# Lower Bounds for the Polynomial Calculus via the "Pigeon Dance"

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### Overview

- Overview of the needed background
- Ideas behind the proof and its structure
- The proof in more detail

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#### Definition

- Similar to sequent calculus, but lines are polynomials
- We use multilinear polynomials  $S_n(\mathbb{K})$  $(xy + xz + v \equiv x^2y + x^3z^5 + v)$
- Addition

$$\frac{f}{af + bg}$$

Multiplication

$$\frac{f}{f \cdot x}$$

#### Motivation

- q is provable from  $f_1, \ldots, f_n$  if and only if it is in the ideal generated by them
- A proof of g=1 exists if and only if  $f_1,\ldots,f_n$  have no common zeroes
- We construct polynomials such that their zeroes correspond to satisfying assignments
- Proving 1 from them is a refutation

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### Example proof

• Try to prove xy + z from x + 1 and z

$$\frac{(y)}{xy+y} \frac{x+1}{xy+y+z} (+)$$

- There is no way to prove y from x+1 and z
- We can not prove xy + z

# Algebraic view of proofs

$$ullet$$
  $V$  are polynomials we can prove

$$ullet$$
  $\Delta$  are leading terms of ones we can not prove

• 
$$S_n(\mathbb{K}) \cong \mathbb{K}\Delta \oplus V$$

$$ullet$$
  $R$  is the projection onto  $\Delta$ 

Similarly:  $V_d, \Delta_d, R_d$  for polynomials up to degree d  $V_I, \Delta_I, R_I$  for polynomials using variables from I

$$xy + y + z$$

$$xy + y + z = -y + xy + y + z$$

$$R(xy + y + z) = \mathbf{y}$$

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# The pigeonhole principle

- ullet If there are m pigeons, n pigeon holes, and m>n then at least two pigeons have to share a hole
- Formally: if m > n there is no injection  $[m] \hookrightarrow [n]$
- Variables:  $x_{i,j}, i \in [m], n \in [n]$
- Assignment of  $x_{3,5}$  corresponds to pigeon 3 being in hole 5

#### Definition $(\neg \mathcal{PHP}_n^m)$

$$Q_i \coloneqq 1 - \sum_{j \in [n]} x_{ij}$$
 for each  $i \in [m]$ 

$$Q_{i_1,i_2,j} \coloneqq x_{i_1j}x_{i_2j}$$
 for each  $i_1 \neq i_2 \in [m], j \in [n]$ 

#### Main result

#### Theorem

For any m > n, every polynomial calculus refutation of  $\neg PHP_n^m$  must have degree at least n/2 + 1.

- Characterize  $R_d$  semantically
- ullet Problem: only works if  $R_I$  agree on their intersections
- ullet Characterize  $R_I$  syntactically
- Show the different operators are identical

#### Idea

- What polynomials are derivable from  $\neg \mathcal{PHP}_n^m$ ?
- Pigeons can not share holes
- Pigeon assignments are variable assignments

$$\delta(x) = \begin{cases} 1, & \text{if } x \in \{x_{1,2}, x_{2,4}, x_{3,1}\} \\ 0, & \text{otherwise} \end{cases}$$

Polynomials are evaluated to 0 if they allow the assignment

$$\delta(1 - x_{1,1} - x_{1,2} - x_{1,3} - x_{1,4}) = 0$$
$$\delta(x_{1,1}x_{3,1}) = \delta(x_{1,2}x_{2,2}) = 0$$

# Characterizing $V_I$

- $I \subseteq [n]$  is a set of pigeons
- ullet  $M_I$  is all assignments corresponding to injections  $I\hookrightarrow [m]$
- ullet Polynomials in  $\neg \mathcal{PHP}_n^m$  are identically zero on  $M_I$
- ullet Polynomials derivable from  $eg \mathcal{PHP}_n^m$  are identically zero on  $M_I$ 
  - $\Rightarrow V_I$  are all polynomials identically zero on  $M_I$

Valid Pigeon Arrangements

# Combining $V_I$

- ullet We want a characterization of  $V_d$ , not  $V_I$
- ullet Set  $V_d\coloneqq igcup_{|I|\leq d} V_I$  and  $R_d$  accordingly
- Does this definition actually work?
- ullet Yes it does! But only if  $R_I(t) = R_{\mathrm{dom}(t)}(t)$

- ullet Define  $R_I(t)$  so that it is independent of pigeons not in t
- Show that the two definitions are identical
- We will only show the broad steps

