MASTER OF TECHNOLOGY (INTELLIGENT SYSTEMS)

PROJECT REPORT

Reasoning Systems

SDOT

Scheduling & Despatch Optimisation Tool

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EXECUTIVE SUMMARY

Today's economy revers competitiveness, agility and sustainability. With key performance indicators of time-to-market and customer responsiveness, many manufacturing companies are adopting Industry 4.0 standards to optimize productivity, efficiency and flexibility via digital manufacturing solutions. While most of the manufacturing systems are operating near optimal productivity under stable conditions, there are certain disturbances such as product changes, dynamic work-in-progress (WIP) and urgent orders, which are inevitably present in the real product environment which can impact the overall manufacturing performance. One possible solutions is providing real-time end-to-end planning and control via an effective scheduling system that aims to handle and provide for the dynamic production environment.

In this project, our group has proposed an integrated solution that is able to address the planning and control issues for the scenario of scheduling in baking of IC (integrated circuit) chips in the semiconductor industry. We implement and propose a solution we call SDOT (Scheduling & Despatch Optimisation Toolkit) and discuss conceivable limitations of the system and suggest possible enhancements that can modify it to handle more complexity in scheduling problems.

1.0) PROBLEM DESCRIPTION

Today's economy revers competitiveness, agility and sustainability. With key performance indicators of time-to-market and customer responsiveness, many manufacturing companies are adopting Industry 4.0 standards to optimize productivity, efficiency and flexibility via digital manufacturing solutions. While most of the manufacturing systems are operating near optimal productivity under stable conditions, certain disturbances such as product changes, dynamic work-in-progress (WIP) and urgent orders, which are inevitably present in the real product environment, can impact the overall manufacturing performance. Such production disturbances can result in an estimated 80% of overall equipment effectiveness (OEE) loss. One of such solutions is providing real-time end-to-end planning and control via an effective scheduling system that aims to handle and optimize the dynamic production environment, thereby reducing cycle time, increase machine utilization and throughput, meet due date and eventually achieving greater customer satisfaction and trust.

The semiconductor backend testing is one of the most highly complex manufacturing systems in terms of equipment, process flows and interdependency relations. Furthermore, the complexity is increased due to a large product variety, complex product to tester relations, multiple level of tester to hardware dependency, sequence dependency, dynamic determination of processing and indexing time, batch processing and rework flow. Besides its complexity, the real production environment is also subjected to many uncertainties and unpredictable events due to continuous arrival of new and unforeseen orders and intermittent occurrence of machine breakdown, and random process and yield variations. Thus, the scheduling is no longer simply a static optimization problem, but an ongoing reactive process. Therefore, the scheduling operations and optimization in semiconductor backend testing are highly demanding and challenging for both researchers and practitioners.

Generally, the basic scheduling approaches formulate schedules computed over a specific time frame assuming all problems and states are known in advance, without consideration of any disturbances in the actual production. In order to handle the dynamic production environment, short-term planning and scheduling can be considered. Schedules can be regenerated based on new events or limitations such as machine breakdown and WIP shortages or urgent order, respectively. A periodic scheduling regenerates periodically over a time period to update the schedules with new job list (based on existing WIP status) while event-driven regenerates schedules upon event occurrence.

Based on the above considerations given to the different facets in the scheduling process, this project aims to introduce a real-time manufacturing scheduler that is capable of helping the production accomplish operation strategies, optimize goals and deal with daily disruptions such as machine breakdown and WIP discontinuity in a more effective way. In the next section, the project objectives and scope are defined, followed by a detailed discussion of the problem definition, requirements and constraints.

1.1) PROJECT OBJECTIVES & SCOPE

In this project, a scheduling problem based on a semiconductor backend test case study is performed via Optaplanner. It is a programmable platform that functions as a constraint satisfaction solver that optimizes business resource planning cases such as scheduling problem in the nature of this project.

The case study mainly focusses on the particular process of the test manufacturing whereby the integrated circuits (ICs) will be placed in an oven at desired high temperature for

prolonged hours (~6-24 hours) to remove any moisture content. This process is also known as the bake-in operation and the lots are processed in the ovens in batches. Therefore, this project aims to schedule and optimize the semiconductor test operation by selecting the most optimal lot to load onto the available oven in future simulated time to:

- Meet due date (improving delivery accuracy; processing the right products and quantity to fulfill delivery demand)
- 2) Minimize cycle time (time taken for lot to complete the bake-in operation)
 The scheduling problem is simplified by making the following assumptions:
- Setup time and breakdown of oven and lot movement time in the operation (i.e. lot moving into and out of oven) are negligible.

