

How two undergrads from the other side of the planet are speeding up your future code

Ken Jin, Jules Poon

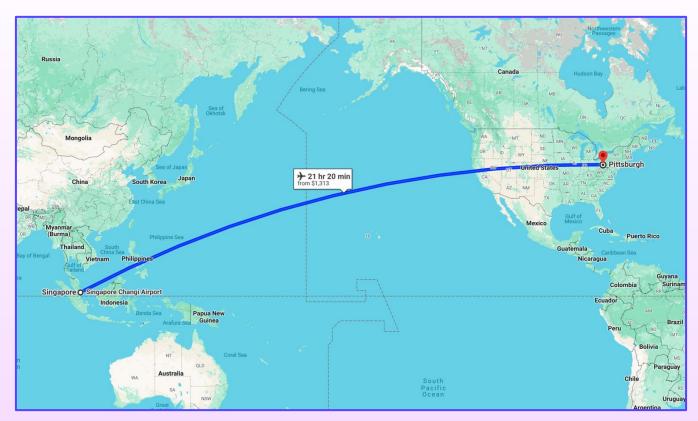


Introduction





The Other Side of the Planet



Introductions

Ken Jin:

- Undergrad student
- CPython Core Developer since 2021
- (Part-time) SWE at Quansight Labs

Jules:

- Math undergrad
- Interested in commutative algebra and programming languages

Goal





^{*} CPython with lazy basic block versioning (Maxime Chevalier-Boisvert and Marc Feeley, 2014).



^{*} CPython with lazy basic block versioning (Maxime Chevalier-Boisvert and Marc Feeley, 2014).



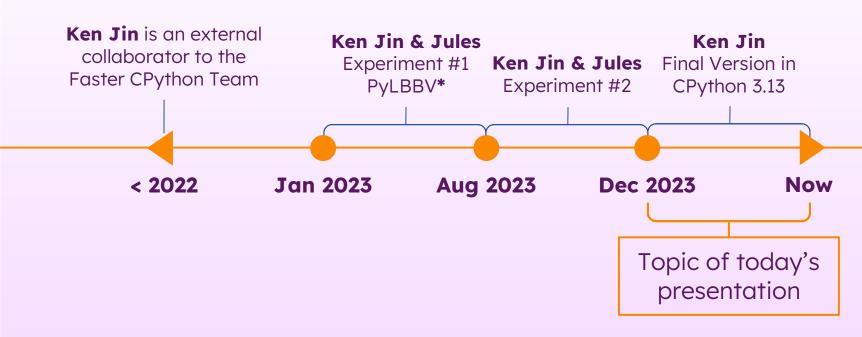
^{*} CPython with lazy basic block versioning (Maxime Chevalier-Boisvert and Marc Feeley, 2014).



^{*} CPython with lazy basic block versioning (Maxime Chevalier-Boisvert and Marc Feeley, 2014).



^{*} CPython with lazy basic block versioning (Maxime Chevalier-Boisvert and Marc Feeley, 2014).



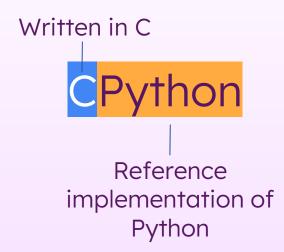
^{*} CPython with lazy basic block versioning (Maxime Chevalier-Boisvert and Marc Feeley, 2014).

Background

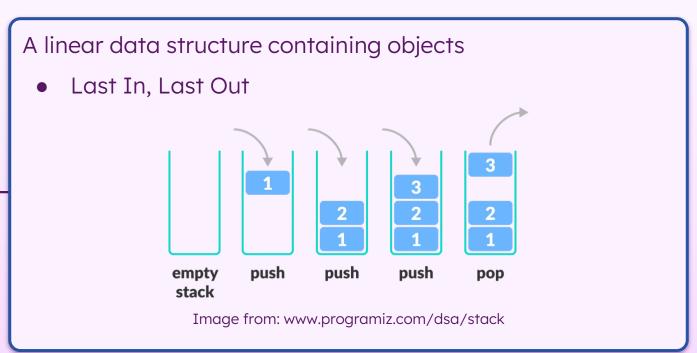




CPython



Bytecode Stack Machine An instruction set for easy interpretation by the *interpreter* Bytecode is easier to interpret compared to Python source A compiler converts Python source to Bytecode **CPython** Python Program Compiler → Bytecode → Interpreter Source output



