MATHEMATICS 1R

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Contents

1	Properties of Numbers	2
	1.1 Introduction	2
2	Complex Numbers	7

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Properties of Numbers

1.1 Introduction

1.1 definition. A set is a finite or infinite collection of objects in which order has no significance, and multiplicity is generally also ignored

1.2 notation. $A = \{a_1, a_2, a_3\}$

1.3 notation. $a \in A$, The element a belongs to the set A

1.4 notation. $A \subseteq B$, A is a subset of B, all elements of A appear in B but the opposite is not necessarily true

1.5 notation. $A \subset B$, the set B has at least one element which does not belong to A

1.6 notation. $a \notin A, A \nsubseteq B, A \notin B$, are negations of the above

1.7 definition. Natural Numbers N, the set of all positive whole numbers

1.8 definition. Whole Numbers \mathbb{N}_0 , $\mathbb{N} + 0$

1.9 definition. Integers \mathbb{Z} , \mathbb{N}_{0} + negative whole numbers

1.10 definition. Rational Numbers \mathbb{Q} , the set of all fractions with denom \neq 0

1.11 definition. Irrational Numbers $\mathbb{R} \setminus \mathbb{Q}$, the set of numbers which can not be expressed as a fraction

1.12 definition. Real Numbers \mathbb{R} , the set of all numbers in the real line

1.13 remark. $\mathbb{N} \subset \mathbb{N}_0 \subset \mathbb{Z} \subset \mathbb{Q} \subset \mathbb{R}$ and $\mathbb{R} \setminus \mathbb{Q} \subset \mathbb{R}$

1.14 example. Prove that $\sqrt{n} \in \mathbb{R} \setminus \mathbb{Q}$, n > 1 for any squarefree integer n, e.g. $\sqrt{20}$

Method

- 1. Proof that any $j \cdot g \in \mathbb{R} \mathbb{Q}$, with $j, g \in \mathbb{Q}$, $\mathbb{R} \mathbb{Q}$
- 2. Decompose $\sqrt{(n)}$ into $k\sqrt{(u)}$, where u is prime
- 3. Prove that $\sqrt{(u)}$ is irrational by Euclid's Lemma Method

Proof.

Step 1 Let $\frac{a}{b}$, $\frac{c}{d} \in \mathbb{Q}$ and $x \in \mathbb{R} - \mathbb{Q}$:

$$(1) x \cdot \frac{a}{b} = \frac{c}{b}$$

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(2) $\iff x = \frac{cb}{da}$

We reach a contradiction at (2) since integers are closed under addition, i.e. x must belong to \mathbb{Q}

performance of that operation on members of the set always produces a member of the same set

A set has closure under an operation if

QED

(Step 2)
$$\sqrt{20} = 2\sqrt{5}$$

Proof.

Step 3 It follows from ?? that $2\sqrt{5}$ is rational *iff* 2 and $\sqrt{5}$ are both rational.

Case 1: 2 is rational is trivially true

Case 2: Let's assume that $\sqrt{5} \in \mathbb{Q}$. This assumption implies that

$$\exists a,b \in \mathbb{Z}, b \neq 0 | \sqrt{5} = a/b$$

We can assume also that a/b is in its most reduced form, i.e. it has no common factors. Then,

(3)
$$\sqrt{5}^2 = \frac{a^2}{b^2}$$

$$(4) 5b^2 = a^2$$

1.15 lemma. Euclid's: For a prime number p if $p \mid mn$ then $p \mid m$ or $p \mid n$

It follows from 1.1 that 5 divides a^2 . And from 1.15 (by setting p =5, m, n = a) that 5 divides a.

We then have that $a = 5k, k \in \mathbb{Z}$, and can therefore rewrite 1.1:

(5)
$$5b^2 = (5k)^2$$

$$(6) \qquad b^2 = 5k^2$$

Similarly as above it follows then, that 5 divides b. Hence we reach the conclusion that 5 divides both a and b. Which contradicts our original assumption that a and b are coprime i.e. that $\frac{a}{b}$ was in its most reduced form.

 $\therefore \sqrt{5}$ is not irrational $\implies \sqrt{20}$ is not irrational

OED

1.16 lemma. *Division Algorithm Given two natural numbers a and b, there integers q and r, such that:*

$$a = bq + r$$
 $q \ge 0$ and $0 \le r < b$

Proof. Let $\frac{a}{b}$ be the positive natural number between q and q + 1.

$$q \le \frac{a}{b} < q + 1$$

$$bq \le a < bq + b$$

$$0 \le a - bq < b$$

Setting
$$r = a - bq$$
, we obtain $a = bq + r, q \ge 0$ and $0 \le r < b$

1.17 lemma. Euclidean Algorithm is obtained by repeated application of 1.1.

$$gcd(a,b) = am + bn$$

Proof.