1.2) PROBLEM STATEMENT AND FORMULATION

The scheduling problem is defined as the assignment of selected lots from the WIP list to be processed on machines in each operation while satisfying constraints and optimizing objective function. The criteria of the lot selection process from the existing WIP is dependent on the delivery plan provided (i.e. delivery accuracy, selecting the right lots to be processed to fulfill delivery demand) (see Appendix A) and its objective function is as follows

Minimize D(s):
$$|\sum_{\substack{i=1\\j=1}} q_{i,j} - Q_j|$$
 (1)

where i refers to the lot identification (ID); j refers to the product type; q refers to the lot quantity size and Q refers to the delivery demand. After the lot selection is performed, the scheduling is carried out with the following requirement and constraints:

Each lot is composed of only one product

- During the bake-in operation, lots of the same product can be processed in the same oven (batch processing)
- However, only some products may be processed in certain ovens because of product requirement
 - Ovens are independent from each other.
 - All lots and ovens are available at initial state.
 - The processing time at each oven is product and oven dependent.
 - At a given time, the oven can only execute one operation.
 - Processing at each oven must be completed without interruption once it starts
 - The oven has a limited capacity.

The scheduling problem sets include the set of selected lots I to be processed, with each lot i corresponding to a specific product j and a set of ovens L with specific temperatures k. The attributes of the sets are shown in Table 1. The planning parameters are also specified in Table 2 which includes the lot quantity q_{ij} , processing time p_k and oven quantity of different temperatures r_k . The oven machines have different operating temperatures k, mainly 125 and 150 degree Celsius and each product j has different processing time p_{jk} . Once the batch processing starts, the process cannot be interrupted. All the product processing requirements are shown in Appendix B. The variable of the scheduling problem will be defined as S_{ni} where S denotes the scheduling solution; n is the ordinal value that measures the order of lot I sequence.

In summary, the scheduling problem considered in this project is a single-stage, parallel batch multi-product job scheduling problem. There are more than 100 different types of products, where same product units are grouped in lot form and assigned to oven in batches for processing.

Table 1. Definition of problem sets

No	Problem sets	Attributes
1	Iii. Set of selected lots that are scheduled	
	for processing. (each lot corresponding	
	to a specific product)	q_{ij} = Quantity
2	L_{lk} : Set of oven units available for	
	processing the lots	<i>k</i> = Oven Temperature Type
		r_k = Oven quantity

Table 2. Planning parameters

Planning Parameters	Definition
q_{ij}	Total number of units within a lot <i>i</i> (each lot corresponding to a specific
	product k)
p_{jk}	Processing time of each product <i>j</i> at specific operating temperature <i>k</i>
r_k	Total number of oven units of specific operating temperature <i>k</i>

Table 3. Oven inventory in the factory

Operations	Machine	Quantity r _k	Temperature Type k
Bake-in	Oven	15	125
		7	150

The appropriate sequencing of lots on the oven in the bake-in operations can be determined via the optimization of the objective function which is derived based on the economic factory performance. The measure of the economic factory performance encompasses of timely delivery fulfillment (minimum makespan or cycle time) and maximum equipment utilization. A set of weightages W are considered to provide prioritization to the two performance measures as in usual practice, these measures often contradict each other. Herein, the objective function f(s) is:

$$CT(s) = \frac{\sum_{l=1}^{L} \sum_{n=1}^{L} t_{ni}}{\sum_{k=1}^{L} r_k}$$
 (2)

where the cycle time CT(s) is measured by averaging the total processing time of all the batch jobs in all ovens.

2.0) APPROACH AND SOLUTION

2.1) SYSTEM ARCHITECTURE

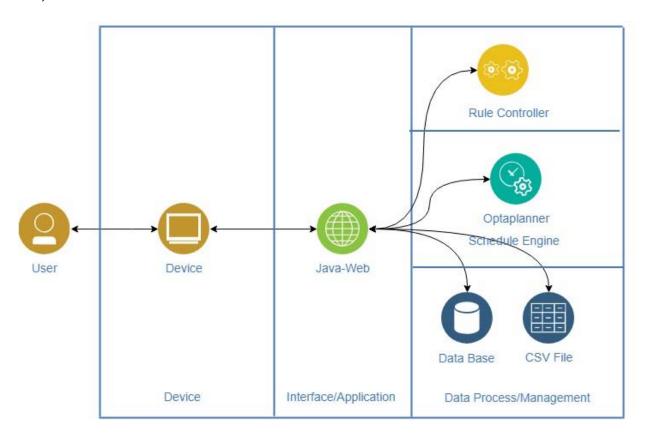


Figure 1. Flow of the effective scheduling system

This system is primarily accessed via a frontend webpage hosted on Apache tomcat locally (at this point). The frontend will take information from CSV and Database to feed to the Optaplanner solver and the Rule engine that does the packaging.

