CSC263 Notes

Data Structures and Analysis

https://github.com/ICPRplshelp/

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1 Data Types, Data Structures

ADTs

- Specification
  - Objects we're working with
  - The operations (WHAT but not how)

Data structures

- Implementation (how)
  - Data
  - Algorithms

Analysis (runtime or complexity)

- Worst case
- Best case
- Upper bounds
  - $\mathcal{O}$
- Lower bounds
  - Ω
- Tight bounds
  - Θ

If we have algorithm A and input x, the runtime $t_A(x) =$ number of constant time operations independent of x.

Ultimately, we want a measure of running time that is a function of the input size. We have lots of inputs for each input size. So if we want to prove an upper bound when looking for the worst case running time:

- I need two functions to prove an upper bound
  - A simple algebraic expression
  - The running time
    * However, the pure runtime function's codomain isn't $\mathbb{R}^{\geq 0}$ – but rather, a list of running times. To turn it into a raw function that outputs $\mathbb{R}^{\geq 0}$, we can take the largest of the list I just described.

Worst case is just us narrowing down a bunch of possible runtimes to the worst one.

My upper bound will always be some value that is larger or equal to the worst case, and the lower bound must be below the worst case but not all the worst cases.

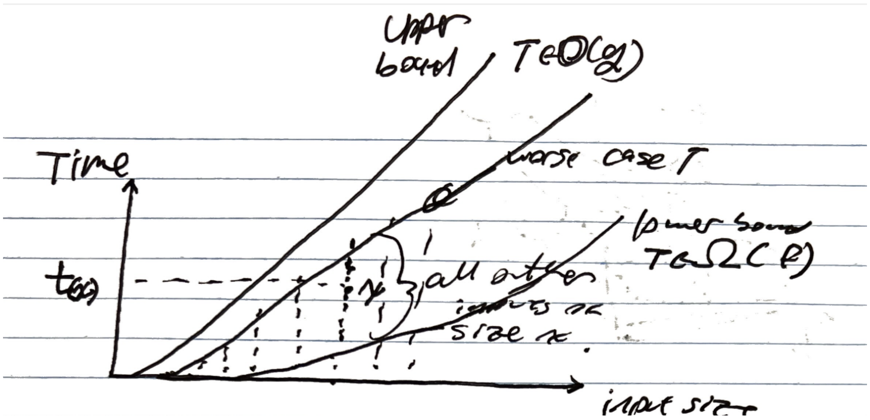

Figure 1: Upper and lower bounds of the worst case