


Computer Arch. Exams



Midterm 2023



Faculty of Computers and Artificial Intelligence
Cairo University



Midterm Exam

Program: Computer Science / Software Engineering
Course Name: Computer Organization and Architecture
Course Code: CS331 / SCS322
Instructor(s): Dr. Amin Allam

Date: 19/11/2022
Duration: 1 hour
Total Marks: 20 marks

تعليمات هامة:

- حيازة التليفون المحمول مفتوحا داخل لجنة الامتحان يعتبر حالة غش تستوجب العقاب وإذا كان ضروري الدخول بالمحمول فيوضع مغلقا في الحقيبة.
- لا يسمح بدخول سماعة الأذن أو تابلت أو توتو.
- لا يسمح بدخول أي كتب أو ملزم أو أوراق داخل اللجنة والمخالفة تعتبر حالة غش.

- Exam consists of 14 equal-weight multiple-choice questions in 2 pages.
- Record in the bubble sheet exactly ONE answer for each question.

Qa

1 The outputs of 16 registers are connected to the inputs of a bus. Each register stores 32 bits. The bus selects the outputs of one of the input registers according to the selection inputs of the bus. The bus must have the following number of multiplexers:

- [A] 3 [B] 4 [C] 5 [D] 16 [E] 32

2 A memory unit contains 2048 words. Each word is 32 bits. The number of input address lines to this memory unit is:

- [A] 5 [B] 10 [C] 11 [D] 16 [E] 32

3 After applying arithmetic shift right to a 6-bit register containing the binary number 101100 it becomes:

- [A] 111100 [B] 010110 [C] 10110 [D] 011001 [E] 011000

4 Consider the ADD instruction which increases the value of the accumulator register by the value located at the effective address. If the instruction uses direct addressing and it contains the address 10 in its address field. The memory word at address 10 contains the value 30, and the memory word at address 30 contains the value 45. After the instruction is executed, the accumulator register increases by:

- [A] 10 [B] 5 [C] 40 [D] 45 [E] 75

5 After addition of unsigned integers, an overflow can be detected by checking:

- [A] last carry [B] last two carries [C] memory buffers [D] last register bit [E] last two register bits

6 The instruction that is being skipped sometimes by its previous instruction is usually:

- [A] Clear [B] Store [C] Branch [D] Halt [E] Increment

7 Assume that 2 clock cycles are required on average to execute any instruction. The time required to execute a program containing 10,000 instructions on a computer with a clock rate of 5,000,000 clock cycles per second is:

- [A] 0.001 seconds [B] 0.002 seconds [C] 0.004 seconds [D] 500 seconds [E] 1000 seconds

$$2 \times 10,000$$

$$5,000,000$$

Qb

- \Rightarrow Let $A = A_2A_1A_0$, $B = B_2B_1B_0$ be input binary numbers in signed two's complement representation.
 \Rightarrow Let $G = G_2G_1G_0$ be an output binary number in signed two's complement representation.
 \Rightarrow Let 2×1 multiplexer M_i selects input M_{i0} when $S = 0$, and selects input M_{i1} when $S = 1$.
 \Rightarrow Let a connection $[a, b]$ be a wire connecting a to b .

For questions 8 to 14, consider electronic circuits with inputs S , A , B and output G such that they contain the **minimum** subset of the following components:

- One or more 2×1 multiplexer M_i with selection input S and inputs M_{i0} , M_{i1} and output M_{if} .
- One or more *Not* gate N_i with input N_{if} and output N_{ig} .
- One or more full adder F_i with inputs F_{i0} , F_{i1} , F_{i2} and outputs F_{is} (sum), F_{ic} (carry).

8 A circuit that outputs $G = A$ or B when $S = 0$ or 1 respectively has the connection:
 [A] $[A_0, F_{00}][N_{ig}, F_{01}]$ [B] $[B_0, M_{01}][M_{0g}, N_{if}]$ [C] $[A_0, M_{00}][M_{0g}, G_0]$ [D] $[A_0, M_{00}][B_0, M_{01}]$
 [E] $[A_0, F_{00}][N_{ig}, M_{01}]$ ☒