2.2) SYSTEM FLOW

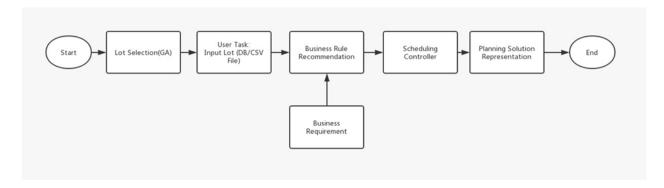


Figure 2. Flow of the effective scheduling system

Figure 2 illustrates the flow of this project. This will include:

- The Lot selector (done by Genetic Algorithm)
- The File Upload
- Packaging Rule
- Scheduling
- Displaying Scheduled Results

Each of these components in the flow is explained below to provide an understanding of how they are integrated for the solution to work.

Lot Selection (done by Genetic Algorithm)

With the lot selection and scheduling problem requirement, constraints and objective functions defined in the previous section, the lot selection process is first performed via Excel Solver and subsequently, the selected lot list is passed to the effective scheduling system which is developed via Optaplanner. The Excel Solver based on the evolutionary algorithm is employed to solve the lot selection problem which is a basic binary assignment problem. When a lot is selected, a value of 1 will be assigned while an unselected lot will assign a value of 0. As the selection list is

relatively large, a population size of 50, mutation rate of 0.075 and convergence of 0.0001 are set to reduce the computational time.



Table 4: Excel GA Solver

<u>User task: Input Lot (DB/CSV File)</u>

From the user end, the process requires the user to upload an input file of the selected lot from the earlier process in the format of a comma separated values (CSV) file. The file has headers in the following format as shown in Table 5 below. The CSV format provides a straightforward information schema in a compact, readable manner that allows any editing manual or otherwise and efficient processing by the scheduling controller based on the business rules. CSV format also allows easy sorting, filtering and manipulation based on any programmable queries.

LOT	STEPNAME	ca der Type	Product Type	QUANTITY
lot1	Bake_In	Trays	LEQUD1106E	3674
lot2	Bake_In	Tubes	UAQAL5502F	27944
lot3	Bake_In	Trays	KILKJ2749Y	8195
lot 4	Raka ta	Travel	1/751 072040	11104

Table 5 : Screenshot of Input file in CSV format

The csv file contains the following lot attribute information which is crucial for the scheduling problem:

- Lot ID
- Process Type

- Product Type
- Quantity

In the next process, we will use these details to generate an input file to feed to the solver.

Batching Rule

In the scheduling problem, the lots are first batched together based on their quantities, process requirement (i.e. processing temperature and time) and the oven capacity before dispatching them to the ovens for processing. The processing requirement and oven inventory are shown in Figure 3 and 4, respectively. This information will be maintained in a database table in mySQL. The rule is designed as follows:

The rule engine will take information provided in the CSV file (Product Type and Quantity), and cross reference with the DB tables. Then, on a first-in first-out (FIFO) basis, it will add cumulative volume percentage, and register one package when the next member of the stack will cause the 100% to overflow. At the end of this process, we will only have data in full packages, ready to be fed into the oven.

As an example, please refer to figure 3. In a scenario,1000000 pieces of 'ACHXA395G' (LOT1), 1000000 pieces of 'ACNYW1625J' (LOT2) and 1000000 pieces of 'AZFGJ6792V' (LOT3), based on the FIFO rule, the rule engine will start with LOT1. 'ACHXA395G' have a unit volume of 0.0000416667 percent out of 100 percent capacity relative to the oven. Therefore, LOT1 will occupy 0.00004167 *1000000 = 41.67% of oven. Using the same rule, LOT2 and LOT3 will both occupy 27.78%. Total: 41.67+27.78+27.78 = 97.23%. This will be batched into one package, when the next lot (with let's say 10% capacity) cause this number to overflow. (97.23%+10% = 107.23%).

productType	volume	temperature	bakeTime	ovenId	temperature	processType
ACHXA3595G	0.0000416667	125	12	BO1	125	BakeIn
ACNYW 1625J	0.0000277778	125	24	BO10	125	BakeIn
ALSRP5590C	0.0000016026	125	24	BO11	125	BakeIn
APYLB9291C	0.0000104167	125	24	BO12	125	BakeIn
ARCYI6496Z	0.0000046296	125	24	BO13	150	BakeIn
ARXZK2314V	0.0000046296	125	24	BO14	150	BakeIn
AYEHS2904M	0.0000055556	150	24	BO15	125	BakeIn
AZFGJ6792V	0.0000277778	125	24	BO16	125	BakeIn
BBKRW9735V	0.0000083333	150	24	BO17	125	BakeIn
BDUFC2928C	0.0000277778	125	24	BO18	150	BakeIn
BGGII2097V	0.0000055556	150	24	BO19	125	BakeIn
BGLWT0100N	0.0000198413	125	24	BO2	125	BakeIn
BHDPO8835E	0.0000104167	125	24	BO20	125	BakeIn
BHKIO5418Q	0.0000055556	150	24	BO21	125	BakeIn
BHLLZ7485H	0.0000416667	125	12	BO22	125	BakeIn
BLCND8992P	0.0000277778	125	24	BO3	150	BakeIn

Figure 3: Product Information DB Table

Figure 4: Oven Information DB Table

For more information on the batching process, please refer to the appendix.

