

# NOTE ON TYPE THEORIES

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**Notation 0.1.** We employ the following notations.

- (2,1)-categories are denoted by bf symbols:  $\mathbf{C}, \mathbf{A}, \mathbf{E}, \dots$  or bb symbols:  $\mathbb{I}, \mathbb{D}, \mathbb{A}, \dots$
- $\mathbf{Set}$  is the category of sets.
- $\mathbf{Cat}$  is the large (2,1)-category of categories.
- There is a fully faithful (2,1)-functor  $\mathbf{disc}: \mathbf{Set} \hookrightarrow \mathbf{Cat}$ .
- $\mathfrak{Cat}$  is the large 2-category of categories.
- $\Delta^1$  is the 1-simplex seen as a category.  $\mathbf{C}^{\Delta^1}$  is the arrow category of  $\mathbf{C}$ .
- $\mathbf{C}_{/A}$  and  $\mathbf{C}_{A/}$  are over and under categories respectively.
- $\mathbf{cod}$  and  $\mathbf{dom}$  mean codomain and domain respectively. They often have as their type  $\mathbf{C}^{\Delta^1} \rightarrow \mathbf{C}, \mathbf{C}_{/A} \rightarrow \mathbf{C}$ , or  $\mathbf{C}_{A/} \rightarrow \mathbf{C}$ .
- By a *replete class of morphisms* of  $\mathbf{C}$ , we mean a replete full subcategory of  $\mathbf{C}^{\Delta^1}$ .
- $\mathbf{WSub}$  is the replete full sub (2,1)-category of  $\mathbf{Cat}^{\Delta^1}$  consisting of wide subcategory inclusions; i.e., faithful functors that induce equivalences on core groupoids.
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## 1. DEPENDENT TYPE THEORIES IN TERMS OF DISPLAY MAPS

**Definition 1.1.** A *display map category*  $\mathbf{C} = (\mathbf{C}, \mathbf{dis}_{\mathbf{C}})$  is a pair of a category  $\mathbf{C}$  and a replete class  $\mathbf{dis}_{\mathbf{C}}$  of morphisms satisfying the following conditions. Arrows in  $\mathbf{dis}_{\mathbf{C}}$  are called *display maps* of  $\mathbf{C}$ .

- $\mathbf{C}$  has a terminal object.
- Let  $h: \Delta \rightarrow \Gamma$  and  $f: A \rightarrow \Gamma$  be morphisms in  $\mathbf{C}$  such that  $f$  is a display map. Then there is a pullback square

$$(1) \quad \begin{array}{ccc} \cdot & \xrightarrow{\quad} & A \\ \downarrow & \lrcorner & \downarrow f \\ \Delta & \xrightarrow{\quad h \quad} & \Gamma \end{array}$$

in  $\mathbf{C}$  such that the left side is also a display map. ■

**Definition 1.2.** Define the 2-category  $\mathfrak{CwD}$  of display map categories as follows.

- 0-cells are display map categories.
- 1-cells are functors preserving the terminal object, display maps, and pullbacks of the form (1).
- A 2-cell  $\alpha: F \Rightarrow G$  is a natural transformation such that naturality squares at display maps are pullback squares. ■

**Definition 1.3.** A display map category is *democratic* if for each object  $\Gamma \in \mathbf{C}$ , there exists a sequence of display maps from  $\Gamma$  to the terminal object. We write  $\mathbf{CwD}^{\mathbf{dm}}$  for the full sub 2-category of  $\mathfrak{CwD}$  spanned by democratic display map categories. ■

**Proposition 1.4.**  $\mathbf{CwD}^{\mathbf{dm}}$  is a (2,1)-category. ◆

**Definition 1.5.** A *clan*  $\mathbf{C} = (\mathbf{C}, \mathbf{fib}_{\mathbf{C}})$  is a display map category satisfying the following conditions. Display maps (i.e., arrows in  $\mathbf{fib}_{\mathbf{C}}$ ) are called *fibrations* of  $\mathbf{C}$ .

- For each object  $A \in \mathbf{C}$ , the unique morphism  $A \rightarrow 1$  towards the terminal object is a fibration.
- $\mathbf{fib}_{\mathbf{C}}$  is closed under composition. ■

We write **Clan** for the full sub 2-category of  $\mathfrak{CwD}$  spanned by clans. Since clans are always democratic, this is a (2,1)-category.

**Definition 1.6.** A *strict category* is a category equipped with its wide subcategory that is a set. In other words, the (2,1)-category **StrCat** of strict categories is defined by the following pullback square.

$$\begin{array}{ccc} \mathbf{StrCat} & \longrightarrow & \mathbf{WSub} \\ \text{Obj} \downarrow & \lrcorner & \downarrow \\ \mathbf{Set} & \xrightarrow{\text{disc}} & \mathbf{Cat}^{\Delta^1} \\ & & \downarrow \text{dom} \\ & & \mathbf{Cat} \end{array}$$

The functor on the left is locally fully faithful and surjective on objects. In particular, **StrCat** is a category. ■

**Definition 1.7.** A *contextual category*  $\mathbf{C} = (\mathbf{C}, \text{dis}_{\mathbf{C}}, \langle \rangle, \ell)$  consists of the following data.

- A strict category  $\mathbf{C}$ .
- A structure  $\text{dis}_{\mathbf{C}}$  of a display map category for the underlying category of  $\mathbf{C}$ .
- An object  $\langle \rangle$  of  $\mathbf{C}$  which is a terminal object in the underlying category of  $\mathbf{C}$ .
- A function  $\ell: \text{Obj}(\mathbf{C}) \rightarrow \mathbb{N}$  satisfying the following conditions.
  - $\ell^{-1}(0) = \{\langle \rangle\}$ .
  - The canonical function

$$(\Gamma, A, f) \mapsto A: \coprod_{\Gamma \in \ell^{-1}(n)} \{(A, f) \mid f: A \twoheadrightarrow \Gamma\} \rightarrow \text{Obj}(\mathbf{C})$$

is a monomorphism and  $\ell^{-1}(n+1)$  coincides with its image.

We write **CwC** for the category of contextual categories and functors preserving those structures. ■

**Theorem 1.8.** There exists an 2-adjunction

$$\begin{array}{ccc} \mathbf{CwC} & \xrightleftharpoons[\text{ctx}]{|-|} & \mathfrak{CwD} \\ & \perp & \end{array}$$

that restricts to a biequivalence

$$\begin{array}{ccc} \mathbf{CwC} & \xrightleftharpoons[\text{ctx}]{|-|} & \mathbf{CwD}^{\text{dm}} \\ & \simeq & \end{array}$$

◆

## 2. UEMURA'S LOGICAL FRAMEWORK

**Definition 2.1.** A *representable map category*  $\mathbf{R} = (\mathbf{R}, \text{rep}_{\mathbf{R}})$  is a display map category satisfying the following conditions. Display maps (i.e., arrows in  $\text{rep}_{\mathbf{R}}$ ) are called *representable maps* of  $\mathbf{R}$ .

- $\mathbf{R}$  is finitely complete.
- For each representable map  $f: X \twoheadrightarrow Y$ , the pullback functor  $f^*: \mathbf{R}_{/Y} \rightarrow \mathbf{R}_{/X}$  has a right adjoint  $f_*$ . ■

**Theorem 2.2.** shoumei no aidani claim ireru yatu ◆

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**Theorem 2.3.** Saigo ni claim kuru taipu. ◆

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