# Multi-objective Optimization

## Features:

- Distributed Data Flow
- Fog nodes mobility
- IoT devices mobility
- Migration support
- $\bullet\,$  Partitioning techniques
- $\bullet\,$  Data placement optimization
- Migration Optimization

## Variables:

- $\bullet$  QoS
  - latency
- $\bullet$  Cost
  - CPU
  - RAM
  - MEM
  - BW
- Energy
  - busyPower
  - idlePower
- Bandwidth

### **Notations:**

- $N \text{fog nodes } S = \{s_1, ..., s_n\}$
- $M \text{modules } U = \{u_1, ..., u_m\}$
- $fP^{Mips}$  matrix  $1 \times N$  representing the MIPS price of each fog node per unit
- $fP^{Ram}$  matrix  $1 \times N$  representing the memory price of each fog node per unit
- $fP^{Strg}$  matrix  $1 \times N$  representing the storage price of each fog node per unit
- $fP^{Bw}$  matrix  $1 \times N$  representing the bandwidth price of each fog node per unit
- $f^{Mips}$  matrix  $N \times 1$  representing the MIPS capacity of each fog node
- $f^{Ram}$  matrix  $N \times 1$  representing the memory capacity of each fog node
- $f^{Strg}$  matrix  $N \times 1$  representing the storage capacity of each fog node
- $f^{bPw} = \text{matrix } N \times 1$  representing the busy power consumption of each fog node per unit
- $f^{iPw}$  matrix  $N \times 1$  representing the idle power consumption of each fog node per unit
- $f^{wPw}$  matrix  $N \times 1$  representing the weight of energy consumption (willing\_to\_waste\_energy^{-1})
- $m^{Mips}$  matrix  $M \times 1$  representing the MIPS needed for each application's module
- $m^{Ram}$  matrix  $M \times 1$  representing the memory needed for each application's module
- $m^{Strg}$  matrix  $M \times 1$  representing the storage needed for each application's module
- $m^{Bw}$  matrix  $M \times 1$  representing the bandwidth needed for each application's module
- $\bullet$  K total number of dependencies/edges
- $e^{Cpu}$  matrix  $1 \times K$  representing the tuple CPU size (MI) needed to be processed
- $e^{Nw}$  matrix  $1 \times K$  representing the tuple network size (MB) needed to be sent
- ullet  $e^{Prob}$  matrix  $1 \times K$  representing the probability of sending the tuple
- $e^{Pe}$  matrix  $1 \times K$  representing the periodicity of the producer (i.e., periodic sources)
- $e^S$  matrix  $1 \times K$  representing the edge source
- $e^D$  matrix  $1 \times K$  representing the edge destination
- $\bullet$  Z total number of dependencies/edges between different pairs of nodes
- $l^S$  matrix  $N \times 1$  representing the starting nodes of each pair of nodes
- $l^D$  matrix  $N \times 1$  representing the ending nodes of each pair of nodes
- mD matrix  $M \times M$  representing the dependencies between modules
- $\bullet$  mB matrix  $M \times M$  representing the bandwidth needed between modules
- fL matrix  $N \times N$  representing the latency between each two direct nodes
- fB matrix  $N \times N$  representing the bandwidth between each two direct nodes
- D matrix  $N \times M$  representing the nodes where each module can be deployed
- $\bullet$  P matrix  $N \times M$  representing the placement mapping between modules and nodes
- R matrix  $Z \times N \times N$  representing the routing map between modules
- e matrix 1xN, with all entries to 1
- $\alpha$  operational weight
- $\beta$  energetic weight
- $\gamma$  processing weight
- $\delta$  latency weight
- $\zeta$  bandwidth weight