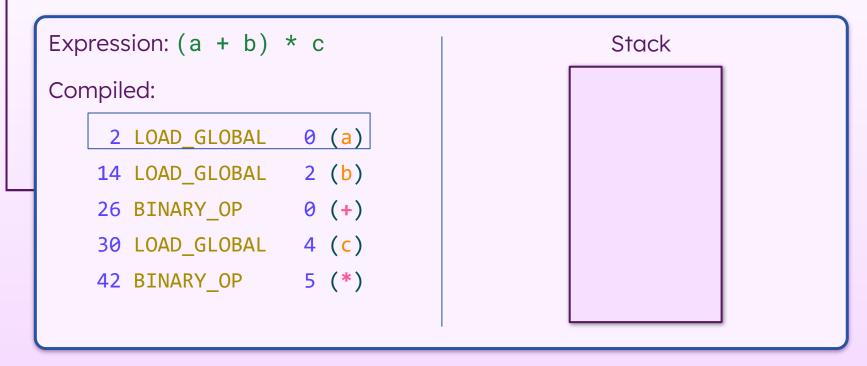
Bytecode Stack Machine

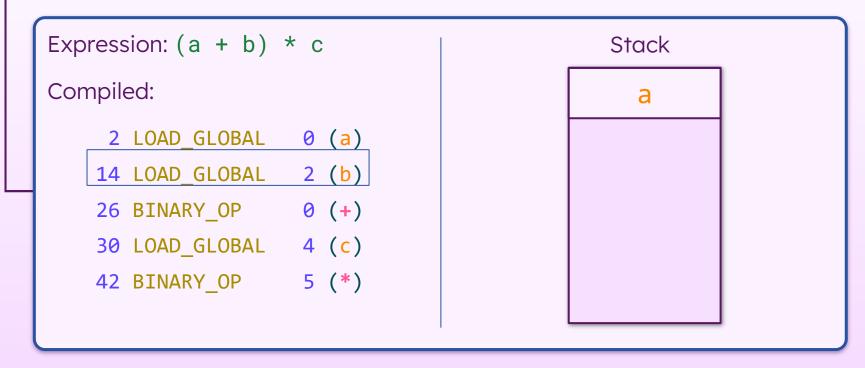
CPython interpreter uses the **stack** to store its intermediate results.

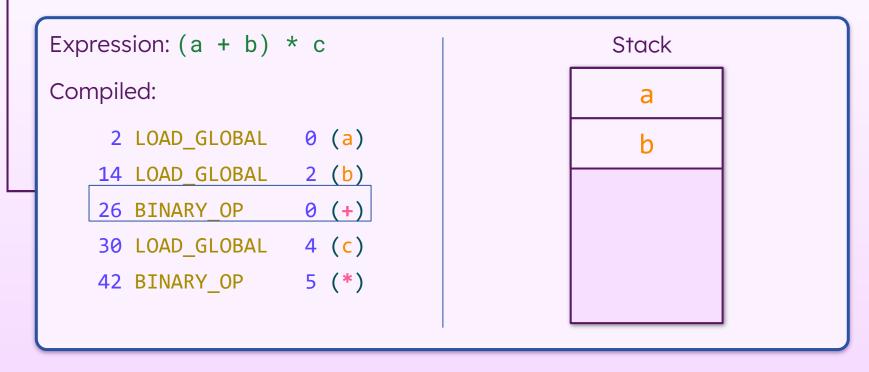
• CPython's Bytecode largely instructs how to manipulate data on the **stack**.