9 A circuit that outputs $G = -A$ or $-B$ when $S = 0$ or 1 respectively has the connections:
 [A] $[A_0, F_{00}][N_{ig}, F_{01}]$ [B] $[B_0, M_{01}][M_{0g}, N_{if}]$ [C] $[A_0, M_{00}][M_{0g}, G_0]$ [D] $[A_0, M_{00}][B_0, M_{01}]$
 [E] $[A_0, F_{00}][N_{ig}, M_{01}]$ ☒

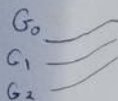
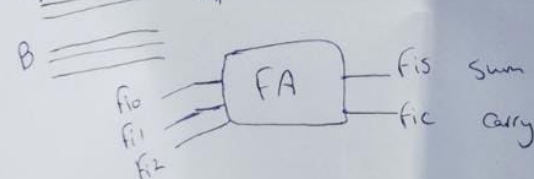
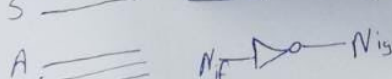
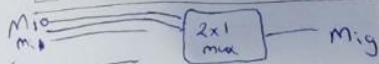
10 A circuit that outputs $G = A + B$ or $A - B$ when $S = 0$ or 1 respectively has the connections:
 [A] $[A_0, F_{00}][N_{ig}, F_{01}]$ [B] $[B_0, M_{01}][M_{0g}, N_{if}]$ [C] $[A_0, M_{00}][M_{0g}, G_0]$ [D] $[A_0, M_{00}][B_0, M_{01}]$
 [E] $[A_0, F_{00}][N_{ig}, M_{01}]$ ☒

11 A circuit that outputs $G = A - B$ or $A - B - 1$ when $S = 0$ or 1 respectively has the connections:
 [A] $[A_0, F_{00}][N_{ig}, F_{01}]$ [B] $[B_0, M_{01}][M_{0g}, N_{if}]$ [C] $[A_0, M_{00}][M_{0g}, G_0]$ [D] $[A_0, M_{00}][B_0, M_{01}]$
 [E] $[A_0, F_{00}][N_{ig}, M_{01}]$ ☒

12 A circuit that outputs $G = A - B$ or $B - A$ when $S = 0$ or 1 respectively has the connections:
 [A] $[A_0, F_{00}][N_{ig}, F_{01}]$ [B] $[A_0, M_{00}][B_0, M_{01}]$ [C] $[A_0, M_{00}][M_{0g}, G_0]$ [D] $[A_0, M_{00}][B_0, M_{01}]$
 [E] $[A_0, F_{00}][N_{ig}, M_{01}]$ ☒

13 A circuit that outputs $G = A$ or $-A$ when $S = 0$ or 1 respectively has the connections:
 [A] $[A_0, F_{00}][N_{ig}, F_{01}]$ [B] $[A_0, M_{01}][M_{0g}, N_{if}]$ [C] $[A_0, M_{00}][N_{ig}, M_{01}]$ [D] $[A_0, M_{00}][M_{0g}, G_0]$
 [E] $[A_0, F_{00}][N_{ig}, M_{01}]$ ☒

14 A circuit that outputs $G = A + 1$ or $A - 1$ when $S = 0$ or 1 respectively has the connections:
 [A] $[A_0, F_{00}][N_{ig}, F_{01}]$ [B] $[A_0, M_{01}][M_{0g}, N_{if}]$ [C] $[A_0, M_{00}][N_{ig}, M_{01}]$ [D] $[A_0, M_{00}][M_{0g}, G_0]$
 [E] $[A_0, F_{00}][N_{ig}, M_{01}]$ ☒



- Shift left : behov zero 31 shenak
wa a shift left

10011
←←←←
10010

- Right : zero small digit

↘ anazul el bazy

10011
— — — —
01001

Final 2022

$$F + \bar{G} + 1 = F - G$$



Faculty of Computers and Artificial Intelligence
Cairo University



Final Exam

Program: Computer Science / Software Engineering
Course Name: Computer Organization and Architecture
Course Code: CS331
Instructor(s): Dr. Amin Allam

Date: 15/1/2022
Duration: 2 hours
Total Marks: 60 marks

تعليمات هامة:

- حيازة التليفون المحمول مفتوحا داخل لجنة الامتحان يعتبر حالة غش تستوجب العقاب وإذا كان ضروري الدخول بالمحمول فيوضع مغلقا في الحقيبة.
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- Exam consists of 40 multiple-choice questions in 6 pages. Each question weights 1.5 marks.
- Record in the bubble sheet exactly ONE answer for each question.