Characterizing  $R_{m{I}}$ 

# Example

#### **Formalization**

- The first pigeon flies to an unoccupied hole to its right
- Repeat until all pigeons have moved once
- If a pigeon can not find an empty hole, the dance is aborted
- ullet Define  $\Delta_I$  to be the set of terms that let pigeons complete the dance
- $R_I(t) = R_{\text{dom}(t)}(t)$  since pigeons not in the dance do not affect it

# Defining $R_I$

- We need  $R_I(t) = f$  with  $\mathrm{LT}(f) \prec t$  and  $t f \in V_I$
- If  $t \in \Delta_I$  then f = t
- Otherwise we use  $Q_{i_1} = 0$  to derive

$$t = x_{i_2 j_2} \cdots x_{i_d j_d} - \sum_{j' < j_1} x_{i_1 j'} x_{i_2 j_2} \cdots x_{i_d j_d}$$
$$- \sum_{j' > j_1} x_{i_1 j'} x_{i_2 j_2} \cdots x_{i_d j_d} \mod V_I.$$

- ullet The first two summands are  $\prec t$  and can be ignored
- Any terms with  $j' \in \{j_2, \dots, j_d\}$  are 0 and can be ignored

# Defining $R_I$ (cont.)

- Repeat the same process with all remaining terms
- ullet At each step the first  $x_{i,j}$  gets replaced with  $x_{i,j'}$  with some unused j'>j
- This is the pigeon dance!
- ullet Since  $t 
  ot\in \Delta_I$  the dance can not be completed
- Process terminates with  $f \prec t$  and  $t = f \mod V_I$

Properties of the dance

### The Kill operator

- The Kill operator kills the first pigeon and moves its hole to the left
- $\text{Kill}(x_{i_1,j_1}\cdots x_{i_d,j_d}) = x_{i_2,j'_2}\cdots x_{i_d,j'_d}$  with

$$j'_k \coloneqq \begin{cases} j_k + 1, & \text{if } j_k < j_1 \\ j_k, & \text{if } j_k > j_1. \end{cases}$$

Properties of the dance

### Properties of Kill

#### Theorem

 $x_{i_1,j_1}\cdots x_{i_d,j_d}\in \Delta_I$  if and only if there is a  $j'>j_1$  such that  $\mathrm{Kill}(x_{i_1,j'}\cdots x_{i_d,j_d})\in \Delta_I$ .

#### Proof

This operator effectively moves the first pigeon to an empty hole and then kills it. This is the same as each step in the dance, where the first pigeon flies to some free hole to its right and then occupies it.  $\hfill\Box$ 

# Properties of Kill (cont.)

#### Theorem

 $\Delta_I$  is closed under Kill.

#### Proof

If  $t \in \Delta_I$  then the pigeons can complete their dance. During this the first pigeon will start at j and fly to j'. Killing the pigeon frees up j' so any other pigeon that wanted to use j can use it instead.

### The lower bound

#### Theorem

If  $|I| \le (n+1)/2, t \in D_I$  and the minimal element i of I is not in dom(t), then there exists a  $j \in [n]$  such that  $Kill(x_{ij}t) \in \Delta_I$ .

#### Proof

At most

$$|\operatorname{dom}(t)| \le |I \setminus \{i\}| \le \frac{n-1}{2}$$

pigeons involved in the dance, each occupying two holes. Thus the total number of holes is n-1 and one hole j remains free. For the purposes of the dance,  $\mathrm{Kill}(x_{ij} \cdot t)$  is the same as t since the only difference is j being moved to the left.

### Putting things together

- We now show that the two operators are identical
- ullet Equivalent to  $\Delta_I$  being linearly independent as functions  $M_I o \mathbb{K}$
- Induction on |I| gives us assignment  $a \in M_{I \setminus \{i\}}$  with  $a(f') \neq 0$
- ullet Pick a j such that  $\operatorname{Kill}(x_{i,j}t)\in\Delta_I$
- $\bullet$  Extend assignment to I with  $a'(f) \neq 0$

### Summary

- ullet If  $d \leq n/2 + 1$ , then definition of  $R_I$  via the pigeon dance and via  $M_I$  are identical
- ullet Pigeons not in the dance do not affect its success, so  $R_I(t)=R_{\mathrm{dom}(t)}(t)$
- ullet  $V_d$  is precisely polynomials identically zero on  $M_I$
- $R_d \neq 0$
- $\bullet$  There is no refutation of  $\neg \mathcal{PHP}_n^m$  with  $d \leq n/2 + 1$