Multi-objective Optimization

Features:

- Distributed Data Flow
- Fog nodes mobility
- IoT devices mobility
- Migration support
- Partitioning techniques
- Data placement optimization
- Migration Optimization

Variables:

- QoS
  - latency
- Cost
  - CPU
  - RAM
  - MEM
  - BW
- Energy
  - busyPower
  - idlePower
- Bandwidth

Notations:

- N – fog nodes $S = \{s_1, ..., s_n\}$
- M – modules $U = \{u_1, ..., u_m\}$

- fP^{Mips} – matrix $1 \times N$ representing the MIPS price of each fog node per unit
- fP^{Ram} – matrix $1 \times N$ representing the memory price of each fog node per unit
- fP^{Strg} – matrix $1 \times N$ representing the storage price of each fog node per unit
- fP^{Bw} – matrix $1 \times N$ representing the bandwidth price of each fog node per unit
- fP^{En} – matrix $1 \times N$ representing the energy price of each fog node per unit

- f^{Mips} – matrix $N \times 1$ representing the MIPS capacity of each fog node
- f^{Ram} – matrix $N \times 1$ representing the memory capacity of each fog node
- f^{Strg} – matrix $N \times 1$ representing the storage capacity of each fog node
- f^{bPw} = matrix $N \times 1$ representing the busy power consumption of each fog node per unit
- f^{iPw} – matrix $N \times 1$ representing the idle power consumption of each fog node per unit

- m^{Mips} – matrix $M \times 1$ representing the MIPS needed for each application's module
- m^{Ram} – matrix $M \times 1$ representing the memory needed for each application's module
- m^{Strg} – matrix $M \times 1$ representing the storage needed for each application's module

- K – total number of dependencies/edges
- e^{Cpu} – matrix $1 \times K$ representing the tuple CPU size (MI) needed to be processed
- e^{Nw} – matrix $1 \times K$ representing the tuple network size (MB) needed to be sent
- e^{Prob} – matrix $1 \times K$ representing the probability of sending the tuple
- e^{Pe} – matrix $1 \times K$ representing the periodicity of the producer (i.e., periodic sources)
- e^{S} – matrix $1 \times K$ representing the edge source
- e^{D} – matrix $1 \times K$ representing the edge destination

- Z – total number of dependencies/edges between different pairs of modules
- l^{S} – matrix $N \times 1$ representing the starting nodes of each pair of modules
- l^{D} – matrix $N \times 1$ representing the ending nodes of each pair of modules

- mD – matrix $M \times M$ representing the dependencies between modules
- mB – matrix $M \times M$ representing the bandwidth needed between modules

- fL – matrix $N \times N$ representing the latency between each two direct nodes
- fB – matrix $N \times N$ representing the bandwidth between each two direct nodes
- f^{Tx} – matrix $N \times N$ representing the transmitter power between each two nodes

- D – matrix $N \times M$ representing the nodes where each module can be deployed
- C – matrix $N \times M$ current module placement

- e – matrix $1 \times N$ with all entries set to 1

- α – operational weight
- β – energetic weight
- γ – processing weight
- δ – transmission weight
- ζ – bandwidth weight
- η – migration weight

Variables:

- P – matrix $N \times M$ representing the placement mapping between modules and nodes
- R – matrix $Z \times N \times N$ representing the tuple routing map between modules
- V – matrix $M \times N \times N$ representing the VM routing map

Preliminary computations:

$$m_i^{Mips} = \sum_{k \in K} \left(\frac{e_k^{Prob} e_k^{Cpu}}{e_k^{Pe}} \right), e_k^D = i$$

$$mB_{i,j} = \sum_{k \in K} \left(\frac{e_k^{Prob} e_k^{Nw}}{e_k^{Pe}} \right), e_k^S = i, \; e_k^D = j$$

$$mD_{i,j} = \sum_{k \in K} \left(\frac{e_k^{Prob}}{e_k^{Pe}} \right), e_k^S = i, \; e_k^D = j$$

Problem formulation:

Operational Cost (C_O)

$$C_O = fP^{Mips} \times P \times m^{Mips} + fP^{Ram} \times P \times m^{Ram} + fP^{Strg} \times P \times m^{Strg} + \left(fP^{Bw} \sum_{z \in Z} mB_{l_z^S, l_z^D} \times R_z \right) e'$$

Power Cost (C_{Pw})

$$\begin{aligned} C_{Pw} = & \sum_{n \in N} \left(f_n^{iPw} + (f_n^{bPw} - f_n^{iPw}) \times \frac{P_n \times m^{Mips}}{f_n^{Mips}} \right) + \\ & \sum_{i \in N} \sum_{j \in N} f_{i,j}^{Tx} \left[fL_{i,j} \sum_{z \in Z} \left(mD_{l_z^S, l_z^D} \times R_{z,i,j} \right) + \frac{\sum_{z \in Z} \left(mB_{l_z^S, l_z^D} \times R_{z,i,j} \right)}{fB_{i,j}} \right] \\ & \sum_{i \in N} \sum_{j \in N} f_{i,j}^{Tx} \left[\frac{\sum_{z \in M} \left(m_z^{Mips} + \times m_z^{Ram} \right) \times V_{z,i,j}}{fB_{i,j}} + fL_{i,j} \sum_{z \in M} \times V_{z,i,j} \right] \end{aligned}$$

Processing Cost (C_P)

$$C_P = \frac{\left(\sum_{n \in N} \frac{P_n \times m^{Mips}}{f_n^{Mips}} \right)^2}{N \times \sum_{n \in N} \left(\frac{P_n \times m^{Mips}}{f_n^{Mips}} \right)^2} \quad -OR- \quad C_P = \sum_{n \in N} \frac{P_n \times m^{Mips}}{f_n^{Mips}}$$

Transmission Cost (C_T)

$$C_T = \sum_{i \in N} \sum_{j \in N} fL_{i,j} \sum_{z \in Z} \left(mD_{l_z^S, l_z^D} \times R_{z,i,j} \right)$$

Bandwidth Cost (C_B)

$$C_B = \sum_{i \in N} \sum_{j \in N} \frac{\sum_{z \in Z} \left(mB_{l_z^S, l_z^D} \times R_{z,i,j} \right)}{fB_{i,j}}$$

Migration Cost (C_M)

$$C_M = \sum_{i \in N} \sum_{j \in N} \left[\frac{\sum_{z \in M} \left(m_z^{Mips} + \times m_z^{Ram} \right) \times V_{z,i,j}}{fB_{i,j}} + fL_{i,j} \sum_{z \in M} \times V_{z,i,j} \right]$$

Final problem:

$$
\begin{aligned}
\underset{P,R}{\text{minimize}} \quad & C = \alpha C_O + \beta C_{Pw} + \gamma C_P + \delta C_T + \zeta C_B + \eta C_M \\
\text{subject to} \quad & P \times m^{Mips} \leq f^{Mips}, \\
& P \times m^{Ram} \leq f^{Ram}, \\
& P \times m^{Strg} \leq f^{Strg}, \\
& P \leq D \\
& P_{i,j} \in \{0,1\}, \forall i \in [0,N], \forall j \in [0,M] \\
& \sum_{i \in N} P_{i,j} = 1, \forall j \in [0,M] \\
& R_{z,i,j} \in \{0,1\}, \forall z \in [0,Z], \forall i \in [0,N], \forall j \in [0,N] \\
& \sum_{j \in N} R_{z,i,j} - \sum_{j \in N} R'_{z,i,j} = P_{i,l_z^S} - P_{i,l_z^D} \quad, \forall z \in [0,Z], \forall i \in [0,N] \\
& \sum_{z \in Z} mB_{l_z^S,l_z^D} \times R_{z,i,j} \leq fB_{i,j} \quad, \forall i \in [0,N], \forall j \in [0,N] \\
& \sum_{j \in N} V_{z,i,j} - \sum_{j \in N} V'_{z,i,j} = C_{j,i} - P_{j,i} \quad, \forall z \in [0,M], \forall i \in [0,N]
\end{aligned}
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