Scheduling Controller (Solver)

Subsequently, the production planner picks up the resulting package list and input it into solver. In this project, the Optaplanner solver engine is used. With proper class parameters (Figure 5) and hard/soft constraints set, these values will be fed to the Optaplanner solver to perform a local search for the best solution.

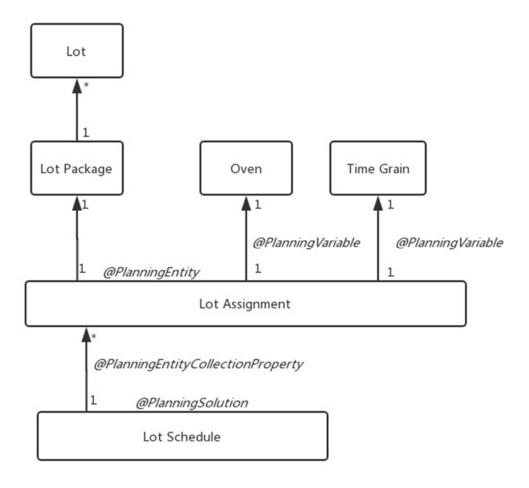


Figure 5: Class Diagram

The objective for this process is the lot scheduling calendar, displayed on a webpage. To fulfil it, we have to set certain hard and soft constraints so that the Optaplanner can perform local search properly. Our hard constraint will be the capacity of the oven, because it cannot be violated. For any instance of time involved, in hours, it will be used as soft constraint.

In this process, we are not setting a algorithm for Optaplanner to use. As such, it will use the default search algorithm (Late Acceptance). This will be for minimizing the amount of idle time in between processes. After doing the local search, Optaplanner will return the data in time-start and time-end format, to be displayed in a chart.

PRESUITE Hard/Score: 2 / 5 oft 5 core: -1570 Oven 8015 Oven 8012 Oven 803 Oven 803 Oven 803 Oven 809 Oven 809 Oven 809 Oven 809 Oven 808 Oven 808 Oven 808 Oven 808 Oven 808 Oven 808 Oven 808

Planning Solution Representation

Figure 6: Screenshot of Planning Solution Representation

After processing, the frontend will take information from the previous process to display the above chart (Figure 6). Engineers can then schedule operators/robots to put the parts in the ovens.

3.0) CONCLUSION

In this project we propose and implement an effective and predictive system, SDOT (Scheduling and Despatch Optimisation Tool), as a solution to address the complexity of problems in predictive scheduling based on software platforms of Excel Solver and Optaplanner. The integrated solution through these two platforms help the test manufacturing

firm to meet its delivery demand in due time. Conventional planning based on static calculation poses a challenging task for the production planner due to the large size and makes lot scheduling prone to human errors. With our scheduling system, SDOT, the complex interrelations between the product requirement, machine parameters, business requirement such as delivery demand fulfilment can be accounted to provide a schedule plan to meet factory performance measure. The solution approach is two fold, where firstly the lot selection is performed using Excel Solver. This relies on application of an Evolutionary computation approach using Genetic Algorithms (GA). Following that, the optimal selected based on the GA is passed to the scheduling system developed using Optaplanner which takes into account the constraints of the system as codified in the form of business rules.

3.1) IMPROVEMENTS FOR THE SDOT SYSTEM

Job shop scheduling problem is a well-researched topic. There are many comprehensive and robust solutions and approaches that have helped to address the manufacturing challenges. The SDOT system we have proposed and implemented is a predictive scheduling system. In reality, scheduling is a more complex problem with multiple states and a myriad of dynamic parameters that needs consideration to achieve an optimal solution. Below, we explain some of the limitations present and potential enhancements that could be integrated into SDOT to improve its functionality.

Predictive systems vs Reactive system

SDOT is currently designed as a predictive scheduling system based on certain premises and conditions. This could be enhanced by a reactive system whereby the tasks and their constraints may become known only incrementally or while the system is already running based

on the current knowledge domain. In reality, there arises dynamic instances such as urgent orders, unpredicted machine breakdown and WIP shortages which may require a better solution. Based on knowledge of the new constraints and limitations, the scheduler can react by recalculating the optimal schedule.

