



# INTRODUCTION TO DATABASE SYSTEMS

## IS211

Dr. Noha Nagy

Lecture 2

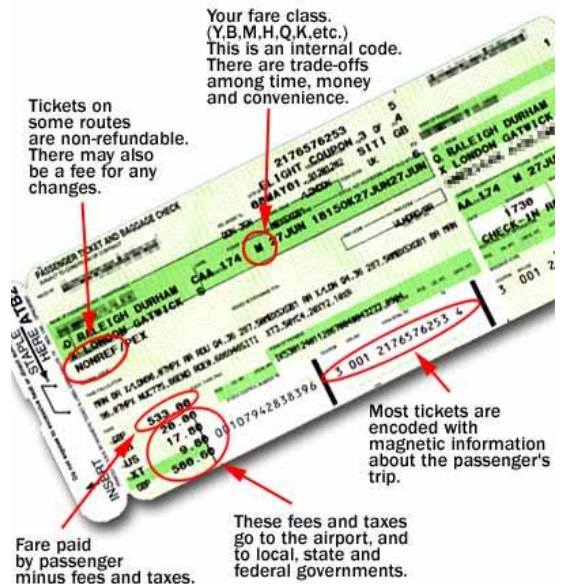
Relational Model

# Databases Everywhere

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amazon.com



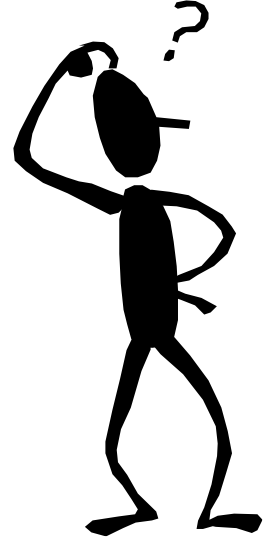
Source: Delta Airlines

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# Definitions

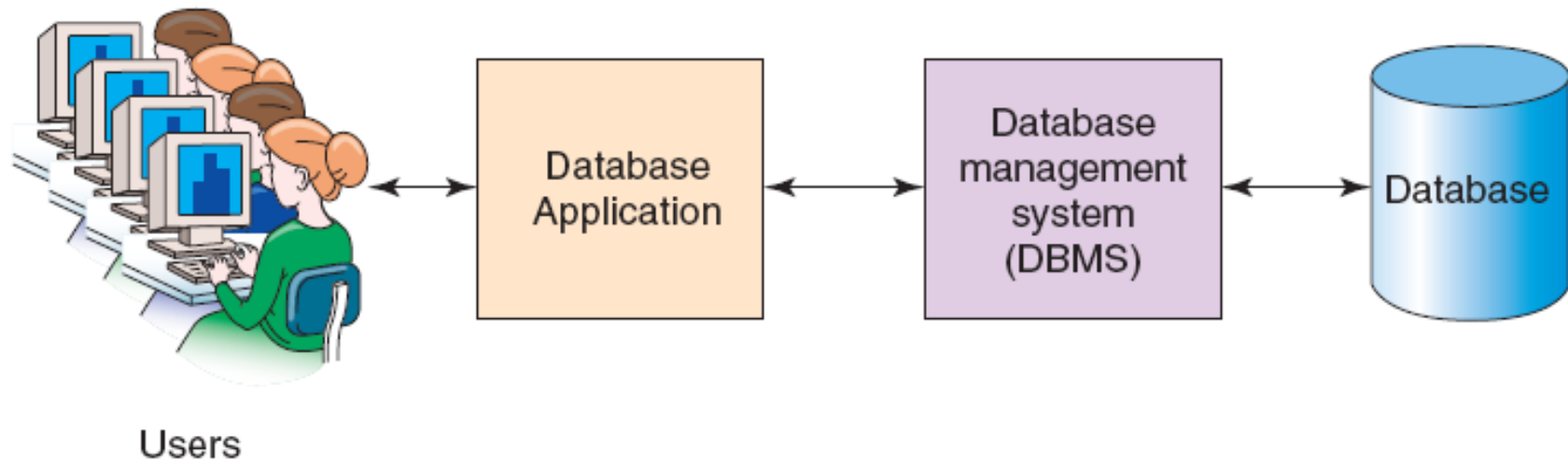
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- **Database:**
  - ▣ Collection of related data.
  - ▣ A database contains a model of something!
- **A Database Management System (DBMS):** is a software system designed to store, manage and facilitate access to the database
- **Database/Application Program:** A computer program that interacts with the database through the DBMS



