

# Introduction to Numerical Analysis notes

**Kasi Reddy Sreeman Reddy**  
Undergraduate Physics Student

IIT Bombay

Spring 2020-2021

11-1-21

MA 214

①

 $x_i$  is the interpolation point.

$P_n$  = set of all polynomials with  
degree  $\leq n$

3-2-21

$\rightarrow$  Joseph-Louis Lagrange's Interpolation theorem

$n+1$  distinct points  $\rightarrow$  Unique polynomial  
interpolant in  $P_n$

Proof: 1) Uniqueness,  $w(x) = p(x) - a(x) \in P_n$

↓

Contradiction  $\in [n+1] \text{ roots} \Leftarrow x_i, i=0, 1, \dots, n \text{ are roots}$

2) Existence

$$L_k^n(x) = \frac{(x-x_0) \cdots (x-x_{k-1})(x-x_{k+1}) \cdots (x-x_n)}{(x_k-x_0) \cdots (x_k-x_{k-1})(x_k-x_{k+1}) \cdots (x_k-x_n)}$$

Lagrange  
polynomials

$$P(x) = \sum_{k=0}^n f_k L_k^n(x)$$

$$L_k^n(x) = \prod_{j=0, j \neq k}^n \frac{(x-x_j)}{(x_{k+1}-x_j)}$$

$\rightarrow$  Similarly Lagrange's Zero One Blocks  $\rightarrow$  for sequences

## Matrix method

$$AX = B$$

$$\begin{bmatrix} 1 & x_0 & x_0^2 & \dots & x_0^n \\ 1 & x_1 & & & \\ 1 & x_n & x_n^2 & \dots & x_n^n \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \\ a_2 \\ \vdots \\ a_n \end{bmatrix} = \begin{bmatrix} f_0 \\ f_1 \\ \vdots \\ f_n \end{bmatrix}$$

→ Distance: If  $f, g$  are defined on  $[a, b]$

$$\max_{x \in [a, b]} |f(x) - g(x)| \text{ is the distance.}$$

→  $C[a, b]$  — set of all continuous functions on  $[a, b]$   
 $C^k[a, b] \rightarrow$  "  $k$  times continuously differentiable"  
 $\|f\| := \max_{x \in [a, b]} |f(x)|$

→ Called Maximum Norm

## Weierstrass Approximation theorem

$f \in C[a, b] \Rightarrow \forall \epsilon > 0, \exists$  a polynomial  $p$  such  
 that  $\|f - p\| < \epsilon$  ( $p$  depends on  $f$  and  $\epsilon$ )

## Error equation theorem

$$f \in C^{n+1}[a, b]$$

Let  $p \in P_n$  s.t.  $p(x_i) = f(x_i) \quad i=0, 1, \dots, n$

then  $\forall x \in [a, b], \exists \xi \in [a, b] \text{ s.t.}$

$$f(x) - p(x) = \frac{1}{(n+1)!} f^{(n+1)}(\xi) \prod_{k=0}^n (x - x_k)$$

depends on  $x$ .

Proof

Fix  $x \in [a, b]$

$$\varphi(s) := (f(s) - p(s)) \prod_{k=0}^n (x - x_k) - (f(x) - p(x)) \prod_{k=0}^n (s - x_k)$$

Rolle's theorem  $\Rightarrow \varphi^{(n+1)}$  has at least 1 root  $\xi$ .

$$\Rightarrow f(x) - p(x) = \frac{1}{(n+1)!} f^{(n+1)}(\xi) \prod_{k=0}^n (x - x_k)$$

$$\max_{x \in [a, b]} |f(x) - p(x)| \leq \frac{1}{(n+1)!} \|f^{(n+1)}\| \max_{x \in [a, b]} \prod_{k=0}^n |x - x_k|$$

For Runge's example ( $f(x) = \frac{1}{1+25x^2}, x \in [-5, 5]$ ) use

Chebyshev interpolation points

$$x_j = 5 \cos\left(\frac{(n-j)\pi}{n}\right) \quad j=0, 1, \dots, n$$

These points imply  $\left\| \prod_{k=0}^n (x - x_k) \right\|_{\min} = \left(\frac{b-a}{2}\right)^{n+1} \times \frac{1}{2^n}$

Newton's divided differences

$$p(x) = f[x_0] + f[x_0, x_1](x - x_0) + \dots + f[x_0, x_1, \dots, x_n] \prod_{j=0}^{n-1} (x - x_j)$$

$$f[x_0] := f(x_0)$$

$$f[x_0, \dots, x_{\theta+j}] := \frac{f[x_{\theta+1}, \dots, x_{\theta+j}] - f[x_0, \dots, x_{\theta}]}{x_{\theta+j} - x_0}$$

Ordering is not important

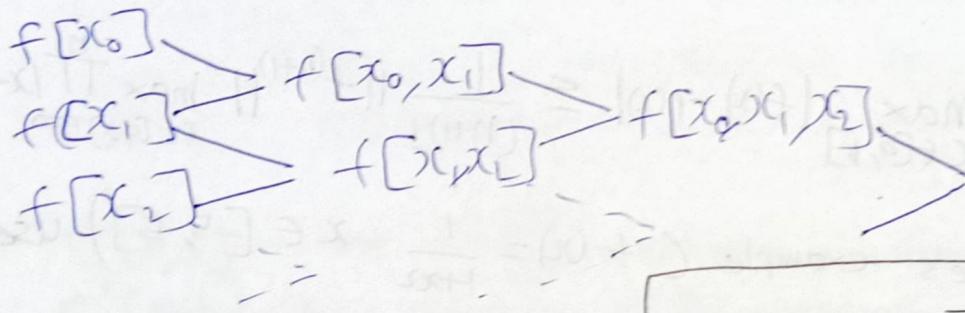
$$(fg)[x_0, \dots, x_n] = \sum_{g=0}^n f[x_0, \dots, x_g] g[x_g, \dots, x_n]$$

$$f[x_0, x_1, \dots, x_n] = \sum_{j=0}^n \frac{f(x_j)}{\prod_{k \in \{0, \dots, n\} \setminus \{j\}} (x_j - x_k)}$$

Divided difference:

$f[x_0, x_1, \dots, x_n]$  = Coefficient of  $x^n$  in  $p \in P_n$

where  $p$  is the interpolant in  $P_n$



Lagrange's recipe

+ - x ÷

$$f[x_0, x_1, x_2, \dots, x_n] = \frac{f^{(n)}(\xi)}{n!}$$

$\xi \in (x_0, x_n)$

Lagrange's	$n$	$2n(n+1)$	$2n^2+n-1$	$n!$
D.D.	$n$	$\frac{3n(n+1)}{2}$	$\frac{n(n+1)}{2}$	

Note: If  $a^{(n)}$  is a sequence of polynomials

s.t.  $\lim_{n \rightarrow \infty} \|f - a^{(n)}\| = 0 \Rightarrow$   $\lim_{n \rightarrow \infty}$  Degree of  $a^{(n)} = \infty$

$\Rightarrow$  But if  $p^{(n)}$  is a sequence of interpolation

points polynomials  $\lim_{n \rightarrow \infty} \|f - p^{(n)}\| \neq 0$ , (They may converge for some points) particular type of

# Splines

Polynomial in  $P_m$

$$x_0 \quad x_1 \quad \dots \quad x_n$$

$$m \leq n$$

$$\psi(x) = a_0^{(1)} + a_1^{(1)}x + a_2^{(1)}x^2 + \dots + a_m^{(1)}x^m \text{ on } [x_0, x_1]$$
$$\vdots$$
$$a_0^{(n)} + a_1^{(n)}x + \dots + a_m^{(n)}x^m \text{ on } [x_{n-1}, x_n]$$

