# Rolling your own Crypto Why you shouldn't

**Jules Poon** 

#### whoami

- 2 Years security experience:
  - Centre for Strategic Infocomm Technologies: Did RE and Tool dev
  - STAR Labs: Died reversing XNU and kernel extensions

- 1 Years playing cyber competitions (Social Engineering Experts)
  - Did Crypto and RE (and a little exploitation)

O Years being a professional cryptographer

#### **Contents**

- 1. Cryptography: A Very Short Intro
- 2. How not to Cryptanalysis
- 3. The **Stack**
- 4. Real Life Attacks
- 5. The Bad API Problem

#### **Contents**

- 1. Cryptography: A Very Short Intro
- 2. How not to Cryptanalysis
- 3. The **Stack**
- 4. Real Life Attacks
- 5. The Bad API Problem

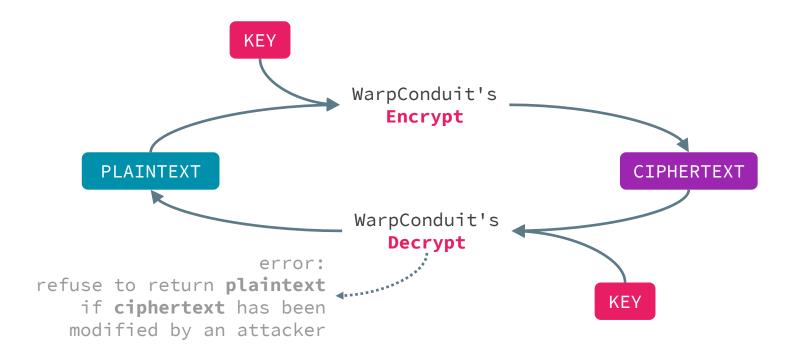


Published: 14/04/2013
Retrieved: 02/10/2022

Encapsulation that does
integrity-authenticated
symmetric encryption

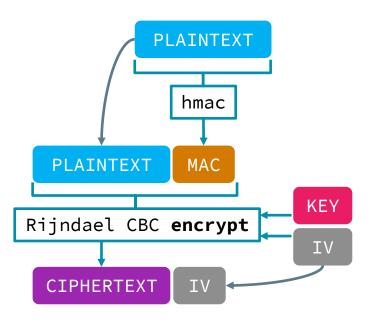
Attackers cannot change the ciphertext to decrypt into a different plaintext

\_\_\_\_

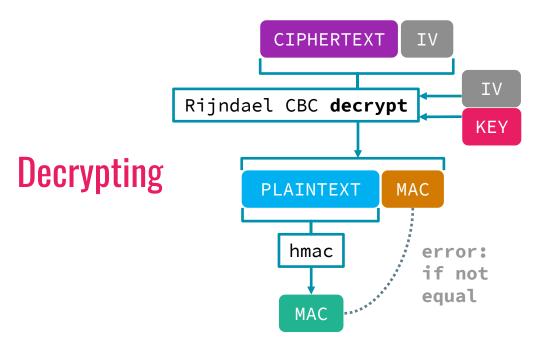


```
function mc encrypt($encrypt, $key){
    $encrypt = serialize($encrypt);
    $iv = mcrypt create iv(
        mcrypt get iv size(
            MCRYPT RIJNDAEL 256,
            MCRYPT MODE CBC),
        MCRYPT DEV URANDOM);
    key = pack('H*', key);
    $mac = hash hmac(
             'sha256', $encrypt,
            substr(bin2hex($key), -32));
    $passcrypt = mcrypt encrypt(
        MCRYPT RIJNDAEL 256,
        $key, $encrypt.$mac,
        MCRYPT MODE CBC, $iv);
    $encoded = base64 encode($passcrypt)
             .base64 encode($iv);
    return $encoded;
```

#### **Encrypting**



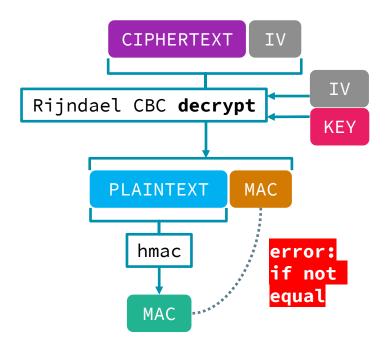
```
function mc decrypt($decrypt, $key){
    $decrypt = explode('|', $decrypt.'|');
    $decoded = base64_decode($decrypt[0]);
    $iv = base64 decode($decrypt[1]);
    if (strlen($iv)
        !== mcrypt get iv size(
            MCRYPT RIJNDAEL 256,
            MCRYPT MODE CBC)){return false;}
    key = pack('H*', key);
    $decrypted = trim(mcrypt_decrypt(
        MCRYPT RIJNDAEL 256, $key,
        $decoded, MCRYPT MODE CBC, $iv));
    $mac = substr($decrypted, -64);
    $decrypted = substr($decrypted, 0, -64);
    $calcmac = hash hmac(
        'sha256', $decrypted, s
        ubstr(bin2hex($key), -32));
    if($calcmac!==$mac){return false;}
    $decrypted = unserialize($decrypted);
    return $decrypted;
```



\_\_\_\_

MAC check prevents
modifying of ciphertext

An attacker that modifies the ciphertext cannot compute a new MAC without knowing the KEY, so the check fails.