Local optimal vs global optimal

Currently SDOT uses GA and search techniques that may result in solutions that are considered feasible for a current frame in time based on current knowledge. This may possibly yield solutions as a local optimum within its search parameters. For a more comprehensive system, that gives consideration to downstream processes, the availability and the states of the machine, a global optimum can be aimed for, by employing 'forward-looking' strategies. However due to the complexity of such techniques, this is not considered for this project.

Evaluation of other techniques

For this project, only one scheduling strategy/optimization/search technique is used. To test and evaluate how different optimization techniques perform and fare against one another based on varying system conditions, they can be run via a simulation model. For example, the schedule results output from the different search/optimization techniques can be then used as an input to a virtual factory simulation model to evaluate other factory performance measures such as operating equipment efficiency (OEE).

Appendix

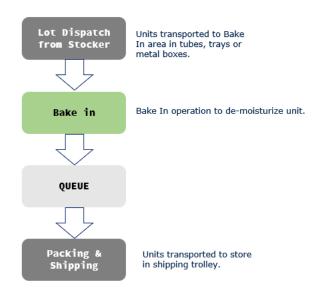
APPENDIX A: DELIVERY PLAN

Screen shot of Delivery Plan and WIP Lot list Document uploaded as "GALotSelection.xlsx"

ည Delivery Plan						
Provide Lot Count(ci) Product Type Requirement Chip Quantity Provide Chip Quantity(xi) Dif						
1	ACNYW1625J	48714	4704	44010	4703	
1	ALSRP5590C	1071	13227	-12156	13226	
2	ARCYI6496Z	19936	17509	2427	17507	
1	ARXZK2314V	9071	7699	1372	7698	
0	AYEHS2904M	3929	0	3929	0	
3	BBKRW9735V	63214	33690	29524	33687	

	WIP Lot List							
LOT	STEPNAME	carrier Type	Product Type	QUANTITY		Α	q'	
lot2	Bake_In	Trays	ACNYW1625J	4704	i	1	4704	
lot87	Bake_In	Trays	ALSRP5590C	13227	i	1	13227	
lot89	Bake_In	Tubes	ARCYI6496Z	8617	i	1	8617	
lot88	Bake_In	Tubes	ARCYI6496Z	8892	i	1	8892	
lot102	Bake_In	Tubes	ARXZK2314V	7699	i	1	7699	
lot136	Bake_In	Tubes	BBKRW9735V	4706	i	1	4706	
1-+124	Dales In	Tudasa	DDVDWOZOEV	11455	1:	1	1115	

APPENDIX B: Product Processing Requirements



APPENDIX C: Batching Rule

	U	U	U	<u> </u>
LOT	STEPNAM	carrier Typ	Product Type	QUANTITY
lot1	Bake_In	Trays	LMCJN3197X	1000000
lot2	Bake_In	Tubes	YGPPA4108I	1000000
lot3	Bake_In	Trays	ACHXA3595G	1000000
lot4	Bake_In	Trays	BHLLZ7485H	1000000
lot5	Bake_In	Trays	EFLFF0922K	1000000
lot6	Bake_In	Trays	EPEOQ3051X	1000000
lot7	Bake_In	Trays	GRKFF2612X	1000000
lot8	Bake_In	Trays	IMBDS8589X	1000000
lot9	Bake_In	Trays	IWFKE5721L	1000000
lot10	Bake_In	Trays	IXHLX5077Z	1000000
lot11	Bake_In	Trays	JILTC2318E	1000000
lot12	Bake_In	Tubes	KTKBX4839E	1000000
lot13	Bake_In	Trays	MHONL8388D	1000000
lot14	Bake_In	Trays	MHZUP0703L	1000000
lot15	Bake_In	Tubes	MTBXW9690F	1000000
lot16	Bake_In	Trays	NLYFD0861S	1000000
lot17	Bake_In	Trays	OPVLI2791R	1000000
lot18	Bake_In	Trays	TBZVN7846Q	1000000
lot19	Bake_In	Trays	UQSFI7003V	1000000
lot20	Bake_In	Tubes	VCBUH0118F	1000000
lot21	Bake_In	Trays	WVFHV5596N	1000000
lot22	Bake_In	Trays	XKVES6845J	1000000
lot23	Bake_In	Trays	ZAJUB9415U	1000000
lot24	Bake_In	Trays	JNCIJ9464N	1000000
lot25	Bake_In	Trays	WPWSV8594A	1000000
lot26	Bake_In	Trays	YLMRS4880I	1000000
lot27	Bake_In	Trays	ACNYW1625J	1000000
lot28	Bake_In	Trays	AZFGJ6792V	1000000
lot29	Bake_In	Trays	BDUFC2928C	1000000
lot30	Bake_In	Tubes	BLCND8992P	1000000
Int31	Rake In	Tuhes	FKCOV7230K	1000000