Qa \Rightarrow A 4×1 multiplexer with selection inputs S_1 and S_0 selects input 0 when $S_1 S_0 = 00$, selects input 1 when $S_1 S_0 = 01$, selects input 2 when $S_1 S_0 = 10$, and selects input 3 when $S_1 S_0 = 11$.

For questions 1 to 6, consider an electronic circuit which has:

- Inputs: S_1, S_0, C_0 , the signed binary integer $F = F_1 F_0$, and the signed binary integer $G = G_1 G_0$.
- output: the signed binary integer $V = V_1 V_0$. (All signed integers are in the 2's complement form).

The circuit consists of the following items:

- A 4×1 multiplexer with selection inputs S_1, S_0 , and 4 inputs: Q_0, Q_1, Q_2, Q_3 , and output: M_0 .
- Another 4×1 multiplexer with selection inputs S_1, S_0 , and 4 inputs: R_0, R_1, R_2, R_3 , and output: M_1 .
- A full-adder with inputs: F_0, M_0, C_0 and outputs: C_1 (carry), V_0 (sum).
- Another full-adder with inputs: F_1, M_1, C_1 and outputs: C_2 (carry), V_1 (sum).
- \bar{G}_0 is connected to Q_0 . • \bar{G}_1 is connected to R_0 . • G_0 is connected to Q_1 . • G_1 is connected to R_1 .
- Logic 0 is connected to both Q_2 and R_2 . • Logic 1 is connected to both Q_3 and R_3 .