Coined by Taylor Hornby @DefuseSec

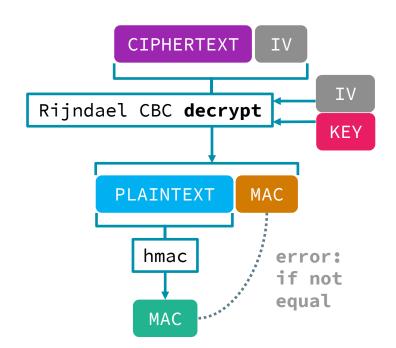


#### Primitives used:

- Rijndael-256
- SHA-256

# Both these **primitives** are secure

- protocol is just a bunch of secure primitives wired together
- protocol is secure





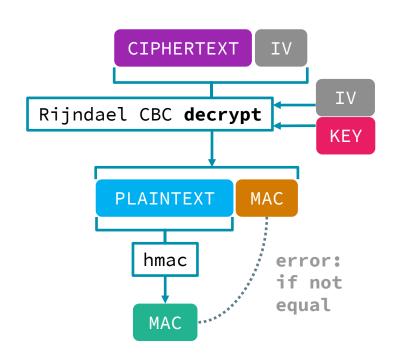
What this protocol should provide:

#### Privacy:

• **Rijndael-256** is secure so no problem

#### Integrity:

SHA-256 is secure so no problem





- Reasoning by LEGO is good for implementation
  - Concept of abstraction in software development



- Reasoning by LEGO is good for implementation
   Concept of abstraction in software development
- Reasoning by LEGO is wrong for cryptanalysis



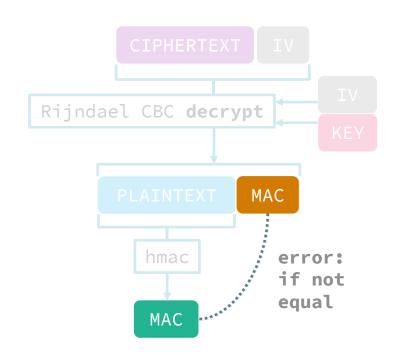
- Reasoning by LEGO is good for implementation
   Concept of abstraction in software development
- Reasoning by LEGO is wrong for cryptanalysis
- One must consider the whole system together

```
function mc decrypt($decrypt, $key){
   $decrypt = explode('|', $decrypt.'|');
   $decoded = base64_decode($decrypt[0]);
   $iv = base64 decode($decrypt[1]);
   if (strlen($iv)
                                                                  Rijndael CBC decrypt
       !== mcrvpt get iv size(
           MCRYPT RIJNDAEL 256,
           MCRYPT_MODE_CBC)){return false;}
   key = pack('H*', key);
   $decrypted = trim(mcrypt_decrypt(
                                                                                            MAC
       MCRYPT RIJNDAEL 256, $key,
                                             Not constant
       $decoded, MCRYPT MODE CBC, $iv));
                                             time compare
   $mac = substr($decrypted, -64);
   $decrypted = substr($decrypted, 0, -6
                                                                                               error:
                                              We can leak
   $calcmac = hash hmac(
                                                                                               if not
       'sha256', $decrypted,
                                               info about
                                                                                               equal
       ubstr(bin2hex($key), -32));
                                                    the
                                                                               MAC
   if($calcmac!==$mac){return false;}
                                               plaintext!
   $decrypted = unserialize($decrypted);
   return $decrypted;
```

\_\_\_\_

#### Reasoning by LEGO:

- A vuln here shouldn't have anything to do with the plaintext (only the MAC).
- Mayyybe integrity is compromised but not privacy??



Not constant-time string compare:

#### Your Code

```
if (!strcmp(secret, attacker_input))
    ERROR("...");
/* ... */
```

#### Glibc source

```
int
strcmp(const char *p1, const char *p2)
   const unsigned char *s1 = (const unsigned char *)p1;
   const unsigned char *s2 = (const unsigned char *)p2;
   unsigned char c1, c2;
   do
       c1 = (unsigned char)*s1++;
       c2 = (unsigned char)*s2++;
       if (c1 == '\0')
                            Early stopping:
           return c1 - c2;
    } while (c1 == c2);
                            First char wrong
    return c1 - c2;
                            -> function returns
                            faster
```

\_\_\_\_

Web Timing Attacks Made Practical\*

Timothy D. Morgan<sup>†</sup> Jason W. Morgan<sup>‡</sup>

August 3, 2015

#### Abstract

This paper addresses the problem of exploiting timing side channels in web applications. To date, differences in execution time have been difficult to detect and to exploit. Very small differences in execution time induced by different security logics, coupled with the fact that these small differences are often lost to significant network noise, make their detection difficult. Additionally, testing for and taking advantage of timing

\_\_\_

Scenario	Classifier		5000	Delta (ns)	200	40
		25000		1000		
lnx						
	midsummary	29  obs	894 obs	17147 obs	16.60% err	38.60% err
	quadsummary	26 obs	894 obs	16289 obs	20.55% err	47.30% err
	septasummary	15 obs	894 obs	17147 obs	22.35% err	45.20% err
	boxtest	146 obs	20.80% err	36.30%  err	47.55% err	49.85% err

Not constant-time string compare: [Aside]

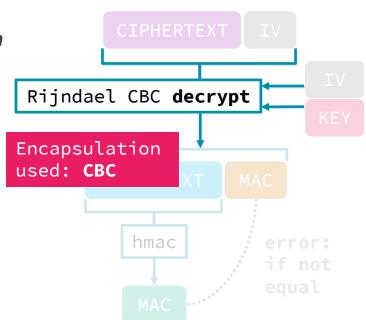
#### Glibc source [x8664 arch optimised impl]

```
L(no_page_cross):
    /* Safe to compare 4x vectors. */
    VMOVU (%rdi), %ymm0
    /* 1s where s1 and s2 equal. Just VPCMPEQ if its not strcasecmp.
        Otherwise converts ymm0 and load from rsi to lower. ymm2 is scratch and ymm1 is the return. */
    CMP_R1_S2_ymm (%ymm0, (%rsi), %ymm2, %ymm1)
    /* 1s at null CHAR. */
    VPCMPEQ %ymm0, %ymmZERO, %ymm2
    /* 1s where s1 and s2 equal AND not null CHAR. */
    vpandn %ymm1, %ymm2, %ymm1
    /* All 1s -> keep going, any 0s -> return. */
    vpmovmskb %ymm1, %ecx
```

