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Private Adresse: XXX  
Telefonnummer: XXX  
E-Mail-Adresse: XXX

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## Keywords

Duplicate Elimination  
Genetic Algorithm  
Greedy Heuristic  
Integer Program  
Meta-Heuristic  
Multi-Dimensional Knapsack Problem  
Multi-Unit Combinatorial Auction  
Phenotypic Distance  
Winner Determination Problem

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## Abbreviations

CA	combinatorial auction
VCG	Vickrey-Clarke-Grooves

# 1. Introduction

Wissenschaftliche Arbeiten weisen standardisierte Strukturen auf, von denen man im Ausnahmefall abweichen kann. Eine Standardstruktur ist die folgende:

- Abstract
- Introduction [Einleitung]
- Related Work [Grundlagen und Verwandte Arbeiten]
- Method/Experimental Design/Implementation
- Results
- Discussion
- Limitations and Future Research [Limitationen und Ausblick]
- Conclusion [Fazit/Schlussfolgerung]
- References [Literaturverzeichnis]
- Appendix [Anhang]

In Bachelor-und Masterthesen können Sie für jeden dieser Abschnitte den Befehl *chapter* wählen. In Seminararbeiten verwenden Sie stattdessen bitte *section*.

Die typische Struktur der Einleitung ist:

1. Problemstellung und Motivation
2. Stand der Forschung, darauf aufbauend Forschungslücke und Forschungsfrage(n) herausarbeiten
3. Ziel der Arbeit und eigener methodischer Ansatz zur Beantwortung der Forschungsfrage(n): Absatz beginnt meist mit: The goal of this thesis/work/manuscript is
4. Ergebnisse der Arbeit [optional]
5. Erwarteter wissenschaftlicher (und praktischer) Beitrag [=Contribution]
6. Manchmal folgt noch eine Gliederung [ausformuliert]

## 2. Figures and Tables

Hier finden Sie nun Beispiele für das Einfügen von Grafiken und Tabellen. Für Tabellen können Sie auch Umgebungen wie `tabularx` und `longtable` verwenden.

We performed experiments for three different product categories ranging from commodity products (energy-saving lamps) over hotel rooms to capital goods (washing machines). A lower average price of the products represents a lower perceived risk. We used energy-saving lamps as rather low priced products (avg. price: 7.57€), hotel rooms as medium priced products (avg. price: 249.50€) and washing machines as rather high priced products (avg. price: 524.33€). For each category, we collected data for 40 products. Each product is described by five attributes. Specifically, we extracted frequently used attributes from Amazon product descriptions (energy-saving lamps, washing machines) or descriptions in the hotel booking platform HRS. Table 2.1 summarizes products and product attributes<sup>1</sup>.

Let us give an example. We assume a choice scenario with five different cameras (see Table 2.2). Product attributes are photo resolution  $ph$ , zoom factor  $zf$ , and price  $pr$ . All consumers have the same preference for the attribute level order: they prefer lower to higher prices, and higher photo resolutions and zoom factors to lower ones. Table 2.2 lists the five exemplary cameras with their corresponding attribute levels as well as single-attribute values  $v_i$  for each attribute  $a_i \in \{ph, zf, pr\}$ . Camera  $E$  has the best price, but the worst photo resolution. In the product domination graph,  $E$  has hence no outgoing edges with respect to price, but four outgoing edges with respect to photo resolution. Figure 2.1 shows the resulting product domination graph.

Eine Tabelle finden Sie in Tabelle 2.2.

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<sup>1</sup>For washing machines and energy-saving lamps, consumer ratings are from Amazon; for hotel rooms, consumer ratings are from hrs.com. The attribute level order for washing machine brands is based on the brands' average sales rank on Amazon.