1 After $S_1 = 0, S_0 = 0, C_0 = 0$, the output V will equal to:

- ☐ A $F + G$ ☐ B $F - G$ ☐ C $F + G + 1$ ☒ D $F - G - 1$ ☐ E $F - G + 1$

$$F - G - 1 = F - G$$

2 After $S_1 = 0, S_0 = 0, C_0 = 1$, the output V will equal to:

- ☐ A $F + G$ ☒ B $F - G$ ☐ C $F + G + 1$ ☐ D $F - G - 1$ ☐ E $F - G + 1$

$$F + \bar{G}$$

3 After $S_1 = 0, S_0 = 1, C_0 = 1$, the output V will equal to:

- ☐ A $F + G$ ☐ B $F - G$ ☒ C $F + G + 1$ ☐ D $F - G - 1$ ☐ E $F - G + 1$

$$G$$

4 After $S_1 = 1, S_0 = 0, C_0 = 1$, the output V will equal to:

- ☐ A $F + G$ ☐ B $F - G$ ☒ C $F + 1$ ☐ D $F - 1$ ☐ E F

$$0$$

$$17.8$$

5 After $S_1 = 1, S_0 = 1, C_0 = 0$, the output V will equal to:

- ☐ A $F + G$ ☐ B $F - G$ ☐ C $F + 1$ ☒ D $F - 1$ ☐ E F

$$1$$

6 After $S_1 = 1, S_0 = 1, C_0 = 1$, the output V will equal to:

- ☐ A $F + G$ ☐ B $F - G$ ☐ C $F + 1$ ☐ D $F - 1$ ☒ E F

3 3 3 3 2 3
B B B C C A +
~ ~ ~ 2.42.4 5.4

- Qb** 7 The outputs of 32 registers are connected to the inputs of a bus. Each register stores 128 bits. The bus selects the outputs of one of the input registers according to the selection inputs of the bus. The bus must have the following number of **selection inputs**:
- A 4 ☒ B 5 ☒ C 7 ☐ D 32 ☐ E 128
- 8 The outputs of 32 registers are connected to the inputs of a bus. Each register stores 7 bits. The bus selects the outputs of one of the input registers according to the selection inputs of the bus. The bus must have the following number of **multiplexers**:
- A 4 ☐ B 5 ☐ C 7 ☒ D 32 ☐ E 128
- 9 A memory unit contains 1024 words. Each word is 32 bits. The number of **input address lines** to this memory unit is:
- A 5 ☒ B 10 ☐ C 11 ☐ D 16 ☐ E 32 $2^x = 1024$
- 10 A memory unit contains 1024 words. Each word is 16 bits. The number of **data output lines** out of this memory unit is:
- A 5 ☐ B 10 ☐ C 11 ☒ D 16 ☐ E 32
- 11 After applying **logical shift right** to a 6-bit register containing the binary number 101100 it becomes:
- ☒ A 010110 ☐ B 110110 ☐ C 011001 ☐ D 011000 ☐ E 111100 010110
- 12 Comparing two numbers to know whether they are equal or not equal, can be done by:
- A bitwise complement ☐ B bitwise and ☐ C bitwise or ☒ D bitwise xor ☐ E arithmetic shift
- 13 Consider the ADD instruction which increases the value of the accumulator register by the value located at the effective address. If the instruction uses **indirect** addressing and it contains the address 10 in its address field. The memory word at address 10 contains the value 30, and the memory word at address 30 contains the value 45. After the instruction is executed, the accumulator register increases by:
- A 10 ☐ B 30 ☐ C 40 ☒ D 45 ☐ E 75
- 14 Consider the ADD instruction which increases the value of the accumulator register by the value located at the effective address. If the instruction uses **direct** addressing and it contains the address 10 in its address field. The memory word at address 10 contains the value 30, and the memory word at address 30 contains the value 45. After the instruction is executed, the accumulator register increases by:
- A 10 ☒ B 30 ☐ C 40 ☐ D 45 ☐ E 75
- 15 The following register contains the instruction that is being executed:
- A Program counter ☐ B Accumulator ☒ C Instruction register ☐ D Input register ☐ E Data register
- 16 Assume that 4 clock cycles are required on average to execute any instruction. The time required to execute a program containing 10,000 instructions on a computer with a clock rate of 5,000,000 clock cycles per second is:
- A 0.002 seconds ☒ B 0.008 seconds ☐ C 125 seconds ☐ D 250 seconds ☐ E 2000 seconds

$$\frac{4 \times 10,000}{5,000,000} =$$

7. ⇒ For questions 17 to 20, consider the following control memory of a microprogrammed control unit. Each line represents a microinstruction stored in the control memory.

Microinstruction	Label	Microoperations	Condition	Branch	Address
0	ADD:	NOP	I	CALL	INDRCT
1		READ	U	JMP	NEXT
2		ADD	U	JMP	FETCH
3	INDRCT:	READ	U	JMP	NEXT
4		DRTAR	U	RET	
5	FETCH:	PCTAR	U	JMP	NEXT
6		READ, INCPC	U	JMP	NEXT
7		DRTAR	U	MAP	

17 After microinstruction 0 is executed, the following microinstruction will be executed:

- ☐ A 1 ☐ B 2 ☒ C 3 ☐ D 1 or 3 ☐ E unknown

18 After microinstruction 2 is executed, the following microinstruction will be executed:

- ☐ A 0 ☐ B 3 ☒ C 5 ☐ D 3 or 5 ☐ E unknown

19 After microinstruction 7 is executed, the following microinstruction will be executed:

- ☐ A 0 ☐ B 4 ☐ C 5 ☐ D 0 or 4 ☒ E unknown

20 The following sequence of microinstruction execution is possible:

- ☐ A 0,1,2,0 ☐ B 0,1,2,3 ☐ C 0,3,4,0 ☒ D 0,3,4,1 ☐ E 0,3,4,5

Qe ⇒ For questions 21 to 24, consider a binary multiplier which multiplies the multiplicand $B=10111$ by the multiplier $Q=10011$ and places the multiplication result into EAQ where A, B, and Q are 5 bits registers and E is a flipflop. Initially, $EAQ=00000010011$. Each stage starts by performing some necessary operations and completes by shifting right EAQ. The multiplication process completes in 6 stages.

21 After the first stage, $EAQ=$

- ☐ A 00000001001 ☒ B 00101111001 ☐ C 00110111001 ☐ D 00110011001 ☐ E 10110011001

22 After the second stage, $EAQ=$

- ☐ A 00010111100 ☐ B 00011011100 ☐ C 10011001100 ☐ D 01011001100 ☒ E 01000101100

23 After the third stage, $EAQ=$

- ☒ A 00100010110 ☐ B 00001011110 ☐ C 00001101110 ☐ D 01001100110 ☐ E 00101100110

24 After the fourth stage, $EAQ=$

- ☒ A 00010001011 ☐ B 00000101111 ☐ C 00000110111 ☐ D 00100110011 ☐ E 00010110011

10111

E	A	Q	sc
0	00000	10011	101
	10111		
0	10111	10011	
0	01011	11001	100
	10111		
1	00010	11001	
0	10001	01100	011
0	10000	10110	