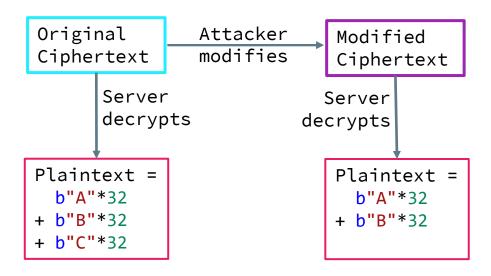
Compares 32 bytes at once:
Timing attack is way harder here

[Aside] Properties of the *Encapsulation* used:

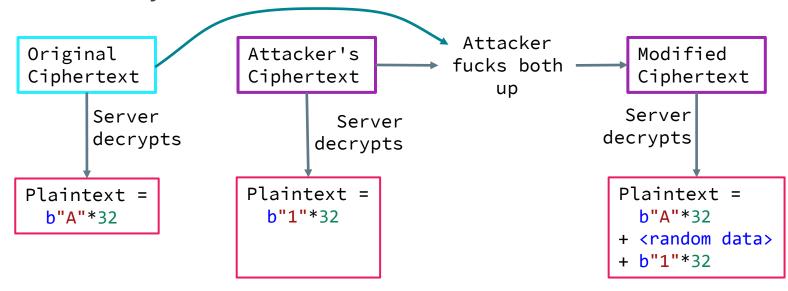
- Rijndael-256 encrypts in blocks of 32 bytes
- To encrypt longer data it uses CBC Encapsulation



CBC: Can truncate plaintext by 32 bytes by modifying ciphertext without the key.



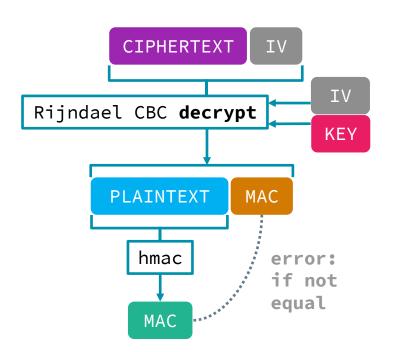
CBC: Can combine two ciphertexts such that the plaintext concats without the key.



\_\_\_\_

#### Attack Scenario:

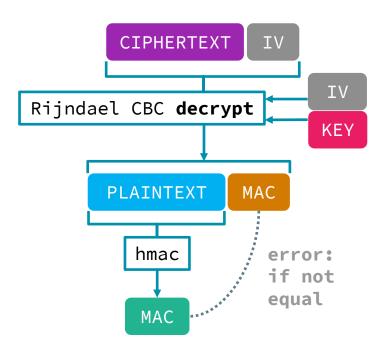
- Server encrypts whatever attacker sends
- Server decrypts whatever attacker sends but doesn't return
- Goal: recover secret given its ciphertext



\_\_\_\_

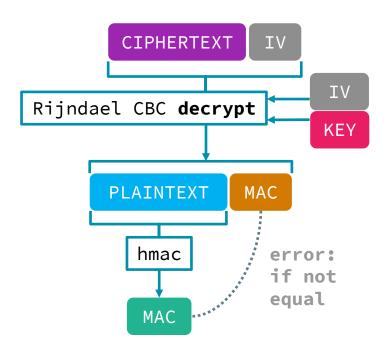
#### Assumptions on attacker:

 Attacker can measure time taken by server to compare MAC to tell how many bytes were matched

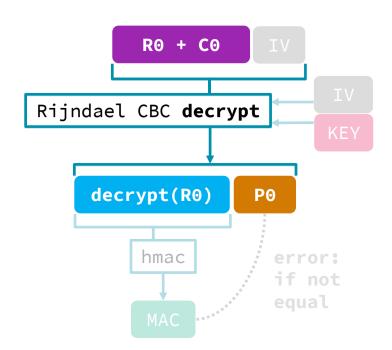


\_\_\_\_

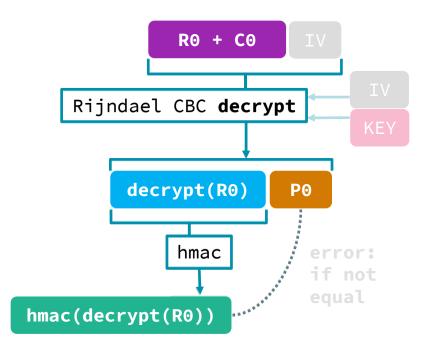
- Not an unrealistic scenario
  - E.g. Secret hidden in session cookie



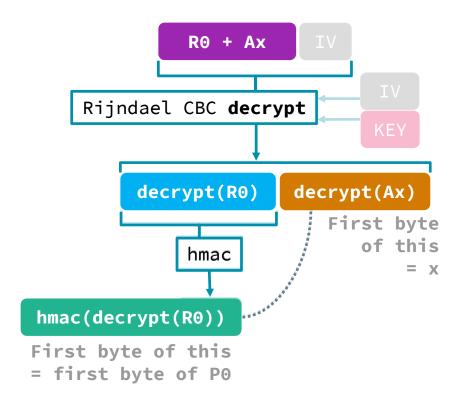
- Get 256 ciphertexts whose plaintext starts with b"\x00" to b"\xff" (A0, A1, ... A255) by requesting the server. decrypt(Ax)[0] = x
- Take first 32 bytes of encrypted secret CO, prepend with randomly generated block RO
  - Tricks server to take first 32
     bytes of secret (P0) as MAC