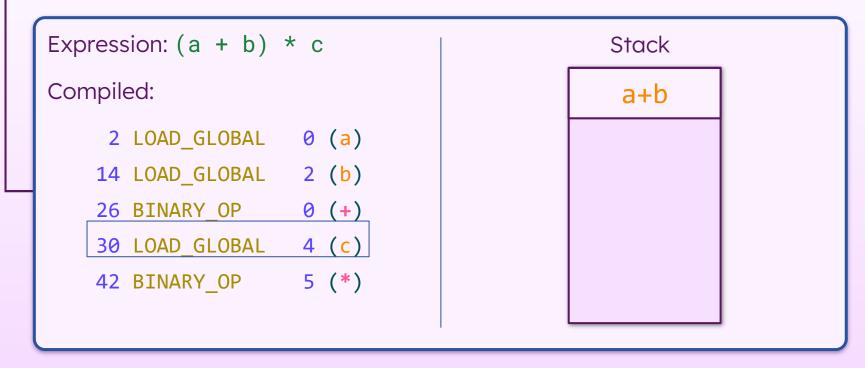
```
Expression: (a + b) * c
```

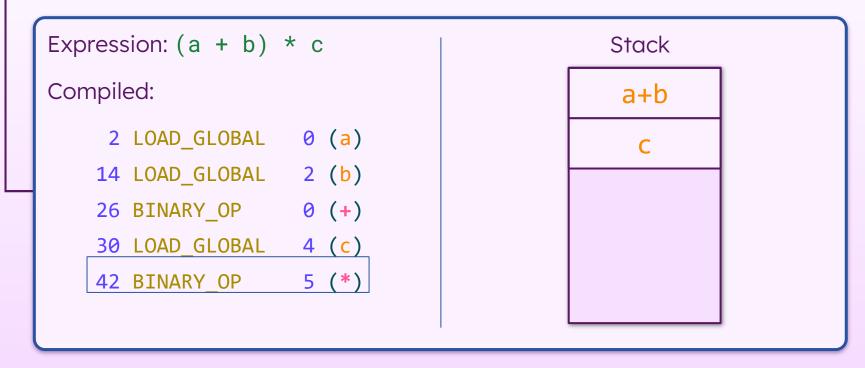
```
Expression: (a + b) * c
Compiled:
    2 LOAD GLOBAL
                   0 (a)
   14 LOAD_GLOBAL 2 (b)
   26 BINARY_OP 0 (+)
   30 LOAD GLOBAL 4 (c)
   42 BINARY OP 5 (*)
```











```
Expression: (a + b) * c
                                            Stack
Compiled:
                                           (a+b)*c
    2 LOAD_GLOBAL 0 (a)
   14 LOAD_GLOBAL 2 (b)
   26 BINARY_OP 0 (+)
   30 LOAD GLOBAL
                   4 (c)
   42 BINARY OP
                    5 (*)
```

CPython: 3.11 Specialising Interpreter

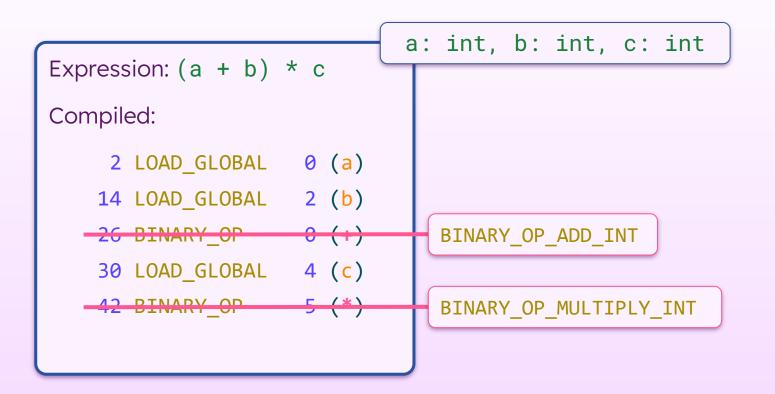
```
Expression: (a + b) * c
Compiled:
     2 LOAD GLOBAL
                     0 (a)
                     2 (b)
    14 LOAD GLOBAL
                     0 (+)
    26 BINARY OP
    30 LOAD GLOBAL
                     4 (c)
    42 BINARY OP
```

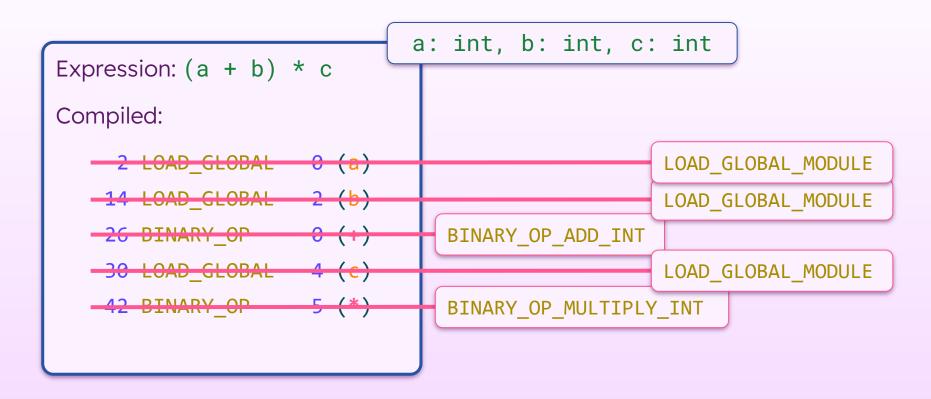
Generic:

- a, b, c can be str, int, float or even objects!
- BINARY_OP has to perform dynamic type dispatch → Slow!

```
Expression: (a + b) * c
Compiled:
    2 LOAD GLOBAL 0 (a)
   14 LOAD GLOBAL 2 (b)
   26 BINARY_OP 0 (+)
   30 LOAD GLOBAL 4 (c)
   42 BINARY_OP 5 (*)
```

```
a: int, b: int, c: int
Expression: (a + b) * c
Compiled:
    2 LOAD GLOBAL 0 (a)
   14 LOAD GLOBAL 2 (b)
   26 BINARY_OP 0 (+)
   30 LOAD GLOBAL 4 (c)
   42 BINARY OP 5 (*)
```







https://youtu.be/shQtrn1v7sQ?si=2BT_V5JiOzwL1wzg

CPython $3.11 \rightarrow 3.13$ and onwards

CPython 3.11:

Specialising Interpreter optimizes across one to two bytecode

CPython 3.13 and onwards:

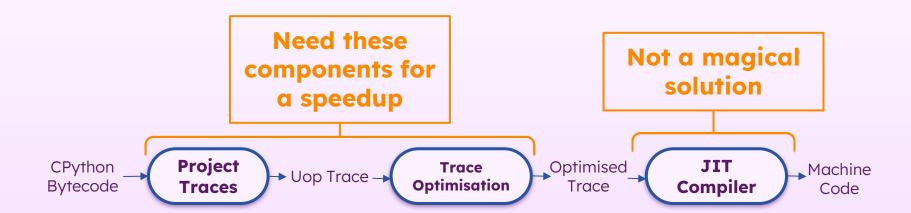
- Learn commonly encountered types at runtime to optimize across larger regions
- Not a new idea, difficulty is implementing correctly and safely and in a maintainable way

CPython 3.13 and onwards

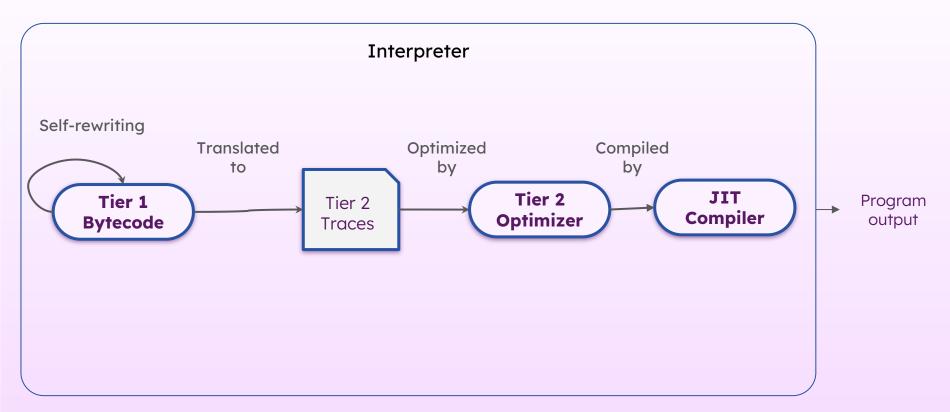




CPython 3.13

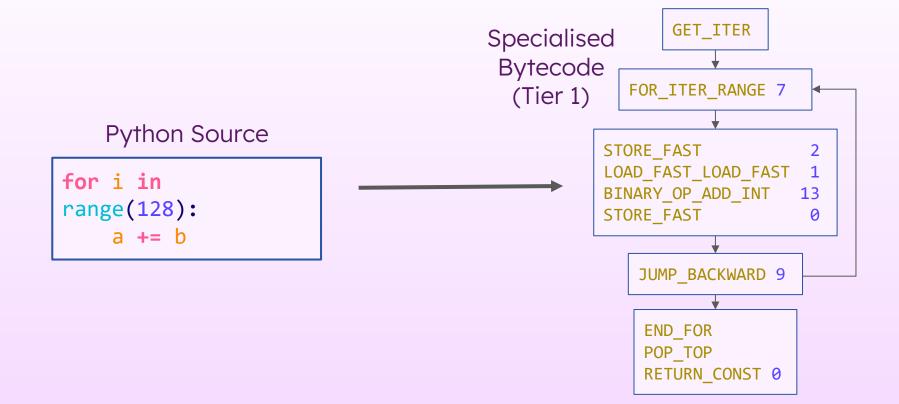


CPython 3.13 Interpretation Pipeline





CPython 3.13 Project Traces

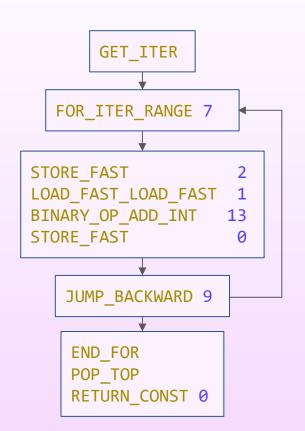




Tracing

```
_START_EXECUTOR
_TIER2_RESUME_CHECK
```

Note: _CHECK_VALIDITY_AND_SET_IP omitted



CPython 3.13 Project Traces

Tracing

_START_EXECUTOR
_TIER2_RESUME_CHECK

Note:

_CHECK_VALIDITY_AND_SET_IP omitted

Trace starts at backwards GET_ITER edge FOR_ITER_RANGE 7 STORE FAST LOAD_FAST_LOAD_FAST BINARY OP ADD INT 13 STORE_FAST JUMP BACKWARD 9 END FOR POP TOP RETURN_CONST 0

