root) {
        if (root!=null) {
            PostOrder(root.left);
            PostOrder(root.right);
            System.out.print(root.data + " ");
        }
    }
}

```

OUTPUT:

```

PS D:\Data Structures\Exp7> cd "d:\Data Structures\
Select 1 operation:
1. Insert      2. Delete
3. PreOrder    4. InOrder    5. PostOrder
1
Enter no. of elements to insert: 12
Enter the elements: 5 8 9 19 7 6 78 20 36 69 4 2
Tree (pre): 5
Tree (pre): 5 8
RR Rotation performed
Tree (pre): 8 5 9
Tree (pre): 8 5 9 19
Tree (pre): 8 5 7 9 19
RL Rotation performed
Tree (pre): 8 6 5 7 9 19
RR Rotation performed
Tree (pre): 8 6 5 7 19 9 78
Tree (pre): 8 6 5 7 19 9 78 20
LR Rotation performed
Tree (pre): 8 6 5 7 19 9 36 20 78
RR Rotation performed
Tree (pre): 8 6 5 7 36 19 9 20 78 69
Tree (pre): 8 6 5 4 7 36 19 9 20 78 69
LL Rotation performed
Tree (pre): 8 6 4 2 5 7 36 19 9 20 78 69
Do you want to continue?
1. Yes  2. No
2
PS D:\Data Structures\Exp7> █

```

CONCLUSION:

In this experiment, we learned how to insert elements into an avl tree while maintaining the balance using single (LL/RR) or double(LR/RL) rotations on the tree.