

Adding Security to File Systems, One Layer at a Time

Erez Zadok

File-systems and Storage Laboratory
Stony Brook University
<http://www.fsl.cs.sunysb.edu/>



Why File Systems?

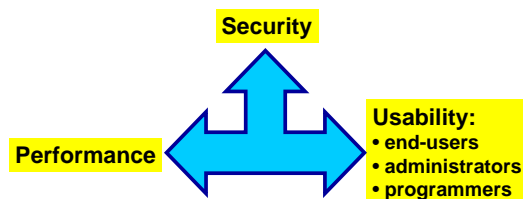
- Hardware can be replaced...
- Networks can be upgraded...
- Even employees come and go...
- ... but a company's data (IP) must **never** be lost.
- The **file system** software is in charge of access to stored data.

1/30/2007

© Erez Zadok – CSE-590



Balancing Conflicting Needs



Our goal: Explore all three aspects for file systems

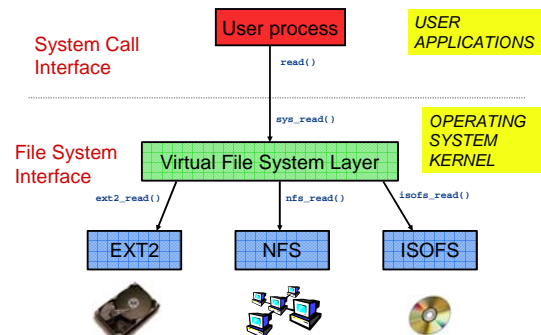
1/30/2007

© Erez Zadok – CSE-590

3



How File Systems Work



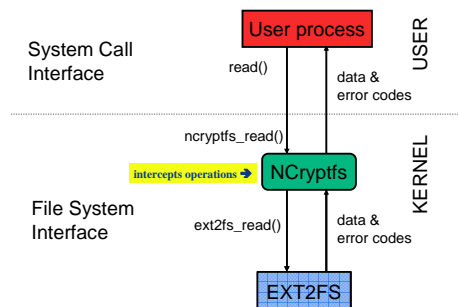
1/30/2007

© Erez Zadok – CSE-590

4



File System Interception



1/30/2007

© Erez Zadok – CSE-590

5



Interception for Security

- Interception a useful security concept
- Network interception
 - ◆ Easy to do without changing clients or servers
 - ◆ Doesn't impact application performance
 - ◆ Hard to infer high-level information from packets
 - ◆ Not all events produce network messages
- System call interception
 - ◆ Easy to do for most OSs
 - ◆ Does not capture all events (mmap)
 - ◆ Does not capture server-side events
 - ◆ May affect performance of every application

1/30/2007

© Erez Zadok – CSE-590

6



File-System Level Interception

- Logically: below system calls and well above network
 - ◆ Uses the Virtual File System (VFS) layer
- Pros:
 - ◆ No need to change clients or servers or applications
 - ◆ Minimal performance overhead, and only for those applications using an intercepted file system
 - ◆ Easy to infer high-level information (users, file names, etc.)
 - ◆ Capture all file accesses (including mmap)
 - ◆ Works equally well on clients, servers, or as a proxy
 - ◆ The file system normalizes data access to individual pages
- Cons:
 - ◆ Kernel development is difficult
 - ◆ Only captures file-related accesses

1/30/2007

© Erez Zadok – CSE-590

7 STONY BROOK

Outline

- NCryptfs [Usenix 2003]
- I3FS [LISA 2004]
- AVFS [Security 2004]
- Other and ongoing projects

1/30/2007

© Erez Zadok – CSE-590

8 STONY BROOK

NCryptfs: Versatile & Convenient

- Multiple ciphers, keys, and key lengths
- Encryption per user, process, group, or session
 - ◆ Ad-hoc groups
- Auth vs. Encryption keys
- Key timeouts and revocations:
 - ◆ Suspend/resume applications
 - ◆ Terminate/sandbox processes
 - ◆ Invoke user-land helper to re-authenticate

1/30/2007

© Erez Zadok – CSE-590

9 STONY BROOK

OS: Cache Cleaning

- Cached data is vulnerable
 - ◆ mmap accesses without consulting file system
 - ◆ Cached data can outlive key
- NCryptfs evicts cleartext pages
 - ◆ When keys become invalid
 - ◆ When authorizations become invalid
- The ciphertext page may still be cached, so I/O is not required to access the data again