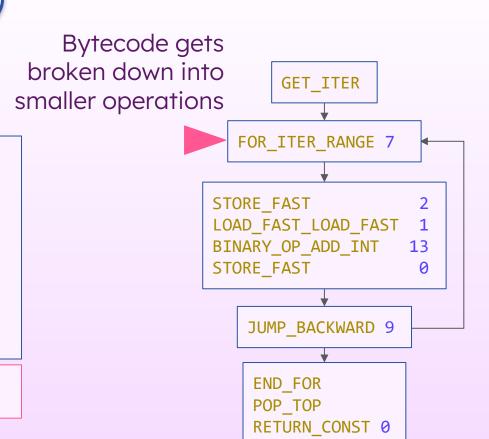
Project Traces

Tracing

_START_EXECUTOR
_TIER2_RESUME_CHECK

Note:

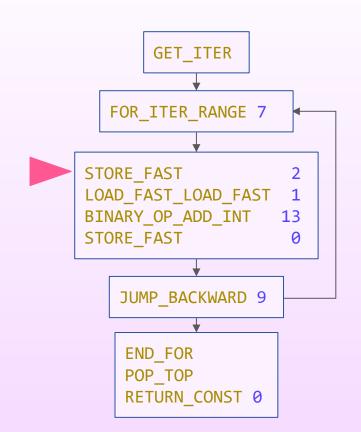
_CHECK_VALIDITY_AND_SET_IP omitted





Tracing

```
_START_EXECUTOR
_TIER2_RESUME_CHECK
_ITER_CHECK_RANGE
_GUARD_NOT_EXHAUSTED_RANGE 7
_ITER_NEXT_RANGE
Note:
_CHECK_VALIDITY_AND_SET_IP omitted
```





Tracing

_CHECK_VALIDITY_AND_SET_IP omitted

GET_ITER FOR_ITER_RANGE 7 STORE FAST LOAD_FAST_LOAD_FAST BINARY_OP_ADD_INT STORE_FAST JUMP BACKWARD 9 END FOR POP TOP RETURN_CONST 0



Tracing

Note:

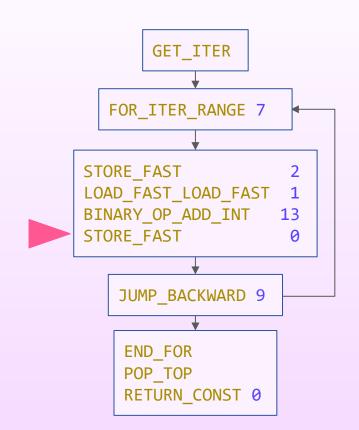
_CHECK_VALIDITY_AND_SET_IP omitted

```
GET_ITER
  FOR_ITER_RANGE 7
STORE FAST
LOAD_FAST_LOAD_FAST
BINARY_OP_ADD_INT
STORE_FAST
   JUMP BACKWARD 9
    END FOR
    POP TOP
    RETURN_CONST 0
```



Tracing

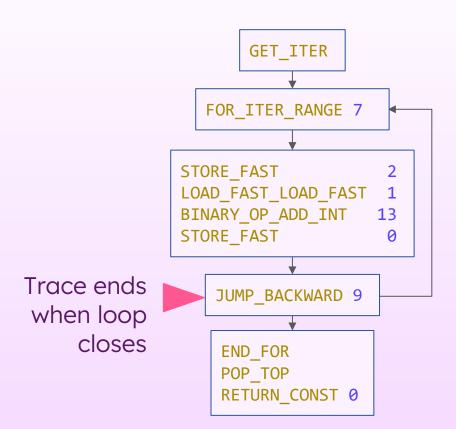
Note: _CHECK_VALIDITY_AND_SET_IP omitted





Tracing

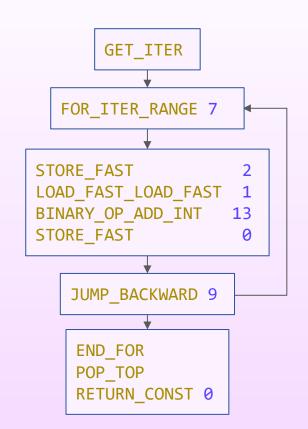
Note: _CHECK_VALIDITY_AND_SET_IP omitted





Tracing

```
Note: 
_CHECK_VALIDITY_AND_SET_IP omitted
```







First pass: By Mark Shannon

Promoting globals to constants (3.13)

Second pass: By Ken Jin, with contributions from Mark Shannon, Guido van Rossum, Peter Lazorchak

- Guard Elimination (3.13 partially implemented)
- True Function Inlining (WIP)
- Deferred Object Creation (WIP)
- Register allocation/TOS caching (WIP)

Analysis via

Abstract

Interpretation (3.13)



Second pass: Abstract Interpretation (3.13)