1.
$$gcd(a,b) = gcd(b,r) = ... = gcd(z,0)$$

Why is this true?

If d|ab, then there exists a k, l s.t. a = dk and b = dl 1.1 also tells us that a = bq + r, replacing dk and dl, we obtain the following:

$$dk = dlq + r$$
$$r = d(k - lq)$$

Hence, there exists a number (represented above by (k-lq)) that multiplied by d gives us r. Therefore, by definition, d must also divide r

2. Assuming that a > b > r, the sequence will decrease until eventually the remainder can not be further divided, i.e. r = 0. Hence,

$$gcd(a,b) = gcd(b,r) = gcd(r,0)$$

By looking at the last term of the equality we find that the gcd must by definition be r (since it divides exactly). Therefore, by back-substituting the result $^{\rm I}$ we find that the gcd of the original expression is equal to the remainder of the penultimate

QED

¹recursively, haskell style

1.18 example. Find integers m and n such that 55m + 7n = 1. Note that by the definition of the division algorithm, if such numbers exist, then we expect gcd(55,7) = 1. So, we start by applying the division algorithm to verify that it exists:

$$55 = 7 \cdot 7 + 6$$
$$7 = 6 \cdot 1 + \mathbf{1}$$
$$6 = 1\dot{6} + 0$$

We can now rearrange the equations above in the form r = a - bq, and backtrack until we reach the original equation in terms of 55 and 7.

$$6 = 55 - 7 \cdot 7$$

$$1 = 7 - 6 \cdot 1$$

$$= 7 - (55 - 7 \cdot 7)(1)$$

$$= 7(8) - 55$$

Hence, m = -1 and n = 8

1.19 definition. Polynomial of degree $n \ge 0$ is an expression of the form:

$$a_n x^n + a_{n-1} x^{n-1} + \cdots + a_1 x + a_0$$

Where a_n is the leading coefficient

1.20 remark. If $a_n = 1$, then the polynomial is said to be monic

1.21 remark. o is considered to be a polynomial

1.22 remark. The Euclidean Algorithm can also be applied to polynomials, since:

$$f(x) = g(x)q(x) + r(x)$$

Where $0 \le r(x) < g(x)$

1.23 theorem. Remainder: Let f(x) be a polynomial in x, and c be a constant:

$$f(x) = (x - c)q(x) + f(c)$$

1.24 theorem. Factor: If f(c) = 0, then (x - c) is a factor of f(x)

1.25 example. Find the highest monic quadratic polynomial which divides both

$$g(x) = x^3 + 6x^2 + 11x + 6$$
 and $f(x) = x^4 + 5x^3 + 10x^2 + 20x + 24$

Following a similar procedure as before, we can apply the Euclidean Algorithm until we find a polynomial which divides g(x) and f(x) exactly (i.e. r(x) = 0). The key point is to keep in mind the fact that **any polynomial wich divides both** f(x) **and** g(x) **must also divide their remainder**.

1. Dividing f(x) by g(x), so as to find q(x) and r(x):

$$\begin{array}{r}
x - 1 \\
x^3 + 6x^2 + 11x + 6) \overline{)x^4 + 5x^3 + 10x^2 + 20x + 24} \\
\underline{-x^4 - 6x^3 - 11x^2 - 6x} \\
-x^3 - x^2 + 14x + 24 \\
\underline{x^3 + 6x^2 + 11x + 6} \\
5x^2 + 25x + 30
\end{array}$$

Hence, we have that:

$$f(x) = g(x)(x-1) + 5(x^2 + 5x + 6)$$

2. Continuing the algorithm, now with g(x) and monic r(x) obtained above:

$$\begin{array}{r}
x+1 \\
x^2+5x+6) \overline{)x^3+6x^2+11x+6} \\
\underline{-x^3-5x^2-6x} \\
x^2+5x+6 \\
\underline{-x^2-5x-6} \\
0
\end{array}$$

Since the remainder is o, we find that $(x^2 + 5x + 6)$ divides exactly into g(x), and therefore also divides into f(x). Lastly, by looking at the remainder of the penultimate iteration, we can assert that, $x^2 + 5x + 6$ is the highest monic polynomial which divides both f(x) and g(x)

2 Complex Numbers