# Components of a Database System

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# Example

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**Welcome to University System**

Register

Close

## Student Registration

ID

Name

Mobile

Address

City

DOB

Submit

Reset

**Application**

## Student Registration

ID 20110021  
Name Mohsan  
Mobile 01110054487  
Address 55 Mohandseen  
City Giza  
DOB 22-2-2002

Submit

Reset

**DBMS**

**DB**

ID	Name	Mobile	Address	City	DOB
20110021	Mohsan	01110054487	55 Mohandseen	Giza	22-2-2002

## Student Registration

ID

20110032

Name

Ahmed

Mobile

01110056577

Address

5 s1 October

City

Giza

DOB

2-9-2002

Submit

Reset

**DBMS**

ID	Name	Mobile	Address	City	DOB
20110021	Mohsan	01110054487	55 Mohandseen	Giza	22-2-2002
20110032	Ahmed	01110056577	5 s1 October	Giza	2-9-2002



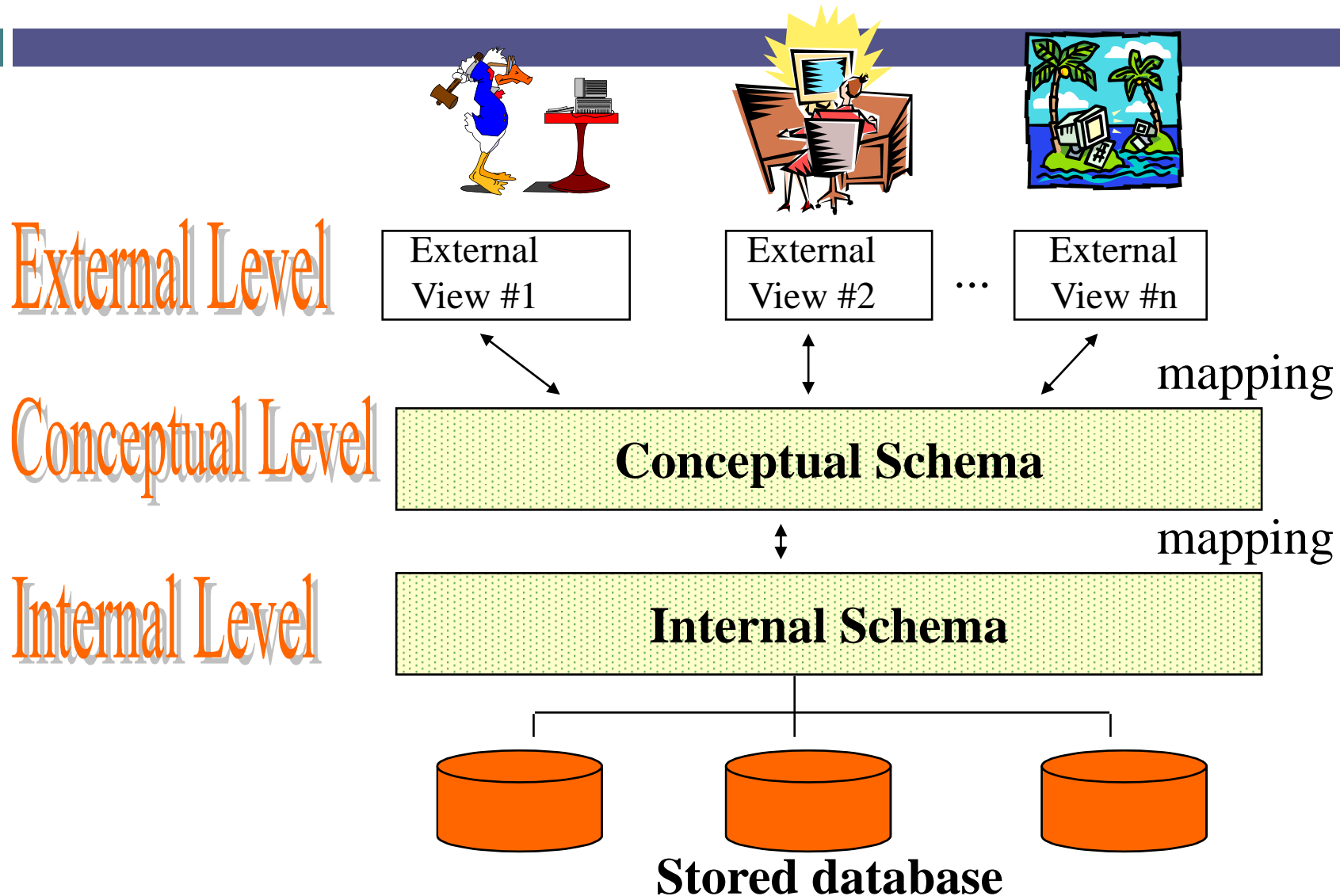
# Data Independence

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- Does data have to be part of programs?
- Do we need to change one if the other changes?
- **Three-tier architecture** of databases

# Three-Schema Architecture

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# Levels of Abstraction

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## □ Views

define how users see the data

(use of data)

## □ Conceptual schema

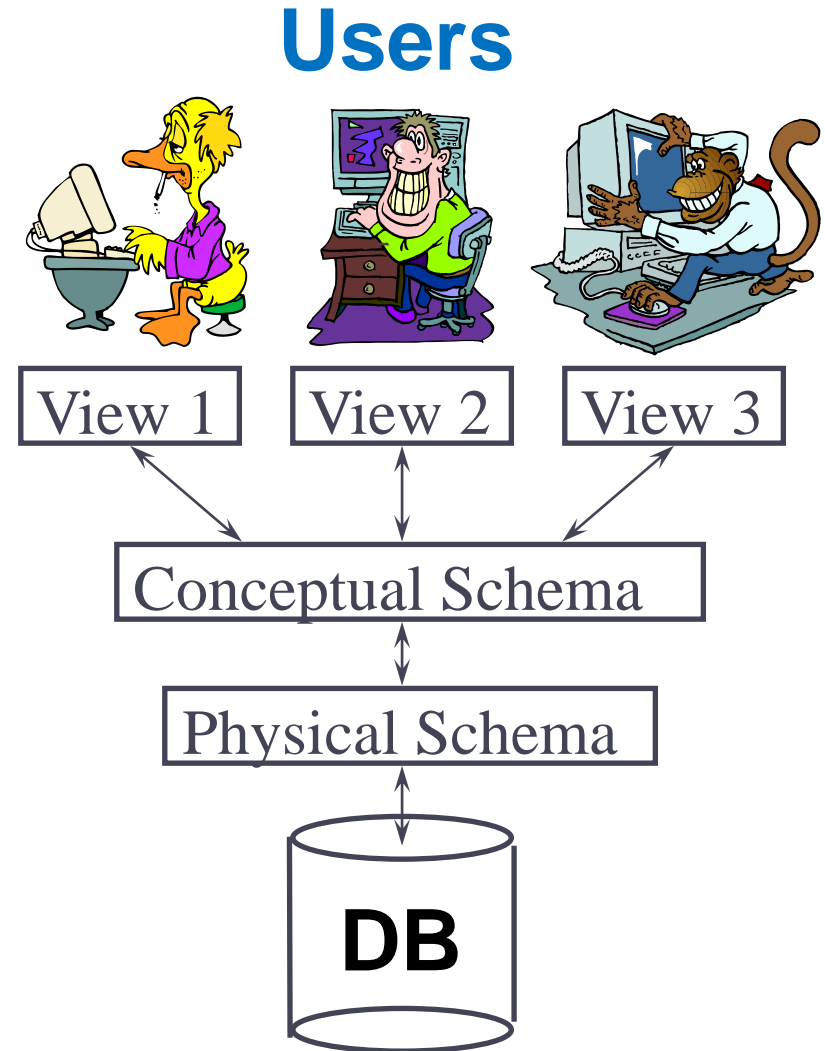
defines logical structure

(meaning of data)

## □ Physical schema

describes the files and indexes used.