Add
Branch
I A E
I A E

Qd ⇒ For questions 25 to 30, consider a three-segment instruction pipeline I, A, E described as follows. The addressing mode used for the Branch instruction is the register indirect mode:

	Load/Store	Add	Branch
I	Fetch and decode the instruction		
A	Evaluate effective address	Perform the addition	Evaluate branch address
E	Transfer operand from/to memory	Put result in destination register	Put branch address in PC

25 The sequence of instructions: (1) Load R1 (2) Load R2 (3) $R3 \leftarrow R1 + R2$ (4) Store R3 has the following pipeline difficulty:

- ☐ A Resource conflict ☒ B Data dependency ☐ C Branch ☐ D Closed pipeline ☐ E Interrupt

26 The pipeline difficulty in the sequence of instructions:

(1) Load R1 (2) Load R2 (3) $R3 \leftarrow R1 + R2$ (4) Store R3 can be solved by:

- ☐ A Adding a no-operation instruction between instructions 1 and 2 ☐ B Swapping instructions 2 and 3
☒ C Adding a no-operation instruction between instructions 2 and 3 ☐ D Swapping instructions 3 and 4
☐ E Adding a no-operation instruction between instructions 3 and 4

27 After solving the pipeline difficulty and pipelining the following sequence of instructions: (1) Load R1 (2) Load R2 (3) $R3 \leftarrow R1 + R2$ (4) Store R3, they will require the following number of clock cycles to be executed:

- ☐ A 4 ☐ B 5 ☐ C 6 ☒ D 7 ☐ E 12

28 After solving the pipeline difficulties by inserting no-operation instructions and/or rearranging instructions, such that the resulting pipeline of the following instruction sequence requires the minimum number of clock cycles: (1) Load R1 (2) $R2 \leftarrow R1 + 8$ (3) Branch R2 (4) $R3 \leftarrow R2 + 8$, they will require the following number of clock cycles to be executed:

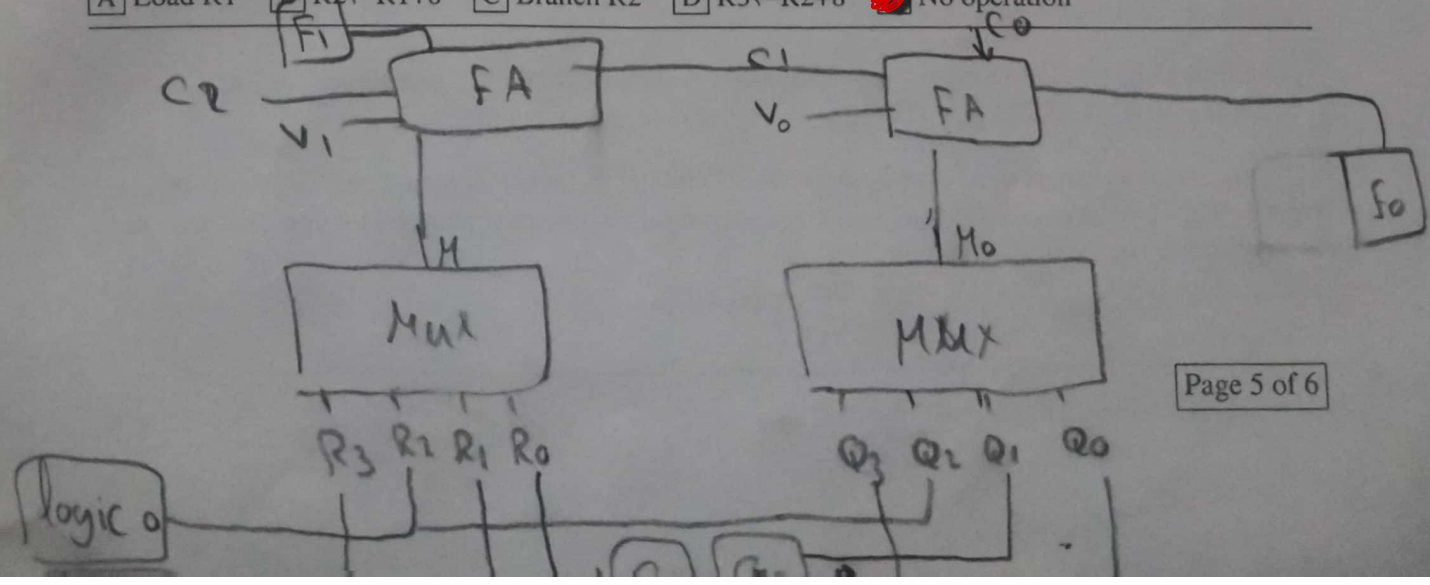
- ☐ A 5 ☐ B 6 ☐ C 7 ☒ D 8 ☐ E 12

29 After solving the pipeline difficulties by inserting no-operation instructions and/or rearranging instructions, such that the resulting pipeline of the following instruction sequence requires the minimum number of clock cycles: (1) Load R1 (2) $R2 \leftarrow R1 + 8$ (3) Branch R2 (4) $R3 \leftarrow R2 + 8$, the third instruction will be:

- ☐ A Load R1 ☐ B $R2 \leftarrow R1 + 8$ ☒ C Branch R2 ☐ D $R3 \leftarrow R2 + 8$ ☐ E No operation

30 After solving the pipeline difficulties by inserting no-operation instructions and/or rearranging instructions, such that the resulting pipeline of the following instruction sequence requires the minimum number of clock cycles: (1) Load R1 (2) $R2 \leftarrow R1 + 8$ (3) Branch R2 (4) $R3 \leftarrow R2 + 8$, the fourth instruction will be:

- ☐ A Load R1 ☐ B $R2 \leftarrow R1 + 8$ ☐ C Branch R2 ☐ D $R3 \leftarrow R2 + 8$ ☒ E No operation



Qe

31 Consider the ADD instruction which increases the value of the accumulator register by the value located at the effective address. If the instruction uses **register** addressing mode and it contains the register R1 in its address field. The register R1 contains the value 10. The memory word at address 10 contains the value 30, and the memory word at address 30 contains the value 45. After the instruction is executed, the accumulator register increases by:

- ☒ A 10 ☐ B 30 ☐ C 40 ☐ D 45 ☐ E 75

32 Consider the ADD instruction which increases the value of the accumulator register by the value located at the effective address. If the instruction uses **register indirect** addressing mode and it contains the register R1 in its address field. The register R1 contains the value 10. The memory word at address 10 contains the value 30, and the memory word at address 30 contains the value 45. After the instruction is executed, the accumulator register increases by:

- ☐ A 10 ☒ B 30 ☐ C 40 ☐ D 45 ☐ E 75

push M[A]

33 Let $M[X]$ = the integer value at memory address X. Consider zero-address instructions: Pu

- (1) **Push X**: pushes $M[X]$ to stack. (2) **Pop X**: pops the top stack element to $M[X]$.
(3) **Add**: pops the two elements on the top of the stack and pushes their sum.

The minimum number of such instructions that compute $M[X] = M[A] + M[B] + M[C]$ is:

- ☐ A 3 ☐ B 4 ☒ C 5 ☐ D 6 ☐ E 7

AC \rightarrow M[A]

34 Let $M[X]$ = the integer value at memory address X. Consider one-address instructions:

- (1) **Load X**: sets AC to $M[X]$. (2) **Store X**: sets $M[X]$ to AC. (3) **Add X**: increases AC by $M[X]$.

The minimum number of such instructions that compute $M[X] = M[A] + M[B] + M[C]$ is:

- ☐ A 3 ☒ B 4 ☐ C 5 ☐ D 6 ☐ E 7

AC \rightarrow AC + M[B]

AC \rightarrow AC + M[C]

M[X]

35 The following event should cause an **internal** interrupt:

- ☐ A keyboard moved a character to input register ☐ B printer ready to take character from CPU
☐ C switching from user mode to administrative mode ☒ D division by zero ☐ E power failure

36 The Zero flipflop becomes 1 if the last operation:

- ☐ A took a zero operand ☐ B result was zero ☐ C was no-operation instruction
☐ D took no operands ☒ E compared two different integers

37 A characteristic of RISC is:

- ☐ A large number of instructions ☐ B variable-length instruction format ☒ C hardwired control unit
☐ D large variety of addressing modes ☐ E small size of control memory