1/30/2007

© Erez Zadok – CSE-590

10 STONY BROOK

OS: On-Exit Callbacks

- Expunge private user info on process exit, advantages over alternatives:
 - ◆ Efficiency: no periodic scans of lists
 - ◆ Security: no gap between process death and cleanup
- NCryptfs uses on-exit callbacks to
 - ◆ purge active sessions and authorizations
 - ◆ challenge-response authentication
 - the task-private data creates a session between a user process and the kernel

1/30/2007

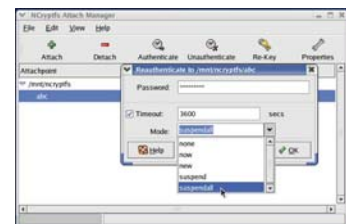
© Erez Zadok – CSE-590

11 STONY BROOK

User Space Callbacks

- Kernel calls user-space program

- NCryptfs GUI to re-key on timeout
- Executes as the user, not as root



1/30/2007

© Erez Zadok – CSE-590

12 STONY BROOK

Outline

- NCryptfs [Usenix 2003]
- **I3FS [LISA 2004]**
- AVFS [Security 2004]
- Other and ongoing projects

1/30/2007

© Erez Zadok - CSE-590

13 STONY BROOK

I³FS: Integrity Checking

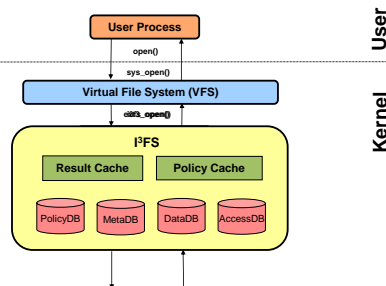
- Intercepts file system calls to protected files
 - ◆ A Stackable file system: inline interception
 - ◆ Performs checksum comparison
 - ◆ Reports checksum mismatch online
 - Optionally prevents access
- Files to be protected
 - ◆ Chosen by the administrator
 - ◆ Valid updates to protected files
 - Requires authentication
- File entities to be checksummed - Policy
 - ◆ Chosen by the administrator
 - ◆ E.g., inode fields, file data

1/30/2007

© Erez Zadok - CSE-590

14 STONY BROOK

I³FS Architecture



1/30/2007

© Erez Zadok - CSE-590

15 STONY BROOK

Outline

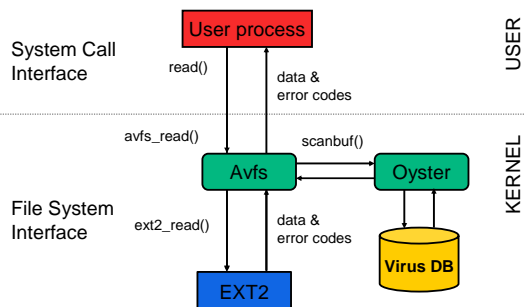
- NCryptfs [Usenix 2003]
- I3FS [LISA 2004]
- **AVFS [Security 2004]**
- Other and ongoing projects

1/30/2007

© Erez Zadok - CSE-590

16 STONY BROOK

Avfs: Anti-Virus File System



1/30/2007

© Erez Zadok - CSE-590

17 STONY BROOK

Avfs Components

1. Oyster virus scanner
 - Based on ClamAV user-mode scanner
 - Modified Aho-Corasick scanner
 - We improved ClamAV's scalability
 - Linear memory consumption with size of virus DB
 - Linux kernel scanning API
2. Use unmodified ClamAV's virus DB
3. Avfs: a stackable file system for inline scanning
 - Forensic modes
 - Sandboxing process's file access
 - Versioning
 - Virus quarantining

1/30/2007

© Erez Zadok - CSE-590

18 STONY BROOK

Outline

- NCryptfs [Usenix 2003]
- I3FS [LISA 2004]
- AVFS [Security 2004]
- **Other and ongoing projects**