Preliminary computations:

$$\begin{split} m_i^{Mips} &= \sum_{k \in K} \left( \frac{e_k^{Prob} e_k^{Cpu}}{e_k^{Pe}} \right), e_k^D = i \\ m_i^{Bw} &= \sum_{k \in K} \left( \frac{e_k^{Prob} e_k^{Nw}}{e_k^{Pe}} \right), e_k^S = i \\ mB_{i,j} &= \sum_{k \in K} \left( \frac{e_k^{Prob} e_k^{Nw}}{e_k^{Pe}} \right), e_k^S = i, \ e_k^D = j \\ mD_{i,j} &= \sum_{k \in K} \left( \frac{e_k^{Prob} e_k^{Nw}}{e_k^{Pe}} \right), e_k^S = i, \ e_k^D = j \end{split}$$

#### Problem formulation:

The cost function is mainly characterized by two components: Operational Cost and Service Quality Cost.

**Operational Cost**  $(C_O)$  is characterized by the resources allocated in each fog node to support all users' computations, namely: CPU, memory, storage, and bandwidth.

$$C_O = fP^{Mips} \times P \times m^{Mips} + fP^{Ram} \times P \times m^{Ram} + fP^{Strg} \times P \times m^{Strg} + \left(fP^{Bw} \sum_{z \in Z} mB_{l_z^S, l_z^D} \times R_z\right) e'$$

**Power Cost**  $(C_{Pw})$  is characterized by the busy/idle power in each fog node to support all users' computations as well as the willing to wast energy to prevent clients to process the whole application.

$$C_{Pw} = \sum_{n \in N} \left( f_n^{iPw} + f_n^{wPw} (f_n^{bPw} - f_n^{iPw}) \times \frac{P_n \times m^{Mips}}{f_n^{Mips}} \right)$$

**Processing Cost**  $(C_P)$  is characterized by the percentage of unused MIPS using the Jain's fairness index.

$$C_P = \frac{\left(\sum_{n \in N} \frac{P_n \times m^{Mips}}{f_n^{Mips}}\right)^2}{N \times \sum_{n \in N} \left(\frac{P_n \times m^{Mips}}{f_n^{Mips}}\right)^2} - OR - C_P = \sum_{n \in N} \frac{P_n \times m^{Mips}}{f_n^{Mips}}$$

**Latency Cost**  $(C_L)$  is characterized by the total latency on the tuple transmission.

$$C_L = e\left(fL. * \sum_{z \in Z} mD_{l_z^S, l_z^D} \times R_z\right) e'$$

**Bandwidth Cost**  $(C_B)$  is characterized by the bandwidth usage on the tuple transmission.

$$C_B = e\left(1./(fB + \epsilon) \cdot * \sum_{z \in Z} mB_{l_z^S, l_z^D} \times R_z\right) e', \epsilon = 1^{-9}$$

### Final problem:

$$\begin{split} & \underset{P,R}{\text{minimize}} & \quad C = \alpha C_O + \beta C_{Pw} + \gamma C_P + \delta C_L + \zeta C_B \\ & \text{subject to} & \quad P \times m^{Mips} \leq f^{Mips}, \\ & \quad P \times m^{Ram} \leq f^{Ram}, \\ & \quad P \times m^{Strg} \leq f^{Strg}, \\ & \quad P \leq D \\ & \quad P_{i,j} \in \{0,1\}, \forall i \in [0,N], \forall j \in [0,M] \\ & \quad \sum_{i \in N} P_{i,j} = 1, \forall j \in [0,M] \\ & \quad R_{z,i,j} \in \{0,1\}, \forall z \in [0,Z], \forall i \in [0,N], \forall j \in [0,N] \\ & \quad \sum_{j \in N} R_{z,i,j} - \sum_{j \in N} R'_{z,i,j} = P_{i,l_z^S} - P_{i,l_z^D} \quad, \forall z \in [0,Z], \forall i \in [0,N] \\ & \quad \sum_{z \in Z} m B_{l_z^S, l_z^D} \times R_{z,i,j} \leq f B_{i,j} \quad, \forall i \in [0,N], \forall j \in [0,N] \end{split}$$