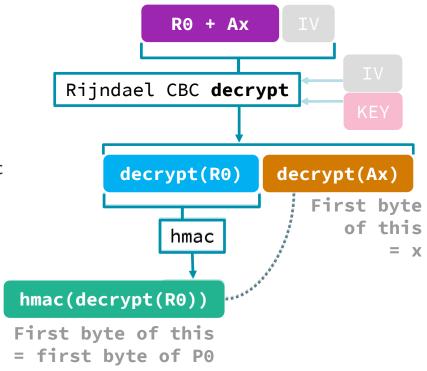
- Try different R0 until
  hmac(decrypt(R0))[0] = P0[0]
  - Attacker can tell via timing attack



- Swap C0 with A0,...,A255 until
  decrypt(Ax)[0] = x =
  decrypt(R0)[0] = P0[0]
  - E.g. If PO's first byte is
     42, using A42 will make hmac
     comparison take longer
- P0[0] recovered!

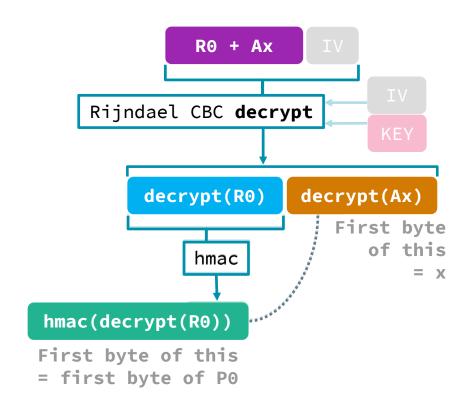


- Can be extended to leak first few bytes of every 32-byte block of plaintext
  - Potentially all bytes of plaintext if one can control block alignment.



#### Take Away:

- Parts of the system that seem totally disconnected can fuck each other up
- There's so much things to consider beyond the primitives
  - Primitives are not a good abstraction of security guarantees

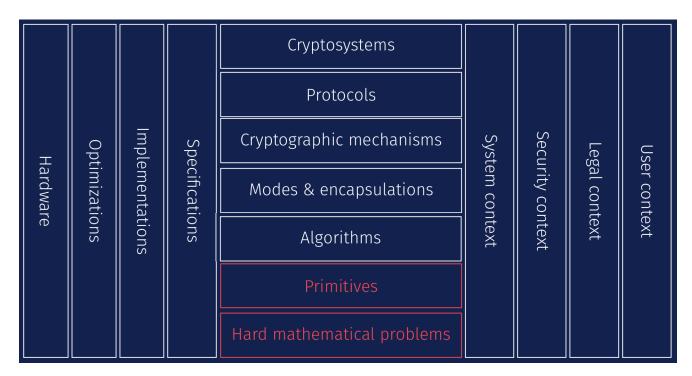


#### **Contents**

- 1. Cryptography: A Very Short Intro
- 2. How not to Cryptanalysis
- 3. The **Stack**
- 4. Real Life Attacks
- 5. The Bad API Problem

#### The Stack

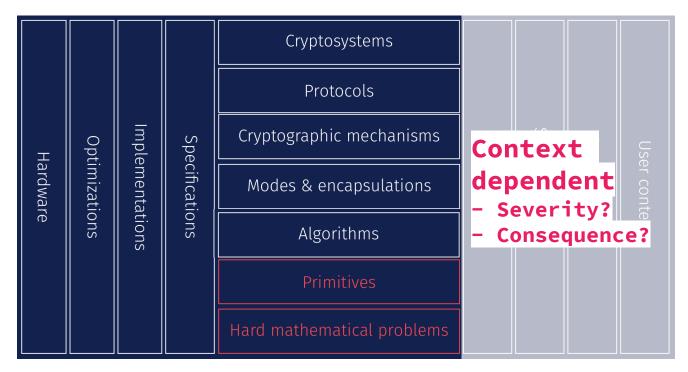
\_\_\_



Stolen from: https://blog.quarkslab.com/status-of-post-quantum-cryptography-implementation.html

#### The Stack

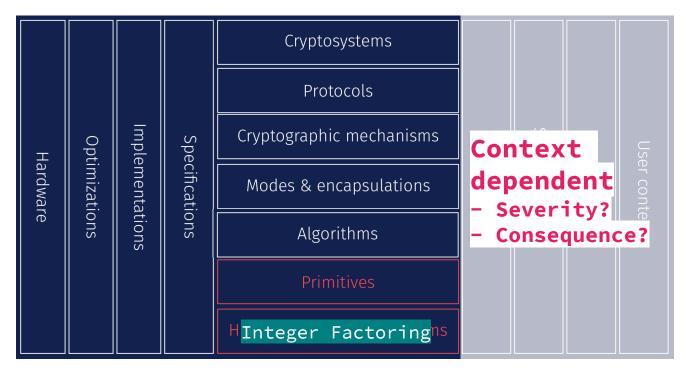
\_ \_ \_



Stolen from: https://blog.quarkslab.com/status-of-post-quantum-cryptography-implementation.html