Table 2.1.: Products and their Attributes

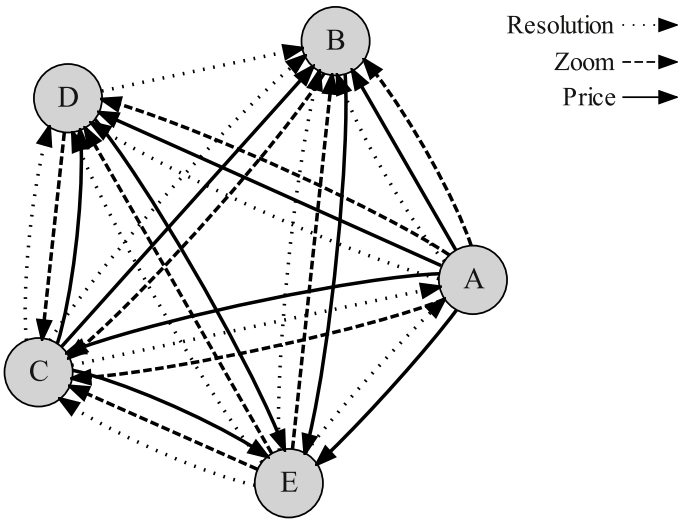
Product	Attribute	Unit	Attribute Level Order
Energy-saving Lamp (n=40)	Price	Euro	Increasing
	Energy Efficiency Grade	–	A+ $\succ$ A $\succ$ B
	Deviation from Day Light	–	None $\succ$ Low $\succ$ Large
	Durability	Hours working time	Decreasing
	Customer Rating	1-5 Stars	Decreasing
Hotel Room (n=40)	Price per Night	Euro	Increasing
	Category	Stars	Decreasing
	Distance from City Center	Kilometers	Increasing
	WLAN availability	–	Available $\succ$ Not available
	Customer Rating	1-5 Stars	Decreasing
Washing Machine (n=40)	Price	Euro	Increasing
	Brand	–	Siemens $\succ$ Bosch $\succ$ AEG $\succ$ Bauknecht $\succ$ Gorenje $\succ$ Blomberg $\succ$ LG
	Energy Consumption	kWh per year	Increasing
	Water Consumption	Liters per year	Increasing
	Customer Rating	1-5 Stars	Decreasing

$\succ$ : is preferred over

Table 2.2.: Example attribute levels and corresponding single-attribute values  $v_i$ 

Camera	Photo Resolution	Zoom Factor	Price
A	$v_{ph}(12MP) = 0.60$	$v_{zf}(10x) = 0.00$	$v_{pr}(610EUR) = 0.00$
B	$v_{ph}(14MP) = 1.00$	$v_{zf}(15x) = 0.63$	$v_{pr}(470EUR) = 0.40$
C	$v_{ph}(10MP) = 0.20$	$v_{zf}(18x) = 1.00$	$v_{pr}(540EUR) = 0.20$
D	$v_{ph}(13MP) = 0.80$	$v_{zf}(15x) = 0.63$	$v_{pr}(470EUR) = 0.40$
E	$v_{ph}(9MP) = 0.00$	$v_{zf}(10x) = 0.00$	$v_{pr}(260EUR) = 1.00$

Figure 2.1.: Example Product Domination Graph



## 3. Combinatorial Auctions

Die ist der Ausschnitt eines Beispielkapitels.

*Combinatorial auctions* (CAs) are a part of electronic market design. Research in electronic market design joins two disciplines: economics and computer science. Economical research focuses on game theoretical aspects by analyzing strategic behavior of self-interested agents. From the viewpoint of computer science, computational problems are addressed, such as finding the optimal allocation in auctions. As this work concentrates on computational aspects, we assume that the reader has a stronger background in computer science than in economics. Thus, in this chapter we will point out the main ideas of the economical perspective to provide some basic knowledge in this area.

### 3.1. Mechanism Design

#### 3.1.1. Definition

Mechanism design was introduced by Hurwicz (1960). It aims at implementing system-wide solutions to problems in non-cooperative environments with multiple self-interested agents. Such problems can be political elections, public projects in which the participants themselves have to invest money, or allocation problems. Given that agents hold only private information about their preferences, a structure has to be chosen in which in equilibrium each agent behaves according to the designer's or principal's intentions. The designer can either act on behalf of the society, for example when collecting taxes for a public project, or she can pursue self-interests when, for instance, being an auctioneer.

Since the agents' information is private, the principal faces the problem that the agents might lie about their real valuations in order to influence the outcome according to their preferences. In most cases, whenever such manipulations occur, they damage the resulting system-wide welfare (Nisan & Ronen, 2000). Thus, simply asking the participants to reveal their preferences is unfavorable. Therefore, the

principal has to define other rules which lead to the desired outcome. The most common solution to this problem is to introduce monetary transfers providing incentives for the agents to behave truthfully.