Continuity

1st diff

2h conditions

$$h-1$$

$\boxed{m-1}$  th diff

$$h-1$$

$$= 2^n + (m-1)(h-1)$$

$$= 2^n + mh - h - m + 1$$

$$= \boxed{mh + h - m + 1}$$

We need  $(m+1)h$

$\Rightarrow$  We need  $(m-1)$  more conditions

$m=1 \Rightarrow$  no conditions

$m=3 \Rightarrow$  2 conditions

Error linear splines

$$f \in C^2[a, b]$$

$$h = \max_{1 \leq i \leq n} (x_i - x_{i-1})$$

$$h_i = x_i - x_{i-1}$$

$$h = \max_{1 \leq i \leq n} h_i$$

$$\|f - s_L\| \leq \frac{h^2}{8} \|f''\|$$

## Cubic splines

2 extra conditions

$$S_0''(x_0) = S_{n-1}''(x_n) = 0$$

$$S_i = a_i(x-x_i)^3 + b_i(x-x_i)^2 + c_i(x-x_i) + d_i \quad x \in [x_i, x_{i+1}]$$

Natural cubic splines with  $h_i = h_j \forall i, j$

$$\sigma_i = S''(x_i)$$

$$d_i = f_i$$

$$b_i = \frac{\sigma_i}{2}$$

$$b_0 = b_n = 0$$

$$a_i = \frac{\sigma_{i+1} - \sigma_i}{6h}$$

$$C_i = \frac{f_{i+1} - f_i}{h} - \frac{h}{6} (2\sigma_i + \sigma_{i+1})$$

$$\sigma_{i-1} + 4\sigma_i + \sigma_{i+1} = \frac{6}{h^2} (f_{i-1} - 2f_i + f_{i+1})$$

$$\begin{matrix} 4 & 1 & 0 & \dots & 0 & 0 & 0 \\ 1 & 4 & 1 & & & & \\ 0 & 1 & 4 & & & & \\ \vdots & & & & & & \\ 0 & 0 & 0 & & & & \\ 0 & 0 & 0 & & & & \end{matrix}$$

$$\begin{pmatrix} 4 & 1 & 0 & & & & \\ & 1 & 4 & 1 & & & \\ & & 1 & 4 & 1 & & \\ & & & 0 & 1 & 4 & \\ & & & & 1 & 4 & 1 \\ & & & & & 1 & 4 \end{pmatrix} \begin{pmatrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \\ \vdots \\ \sigma_{n-1} \\ \sigma_{n-1} \end{pmatrix} = \frac{6}{h^2} \begin{pmatrix} f_0 - 2f_1 + f_2 \\ f_1 - 2f_2 + f_3 \\ \vdots \\ f_{n-2} - 2f_{n-1} + f_n \end{pmatrix}$$

~~Diagonally dominant~~ Invertible

## Natural cubic spline error

7

$$\|f - s\| \leq \frac{h^4}{8} \|f^{(IV)}\|$$

$$\boxed{\|f - s\| \leq \|f'' - s''\| \left\| \frac{h^2}{8} \right\|}$$

Proof  $g := f - s$  on  $[x_i, x_{i+1}]$

$$\Rightarrow g(x_i) = 0, g(x_{i+1}) = 0 \Rightarrow$$

Linear spline for  $g$  is just 0 function.

Using linear error  $\|g - 0\| \leq \frac{h^2}{8} \|g''\|$

$$\Rightarrow \|f - s_i\| \leq \|f'' - s_i''\| \left\| \frac{h^2}{8} \right\|$$

$$\Rightarrow \boxed{\|f - s\| \leq \frac{h^2}{8} \|f'' - s''\|}$$

## 2A Numerical Integration

$$\int_a^b f(x) dx \approx \int_a^b p(x) dx$$

Interpolation for

$$a = x_0 < x_1 < x_2 < \dots < x_n = b$$

$$x_i = a + ih, i=0, 1, \dots$$

Newton and Cotes integration formulas  
Take

# Newton and Cotes formula for integration

8

Take equally spaced points  $x_i = a + ih$   $[0 \leq i \leq h]$

$$h = \frac{b-a}{n} \text{ or } nh = b-a$$

$$P(x) = \sum_{i=0}^n f(x_i) L_i(x), \quad L_i(x) = \prod_{k=0, k \neq i}^n \frac{(x - x_k)}{(x_i - x_k)}$$

$$t = \frac{x-a}{h} = n \left( \frac{x-a}{b-a} \right) \in [0, n]$$

$$L_i(x) = \prod_{k=0, k \neq i}^n \left( \frac{t-k}{i-k} \right) := \ell_i(t)$$

$$\int_a^b f(x) dx \approx \int_a^b \sum_{i=0}^n f(x_i) L_i(x) dx$$

$$= \sum_{i=0}^n f(x_i) \int_a^b L_i(x) dx$$

$$\int_a^b f(x) dx \approx h \sum_{i=0}^n w_i f(x_i)$$

$$w_i := \int_0^n \ell_i(t) dt$$

depends only  
on n.

$$\sum_{i=0}^n w_i = n$$

$$w_k = w_{n-k}$$

$n=1$  Trapezoidal rule

$$\frac{h}{2} (f_0 + f_1)$$

$n=2$  Simpson's rule

$$\frac{h}{3} (f_0 + 4f_1 + f_2)$$

$n=3$  Simpson's  $\frac{3}{8}$  rule

$$\frac{3h}{8} (f_0 + 3f_1 + 3f_2 + f_3)$$

$n=4$  Boole's rule  
Milne's

$$\frac{2h}{45} (7f_0 + 32f_1 + 12f_2 + 32f_3 + 7f_4)$$

## 2B) Error in Newton-Cotes formula

9

$$\mathcal{I}_f = \int_a^b f(x) dx \quad \mathcal{I}_{P_n} = \int_a^b P_n(x) dx$$

If  $f \in C^{n+1}[a, b]$ ,  $f(x) - P_n(x) = \frac{1}{(n+1)!} f^{(n+1)}(\xi) \prod_{i=0}^n (x-x_i)$   
 $\xi \in [a, b]$

$$|\mathcal{I}_f - \mathcal{I}_{P_n}| = \left| \int_a^b (f(x) - P_n(x)) dx \right| \leq \int_a^b |f(x) - P_n(x)| dx$$

$$|\mathcal{I}_f - \mathcal{I}_{P_n}| \leq \frac{\|f^{(n+1)}\|}{(n+1)!} \int_a^b \prod_{i=0}^n |x-x_i| dx$$

For  $n=1$

$$|\mathcal{I}_f - \mathcal{I}_{P_1}| \leq \frac{\|f''\| (b-a)^3}{2}$$

$$\text{For } n=2 \quad |\mathcal{I}_f - \mathcal{I}_{P_2}| \leq \frac{\|f'''(x)\|}{192} (b-a)^4$$

2C In Runge example, Newton-Cotes method  
 does not converge as  $n \rightarrow \infty$

## Gaussian quadrature

$$G_n(f) = \sum_{i=0}^n w_i f(x_i)$$

$$w_i = \int_a^b (L_i(x))^2 dx = \int_a^b \prod_{k=0, k \neq i}^n \frac{(x-x_k)^2}{(x_i-x_k)^2} dx$$

Not equally spaced

Weights non negative

$$\text{Let } f \in C[a, b] \Rightarrow \lim_{n \rightarrow \infty} |G_n(f) - \mathcal{I}_f| = 0$$

# Composite Newton Cotes (converges)

10

## Composite trapezoidal

$$C_{P_1}(f) = h \left( \frac{1}{2} f(a) + f(a+h) + \dots + f(a+(m-1)h) + \frac{1}{2} f(b) \right)$$

$h = \frac{b-a}{m}$

Error:

$$|C_{P_1}(f) - I_f| \leq m \times \frac{1}{12} \|f''\| h^3 = \frac{b-a}{12} \|f''\| h^3$$

## Composite Simpson's

$$C_{P_2}(f) = \frac{h}{3} \left( f(a) + 4f(a+h) + 2f(a+2h) + \dots + 2f(a+(2m-2)h) + 4f(a+(2m-1)h) + f(b) \right)$$

$h = \frac{b-a}{2m}$

$$|C_{P_2}(f) - I_f| \leq \frac{(b-a)}{180} \|f''\| h^3$$

## 3A ODEs

$$y' = \sqrt{y} \quad y(0) = 0$$

Solution:  $y(t) = \begin{cases} \frac{(t-k)^2}{4} & t \geq k \\ 0 & t < k \end{cases}$

or  $y(t) = 0 \quad \forall y$   
 $y' = \sqrt{y}, \quad y(0) \neq 0 \Rightarrow$  Unique solution.