(storage of data)



# Three-Schema Architecture

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- **Defines DBMS schemas at *three* levels:**
  - ▣ **Internal schema** is used to describe physical storage structures and access paths (e.g indexes).
    - Typically uses a physical data model.
  - ▣ **Conceptual schema** is used to describe the structure and constraints for the whole database for a community of users.
  - ▣ **External schemas** is used to describe the various user views.

# Data Models

# A Map is a Model of Reality

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# Data Models

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- A collection of tools for describing
  - ▣ data
  - ▣ data relationships
  - ▣ data constraints
- Models:
  - ▣ Relational Model
  - ▣ object-oriented model
  - ▣ semi-structured data models
  - ▣ Older models: network model and hierarchical model

# History of Relational Model

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- Introduced by **Ted Codd** in 1970 in a classic paper
- Ted Codd was an **IBM** Researcher
- Many database concepts & products based on this model





# Relational Model

# Relations

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- A relational database is a set of **relations**
- Relations are basically **tables** of data
- Each row represents a **record** in the relation
- Each relation has a **unique name** in the database
- Each row in the table specifies a relationship between the values in that row

acct_id	branch_name	balance
A-301	New York	350
A-307	Seattle	275
A-318	Los Angeles	550
...	...	...

The *account* relation

- Example :

The account ID “A-307”, branch name “Seattle”, and balance “275” are all related to each other

# Relation in RDB

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**Student**

**Attribute**

**Table Heading**

SID	Name	City	Phone
012520	Rana	Cairo	01225928765
875466	Nawal	Giza	01228796457
897555	Omar	Cairo	01110097564
987458	Ali	Cairo	01110968874