Normal interpretation: Operate on values

Abstract interpretation: Operate on abstractions of values

In CPython 3.13, the abstraction is (mostly) the type of the value/object

Problem:

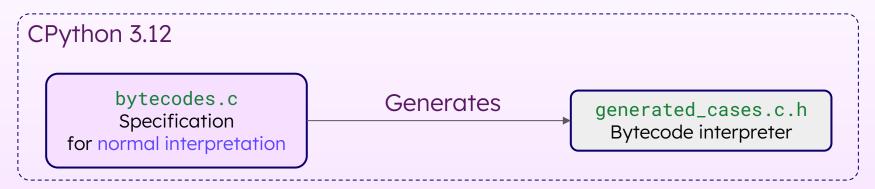
Need to maintain two largely disconnected interpretation specifications

Solution:

CPython 3.12 introduced a DSL to specify the operations of bytecode

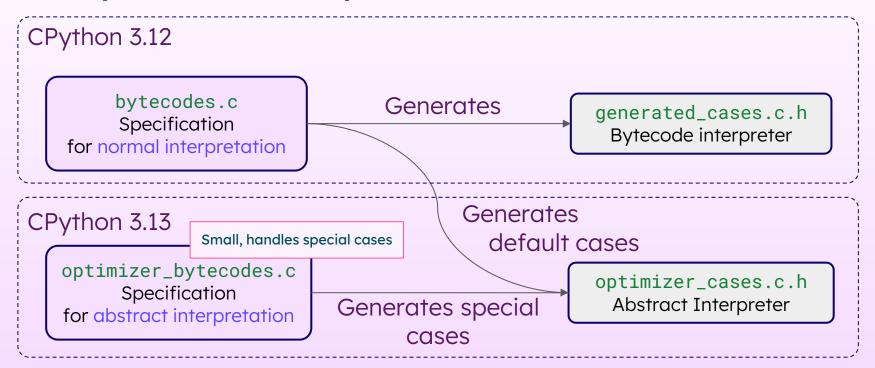


Second pass: Abstract Interpretation



Trace Optimisation

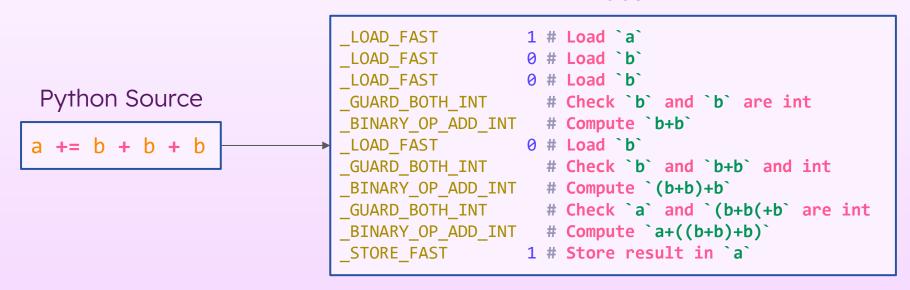
Second pass: Abstract Interpretation





Second pass: Guard Elimination (3.13 partially implemented)

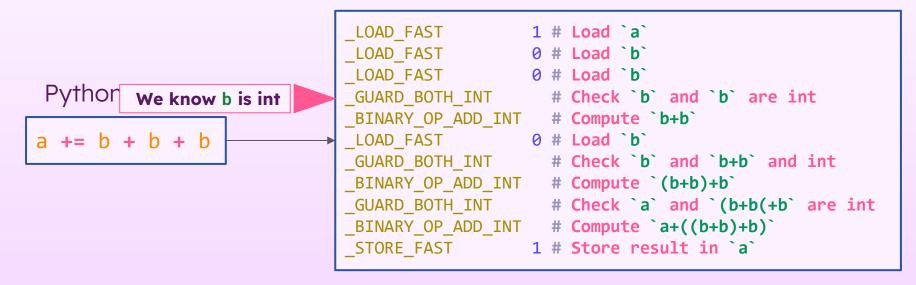
Abstract Interpretation learns the types of each variable





Second pass: Guard Elimination (3.13 partially implemented)

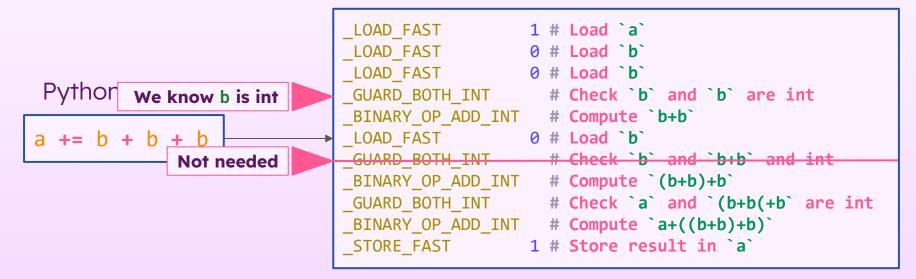
Abstract Interpretation learns the types of each variable