Notation:

$$y' = f(t, y) \Rightarrow 1^{\text{st}} \text{ order ODE}$$

## n<sup>th</sup> order ODE

$$F(t, y, y', \dots, y^{(n)}) = 0$$

$$\vec{y} = \begin{pmatrix} y_1 \\ y_2 \\ \vdots \\ y_n \end{pmatrix} = \begin{pmatrix} y \\ y_1 \\ \vdots \\ y^{(n-1)} \end{pmatrix}$$

$$\Rightarrow \boxed{\vec{y}' = f(t, \vec{y})}$$

### Existence - Peano

$f: \mathbb{R} \times \mathbb{R} \rightarrow \mathbb{R}$  s.t it is continuous in ~~( $s, t$ )~~

$(-s, s) \times (y_0 - n, y_0 + n)$  for some  $s > 0, n > 0$   
 $\Rightarrow \exists \epsilon > 0$  and function  $y \in C^1(\epsilon, \epsilon)$  s.t  
 $y' = f(t, y); y(0) = y_0$

Lipschitz continuity:  $g: [a, b] \rightarrow \mathbb{R}$

$$L > 0 \quad |g(x) - g(y)| \leq L|x - y| \quad \forall x, y \in [a, b]$$

### Uniqueness - Cauchy-Lipschitz

$f(t, y) \rightarrow$  continuous in both variables and  
 Lipschitz continuous in  $y$  variable.

$\exists \epsilon > 0$  and a unique function  $y \in C^1(-\epsilon, \epsilon)$   
 s.t  $y' = f(t, y); y(0) = y_0$

Assume Lipschitz from now

$$y' = f(t, y); y(0) = y_0 \Leftrightarrow y(t) = y_0 + \int_0^t f(s, y(s)) ds$$

Rather than solving integration at all time points we do at some points.

### Mesh points

$$N > 0$$

$$h = \frac{T}{N}$$

$$t_n = nh$$

$$h = 0, 1, \dots, N$$

$$y_n \approx y(t_n) \quad h > 0$$

$$y_0 = y(t_0)$$

$$y(t_{n+1}) = y(t_n) + \int_{t_n}^{t_{n+1}} f(s, y(s)) ds$$

### Eulers method

approximate  $f(s, y(s))$  on  $[t_n, t_{n+1}]$

$$f(s, y(s)) \approx f(t_n, y_n) \quad s \in [t_n, t_{n+1}]$$

$$y_{n+1} = y_n + hf(t_n, y_n)$$

### Explicit

### Trapezoidal method

### Implicit

$$y_{n+1} = y_n + \frac{h}{2} (f(t_n, y_n) + f(t_{n+1}, y_{n+1}))$$

$$y_{n+1} - \frac{h}{2} f(t_{n+1}, y_{n+1}) = y_n + \frac{h}{2} f(t_n, y_n)$$

3B

Order

Let  $y_{n+1} = F(t, f, y_0, y_1, \dots, y_n, y_{n+1})$  be a recurrence relation, the order is p if

$$y(t_{n+1}) - F(t, f, y(t_0), y(t_1), \dots, y(t_n), y(t_{n+1})) = O(h^{p+1})$$

Euler method

$$\begin{aligned} y(t_{n+1}) - y(t_n) - hf(t_n, y(t_n)) &= y(t_n + h) - y(t_n) - hf(t_n, y(t_n)) \\ &= y(t_n) + hy'(t_n) + \frac{h^2}{2} y''(\xi) - y(t_n) - hf(t_n, y(t_n)) \\ &= \frac{h^2}{2} y''(\xi) = O(h^2) \end{aligned}$$

Order = 1Trapezoidal method

$$\begin{aligned} y(t_{n+1}) - y(t_n) - \frac{h}{2} (f(t_n, y(t_n)) + f(t_{n+1}, y(t_{n+1}))) \\ &= y(t_n + h) - y(t_n) - \frac{h}{2} (y'(t_n) + y'(t_n + h)) \\ &= y(t_n) + hy'(t_n) + \frac{h^2}{2} y''(t_n) + \frac{h^3}{6} y'''(\xi) \\ &\quad - y(t_n) - \frac{h}{2} (2y'(t_n) + hy''(t_n) + \frac{h^2}{2} y'''(t_n)) \\ &= \frac{h^3}{6} y'''(\xi) - \frac{h^3}{4} y'''(h) = O(h^3) \end{aligned}$$

Order = 2

Method of order  $p$  recoveres exactly every polynomial with degreee  $\leq p$ .

Global error:  $y_{n,h} = y_n$   $E_n = \max_{i=0,1,\dots,[\frac{T}{h}]} |e_{i,h}|$  Global error

New notation

$$e_{n,h} = y_{n,h} - y(t_n)$$

In  $[0, T]$  there are  $[\frac{T}{h}] + 1$  equally spaced mesh points

Convergence: A numerical method is convergent if

$$\lim_{h \rightarrow 0^+} \max_{n=0,1,\dots,[\frac{T}{h}]} |e_{n,h}| = 0$$

Euler method

If  $|e_{n,h}| \leq Ch^p$  then the convergence order is  $p$ . (also called order of Global error)

For Euler's method

$$\begin{aligned} e_{n+1,h} &= y_{n+1,h} - y(t_{n+1}) \\ &= y_{n,h} - h f(t_n, y_{n,h}) - y(t_n) - h f(t_n, y(t_n)) + o(h) \\ &= y_{n,h} + h f(t_n, y_{n,h}) - y(t_n) - h f(t_n, y(t_n)) + o(h) \\ &= e_{n,h} + h (f(t_n, y_{n,h}) - f(t_n, y(t_n))) + o(h^2) \\ \Rightarrow |e_{n+1,h}| &\leq |e_{n,h}| + h |f(t_n, y_{n,h}) - f(t_n, y(t_n))| + o(h^2) \end{aligned}$$

$$\Rightarrow |e_{n+1,h}| \leq |e_{n,h}| (1 + Lh) + ch^2 \quad (L = \text{Lipschitz constant})$$

$$\Rightarrow |e_{n,h}| \leq \frac{Ch}{L} ((1 + Lh)^n - 1) \quad (\text{By induction})$$

$$\Rightarrow |e_{n,h}| \leq \frac{ch}{L} ((1+Lh)^n - 1) \leq \frac{ch}{L} (e^{LT} - 1)$$

Order=1

$$\Rightarrow \lim_{h \rightarrow 0} |e_{n,h}| = 0 \quad \left( (1+Lh)^n < e^{nLh} \leq e^{\left[\frac{L}{h}\right]Lh} \leq e^{LT} \right)$$