In mechanism design two economic areas are joined: *game theory* and *social choice theory*. In game theory the agents' strategies are analyzed, and in social choice theory an outcome is selected according to a set of agents' preferences. The outcome in social choice theory is determined by a social choice function, which is to be implemented by a mechanism. Formally we have a set of possible outcomes  $O$  and agents  $i \in I$ ,  $|I| = n$ . Each agent  $i$  has a type  $\theta_i \in \Theta_i$  reflecting the possible preference sequences the agent can have. The type captures all of the agent's private information relevant to her decision. The agent's utility  $u_i(o, \theta_i)$  over each outcome depends on her type; while  $u_i(o_1, \theta_i) > u_i(o_2, \theta_i)$  means that the outcome  $o_1$  is preferred over the outcome  $o_2$ . The social choice function maps from the space of all types  $\Theta$  to the space of all outcomes  $O$ ,

$$f : \Theta_1 \times \Theta_2 \times \dots \times \Theta_n \rightarrow O. \quad (3.1)$$

Examples for such social choice functions are allocation problems or political voting protocols in which a candidate or a party is chosen. The most common objective of a social choice function is the maximization of the social welfare, the so called *allocative-efficiency* (Parkes, 2001). It maximizes the sum of all utilities over all agents:

$$f(\theta) = \arg \max_{o \in O} \sum_{i \in I} u_i(o, \theta_i). \quad (3.2)$$

Another objective is *individual rationality*; the agent's payoff is never less when participating in the mechanism than her payoff without participating. Additionally there is *Pareto optimality*. An outcome is Pareto optimal whenever none of the agents could perform better without causing another agent to perform worse than in the current situation.

So far, we have learned what a social choice function is, and what typical objectives for the choices of outcomes are. Now, a mechanism has to be found which implements a given social choice function with one or several of these objectives. For this purpose, the agents' possible strategies have to be specified together with

an outcome function based on these strategies. The mechanism should guarantee an implementation despite the self-interest of the agents (Parkes, 2001). Mathematically, a mechanism  $M$  is defined on the strategy spaces  $S_i$  of the agents:

$$M = ((S_1, \dots, S_n), g(\cdot)) g : S_1 \times \dots \times S_n \rightarrow O, \quad (3.3)$$

where  $g$  is an outcome function and  $S_i$  denotes all strategies or actions an agent  $i$  is allowed to take. A mechanism implements a social choice function if there is an equilibrium strategy profile  $s^*(\cdot) = (s_1^*(\cdot), \dots, s_n^*(\cdot))$  of the game induced by  $M$  so that

$$g(s_1^*(\theta_1), \dots, s_n^*(\theta_n)) = f(\theta_1, \dots, \theta_n), \quad \forall (\theta_1, \dots, \theta_n) \in (\Theta_1, \dots, \Theta_n), \quad (3.4)$$

where  $s_i^*(\theta_i)$  is the strategy agent  $i$  with type  $\theta_i$  plays in the equilibrium. Please note that the equilibrium concept is not specified in this definition. It could, for example, be a Nash equilibrium. In this case, given the other players  $j$ ,  $j \neq i$ , conform to the equilibrium strategies  $s_j^*(\theta_j)$ , no other player  $i$  has an incentive to unilaterally deviate from her equilibrium strategy. Other examples are the dominant strategy or the Bayes-Nash strategy equilibrium. The dominant strategy equilibrium facilitates it for the agents since the optimal strategy for an agent is independent of any strategies the other agents could play. Thus, the agents do not need to speculate about the way the others might behave. Informally, we could say that the concept of dominant strategies "removes game theory from the problem" Parkes (2001, p. 5). The Bayes-Nash equilibrium is similar to Nash equilibriums, but assumes that agents have incomplete information about the opponents' types. Therefore, agents use probability functions to speculate about the other agents' preferences (Osborne & Rubinstein, 1994).

### 3.1.2. Revelation Principle and Gibbard-Satterthwaite Theorem

In equation 3.3, we see that a mechanism defines the available strategies and the function for selecting an outcome. It is necessary that these strategies are kept simple so that they can be applied by the agents. The easiest strategies occur when choosing a *direct mechanism* asking the agents to report their types directly to the principal,  $S_i = \Theta_i$ . Direct mechanisms lead to a centralization of the problem as agents report their types to a center that determines the outcome and reports it back to the agents. On the contrary, when applying *indirect mechanisms* agents

have to think about how to transform their type into a strategy and the latter is reported to the mechanism. In other words,

”the computations that go on within the mind of any bidder in the non-direct mechanism are shifted to become part of the mechanism in the direct mechanism”. McAfee and McMillan (1987, p. 712)

When applying these direct mechanisms agents may still lie about their true types. Mechanisms which, in contrast, succeed in establishing an equilibrium in which all agents tell the truth, are called *incentive-compatible*. In this case, it is in the interest of all agents to report their true types,  $s_i^*(\theta_i) = \theta_i$ ,  $\forall \theta_i \in \Theta_i$ . Further, if telling the truth is a dominant strategy, the mechanism is called *strategy-proof*. As will be shown later on, this can be achieved by the *Vickrey-Clarke-Grooves* (VCG) mechanism.