Second pass: Guard Elimination (3.13 partially implemented)

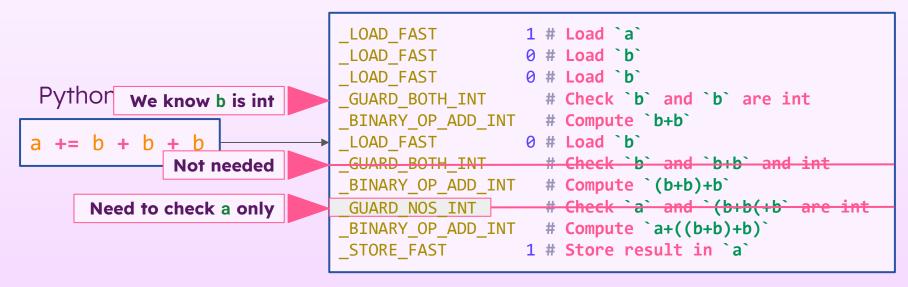
Abstract Interpretation learns the types of each variable





Second pass: Guard Elimination (3.13 partially implemented)

Abstract Interpretation learns the types of each variable





Second pass: True Function Inlining (WIP)

Currently worked on by Ken Jin.

Problem:

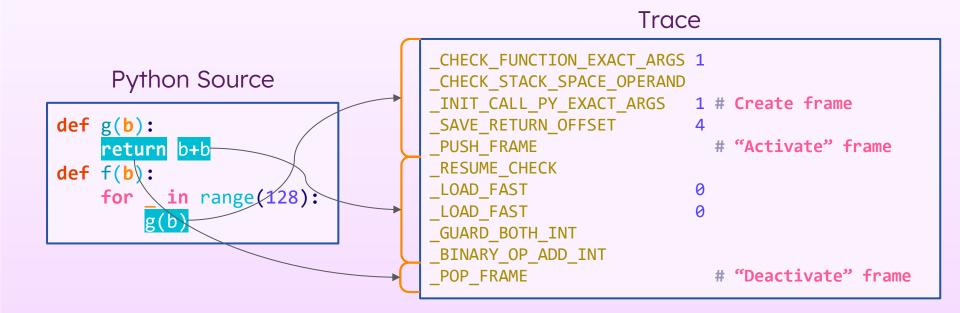
Function calls have some overhead (E.g., Creating a new frame)

Idea:

Inline the function as if it is one big function

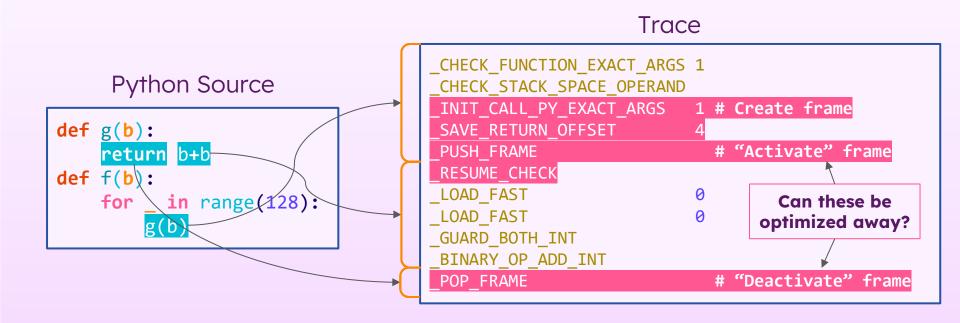
Trace Optimisation

Second pass: True Function Inlining (WIP)





Second pass: True Function Inlining (WIP)





Second pass: Deferred Object Creation/Scalar Replacement (WIP)

Idea from Mark Shannon: Defer or avoid entirely object creation if possible

E.g.,

- [0,1,2,3][3] returns 3 without creating a list
- filter(lambda x: x%2, [1,2,3,4,5]) returns filter without creating
 an intermediate list literal



Second pass: Register allocator/top of stack caching (WIP)

Currently worked on by Brandt Bucher.

Based on:

Ertl, M. A. (1995b). Stack caching for interpreters. SIGPLAN Not., 30(6), 315–327. doi:10.1145/223428.207165

Idea: Cache the top few items on the stack in registers

- Memory access is slow, register access is fast





Talks: Building a JIT compiler for CPython

Sunday - May 19th, 2024 1 p.m.-1:30 p.m. in Ballroom A

Presented by:

***** Brandt Bucher

Description

CPython is a programming language implementation that is mostly maintained by volunteers, but has a huge, diverse user base spread across a wide variety of platforms. These factors present a difficult set of challenges and tradeoffs when making design decisions, especially those related to just-in-time machine code generation.