Trapezoidal rule

$$e_{n+1,h} = e_{n,h} + \frac{h}{2} (f(t_n, y_{n,h}) - f(t_n, y(t_n))) + \frac{h}{2} (f(t_{n+1}, y_{n+1,h}) - f(t_{n+1}, y(t_{n+1}))) + O(h^3)$$

$$\Rightarrow |e_{n+1,h}| \leq |e_{n,h}| + \frac{h}{2} L |e_{n,h}| + \frac{h}{2} L |e_{n+1,h}| + ch^3$$

$$\Rightarrow |e_{n+1,h}| \leq \left(\frac{1+Lh}{1-Lh}\right) |e_{n,h}| + \left(\frac{1}{1-Lh}\right) ch^3$$

$$\Rightarrow |e_{n,h}| \leq \frac{ch^2}{L} \left[ \left(\frac{1+Lh}{1-Lh}\right)^n - 1 \right] \quad n=0, 1, \dots, \left[\frac{T}{h}\right]$$

$$\Rightarrow |e_{n,h}| \leq \frac{ch^2}{L} e^{\frac{nLh}{1-Lh}} \leq \frac{ch^2}{L} e^{\frac{TL}{1-\frac{Lh}{2}}}$$

$$\Rightarrow \lim_{h \rightarrow 0} |e_{n,h}| = 0 \quad (\text{convergence order}=2)$$

3c

One Step method: approximation at  $n$ th step depends on approx at  $(n-1)$ th step

General S-step method:  $S \geq 1$

$$\sum_{k=0}^S \alpha_k y_{n+k} = h \sum_{k=0}^S \beta_k f_{n+k}$$

$$f_{n+k} = f(t_{n+k}, y_{n+k})$$

$\alpha_i, \beta_j$  are constants.

$\beta_S = 0$	Explicit
$\beta_S \neq 0$	Implicit

Need  $y_0, \dots, y_{n-1}$  to start

Sometimes these are generated using a 1-step method like Euler.

Adams' s-step method or (Adams - Bashforth) (Explicit)  
(Order  $s$ )

$$t_n = nh$$

Let  $\Psi(t)$  be the polynomial interpolation with

$$\Psi_i = f_i = f(t_i, y_i) \quad i = n, n+1, \dots, n+s-1$$

$$\Psi(t) = \sum_{k=0}^{s-1} \Psi_k(t) f(t_{n+k}, y_{n+k}) \quad (\text{degree } s-1)$$

$$\Psi_k(t) = \prod_{\lambda=0, \lambda \neq k}^{s-1} \frac{(t - t_{n+\lambda})}{(t_{n+k} - t_{n+\lambda})}$$

Substitute  $\Psi(t)$  for  $f(t, y(t))$  in  $(t_{n+s}, t_{n+s})$

$$y(t_{n+s}) - y(t_{n+s-1}) = \int_{t_{n+s-1}}^{t_{n+s}} f(\tau, y(\tau)) d\tau$$

$$\Rightarrow y_{n+s} - y_{n+s-1} = h \sum_{k=0}^{s-1} \beta_k f(t_{n+k}, y_{n+k})$$

$$\beta_k = \frac{1}{h} \int_0^h \Psi_k(t_{n+s-1} + \tau) d\tau \quad (\text{Independent of } n)$$

2 step  $y_{n+2} - y_{n+1} = \frac{h}{2} (3f_{n+1} - f_n)$

Adams Bashforth 3 step 4 step

$$\boxed{3} \quad y_{n+3} = y_{n+2} + h \left( \frac{23}{12} f_{n+2} - \frac{16}{12} f_{n+1} + \frac{5}{12} f_n \right)$$

$$\boxed{4} \quad y_{n+4} = y_{n+3} + h \left( \frac{55}{24} f_{n+3} - \frac{59}{24} f_{n+2} + \frac{37}{24} f_{n+1} - \frac{9}{24} f_n \right)$$

$$\boxed{1} \quad y_{n+1} = y_n + h f_n$$

(depends on s)  ~~$\beta_{s-j}$~~   $\beta_{s-j-1} = \frac{(-1)^j}{\prod_{i=0}^{s-j-1} i!} \int_0^s f(t) dt$

$$\beta_{s-j-1} = \frac{(-1)^j}{\prod_{i=0}^{s-j-1} i!} \int_0^s f(t) dt$$

$$j = 0, \dots, s-1$$

Adams-Moulton method ( $s$  step =  $s+1$  order) (Implicit)

$$y_{n+s} = y_{n+s-1} + h \sum_{k=0}^s \beta_k f_{n+k}$$

$$\beta_{s-j} = \beta_{s-j} = \frac{(-1)^j}{\prod_{i=0}^{s-j} i!} \int_0^s f(t) dt \quad j = 0, \dots, s$$

$$\Psi_i = f_i \quad i = n, n+1, \dots, n+s$$

$$y_n = y_{n-1} + h f(t_n, y_n)$$

$$1 \quad y_{n+1} = y_{n+1} + h (f_{n+1} + f_n)$$

$$2 \quad y_{n+2} = y_{n+1} + h \left( \frac{9}{12} f_{n+2} + \frac{2}{3} f_{n+1} - \frac{1}{12} f_n \right)$$

$$3 \quad y_{n+3} = y_{n+2} + h \left( \frac{9}{24} f_{n+3} + \frac{19}{24} f_{n+2} - \frac{5}{24} f_{n+1} + \frac{1}{24} f_n \right)$$

Order of multistep Adams-Basforth method

$$\sum_{k=0}^s \alpha_k y_{n+k} = h \sum_{k=0}^s \beta_k f_{n+k} \text{ is } s\text{-step and order } p$$

$$\sum_{k=0}^s \alpha_k = 0$$

$$\Rightarrow \sum_{k=0}^s k^m \alpha_k = m \sum_{k=0}^{m-1} k^{m-1} \beta_k \quad m = 1, 2, \dots, p$$

$$\sum_{k=0}^s k^{p+1} \alpha_k \neq (p+1) \sum_{k=0}^s k^p \beta_k$$

Proofs:

$$\sum_{k=0}^s \alpha_k y(t_{n+k}) - h \sum_{k=0}^s \beta_k f(t_{n+k}, y(t_{n+k}))$$

$$= \sum_{k=0}^s \alpha_k \left( y(t_n) + \sum_{m=1}^{\infty} \frac{k^m h^m}{m!} y^{(m)}(t_n) \right) - h \sum_{k=0}^s \beta_k y'(t_n + kh)$$

$$y'(t_n + kh) = \sum_{m=0}^{\infty} (m+1) \frac{k^m h^m}{(m+1)!} y^{(m+1)}(t_n)$$

$$= \sum_{m=1}^{\infty} m \frac{k^{m-1} h^{m-1}}{m!} y^{(m)}(t_n)$$

$$= \left( \sum_{k=0}^s \alpha_k \right) + \sum_{m=1}^{\infty} \left( \sum_{k=0}^s k^m \alpha_k - m \sum_{k=0}^s k^{m-1} \beta_k \right) \frac{h^m}{m!} y^{(m)}(t_n)$$

$$= O(h^{p+1})$$

Convergence

1st characteristic polynomial:  $e(z) = \sum_{k=0}^s \alpha_k z^k$

2nd characteristic polynomial:  $\sigma(z) = \sum_{k=0}^s \beta_k z^k$

Dahlquist Equivalence: An  $s$ -step method is convergent

iff 1) is of order  $p \geq n$

2) roots of  $e(z)$  lie in the closed disk

$$|z| \leq 1$$

with any which lie on the unit circle being SIMPLE.

