

course_content



Data Structures and Algorithms

Graph Algorithms

Week 10

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W10: Graphs

W11: Greedy Algorithms

Graph Algorithms - Activity

- Once we have a graph, we want to do things with graphs.
- Try to think about/imagine how you would do the following graph algorithms.
 - Find the node with the largest value at the node
 - Find if there is a node with a specific value
 - Given two nodes, find a path between them
- This should be hard. If it's too much, try it with a tree first.
- What is the issue when we introduce the possibility of 'cycles'?
- In general searching is very important in graphs!

Breadth First Search

- Graph G and starting vertex s
- Finds all vertices 1 away from s , 2 away from s , 3 away from s
- Puts all vertices 1 away from s into a queue e.g. a, b, c, d, e
- Then runs through queue again finding vertices 1 away from each of a, b, c, d, e putting those into the queue and then repeats
- If we reach a vertex previously visited, don't re-explore (e.g. check stored set)
- Through this, we can store the max, look for a specific value, find path to vertex t

Look at it [here](#), visualise it [here](#), read about it [here](#)

Breadth First Search - Applications

- Web Crawlers (in Search Engines) - to index popularity of websites, crawler needs to search through all possible paths
- [Torrenting](#) (P2P sharing) - Checks peers for files already downloaded
- Basic shortest path algorithms (e.g. Google Maps, Flights)
- Social Networks (E.g. In LinkedIn, determine if connection is 1st/2nd/3rd+ degree)
- Spell check - looks for similar words with slight changes (related to [this](#))
- Read more [here](#)

Breadth First Search - Complexity

- 2 different inputs here - V is the number of vertices, E is the number of edges.
- **Adjacency List** - $O(V+E)$ as we have to visit every vertex exactly once and for each vertex we have to get all of its neighbours each of which is through an edge.
- **Adjacency Matrix** - $O(V^2)$ since for each vertex, we have to go through the array/list for that vertex which is length V .
- **Edge List** - An edge list would be very poor here as you wouldn't have a nice way of getting which edges are adjacent to any given vertex.

E can be $O(V^2)$ if there are many edges (since each edge connects 2 vertices so at most “ n choose 2” edges but for most graphs E is much smaller than V^2).

Depth First Search

- Graph G and starting vertex s
- Picks a neighbour of s and adds it to a stack. E.g. a
- Picks an unvisited neighbour of a and adds it to a stack. E.g. b
- Visited neighbours are checked through a set (hash map) - $O(1)$ for each check
- Eventually we reach a vertex where all neighbours have already been visited.
- Remove this from the stack and again check for an unvisited neighbour of the top
- Keep going until the entire stack is empty (it will grow again)

Look at it [here](#), visualise it [here](#), read about it [here](#)

Depth First Search - Applications

- Testing if a network is 'connected'. (E.g. will message reach everyone in network)
- Shortest Path Algorithm Alternative (Flights, Google maps etc) but this generally performs poorly compared to other algorithms
- Detecting cycles (e.g. in a schedule, a cycle would make it impossible to complete)
- [Topological Sorting](#) - Creates an ordering of vertices (in a directed graph) where vertices that come before lead to vertices that come after. [More here](#).
- [A variety of interview questions](#)

Depth First Search - Complexity

- **Adjacency List** - $O(V+E)$ as we again have to visit every vertex exactly once and for each vertex we have to access/check each of its neighbours until it is complete
- **Adjacency Matrix** - $O(V^2)$ since again for each vertex, we have to go through the array/list for that vertex which is length V to get the adjacent vertices.
- **Edge List** - An edge list would be very poor here as you wouldn't have a nice way of getting which edges are adjacent to any given vertex.

Can't do parallel computing as we shouldn't take next item in stack without finish top

Graph Algorithms

- There are many other graph algorithms. You can explore some of them:
 - [Dijkstra's Algorithm](#) is the standard 'shortest path' algorithm that uses a priority queue (built using a heap)
 - [Prim's Spanning Tree Algorithm](#) finds the tree from a central root to every other vertex that minimises the time taken to 'broadcast a message out'
 - [Ford-Fulkerson Algorithm](#) for sending the maximal 'flow' from one node through to another restricted by the weights on the edges.
- A lot of graph algorithms are incredibly useful for computer science. They are explored more in the 4th-year elective 'Graph Theory'.

Time:
1 min

Questions?....Office Hours!