As one of the engineers working on Microsoft's ambitious "Faster CPython" project, I'll introduce our prototype of "copy-and-patch", an interesting technique for generating high-quality template JIT compilers. Along the way, I'll also cover some of the important work in recent CPython releases that this approach builds upon, and how copy-and-patch promises to be an incredibly attractive tool for pushing Python's performance forward in a scalable, maintainable way.

Side Exits + De-Opts:

When

- The assumptions a trace made is invalid, or
 - (e.g., invalid cache/different runtime type encountered)
- Control-flow exits the trace,

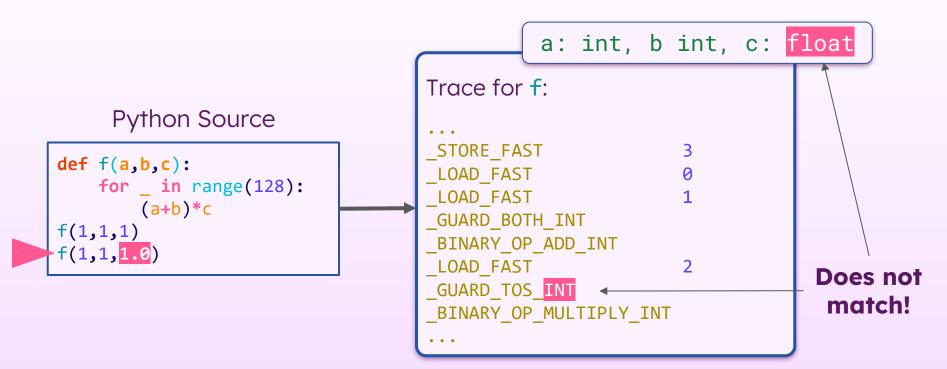
CPython performs one of two exits:

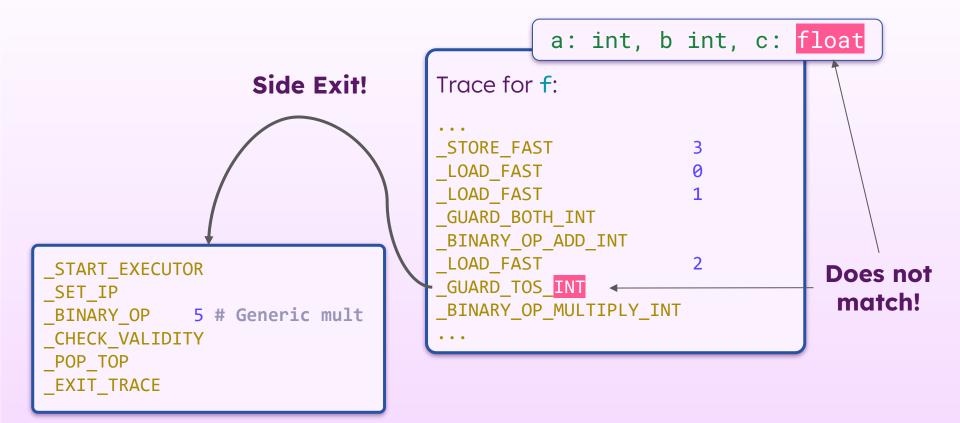
- 1. Side Exits
- 2. De-Opts (De-Optimisation)

Side Exits:

If current progress of execution of the trace is still valid, a side exit either:

- 1. Jumps back to Tier 1
- 2. Creates a new trace corresponding to the side exit (if the side exit is taken enough times)
- 3. Jumps to an existing trace (if 2 has already happened)





De-Opts:

If continued execution of the trace is no longer valid (rare)

• Drop back to Tier 1

Metadata: Open Source as a University Student





How did we come to all this?

Before University:

- Ken Jin started on CPython before university.

University Sem 1:

- We took Programming Language Implementation under Martin Henz.
- Martin Henz supervised our first Experiment #1

Year 1 Summer:

Continued Experiment #1 as a credit bearing software engineering project

Subsequent Work:

Manuel Rigger became our advisor for the next year

Lesson 1: Make full use of your university's mentorship as a student.





How did we come to all this?

Classes we took:

- Programming Language Implementation
- Compiler Design
- Principles of Program Analysis (Ken Jin)

Lesson 2: Take classes that help your journey





How did we come to all this?

 Jules and Ken Jin wrote 2 versions of the optimizer, Ken Jin then rewrote that 2 more times before it went into CPython.

Lesson 3: Be patient with yourself





That's a wrap!

- Ken Jin will try to graduate.
- Jules has no plans.

Thank You!

@Fidget-Spinner | kenjin@python.org @JuliaPoo | juliapoopoopoo@gmail.com



