

## course\_content

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### **S2 - Complexity Analysis**

W2: In-built Python

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W9: Graphs

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# Data Structures and Algorithms

## Complexity Analysis

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# Week 1

# Counting Operations

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- Every algorithm takes a certain amount of time
- A proxy for the amount of time is the number of 'elementary operations'
- An elementary operation is a computer instruction executed in a single time unit.
  - `+, -, x, /, %, >, <, |, &, ==, =`
- The running time of an algorithm is a linear function of the number of operations

# Example 1

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**Algorithm** contains(a, x):

```
for i = 0 to len(a)-1
```

```
    if x == a[i]
```

```
        return true
```

```
return false
```

- Worst-case
- Best-case
- Average-case

Generally, we look at worst-case  
when analysing algorithms.

# Example 2

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**Algorithm** max(a):

```
max = a[0]
```

```
for i = 1 to len(a)-1
```

```
    if a[i] > max
```

```
        max = a[i]
```

```
return max
```

What is the runtime?

How much space is used?

# Asymptotic Analysis

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- Asymptotic - what is the 'limiting behaviour'. What happens as  $n \rightarrow \infty$ ?
- We wish to classify algorithms into smaller groups rather than computing the exact number of operations each time.
- There are 3 important notations here - **Big-O**, **Big-Omega**, **Big-Theta**.
- **Big-O** (upper bound for the algorithm) - it will take at most this time
- **Big- $\Omega$**  (lower bound for the algorithm) - it will take at least this time
- **Big- $\theta$**  (lower and upper bound for the algorithm) - it will take roughly this time

# Big-O Definition

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- Let's call the runtime function  $T(n)$
- We can say that the runtime  $T(n)$  is  $O(f(n))$  for a function  $f(n)$  (e.g.  $n$ ,  $\log n$ ,  $n^2$ ,  $2^n$ ,  $n!$ ) if:
  - In **maths**:  $T(n) \leq Mf(n)$  for all  $x > c$  for some constants  $M$  and  $c$
  - In **words**:  $T(n)$  is eventually bounded above by a multiple of the function  $f(n)$ .
- For example  $T(n) = 5n^2 + 2n$  is  $O(n^2)$  since  $5n^2 + 2n \leq 7n^2$  for all  $n > 0$
- Or  $T(n) = 18n \log n + 100n$  is  $O(n \log n)$  since  $18n \log n + 100n \leq 200n \log n$  for all  $n > 1$

**Intuition** What is the dominant term in the expression? Which term grows the fastest?

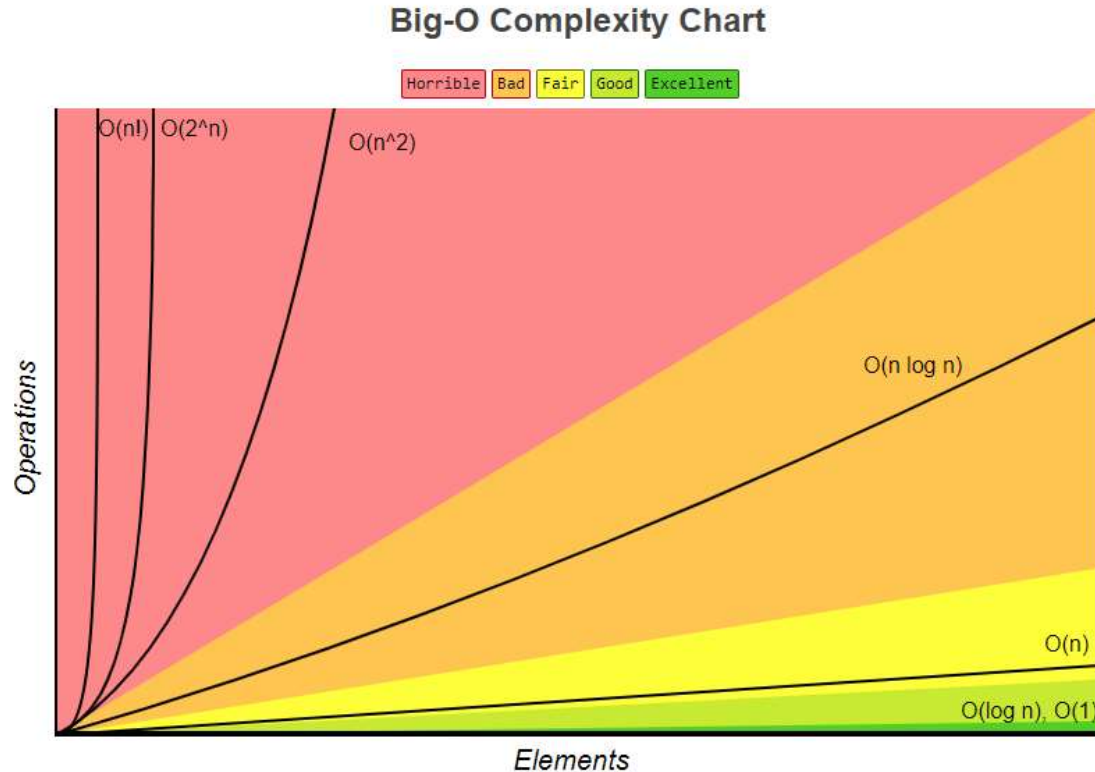
# Big-O Categories

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- There are various **Big-O** categories but the most used ones are
  - $O(1)$  - constant time
  - $O(\log n)$  - log time
  - $O(n)$  - linear time
  - $O(n \log n)$  - loglinear time
  - $O(n^2)$  - quadratic time
  - $O(n^3)$  - cubic time
  - $O(2^n)$  - exponential time
  - $O(n!)$  - factorial time

# Big-O Complexity Chart

Time:  
5 min



Source: [Medium](#)



# Big O Notation - Exercise

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- Get into groups of 3.
- Each group should have only one laptop open.
- Follow the link below to access the instructions and prompt for the exercise.
- Make sure to discuss the asymptotic runtime and space complexity of the algorithms provided.

[bit.ly/DSAW1BigO](https://bit.ly/DSAW1BigO)

# Big O Notation - Exercise Solutions

index	Time Complexity	Space Complexity
Alg1	$O(n^2)$	$O(n)$
Alg2	$O(n)$	$O(n)$
Alg3	$O(n^2)$	$O(n)$
Alg4	$O(\log(n))$	$O(n)$
Alg5	$O(ab)$	$O(a + b)$

Look up **Auxiliary Space** and see how that differs from space complexity

Time:  
7 min

# Questions

# Next Steps

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1. Commit implementations for the week
2. Readings for Week 1 - Make sure you understand time and space complexity
3. Review Python Basic Data Structures -
4. Brainstorm Project Topics
5. Post questions on Piazza