Dahlquist's Barrier theorem: For  $s$ -step order cannot exceed

1) S<sub>1</sub> if  $s$  is ODD

2) S<sub>2</sub> if  $s$  is EVEN

# Runge - Kutta Method

Explicit if

A is  
strictly  
lower  
triangular

4-stage

$$y_{n+1} = y_n + h \sum_{i=1}^4 b_i f(t_n + c_i h, \tilde{y}_i)$$

$$\tilde{y}_i = y_n + h \sum_{j=1}^{i-1} a_{ij} f(t_n + c_j h, \tilde{y}_j)$$

$A = (a_{i,j})$  is called RK Matrix

$b_i$  = RK weights     $c_i$  = RK Nodes

$$c_1 = 0$$

## RK Tableau

$$\begin{array}{c|ccccc} c & A & & & & \\ \hline & b_1 & b_2 & \dots & b_{s-1} & b_s \end{array} = \begin{array}{c|ccccc} 0 & & & & & \\ \hline c_2 & a_{21} & & & & \\ c_3 & a_{31} & a_{32} & & & \\ \hline c_s & a_{s1} & a_{s2} & \dots & a_{ss-1} & \end{array} \quad S = 4$$

Consistent

$$\text{iff } \sum_{i=1}^s b_i = 1$$

Popular:  $\sum_{j=1}^{i-1} a_{ij} = c_i$

## Standard RK or RK4

$$\begin{array}{c|ccccc} 0 & & & & & \\ \frac{1}{2} & \frac{1}{2} & & & & \\ \frac{1}{2} & 0 & \frac{1}{2} & & & \\ \hline 1 & 0 & 0 & 1 & & \\ \hline & \frac{1}{6} & \frac{1}{3} & \frac{1}{3} & \frac{1}{6} & \end{array}$$

$$y_{n+1} = y_n + \frac{1}{6} h (k_1 + 2k_2 + 2k_3 + k_4)$$

$$t_{n+1} = t_n + h$$

$$k_1 = f(t_n, y_n)$$

$$k_2 = f\left(t_n + \frac{h}{2}, y_n + h \frac{k_1}{2}\right)$$

$$k_3 = f\left(t_n + \frac{h}{2}, y_n + h \frac{k_2}{2}\right)$$

$$k_4 = f(t_n + h, y_n + h k_3)$$

4A

 $M_{m,n}(\mathbb{R}) = \text{All } m \times n \text{ matrices.}$ 
 $M_{m,n}(\mathbb{R}) = \text{All } m \times n \text{ matrices}$ 

$A = (A_{ij})_{1 \leq i \leq m, 1 \leq j \leq n}$

 $\in M_{m,n}(\mathbb{R})$ 
Objective

$AX = b$

$A \in M_{m,n}(\mathbb{R})$

Underdetermined

Overdetermined

$1] m < n$

more unknowns

→ No solution

→ Infinitely many

$2] m > n$

more equations

→ No solution

→ Single solution

→ Infinite

$3] m = n$

No solution

Single solution

Infinite

$A = A^i_j, X = x^j$

$A^i_j x^j = b^i$

$c^i_k = A^i_j B^j_k$

 $A$  is not invertible
 $\Leftrightarrow \text{Ker}(A) \neq \{0\} \Leftrightarrow A \text{ is singular}$ 

$\text{Ker}(A) = \{v \in V \mid Av = 0\}$

No solution or  $\infty$ If  $Ay = b$  then  $Ax = b$ 

$x = y + cg \quad c \in \mathbb{R}, g \in \text{Ker}(A) - \{0\}$

For square  
matrices
 $\text{Ran } k(A) + \text{nullity } (A) = n$ 

dimension of rowspace

dimension of kernel.

null or space

 $A$  is invertible

$x = A^{-1} b$

(Hard)

Cramer's rule

$x_i = \frac{\det(A_i)}{\det(A)}$

$A_i$  is obtained  
by replacing  $i$ th  
column of  $A$  by  $b$ .

Diagonal matrices

$$x_i = \frac{b_i}{A_{ii}}$$

Unitary matrix ( $A^{-1} = A^T$ )

$$|A|=1$$

Upper triangular Matrix

$$x_n = \frac{b_n}{A_{nn}}$$

$$x_i = \sum_{j=1}^n A_{ji} b_j$$

(Orthogonal  
 $\Rightarrow A^T = A^{-1}$ )

(Unitary  
 $\Rightarrow A^* = A^{-1}$ )

$$(A^i_j = 0 \forall i > j)$$

$$x_i = \frac{b_i - \sum_{j=i+1}^n a_{ij} x_j}{a_{ii}} \quad (\text{Back substitution algorithm})$$

(similarly for lower)

Gauss method (Only for invertible matrices)

$$\{A^{(1)}, \dots, A^{(n)}\} \quad \{b^{(1)}, \dots, b^{(n)}\}$$

$$A^{(1)} = A \quad A^{(k)} = \text{Upper triangular}$$

$$A_{ij}^{(k+1)} = A_{ij}^{(k)} - m_{ik} A_{kj}^{(k)}$$

for  $i=k+1, \dots, n$   
 $j=1, \dots, n$

$$b_i^{(k+1)} = b_i^{(k)} - m_{ik} b_k^{(k)}$$

for  $i=k+1, \dots, n$

$$m_{ik} = \frac{A_{ik}^{(k)}}{A_{kk}^{(k)}}$$

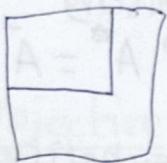
Theorem:  
Let  $A$  be a  $n \times n$  matrix. There exists at least one non-singular matrix  $B$  such that the product  $BA$  is upper-triangular.

$$BAx = Bb$$

For non-invertible matrices the diagonal contains at least 1 zero element.

Pivoting: If matrix is invertible but  $A_{ii} = 0$  for some  $i$ , we need to do pivoting.

Diagonal submatrices



$$\Delta^k = \begin{pmatrix} A_{11} & \cdots & A_{1k} \\ \vdots & \ddots & \vdots \\ A_{k1} & \cdots & A_{kk} \end{pmatrix}$$

Theorem: Let all diagonal submatrices of order  $k=1, 2, \dots, n$  [are invertible] for  $A$ . Then Gaussian elimination doesn't need any pivoting strategy.

GB

Factorization

$$A = BC$$

$$B(Cx) = b$$

$$B y = b \quad \text{and} \quad y = cx$$

$$L_i i = 1 \Rightarrow \text{Unique}$$

LU factorization

$(A = LU)$  It should be

LU factorization  $\Leftrightarrow$  Gaussian elimination without pivoting.

$$U = A^{(n)}$$

$$L_{ij} = \begin{cases} 1 & i=j \\ m_{ij} & i>j \\ 0 & i<j \end{cases}$$

$$U_{1j} = A_{1j}$$

$$U_{ij} = A_{ij} - \sum_{k=1}^{j-1} L_{ik} U_{kj}$$

$$L_{j+1,j} = A_{j+1,j} - \sum_{k=1}^{j-1} L_{j+1,k} U_{kj}$$

$$\begin{array}{c} i=1, \dots, n \\ \cancel{j=i, \dots, n} \end{array}$$

$$L_{ij} = A_{ij} - \sum_{k=1}^{j-1} L_{ik} U_{kj} \quad i=2, \dots, n$$

23

$$U_{jj}$$

$$j=1, \dots, i-1$$

$$A_{ij} = \sum_{k=1}^{\min(i,j)} L_{ik} U_{kj}$$

Cholesky Factorisation ( $A = BB^T$ )      Uniax

Only for Hermitian and positive definite matrices.  
 ↓  
 symmetric

Positive definite  $\Leftrightarrow$  all  $\lambda_i > 0$

Positive semidefinite  $\Leftrightarrow \lambda_i \geq 0$

$$\begin{aligned} & \forall y \neq 0 \\ & \langle Ay, y \rangle > 0 \\ & \langle Ay, y \rangle \geq 0 \end{aligned}$$

$$B_{jj} = \sqrt{A_{jj} - \sum_{k=1}^{j-1} B_{j,k} B_{j,k}^T}$$

$$B_{ij} = \frac{1}{B_{jj}} \left( A_{ij} - \sum_{k=1}^{j-1} B_{i,k} B_{j,k}^T \right) \text{ for } i > j$$

for strictly positive semidefinite  $B_{jj} = 0$  for some  $j$   
 and the method fails.