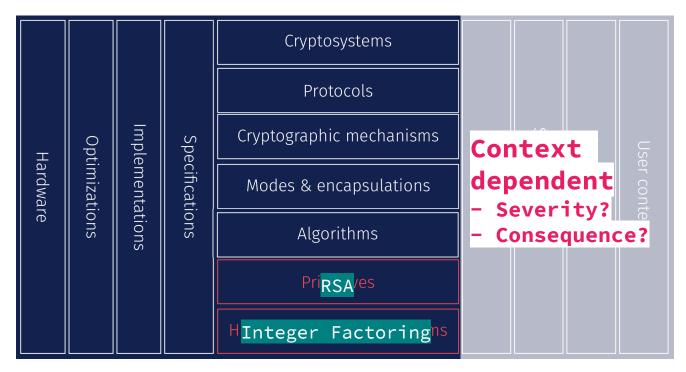
#### The Stack

\_\_\_\_

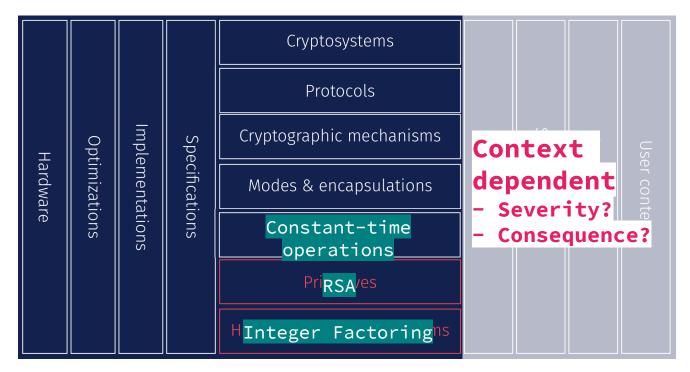


Stolen from: https://blog.quarkslab.com/status-of-post-quantum-cryptography-implementation.html