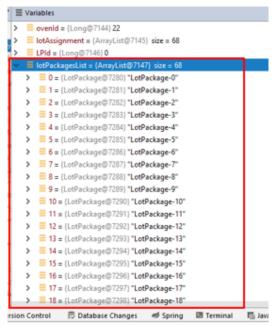
Appendix C1: Lot List (test_data.csv)

4	А	D	C
1	ProductType	Volume	
2	ACHXA3595G	0.00004167	
3	ACNYW1625J	0.00002778	
1	ALSRP5590C	0.00000160	
5	APYLB9291C	0.00001042	
5	ARCYI6496Z	0.00000463	
7	ARXZK2314V	0.00000463	
3	AYEHS2904M	0.00000556	
9	AZFGJ6792V	0.00002778	
0	BBKRW9735V	0.00000833	
1	BDUFC2928C	0.00002778	
2	BGGII2097V	0.00000556	
3	BGLWT0100N	0.00001984	
4	BHDPO8835E	0.00001042	
5	BHKIO5418Q	0.00000556	
6	BHLLZ7485H	0.00004167	
7	BLCND8992P	0.00002778	
8	BLXHK2495Y	0.00000641	
9	BQJOE1724N	0.00001042	
0	CEBNQ7448Y	0.00001042	
1	CFMVM3542V	0.00000340	
2	CGAQM8807C	0.00000868	
3	COWVG8792S	0.00001984	
4	CREGF0290P	0.00000667	
5	CRTCK3425B	0.00000397	
6	DFTHC1126O	0.00001984	
7	DGDKI8598E	0.00000556	

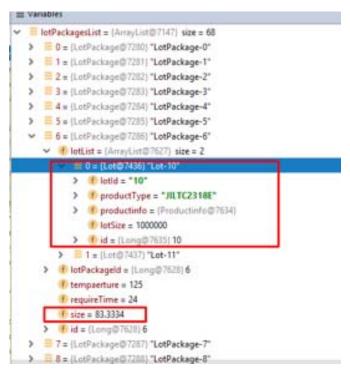
Appendix C2: ProductType (Extracted from DB)

```
IotList = {ArrayList@7138} size = 212
   = 0 = {Lot@7418} "Lot-0"
       f) lotld = "0"
      f productType = "LMCJN3197X"
   productinfo = {Productinfo@7422}
      f) lotSize = 1000000
   > fid = {Long@7146} 0
> = 1 = {Lot@7427} "Lot-1"
> = 2 = {Lot@7428} "Lot-2"
> = 3 = {Lot@7429} "Lot-3"
> = 4 = {Lot@7430} "Lot-4"
> = 5 = {Lot@7431} "Lot-5"
> = 6 = {Lot@7432} "Lot-6"
> = 7 = {Lot@7433} "Lot-7"
> = 8 = {Lot@7434} "Lot-8"
> = 9 = {Lot@7435} "Lot-9"
> = 10 = {Lot@7436} "Lot-10"
> = 11 = {Lot@7437} "Lot-11"
> = 12 = {Lot@7438} "Lot-12"
> = 13 = {Lot@7439} "Lot-13"
> = 14 = {Lot@7440} "Lot-14"
> = 15 = {Lot@7441} "Lot-15"
```

Appendix C3: Initial State (212 lots)



Appendix C4a: 68 Assigned Packages



Appendix C4: Assigned Package 10 (Out of 68 Package)

Appendix D: Batching Rule Code

```
public void generatePackageList (List<Lot> lotList, List<LotPackage> lotPackageSList,
List<LotAssignment> lotAssignment) {
Long LPId = OL;
HashMap<String, Double> filter = new HashMap<>();
HashMap<String, List<Lot>> container = new HashMap<>();
       //Traversing the lot list array
       for (Lot 1 : lotList) {
               int time = l.getProductinfo().getBakeTime();
               int temp = 1.getProductinfo().getTemperature();
              String s = Integer.toString(time) + temp;
       //Add create the package by unique temperate and time combination and add into map
               if (filter.get(s) == null) {
                      filter.put(s, 1.getRequiredCapacity());
                      List<Lot> llist = new ArrayList<>();
                      llist.add(1);
                      container.put(s, llist);
       //The package size exceed the oven volume, create a new package
              else {
                      if (filter.get(s) + l.getRequiredCapacity() > 100.0) {
                             LotPackage lp = new LotPackage(container.get(s), LPId++, temp,
                             time, filter.get(s));
                              lp.setId(lp.getLotPackageId());
                             lotPackagesList.add(lp);
                             filter.remove(s);
                             container.remove(s);
                             filter.put(s, l.getRequiredCapacity());
                             List<Lot> llist = new ArrayList<>();
                             llist.add(1);
                             container.put(s, llist);
                      }
       //The package exist in oven, add the capacity
              else {
                      double n = filter.get(s) + l.getRequiredCapacity();
                      List<Lot> llist = container.get(s);
                      llist.add(1);
                      filter.put(s, n);
                      container.put(s, llist);
                      }
              }
       }
```

```
//Traversing the map add the package into the package list
       for (String key : filter.keySet()) {
              List<Lot> llist = container.get(key);
              LotPackage lp = new LotPackage(container.get(key), LPId++,
              container.get(key).get(0).getProductinfo().getTemperature(),
              container.get(key).get(0).getProductinfo().getBakeTime(), filter.get(key));
              lp.setId(lp.getLotPackageId());
              lotPackagesList.add(lp);
       }
       //Traversing the package list, assign the unique id
       for (LotPackage 1 : lotPackagesList) {
              LotAssignment la = new LotAssignment();
              la.setLotPackage(1);
              la.setId(l.getId());
              lotAssignment.add(la);
       }
}
```