for indefinite matrix  $\Rightarrow B_{jj} < 0$  for some  $j$

Test for +Ve definiteness: Fails if  
 $B_{jj} = 0$  or  $B_{jj}^2 < 0$

24

QR Factorization ( $A = QR$ ) or  $QV$  (for any ~~method~~ matrix)

$Q = \text{Orthogonal}$  ( $Q^T = Q^{-1}$ )

$R = \text{Upper triangular}$

Theorem: If  $\det(A) \neq 0 \Rightarrow A = QR$  such that  $\det(R) \neq 0$ .

If  $R$  is assumed to have the diagonal entries then  $(Q, R)$  is unique.

Algorithm Let  $A_1, A_2, \dots, A_n$  be columns of  $A$ .

$$q_i = A_i - \sum_{k=1}^{i-1} \langle q_k, A_i \rangle q_k$$

$\|A_i - \sum_{k=1}^{i-1} \langle q_k, A_i \rangle q_k\|$

$q_i$  is a column

$$A_i = \sum_{k=1}^n r_{ki} q_k$$
$$R_{ki} = \begin{cases} \langle q_k, A_i \rangle & \text{if } 1 \leq k \leq i-1 \\ \|A_i - \sum_{k=1}^{i-1} \langle q_k, A_i \rangle q_k\| & \text{if } k=i \\ 0 & \text{if } k > i \end{cases}$$

or

$$R_{ki} = \begin{cases} \langle q_k, A_i \rangle & \text{if } 1 \leq k \leq i \\ 0 & \text{if } k > i \end{cases}$$

#### AC Norm on $\mathbb{R}^n$

$$1) \|x\| > 0 \quad 2) \|\alpha x\| = |\alpha| \|x\| \quad 3) \|x+y\| \leq \|x\| + \|y\|$$

25

$\ell^p$ -norm for  $p \geq 1$

$$\|x\|_{\ell^p} = \left( \sum_{i=1}^n |x_i|^p \right)^{\frac{1}{p}}$$

$\ell^\infty$  norm

$$\|x\|_{\ell^\infty} = \max_{1 \leq i \leq n} |x_i|$$

$$\boxed{\|x\|_{\ell^2} = \|x\|_2 = \sqrt{\langle x, x \rangle}}$$

Equivalence

$$m\|x\| \leq \|x\|' \leq M\|x\| \quad \forall x \in \mathbb{R}^n$$

$$\downarrow \quad \downarrow$$

$$\text{Independent of } x$$

$$\frac{1}{M}\|x\| \leq \|x\|' \leq \frac{1}{m}\|x\|'$$

$\rightarrow$  All norms are equivalent on  $\mathbb{R}^n$

$$\|x\|_{\ell^\infty} \leq \|x\|_{\ell^p} \leq \sqrt[p]{n} \|x\|_{\ell^\infty}$$

$$\|x\|_{\ell^2} \leq \|x\|_{\ell^1} \leq \sqrt{n} \|x\|_{\ell^2}$$

Matrices

$$\|A\|_{\ell^p} = \left( \sum_{j=1}^n \sum_{m=1}^m |A_{ij}|^p \right)^{\frac{1}{p}}$$

$\boxed{p=2 = \text{Frobenius norm}}$

$$\|A\|_{\ell^\infty} = \max_{1 \leq i \leq m, 1 \leq j \leq n} |A_{ij}|$$

$$\|A\|_{L_{p,q}} = \left( \sum_{i=1}^n \left( \sum_{j=1}^m |A_{ij}|^p \right)^{\frac{q}{p}} \right)^{\frac{1}{q}}$$

$$\|AB\|_{\ell^2} \leq \|A\|_2 \|B\|_2$$

Matrix Norm

or Submultiplicative norm

$$\boxed{3) + \|AB\| \leq \|A\| \|B\| \quad \forall A, B \in M_n(\mathbb{R})}$$

Ex:  $\|AB\|_2$  or  $\|A\|_2$

$\boxed{1 \leq \|I_n\|}$

$\| \cdot \|_{\ell^\infty}$  is not a matrix norm.

Subordinate matrix norm:  $\|\cdot\|$  be a vector norm, 26

$$\|A\| = \sup_{x \in \mathbb{R}^n \setminus \{0\}} \frac{\|Ax\|}{\|x\|} = \sup_{\substack{x \in \mathbb{R}^n \\ \|x\|=1}} \|Ax\|$$

$$\Rightarrow \|Ax\| \leq \|A\| \|x\| \quad \& \quad \|AB\| \leq \|A\| \|B\|$$

Proof:  $\|AB\| \leq \|A\| \sup_{\|x\|=1} \frac{\|Bx\|}{\|x\|} \leq \|A\| \|B\|$

Subordinate matrix norm  $\Rightarrow \|I_n\| = 1$

Forbenius norm is not subordinate.  $\|I_n\|_F = \sqrt{n}$  But vector norm.

$$\|A\|_F = \sup_{x \in \mathbb{R}^n} \frac{\|Ax\|_F}{\|x\|_F}$$

$\|\cdot\|_F$  forbenius subordinate

1)  $A$  is unitary ( $A^{-1} = A^T$ )  $\Rightarrow \|A\|_F = 1$

2)  $A = \text{Unitary} \Rightarrow \|AB\|_F = \|BA\|_F = \|B\|_F$

due to  $\|x\|_F = \sqrt{\langle x, x \rangle}$

Diagonal matrix  $\Rightarrow \|A\|_F = \max_{1 \leq i \leq n} |a_{ii}|$

4D

Normal Matrix:  $AA^* = A^*A$

$\Rightarrow \exists U$  st  $U$  is unitary and

$$A = U \text{diag}(\lambda_1, \dots, \lambda_n) U^*$$

If  $A = \underline{\text{normal}}$   $\Rightarrow \|A\|_F = \|\text{diag}(\lambda_1, \dots, \lambda_n)\|_F$

$$\|A\|_F \geq 1$$

$$= \max_{1 \leq i \leq n} |\lambda_i|$$

$$= e(A)$$

Spectral radius: For any matrix  $A$

↓  
not a matrix  
norm.

$$\rho(A) = \max_{1 \leq i \leq n} |\lambda_i|$$

→ For any matrix norm  $\rho(A) \leq \|A\|$

Converse: For any  $A \exists$  a subordinate matrix such that  $\|A\| \leq \rho(A) + \epsilon$   
 $\|.\|$  depends on  $A$  ~~& E.~~.

$$\|A\|_1 = \max_{1 \leq j \leq n} \left( \sum_{i=1}^n |a_{ij}| \right) \quad (\text{column sum})$$

$$\|A\|_\infty = \max_{1 \leq i \leq n} \left( \sum_{j=1}^n |a_{ij}| \right) \quad (\text{Row sum})$$

$$\Rightarrow \rho(A) \leq \min(\|A\|_1, \|A\|_\infty)$$

Hermitian: ift  $A = U \text{diag}(\lambda_1, \dots, \lambda_n) U^*$   
 $(A = A^*)$  & all  $\lambda$ 's are real.