We learned that the equilibrium strategy profile  $s^*(\cdot)$  does not determine the concept of equilibrium. Some equilibrium concept must be chosen and implemented together with the mechanism. In the worst case, in order to find out if a certain social choice function can be implemented by a certain mechanism with, for instance, dominant strategies, one would have to consider all possible mechanisms. However, research on mechanism design led to the *revelation principle* as a solution to this. It states that for any mechanism, there is a direct, incentive-compatible mechanism with the same outcome (McAfee & McMillan, 1987). An intuitive explanation for this principle consists in: the transformation from types into strategies, which occurs in the agents’ minds in indirect mechanisms, and which is used as a filter in the direct mechanism. That is, the direct mechanism first filters all reports of the agents and simulates the indirect mechanism with the filtered input. This principle is valid for the optimal mechanism as well. Thus, the search for a mechanism can focus on direct mechanisms. Therefore, if no direct mechanism can implement a given social choice function, then no indirect mechanism will do so.

In contrast to the positive result of the revelation principle, there also exists a negative result, the *Gibbard-Satterthwaite* theorem. According to it, it is impossible to find a mechanism with certain positive characteristics. To understand the theorem, first note that a social choice function is *truthfully implementable* if and only if the dominant strategy is to reveal the truth. Furthermore, a social choice function  $f$  is *onto* if for each  $o \in O$  at least one element in  $\Theta$  exists so that  $f$  maps to  $o$ . Finally, a social choice function  $f$  is *dictatorial* whenever there is a dictator  $j$  among

the agents so that for all outcomes,  $o_j$  is strictly preferred to another outcome  $o_k$  whenever the dictator  $j$  strictly prefers  $o_j$  to  $o_k$ . Obviously, this is an unwanted characteristic. It turns the Gibbard-Satterthwaite theorem impractical for real-life mechanisms since they allow manipulation.

**Gibbard-Satterthwaite Theorem:** *Given  $O$  is finite,  $|O| \geq 3$ , and the social choice function  $f$  is onto, then  $f$  is truthfully implementable in dominant strategies if and only if  $f$  is dictatorial.*

According to the theorem it is impossible to elicit the truth if dominant strategies exist. However, despite this result, the theorem can be circumvented by placing restrictions on the agents' preferences, the way it is done in the VCG mechanism.

### 3.1.3. Vickrey-Clarke-Grooves Mechanism

The VCG mechanism combines the following important virtues by introducing a special payment scheme. First, it implements social choice functions in dominant strategies. Thus, agents do not have to speculate which strategies the other agents might play, and they do not need to waste resources on learning about their competitors' strategies. Second, the mechanism does not have to make any assumptions about the information agents have on each other. And, third, the VCG mechanism is allocative-efficient (see equation 3.2), strategy-proof and non-dictatorial.

AND SO ON...

## **4. Conclusions**

Auf Deutsch: Fazit

### **4.1. Summary**

Auf Deutsch: Zusammenfassung

Die Zusammenfassung der Arbeit ist optional. Sollten Sie die Arbeit noch einmal zusammenfassen wollen, so halten Sie dies bitte eher kurz und wiederholen Sie sich nicht zu sehr.

### **4.2. Limitations and Future Research**

Auf Deutsch: Limitationen und Ausblick

Es ist sinnvoll, jede Limitation an eine Idee zu knüpfen, wie diese in zukünftigen Arbeit zu adressieren wäre.

### **4.3. Contribution**

Auf Deutsch: Beiträge

Was sind die Beiträge Ihrer Arbeit sowohl für die Wissenschaft (und Theorie) als auch für die Praxis? Hier sollten Sie versuchen über den Tellerrand hinauszuschauen und einen eher weiten Blick einnehmen.



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## A. Appendix

Table A.1.: Size of the search space for BASIC strategies ( $n = 4$ ).

m	s	search space size
4	1	64
	2	4096
	3	262,144
	4	16,777,216
	5	1,073,741,824
7	1	262144
	2	16777216
	3	68,719,476,736
	4	2.81E+14
	5	1.15E+18

## Eidesstattliche Erklärung

Ich versichere, dass ich meine Thesis ohne Hilfe Dritter und ohne Benutzung anderer als der angegebenen Quellen und Hilfsmittel angefertigt und die den benutzten Quellen wörtlich oder inhaltlich entnommenen Stellen als solche kenntlich gemacht habe. Diese Arbeit hat in gleicher oder ähnlicher Form noch keiner Prüfungsbehörde vorgelegen.

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