**Relation Degree**

**Relation Cardinality**

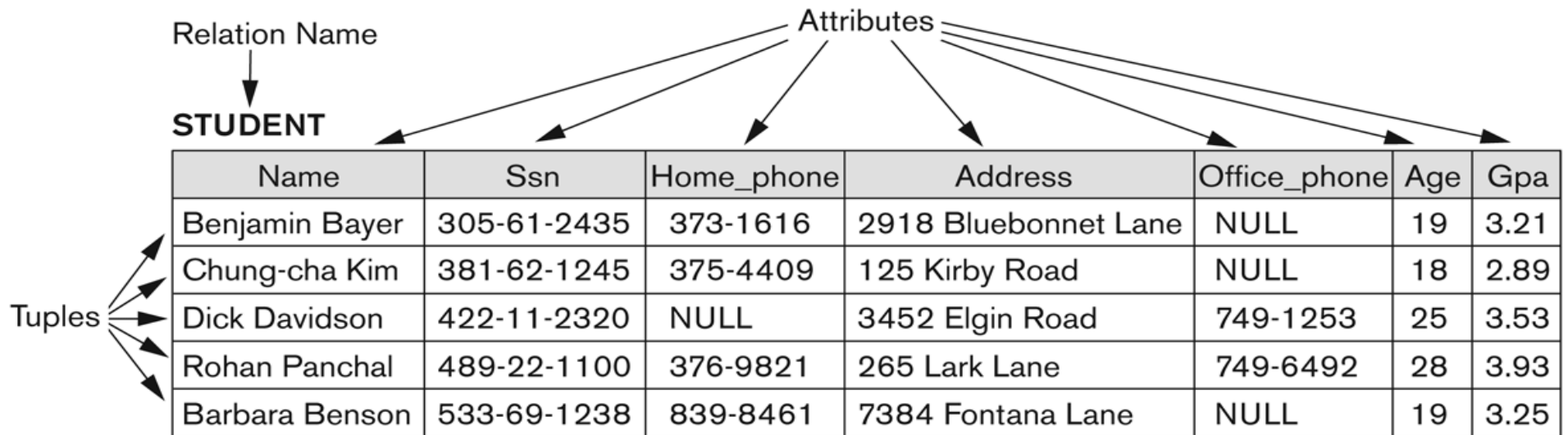
# Definition Summary

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<u><b>Informal Terms</b></u>		<u><b>Formal Terms</b></u>
<b>Table</b>		<b>Relation</b>
<b>Column Header</b>		<b>Attribute</b>
<b>All possible Column Values</b>		<b>Domain</b>
<b>Row</b>		<b>Tuple</b>
<b>Table Definition</b>		<b>Schema of a Relation</b>
<b>Populated Table</b>		<b>State of the Relation</b>

# Relation Example

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The attributes and tuples of a relation STUDENT.

**Degree =7**  
**Cardinality=5**

# Relations and Attributes

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- Each relation has some number of attributes Sometimes called “**columns**”
- Each attribute has a **domain** Specifies the set of valid values for the attribute

acct_id	branch_name	balance
A-301	New York	350
A-307	Seattle	275
A-318	Los Angeles	550
...	...	...

The *account* relation

- The *account* relation:
  - 3 attributes
  - Domain of *balance* is the set of nonnegative integers
  - Domain of *branch\_name* is the set of all valid branch names in the bank

# DOMAINS

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- A domain **D** is a set (pool) of values, from which one or more attributes takes their values.

- Example

CITY = {London, Paris, Doha, Cairo, Athens, Rome, Dobai, Madrid}

CITY is a pool of cities from which The attributes Supplier.City, Customer.City take their own values.

DATE = (DAY, MONTH, YEAR)

Where:

DAY = {1..31}, MONTH = {1..12}, YEAR = {1990..2100}

CITY is a simple domain, but DATE is a **composite** Domain

# Tuples and Attributes

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- Each row is called a tuple
  - ▣ A fixed-size, ordered set of name-value pairs
- Each attribute in the tuple has a **unique name**

acct_id	branch_name	balance
A-301	New York	350
A-307	Seattle	275
A-318	Los Angeles	550
...	...	...

The *account* relation



# Tuples and Relations

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- A **relation** is a set of tuples
- Each tuple appears **exactly once**
- The **order of tuples** in a relation is not relevant

Instructor

<i>ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>
22222	Einstein	Physics	95000
12121	Wu	Finance	90000
32343	El Said	History	60000
45565	Katz	Comp. Sci.	75000
98345	Kim	Elec. Eng.	80000
76766	Crick	Biology	72000
10101	Srinivasan	Comp. Sci.	65000
58583	Califieri	History	62000
83821	Brandt	Comp. Sci.	92000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
76543	Singh	Finance	80000

# Schema VS Instance

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- the name of the relation and the set of attributes is called the **schema**
- the current values contained in the relation represent an **instance**

# Relation Schemas

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- Every relation has a schema
- A relation schema includes:
  - ▣ an ordered set of **attributes**
  - ▣ the **domain** of each attribute

acct_id	branch_name	balance
A-301	New York	350
A-307	Seattle	275
A-318	Los Angeles	550
...	...	...

The *account* relation

- The relation schema of *account* is:  
 $Account\_schema = (acct\_id, branch\_name, balance)$
- Database schema is a collection of relation schemas

# COMPANY Database Schema

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## EMPLOYEE

Fname	Minit	Lname	<u>Ssn</u>	Bdate	Address	Sex	Salary	Super_ssn	Dno
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## DEPARTMENT

Dname	<u>Dnumber</u>	Mgr_ssn	Mgr_start_date
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## DEPT\_LOCATIONS

<u>Dnumber</u>	<u>Dlocation</u>
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## PROJECT

Pname	<u>Pnumber</u>	Plocation	Dnum
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## WORKS\_ON

<u>Essn</u>	<u>Pno</u>	Hours
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## DEPENDENT

<u>Essn</u>	<u>Dependent_name</u>	Sex	Bdate	Relationship
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Schema diagram for  
the COMPANY  
relational database  
schema.