→ For any  $A$ ,  $A^*$  is Hermitian.

Singular values: of  $A$  are the non negative square roots of  $AA^*$  eigenvalues.

→ remove zeroes.

→ For  $A$  normal  $\Rightarrow A = U D U^* \Rightarrow A^* A = U D^* D U^*$   
So singular values are  $\{\sqrt{|\lambda_1|}, \sqrt{|\lambda_2|}, \dots, \sqrt{|\lambda_n|}\}$

Subordinate  $\|A\|_2$

For any matrix  $A$

$\lambda_i$  are eigenvalues of  $D$ .

$$\|A\|_2 = \max_{1 \leq i \leq n} \sqrt{|\lambda_i|} = \text{largest singular value of } A.$$

$$A = UDU^*$$

Both have  $\lambda_1, \lambda_2, \dots, \lambda_n$  eigenvalues.

4E

Convergence for any Norm (since all are equivalent)

$$\lim_{k \rightarrow \infty} A^{(k)} \rightarrow A \iff \lim_{k \rightarrow \infty} \|A^{(k)} - A\| = 0$$

Theorem i), ii) iii) & iv) are equivalent.

i)  $\lim_{k \rightarrow \infty} A^k = 0$  ii)  $\lim_{k \rightarrow \infty} A^k x = 0 \forall x \in C^n$

iii)  $e(A) < 1$  iv)  $\exists$  a subordinate matrix norm such that  $\|A\| < 1$

Geometric series:  $\sum_{k=0}^{\infty} A^k$  converges iff  $e(A) < 1$ .

~~Any invertible matrix has a neighbourhood in which there are invertible matrices.~~

$$\|A - B\| < \frac{1}{\|A^{-1}B\|}$$

$$\Rightarrow e(A^{-1}(A - B)) < 1$$

$\Rightarrow B = A(I - (A^{-1}(A - B)))$  is invertible.

Relative error

$$\frac{\|\tilde{y} - y\|}{\|y\|} \quad (\text{depends on } \|.\|)$$

Condition Number

$$\text{cond}(A) := \|A\| \|A^{-1}\|$$

relative to subordinate matrix norm  $\|.\|$

$$\text{cond}(A) \geq 1$$

$$(A+B)^{-1} = (I + A^{-1}B) A^{-1}$$

Error

$$x = A^{-1}b$$

$$x_\epsilon = A_\epsilon^{-1} b_\epsilon$$

$$A_\epsilon = A + \epsilon B$$

$$b_\epsilon = b + \epsilon c$$

$$\frac{\|x_\epsilon - x\|}{\|x\|} \leq \text{cond}(A) \left\{ \frac{\|b_\epsilon - b\|}{\|b\|} + \frac{\|A_\epsilon - A\|}{\|A\|} \right\}$$

Well conditioned  $\Rightarrow \text{cond}(A) \approx 1$  ILL conditioned  $\Rightarrow \text{cond}(A) \gg 1$

$\text{cond}_2(A) = 1$  for any Unitary A

$$\text{Cond}(A) = \|A\| \|A^{-1}\| \geq e(A) e(A^{-1})$$

If A = normal

$$\Rightarrow \text{cond}_2(A) = \frac{\|\lambda_{\max}\|}{\|\lambda_{\min}\|}$$

If A = any invertible

$$\Rightarrow \text{cond}_2(A) = \frac{\mu_{\max}}{\mu_{\min}}$$

4F

Iterative method converges if for any choice of  $x^{(0)} \in \mathbb{R}^n$ ,  $x^{(k)}$  converges to  $x$ .

$$g^{(k)} = A x^{(k)} - b = A e^{(k)}$$

$$e^{(k)} = x^{(k)} - x$$

$$\lim_{k \rightarrow \infty} x^{(k)} - x = \lim_{k \rightarrow \infty} g^{(k)} = 0$$

Splitting: ~~M~~ Invertible  $\Leftrightarrow A = M - N$

30

$$x^{(k+1)} = M^{-1}N x^{(k)} + M^{-1}b$$

Converges

$$x^{(k+1)} = M^{-1}N x^{(k)} + M^{-1}b$$

converges iff  $e(M^{-1}N) < 1$

$$e^{(k)} = M^{-1}N e^{(k-1)}$$

Estimate

$$\|e^{(k)}\| \sim [e(M^{-1}N)]^k \|e^{(0)}\| \text{ for } k \gg 1$$

Richard's iterative method (steepest descent/gradient)

$$M = \frac{1}{\alpha} I_n \quad N = \frac{1}{\alpha} I_n - A$$

$$M^{-1}N = I_n - \alpha A$$

$$\boxed{\|I - \alpha x_i\| < 1 \forall i}$$

If  $A$  is real & symmetric  $\alpha x_1 \leq x_2 \leq \dots \leq x_n$   
 $\rightarrow x_i x_j < 0$  for some  $i, j \Rightarrow$  Richardson's method  
doesn't converge.

$\rightarrow x_i > 0 \forall i$ , then it converges iff  $0 < \alpha < \frac{2}{\lambda_1}$

$\rightarrow x_i < 0 \forall i$  then  $\| \cdot \|$  iff  $\frac{2}{\lambda_1} < \alpha < 0$

Optimal  $\alpha$ : Let  $A$  be a matrix with  $n$  eigenvalues

$$\Rightarrow \min_{\alpha} e(I - \alpha A) \Rightarrow \tilde{\alpha} = \frac{2}{\lambda_1 + \lambda_n}$$
$$e(I - \tilde{\alpha} A) = \frac{\lambda_n - \lambda_1}{\lambda_n + \lambda_1}$$

$$\epsilon(I_n - \alpha A) = \frac{\text{cond}_2(A) - 1}{\text{cond}_2(A) + 1}$$

Well conditioned  $\Rightarrow$  fast Richardson convergence.

4G

Jacobi method:  $M = D = \text{diag}(A_{11}, \dots, A_{nn})$

$$N = D - A \quad M^{-1}N = I_n - D^{-1}A$$

$|A_{ii}| > \sum_{k=1, k \neq i}^n |A_{ik}| \Leftrightarrow$  strictly row-diagonally dominant

$|A_{ii}| > \sum_{k=1, k \neq i}^n |A_{ki}| \Leftrightarrow$  strictly column-diagonally dominant

If  $A$  is either strictly row or column diagonal dominant then Jacobi converges.

Gauss-Seidel method

$$D = \text{diag}(A_{11}, A_{22}, \dots, A_{nn})$$

$$F_{ij} = \begin{cases} -A_{ij} & i < j \\ 0 & i \geq j \end{cases} \quad E_{ij} = \begin{cases} -A_{ij} & \text{for } i > j \\ 0 & \text{otherwise} \end{cases}$$

$$A = D - E - F$$

$$M = D - E \quad N = F$$

$$M^{-1}N = (D(I_n - D^{-1}E))^{-1}F = (I_n - D^{-1}E)^{-1}D^{-1}F$$

$$L = D^{-1}E \quad U = D^{-1}F$$

$$G = M^{-1}N = (I_n - L)^{-1}U$$

$$\text{For Jacobi method } J = I_n - D^{-1}A = I_n - D^{-1}(D - E - F)$$

$$J = L + U$$

32

If  $A$  is strictly row-diagonally dominant matrix

then Gauss-Seidel method converges.

$$\|G\|_{\infty} \leq \|J\|_{\infty} < 1$$

Theorem If  $J_{ij} \geq 0$ , then one of the following is true

- 1)  $e(G) = e(J) = 0$
- 2)  $0 < e(G) < e(J) < 1$
- 3)  $e(G) = e(J) = 1$
- 4)  $e(G) > e(J) > 1$

i.e if  $J_{ij} \geq 0$  either both converge or diverge.  
when both converge Gauss-Seidel is better.