Appendix E: ProductType (Extracted from DB)

	PP		<u> </u>	<u> </u>	
ProductType	Volume	ProductType	Volume	ProductType	Volume
ACHXA3595G	0.00004167	ELXUA2297T	0.00000868	JNCIJ9464N	0.00003788
ACNYW1625J	0.00002778	EPEOQ3051X	0.00004167	JSVFF0225E	0.00002778
ALSRP5590C	0.00000160	EZMHH6514U	0.00000694	KFUST0163P	0.00002778
APYLB9291C	0.00001042	FBGGJ2747C	0.00001042	KGMOF0693Q	0.00002525
ARCYI6496Z	0.00000463	FHQHB3513Q	0.00001736	KILKJ2749Y	0.0000160
ARXZK2314V	0.00000463	FJYSH3065J	0.00000160	KPUGN4232I	0.00000309
AYEHS2904M	0.00000556	FOHRG5276H	0.00000556	KTKBX4839E	0.00004167
AZFGJ6792V	0.00002778	FUJLR9745W	0.00000099	KUQMO4240H	0.00000833
BBKRW9735V	0.00000833	GDXAE0024A	0.00001042	KZVRW4482C	0.00000309
BDUFC2928C	0.00002778	GEIGG3364J	0.00000556	LEQUD1106E	0.00001401
BGGII2097V	0.00000556	GJEYB5591K	0.00000556	LGOQD2787P	0.00001852
BGLWT0100N	0.00001984	GMLCG7480V	0.00000425	LKNAX8537D	0.00000868
BHDPO8835E	0.00001042	GOHUP4823F	0.00000868	LLLCH5198U	0.00000556
BHKIO5418Q	0.00000556	GRKFF2612X	0.00004167	LMCJN3197X	0.00006944
BHLLZ7485H	0.00004167	GSJNU1904S	0.00000926	LQKJX0484R	0.00001984
BLCND8992P	0.00002778	GUJLK0157P	0.00000694	MACNC8175W	0.00001401
BLXHK2495Y	0.00000641	GXOCL6694O	0.00001984	MALXT4126W	0.00000463
BQJOE1724N	0.00001042	HAOWD4659D	0.00000926	MEVBJ7623P	0.00000425
CEBNQ7448Y	0.00001042	HAXFJ6054Z	0.00001401	MGMOY2487X	0.00000463
CFMVM3542V	0.00000340	HCDIE3287X	0.00000556	MHONL8388D	0.00004167
CGAQM8807Q	0.00000868	HPHWC7201Y	0.00000667	MHSTD9087W	0.00001984
COWVG8792S	0.00001984	HRPGV1601E	0.00001852	MHZUP0703L	0.00004167
CREGF0290P	0.00000667	HWGWC8123C	0.00000309	MNREI7200U	0.00001984
CRTCK3425B	0.00000397	HXCVA2371C	0.00000556	MSBFN3632E	0.00000463
DFTHC1126O	0.00001984	HXSEL0087W	0.00000833	MSXWB6402C	0.00000490
DGDKI8598E	0.00000556	HYTHB7799D	0.00000309	MTBXW9690F	0.00004167
DGEPX1320V	0.00000667	ICICQ2140V	0.00001667	MWYBG8201W	0.00000309
DHHSP8518L	0.00000833	IMBDS8589X	0.00004167	MWYUR7713Y	0.00001852
DHVER2972M	0.00001852	ITWGD6876S	0.00000463	MZXTQ7065C	0.00000868
DOPHU4614Q	0.00000160	IWFKE5721L	0.00004167	NHJWD4133V	0.00000340
DQAYA2429F	0.00000694	IXHLX5077Z	0.00004167	NKXIQ7419H	0.00001042
DQOLT4151J	0.00000309	IYBXO3301K	0.00000309	NLIJA20910	0.00001852
DSXXJ1471E	0.00000556	IZQFY5290K	0.00000309	NLYFD0861S	0.00004167
DUMTB6548U	0.00000556	IZXCE0967W	0.00001984	NNWVS3437K	0.00001852
DVYYW0521A	0.00000868	JAPNC71290	0.00001042	NOXPZ4596E	0.00000641
DXYNT3924J	0.00001042	JDNKS5460I	0.00000694	NXAYF1233E	0.00000309
EFLFF0922K	0.00004167	JILTC2318E	0.00004167	NXPZQ2196V	0.00000340
EGZPY1599G	0.00000694	JLIWY4668I	0.00000160	OATAD7946W	0.00000694
EKCOV7230K	0.00002778	JMIHP3141T	0.00002778	OEWVV1941W	0.00000833