1/30/2007

© Erez Zadok – CSE-590

19



Tracefs

- **Flexibility**
 - ◆ Fine-grained control of what should be traced
 - **Performance**
 - ◆ Low overheads. Trade-offs
 - **Convenience**
 - ◆ Simple and self-contained trace file format
 - ◆ User-level utilities to configure and manipulate traces
 - **Security**
 - ◆ Keyed checksums and encryption
 - ◆ Tracing to remote location
 - ◆ Anonymization
- [FAST 2004]

1/30/2007

© Erez Zadok – CSE-590

20



Replays

- Select subsets of traces to replay
 - ◆ e.g., IDS can pick parts to replay
- Compile portable traces into OS-specific format
 - ◆ even for another OS
- Replay traces at different speeds
- Single step undo/redo
 - ◆ Forensics

[FAST 2005]

→ coming to a SNIA WG near you...

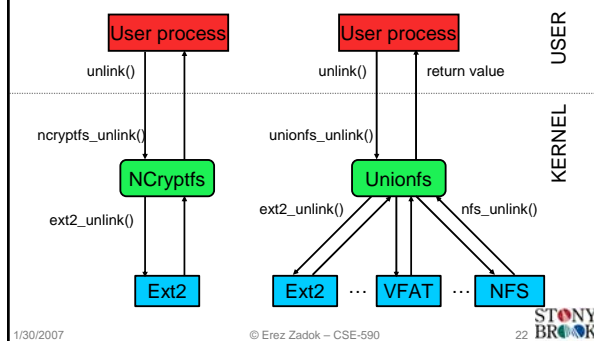
1/30/2007

© Erez Zadok – CSE-590

21



Stacking and Fan-Out



1/30/2007

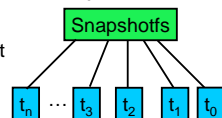
© Erez Zadok – CSE-590

22



Snapshots

- Create a consistent image of an entire file system
- When a snapshot is taken at time n :
 - ◆ Mark existing branches file system read-only
 - ◆ Add a new high-priority branch (t_n)
 - ◆ All writes take place on t_n
 - ◆ Transparently migrate open files on t_{i-1} to t_n
- To view snapshot i : union t_i through t_0
- Important fan-out features:
 - ◆ Dynamic branch management
 - ◆ File revalidation
 - ◆ Scalable to many branches



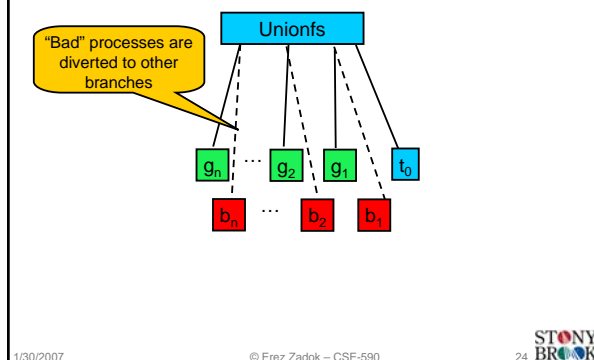
1/30/2007

© Erez Zadok – CSE-590

23



File System Sandboxing



1/30/2007

© Erez Zadok – CSE-590

24



Ongoing

- **RAIF**: Distributed storage survivability
 - ◆ code released in Jan 2007
 - ◆ [Cluster Sec 2005]
- **SDFS**: Secure Deletion of files
 - ◆ code released in Jan 2007
 - ◆ [SISW 2005, StorageSS 2006]
- **ACIDfs**: Transactions (TOCTOU)
 - ◆ code released in mid 2006
 - ◆ [ACM TOS 2007]

1/30/2007

© Erez Zadok – CSE-590



Take Home Messages

1. Security work must address *usability*
2. Using *layered* file systems to add security

1/30/2007

© Erez Zadok – CSE-590



Future

- **eCryptfs** (IBM)
 - ◆ already in Linux 2.6.19
 - ◆ More stacking support in 2007
- **Unionfs**:
 - ◆ basis for fan-out stackable file systems
 - ◆ code stability & cleanup
 - ◆ ETA 2007?

→ *expect more*

1/30/2007

© Erez Zadok – CSE-590



Adding Security to File Systems, One Layer at a Time

Erez Zadok

Filesystems and Storage Laboratory
Stony Brook University
<http://www.fsl.cs.sunysb.edu/>