4H

Relaxed Gauss-Seidel

$$A = \frac{1}{\alpha} D - E + \left(1 - \frac{1}{\alpha}\right) D - F$$

$$M = \frac{1}{\alpha} D - E, \quad N = \left(\frac{1}{\alpha} - 1\right) D + F$$

$$G_{\alpha} = \left(\frac{1}{\alpha} I - L\right)^{-1} \left(\left(\frac{1-\alpha}{\alpha}\right) I + U\right)$$

$$e(G_{\alpha}) \geq |\alpha - 1|$$

convergence  $\rightarrow \alpha \in (0, 2)$

Let  $A$  be Hermitian ( $A = A^*$ ) positive definite matrix. Also assume  $M + M^* - A$  is positive definite.

$\lambda$  is eigenvalue of

$$H = A^{-1}(2M - A)$$

$$\operatorname{Re}(\lambda) > 0$$

Also

$$M^{-1}N = (H - I)(H + I)^{-1}$$

$$\text{and } e(M^{-1}N) < 1$$

Proof

$$\text{Let } M^{-1}N Z = MZ$$

$$\text{Let } Y = (H + I)^{-1}Z$$

$$\Rightarrow H Y = \left(\frac{I+M}{I-M}\right)Y$$

$$\Rightarrow M = \frac{\lambda - 1}{\lambda + 1} \Rightarrow |M|^2 = \frac{|\lambda|^2 + 1 - 2\operatorname{Re}(\lambda)}{|\lambda|^2 + 1 + 2\operatorname{Re}(\lambda)} < 1$$

→ For symmetric positive definite matrices

$$A \in M_n(\mathbb{R}) \quad e(G_\alpha) < 1 \quad \forall \alpha \in (0, 2)$$

Proof:

$$M + M^T - A = \left(\frac{2}{\alpha} - 1\right) D$$

$$\operatorname{Trace}(A) = \sum_{i=1}^n \alpha_i$$

$$\det(A) = \prod_{i=1}^n \alpha_i$$

Similar matrices

$$B = P^{-1}AP \quad (A, B \text{ same eigenvalues})$$

For orthogonal  $|\alpha_i| = 1$

Let  $Ax = \alpha x$  and  $|x_i| \geq |x_j| \quad \forall j$

$$\Rightarrow \sum_{j \neq i} A_{ij} x_j = (\alpha - \alpha_{ii}) x_i$$

$$\Rightarrow |\alpha - \alpha_{ii}| \leq \sum_{j \neq i} |A_{ij}|$$

$$D_i = \{z : |z - a_{ii}| \leq r_i\}$$

$$R_i = \sum_{j \neq i} |a_{ij}|$$

Any

$$z \in \bigcup_{i=1}^n D_i$$

SA

No formula for polynomial with degree  $\geq 5$

Fixed point if  $\psi(y) = y$

Write  $f(x) = 0$  as  $x - g(x) = 0$

Fixed point theorem: continuous  $g: [a, b] \rightarrow [a, b]$

$\exists \xi \in [a, b]$  s.t.  $\xi = g(\xi)$

Proof:  $g(a) - a \geq 0$  &  $g(b) - b \leq 0$

Iteration

$$x^{(k+1)} = g(x^{(k)})$$

Contraction continuous  $g: [a, b] \rightarrow \mathbb{R}$  is a contraction

on  $[a, b]$  if  $L \in (0, 1)$   
 $|g(x) - g(y)| \leq L |x - y|$

$\rightarrow$  If  $g$  is a contraction on  $[a, b]$  it has a

unique fixed point on  $[a, b]$  & any  $x^{(0)} \in [a, b]$

converge. (order =  $< 1$ )

$$|x^{(k)} - \xi| \leq L^k |x^{(0)} - \xi|$$

→ If  $g'$  is continuous on  $[a-h, a+h]$   
 and  $|g'(x)| < 1$  then converges for  
 $x^{(0)} \in (a-\delta, a+\delta)$  s.t.  $|g(x)| \leq L < 1 \forall x \in [a-\delta, a+\delta]$

Relaxation

~~zero iteration~~

$$x^{(k+1)} = x^{(k)} - \lambda f(x^{(k)})$$

$$g(x) = x - \lambda f(x)$$

$$f(x)=0 \Leftrightarrow g(x)=x$$

$\exists \lambda > 0 \text{ s.t. it converges if } f'(x) \neq 0$   
 $x^{(0)} \in [a-\delta, a+\delta] \text{ where } |f'(x)| < 1$   
 $\forall x \in [a-\delta, a+\delta]$

SB Error:  $e^{(k)} = x^{(k)} - a$

Order of convergence If  $\exists c > 0 \text{ s.t. } p \geq 1$   
 (with  $c < 1$  if  $p=1$ ) and ~~an~~ integer  
 $N$  s.t. for all  $k \geq N$

$$|e^{(k+1)}| \leq c |e^{(k)}|^p \quad (c < 1 \text{ if } p=1)$$

Generalised relaxation iteration

$$x^{(k+1)} = x^{(k)} - \lambda f(x^{(k)})$$

$$g(x) = x - \lambda(x) f(x)$$

Convergence if  $|1 - \lambda(a) f'(a)| < 1$

Newton's method:

$$x^{(k+1)} = x^{(k)} - \frac{f(x^{(k)})}{f'(x^{(k)})}$$

36

$$g(x) = x - \frac{f(x)}{f'(x)}$$

$$f'(x^{(k)}) \neq 0 \forall k$$

Convergence

If  $f: [a, b] \rightarrow \mathbb{R}$  is  $C^2$  on  $I_S = [\xi - \delta, \xi + \delta]$

for  $\delta > 0$ , let  $f(\xi) = 0$ ,  $f'(\xi) \neq 0$  &  $f''(\xi) \neq 0$ .

If

$$\frac{|f''(x)|}{|f'(y)|} \leq M \quad \forall x, y \in I_S$$

and

if  $|x^{(0)} - \xi| \leq h$  with  $h = \min \{\delta, \frac{1}{M}\}$

then it converges with order 2.

$$\lim_{k \rightarrow \infty} \left| \frac{\xi - x^{(k+1)}}{(\xi - x^{(k)})^2} \right| \leq \frac{M}{2} \quad (p=2)$$

Zeroes with multiplicities

If  $m > 1$  is the multiplicity,

$$x^{(k+1)} = x^{(k)} - \frac{f(x^{(k)})}{f'(x^{(k)})}$$

still converges  
but only  
linearly ( $p=1$ )

$$x^{(k+1)} = x^{(k)} - m \frac{f(x^{(k)})}{f'(x^{(k)})}$$

$$x^{(k+1)} = x^{(k)} - \frac{f(x^{(k)})}{f'(x^{(k)})}$$

$$f(x) = f(x)$$

converge with  $p=2$

$$g(x) = x - \frac{f(x)}{f'(x)} \quad |(x_i) = 1 - \frac{1}{m} < 1$$

Linear convergence

$$f(x) = (x - x_i)^m \psi(x)$$

Bisection method

Slow

Secant method

$$\& x_{(k+1)} = x_k - f(x_k) \left( \frac{x_k - x_{k-1}}{f(x_k) - f(x_{k-1})} \right)$$

$$f'(x_i) \neq 0 \Rightarrow \frac{|x_i - x_{(k+1)}|}{|x_i - x_k|} \leq \frac{2}{3}$$

P=1 for convergence

End