ProductType	Volume	ProductType	Volume	ProductType	Volume
OGNPW3564G	0.00001042	TLFTZ7685U	0.00000463	XRVIV5928G	0.00001852
OLTPG9765V	0.00001852	TNMTZ9536S	0.00000641	XTMEP04340	0.00000641
OPVLI2791R	0.00004167	TQPDI3370U	0.00002778	XUEWP9238F	0.00002778
OSIFO7134I	0.00000417	TVLEI1958Q	0.00000667	XZABM9508E	0.00000556
OXPNF3175U	0.00001852	TZXME3979T	0.00000556	YBJAY1812F	0.00002778
OXWYW5999S	0.00000833	UAQAL5502F	0.00000309	YEQSY5057S	0.00000309
PBZUJ0193U	0.00000309	UBAOQ0964L	0.00000926	YGPPA4108I	0.00006944
PCBDV9042T	0.00000309	UEZQO1285W	0.00000463	YLMRS4880I	0.00003788
PCPWW6097L	0.00000309	UGPNC8214R	0.00000694	YMDGT6550X	0.00000556
PIKIT7942X	0.00001042	UGXYU1286G	0.00000833	YNFVG0559N	0.00000868
POVWR0516E	0.00000868	UJKEZ2243J	0.00001984	YOCFE9461J	0.00001042
PPJSS3533C	0.00000926	UKNSR1520H	0.00000160	YRKOG7073J	0.00000868
PSTCF2893D	0.00000833	UNLHK2641P	0.00000556	ZAJUB9415U	0.00004167
PZHGH2590B	0.00000309	UQSFI7003V	0.00004167	ZHSLS1053Q	0.00000833
QBXAQ7463A	0.00000556	UXZWY6716Q	0.00001042	ZHWKP7472G	0.00000694
QCZRA2143O	0.00000556	VAHUK4771M	0.00000309	ZKLJW9348H	0.00001852
QFZAE7569C	0.00000439	VCBUH0118F	0.00004167	ZNUCC9489F	0.00000463
QMRCC1228A	0.00001984	VCDWK7350G	0.00000694	ZOHLK8721D	0.00001852
QMTNJ07210	0.00001042	VDOQD63610	0.00000309	ZQEXD8640N	0.00000833
QQNAB1200C	0.0000160	VHRFM7199D	0.00001042	ZYTGD0971J	0.0000160
QVQBO0440I	0.00001401	VKWRZ0150F	0.00000309		
RBMHS8189G	0.00000490	VOTDJ2592Y	0.00000868		
RCFNP7766A	0.00000099	VVBTG9552O	0.00000463		
RGWMD6191Y	0.00000463	VYCGE5551R	0.00001984		
RMAQD0021E	0.00000833	VYYEL0200D	0.00000278		
RMNKR9893J	0.00000833	VZELO7294C	0.00000160		
RVRSE4446O	0.00000463	WDJHD1887C	0.00000926		
SCDWH9206T	0.00000556	WPWSV8594A	0.00003788		
SDPCN0119Q	0.00000868	WSVLW7887T	0.00001126		
SGVYA6569T	0.00000556	WTIRB5225H	0.0000340		
SHAED4090K	0.00000556	WVFHV5596N	0.00004167		
SJYTK2177B	0.00002778	WWACH5234I	0.00001042		
SMPDM6000H	0.00000833	WYQZK3313M	0.00000556		
SPIPQ7462M	0.00000309	WYSAA6638W	0.00000641		
SPJJE8376D	0.00000309	XATEW2281B	0.00000833		
SRMKK58150	0.00000868	XFYKP7730H	0.00002778		
SVFJY5019T	0.00001852	XKVES6845J	0.00004167		
TBZVN7846Q					
IBZVIV/040Q	0.00004167	XLABU6084F	0.00000309		

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