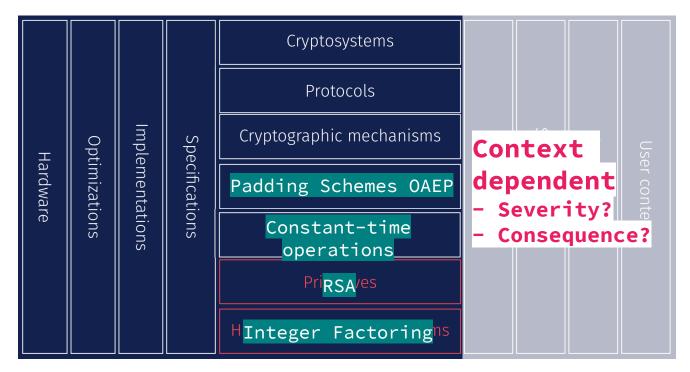
\_\_\_\_



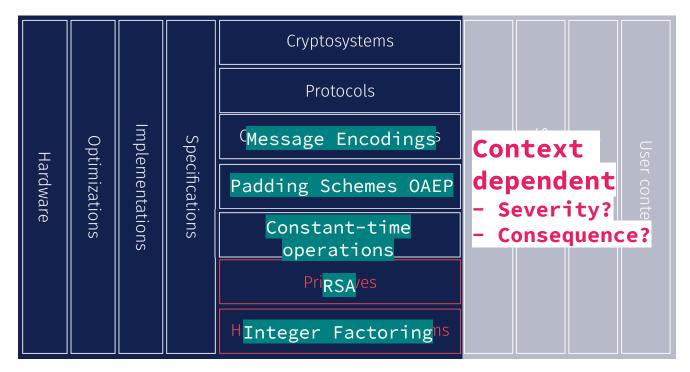
\_ \_ \_



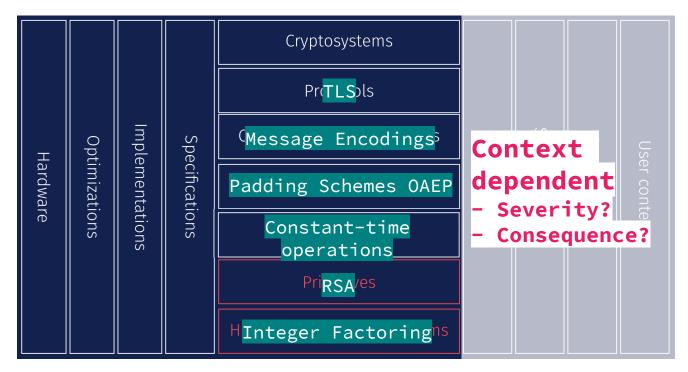
\_\_\_\_\_



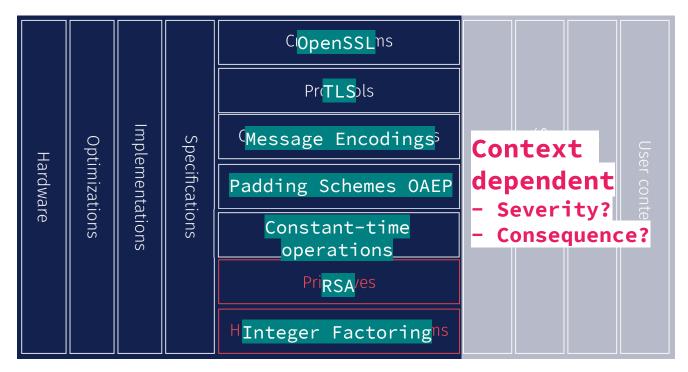
\_ \_ \_



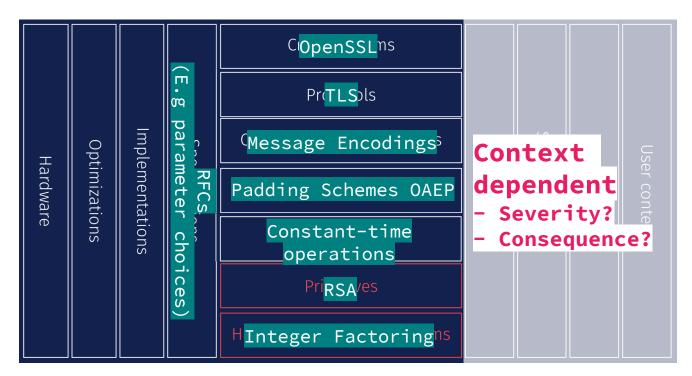
\_\_\_\_



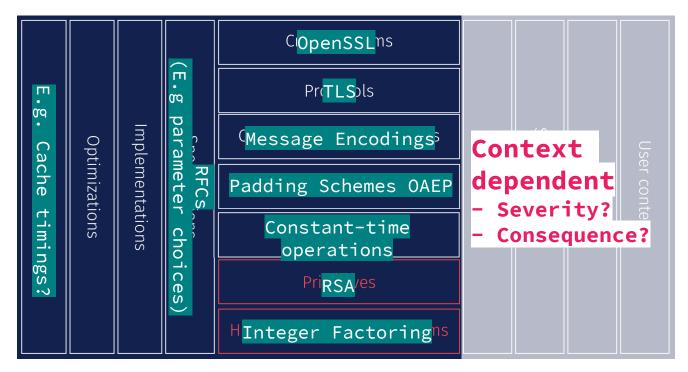
\_\_\_\_



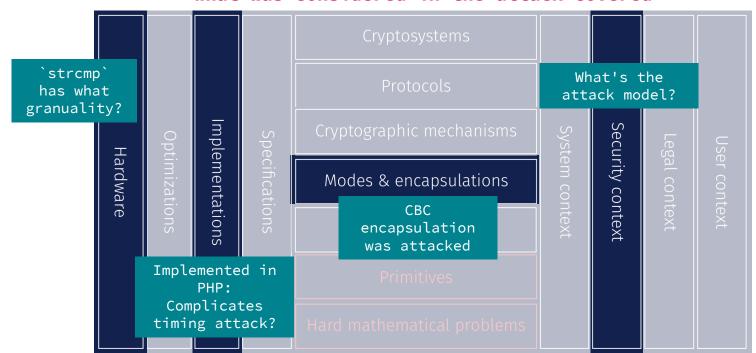
\_\_\_\_

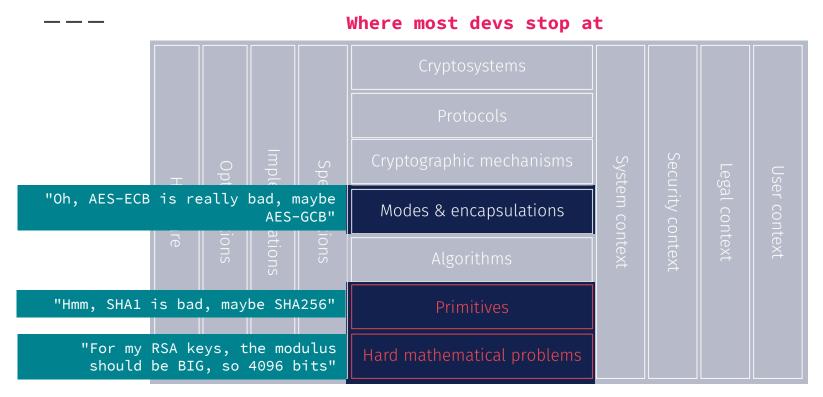


\_\_\_\_\_



What was considered in the attack covered





### Unknown Unknowns

```
"There are things that we don't know we don't know." ~ Donald Rumsfeld
```

- Cryptography is all about unknown unknowns.
  - MACs have to be compared constant time?
  - Nonce shouldn't be used more than once?
  - vou need to properly pad RSA messages?

### **Contents**

- 1. Cryptography: A Very Short Intro
- 2. How not to Cryptanalysis
- 3. The **Stack**
- 4. Real Life Attacks
- 5. The Bad API Problem

- Firmware updates are often signed
  - Ensure updates are from a trusted source (not an attacker)
- Firmware is also often encrypted
  - Prevents reverse engineering

#### D-Audio2V Firmware Update: A Disaster

Research done by programmingwithstyle.com/posts/howihackedmycar

- Firmware signed with RSA
  - Public Key is an example key -> Private key is online

#### D-Audio2V Firmware Update: A Disaster

Research done by programmingwithstyle.com/posts/howihackedmycar

- Firmware signed with RSA
  - Public Key is an example key -> Private key is online
- Firmware update comes in encrypted ZIP
  - PKZIP vulnerable to a Known Plaintext Attack

#### D-Audio2V Firmware Update: A Disaster

Research done by programmingwithstyle.com/posts/howihackedmycar

- Firmware signed with RSA
  - Public Key is an example key -> Private key is online
- Firmware update comes in encrypted ZIP
  - PKZIP vulnerable to a Known Plaintext Attack
- System image in update encrypted with AES-CBC
  - Key (same for enc and dec) publicly available in opensourced code

### **Contents**

- 1. Cryptography: A Very Short Intro
- 2. How not to Cryptanalysis
- 3. The **Stack**
- 4. Real Life Attacks
- 5. The Bad API Problem

- Developers not aware of what cryptographic operations to do
   Sign? Encrypt??
- Developers no idea how to implement operations

Javascript Object Signing and Encryption (JOSE)

A set of standards to sign and encrypt json (essentially)

**JWS:** Anybody can see the data, and there's a signature to verify the legitimacy of data

JWE: Data is encrypted

# JWT Web Tokens:

#### Encoded PASTE A TOKEN HERE

eyJhbGciOiJIUzI1NiIsInR5cCI6IkpXVCJ9.ey
JzdWIiOiIxMjMONTY3ODkwIiwibmFtZSI6Ikpva
G4gRG91IiwiaWF0IjoxNTE2MjM5MDIyfQ.SflKx
wRJSMeKKF2QT4fwpMeJf36P0k6yJV\_adQssw5c

#### Decoded EDIT THE PAYLOAD AND SECRET

```
HEADER: ALGORITHM & TOKEN TYPE
                     Alg used
  "alg": "HS256",
  "typ": "JWT"
                      (not encrypted)
PAYLOAD: DATA
                     Data
   "sub": "1234567890".
  "name": "John Doe",
                      (not encrypted)
  "iat": 1516239022
VERIFY SIGNATURE
                     Signature
HMACSHA256(
  base64Ur1Encode(header
                      (verifies
  base64UrlEncode(payload
                      `HEADER` and
  your-256-bit-secret
                       PAYLOAD` using
 ) secret base64 encoded
                       alg`)
```

```
Encoded PASTE A TOKEN HERE
                                                     Decoded EDIT THE PAYLOAD AND SECRET
                                                      HEADER: ALGORITHM & TOKEN TYPE
  eyJhbGciOi Problem 1:
  JzdWIi0iIx
              `alg` can be tampered
                                                                          Alg used
                                                         "alg": "HS256",
  G4gRG91Iiv
                                                         "typ": "JWT"
              by attacker BEFORE it
                                                                           (not encrypted)
  wRJSMeKKF2
              is verified
                                                      PAYLOAD: DATA
                                                                          Data
                                                         "sub": "1234567890",
                                                         "name": "John Doe".
                                                                           (not encrypted)
                                                         "iat": 1516239022
                                                      VERIFY SIGNATURE
                                                                          Signature
                                                       HMACSHA256(
                                                                           (verifies
                                                        base64Ur1Encode(header
                                                                            `HEADER` and
                                                        base64UrlEncode(payload
                                                        vour-256-bit-secret
                                                                           `PAYLOAD` using
                                                        ☐ secret base64 encode
                                                                            alg`)
```

```
Possible `alg` values:
```

- HS256 <-- Symmetric signing
- RS256 <-- Asymmetric signing
- None <-- Lmao</li>

#### Trivial Attack 1:

Attacker specifies `alg: None`. Server doesn't verify signature

#### **Trivial Attack 2:**

Attacker changes `alg: RS256` to `alg: HS256` and signs token with PUBLIC key.

```
Possible `alg` values:
```

- HS256 <-- Symmetric signing
- RS256 <-- Asymmetric signing
- None <-- Lmao</li>

The "alg" value [...] MUST be present and MUST be understood and processed by implementations.

~ RFC7515 Section 4.1.1

Server **HAVE** to use the `alg` given in JWT to be RFC7515 compliant

Signing: JWS

```
import jose
                                            jwt = jose.serialize compact(jws)
                                            'eyJhbGciOiAiSFMyNTYifQ.eyJpc3MiOiAiaHR0cDov
                                            L3d3dy5leGFtcGxlLmNvbSIsICJzdWIiOiA0MiwgImV4
claims = {
    'iss': 'http://www.example.com',
                                            cCI6IDEzOTU2NzQ0Mjd9.WYApAiwiKd-
    'exp': int(time()) + 3600,
                                            eDClA1fg7XFrnfHzUTgrmdRQY4M19Vr8'
    'sub': 42,
                                            jws = jose.deserialize compact(jwt)
                                            jose.verify(jws, jwk, jsw.alg)
jwk = {'k': 'password'}
                                            JWT(header={u'alg': u'HS256'},
                                            claims={u'iss': u'http://www.example.com',
jws = jose.sign(claims, jwk, alg='HS256')
                                            u'sub': 42, u'exp': 1395674427})
```

Signing: JWS

```
import jose
                                             jwt = jose.serialize compact(jws)
                                             'eyJhbGciOiAiSFMyNTYifQ.eyJpc3MiOiAiaHR0cDov
                                             L3d3dy5leGFtcGxlLmNvbSIsICJzdWIiOiA0MiwgImV4
claims = {
    'iss': 'http://www.example.com',
                                             cCI6IDEzOTU2NzQ0Mjd9.WYApAiwiKd-
    'exp': int(time()) + 3600,
                                             eDClA1fg7XFrnfHzUTgrmdRQY4M19Vr8'
    'sub': 42,
                                             jws = jose.deserialize compact(jwt)

√(jws, jwk, jsw.alg)

                         Problem 2:
jwk = {'k': 'password'}
                          Potential for key reuse iss': u'HS256'},

iss': u'http://www.example.com',
jws = jose.sign(claims, jwk, alg='HS256') u'sub': 42, u'exp': 1395674427})
```

```
A JSON Web Key (JWK) [...] represents a cryptographic key. [...] allows for a generalized key as input that can be applied to a number of different algorithms that may expect a different number of inputs.

~ JOSE Official Documentation
```

JWK designed with reuse in different algos in mind

Sophie, indistinguishable from random noise for a gen @SchmiegSophie different cumentation Replying to @str4d The even larger point is that you should never use a key without context outside of the actual cryptographic implementation. If I manage to change your AES-GCM mode into AES-CTR on a decryption oracle, I get your auth key, not disimilar to how changing the curve parameters leaks

- 1. Algo to encrypt the key (CEK Encryption, `alg` param)
  - 1. RSA with PKCS #1v1.5 padding
  - 2. RSA with OAEP padding
  - 3. ECDH
  - 4. AES-GCM
- 2. Algo to encrypt the data (Claims Encryption, 'enc' param)
  - 1. AES-CBC + HMAC
  - 2. AES-GCM

### JWE: Developers specify two algos

- 1. Algo to encrypt the key (CEK Encryption, `alg` param)
  - 1. RSA with PKCS #1v1.5 padding
  - 2. RSA with OAEP padding
  - 3. ECDH
  - 4. AES-GCM

Problem 3:

**INSECURE CONFIGURATIONS** 

- 2. Algo to encrypt the data (Claims Encryption, 'enc' param)
  - 1. AES-CBC + HMAC
  - 2. AES-GCM

- 1. Algo to encrypt the key (CEK Encryption, `alg` param)
  - 1. RSA with PKCS #1v1.5 padding
  - 2. RSA with OAEP padding
  - 3. ECDH
  - 4. AES-GCM
- 2. Algo to encrypt the data (Claims Encryption, `enc` param)
  - 1. AES-CBC + HMAC
  - 2. AES-GCM

- 1. Algo to encrypt the key (CEK Encryption, `alg` param)
  - 1. RSA with PKCS #1v1.5 padding Very famous padding oracle attack
  - 2. RSA with OAEP padding
  - 3. ECDH
  - 4. AES-GCM
- 2. Algo to encrypt the data (Claims Encryption, 'enc' param)
  - 1. AES-CBC + HMAC
  - 2. AES-GCM

- 1. Algo to encrypt the key (CEK Encryption, `alg` param)
  - 1. RSA with PKCS #1v1.5 padding Very famous padding oracle attack
  - 2. RSA with OAEP padding
  - 3. **ECDH** Weierstrass form: Prone to invalid curve attacks
  - 4. AES-GCM
- 2. Algo to encrypt the data (Claims Encryption, 'enc' param)
  - 1. AES-CBC + HMAC
  - 2. AES-GCM

- 1. Algo to encrypt the key (CEK Encryption, `alg` param)
  - 1. RSA with PKCS #1v1.5 padding Very famous padding oracle attack
  - 2. RSA with OAEP padding
  - 3. ECDH Weierstrass form: Prone to invalid curve attacks
  - 4. AES-GCM Shared key?????
- 2. Algo to encrypt the data (Claims Encryption, 'enc' param)
  - 1. AES-CBC + HMAC
  - 2. AES-GCM

- 1. Algo to encrypt the key (CEK Encryption, `alg` param)
  - 1. RSA with PKCS #1v1.5 padding Very famous padding oracle attack
  - 2. RSA with OAEP padding
  - 3. ECDH Weierstrass form: Prone to invalid curve attacks
  - 4. AES-GCM Shared key?????
- 2. Algo to encrypt the data (Claims Encryption, `enc` param)
  - 1. AES-CBC + HMAC CBC: Malleable (as we have seen)
  - 2. AES-GCM

#### JOSE: Takeaway

- 1. Insecure standard
  - Secure usage breaks standard
- 2. "Sign or encrypt? Or both?? Wtf is none??"
- 3. "What **alg** for key/msg enc?"
  - More insecure configs than secure configs
- 4. API leaves too much room for error (e.g. usage of key)

### The Bad API Problem

Devs trying to implement crypto properly go down deep rabbitholes

• "developers lacked the confidence to choose the best algorithm or parameter" (Hazhirpasand et al, 2021)

#### Solution?

- Developers should git gud
- Crypto libraries should be proper abstractions
  - easy to use, hard to misuse, and secure by default

Libsodium: Password hashing

Libsodium: Password hashing

Libsodium: Password hashing

Libsodium: Password hashing #define PASSWORD "Correct Horse Battery Staple" #define KEY LEN crypto box SEEDBYTES unsigned char salt[crypto\_pwhash\_SALTBYTES]; unsigned char key[KEY LEN]; Generates salt for you randombytes buf(salt, sizeof salt); if (crypto pwhash **Dev-relevant parameters** (key, sizeof key, PASSWORD, strlen(PASSWORD), salt, crypto pwhash OPSLIMIT INTERACTIVE, crypto pwhash MEMLIMIT INTERACTIVE, crypto pwhash ALG DEFAULT) != 0) { /\* out of memory \*/ Never touch

the primitives

Libsodium: Authenticated Encryption Auxillary Data (AEAD)

### **AEAD** constructions

Documentation tells you exactly what AEAD is

#### This operation:

- Encrypts a message with a key and a nonce to keep it confidential
- Computes an authentication tag. This tag is used to make sure that the message, as well as optional, non-confidential (non-encrypted) data, haven't been tampered with.

A typical use case for additional data is to authenticate protocol-specific metadata about the message, such as its length and encoding.

### Libsodium: Authenticated Encryption Auxillary Data (AEAD)

vailability and in	teroperability #			
Construction	Key size	Nonce size	Block size	
AES256-GCM	256 bits	96 bits	128 bits	Multiple choices?
ChaCha20-Poly1305	256 bits	64 bits	512 bits	All of them are secur
ChaCha20-Poly1305- IETF	256 bits	96 bits	512 bits	
XChaCha20-Poly1305- IETF	256 bits	192 bits	512 bits	

Libsodium: Authenticated Encryption Auxillary Data (AEAD)

#### Limitations

Construction	Max bytes for a single (key,nonce)	Max bytes for a single key
AES256-GCM	~ 64 GB	~ 350 GB (for ~16 KB long messages)
ChaCha20-Poly1305	No practical limits (~ 2^64 bytes)	Up to 2^64* messages, no practical total size limits
ChaCha20-Poly1305-IETF	256 GB	Up to 2^64* messages, no practical total size limits
XChaCha20-Poly1305-IETF	No practical limits (~ 2^64 bytes)	Up to 2^64* messages, no practical total size limits

Developers are given dev-relevant info about each choice

\_\_\_\_

Libsodium: Authenticated Encryption Auxillary Data (AEAD)

#### TL;DR: which one should I use?

XChaCha20-Poly1305-IETF is the safest choice.

if you don't have any
special considerations

Still unsure?

There's recommendations

Other choices are only present for interoperability with other libraries that don't implement XChaCha20-Poly1305-IETF yet.

# QnA