38 The following step is **not** required in multiplication of floating point numbers:

- ☐ A check for zeros ☐ B add the exponents ☒ C align the mantissas ☐ D multiply the mantissas
☐ E normalize the product

10-10

39 The following associative memory word matches Argument = 1010011 and Key = 11011000:

- ☒ A 10010111 ☐ B 11100111 ☐ C 11110111 ☐ D 10111111 ☐ E 11111111

40 A main memory contains 2^{15} words. Each word is 9 bits. A direct mapping cache memory contains 64 words. The value at address $(34567)_8$ in main memory (8 means octal number) may be cached at the following address in the cache memory:

- ☐ A $(34)_8$ ☐ B $(345)_8$ ☐ C $(456)_8$ ☒ D $(567)_8$ ☐ E $(67)_8$

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20196022

رئيسية رسمى كورس مسجل



Midterm 2021

Faculty of Computers and Artificial Intelligence
Cairo University

Midterm Exam



Program: Computer Science / Software Engineering
Course Name: Computer Organization and Architecture
Course Code: CS331
Instructor(s): Dr. Amin Allam

Date: 27/11/2021
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- Exam consists of 20 multiple-choice questions in 3 pages. Each question weights 1 mark.
- Record in the bubble sheet exactly ONE answer for each question.

Qa ⇒ Given the characteristic table of the JK flip-flop:

J	K	$Q(t+1)$
0	0	$Q(t)$
0	1	0
1	0	1
1	1	$\overline{Q(t)}$

For questions 1 to 4, consider an electronic circuit which consists of the following items:

- A positive edge-triggered JK flip-flop with inputs: clock, J_0 , K_0 , and output: A_0 .
- Another positive edge-triggered JK flip-flop with inputs: clock, J_1 , K_1 , and output: A_1 .
- Logic 1 is connected to both J_0 and K_0 .
- A_0 is connected to both J_1 and K_1 .

1 Assuming $A_1 = 0$ and $A_0 = 0$ while clock input is logic 0. After the clock input becomes logic 1, the values of A_1 and A_0 become:

[A] $A_1 = 0, A_0 = 0$ [B] $A_1 = 0, A_0 = 1$ [C] $A_1 = 1, A_0 = 0$ [D] $A_1 = 1, A_0 = 1$ [E] unknown

2 Assuming $A_1 = 0$ and $A_0 = 1$ while clock input is logic 0. After the clock input becomes logic 1, the values of A_1 and A_0 become:

[A] $A_1 = 0, A_0 = 0$ [B] $A_1 = 0, A_0 = 1$ [C] $A_1 = 1, A_0 = 0$ [D] $A_1 = 1, A_0 = 1$ [E] unknown

3 Assuming $A_1 = 1$ and $A_0 = 1$ while clock input is logic 0. After the clock input becomes logic 1, the values of A_1 and A_0 become:

[A] $A_1 = 0, A_0 = 0$ [B] $A_1 = 0, A_0 = 1$ [C] $A_1 = 1, A_0 = 0$ [D] $A_1 = 1, A_0 = 1$ [E] unknown

4 The above circuit is:

[A] half adder [B] shift register [C] binary counter [D] bus [E] other

Qb \Rightarrow A 4×1 multiplexer with selection inputs S_1 and S_0 selects input 0 when $S_1S_0 = 00$, selects input 1 when $S_1S_0 = 01$, selects input 2 when $S_1S_0 = 10$, and selects input 3 when $S_1S_0 = 11$.

For questions 5 to 10, consider an electronic circuit which has:

• Inputs: S_1, S_0, C_0 , the binary number $F = F_1F_0$, and the binary number $G = G_1G_0$.

• output: the binary number $V = V_1V_0$.

The circuit consists of the following items:

• A 4×1 multiplexer with selection inputs S_1, S_0 , and 4 inputs: Q_0, Q_1, Q_2, Q_3 , and output: M_0 .

• Another 4×1 multiplexer with selection inputs S_1, S_0 , and 4 inputs: R_0, R_1, R_2, R_3 , and output: M_1 .

• A full-adder with inputs: F_0, M_0, C_0 and outputs: C_1 (carry), V_0 (sum).

• Another full-adder with inputs: F_1, M_1, C_1 and outputs: C_2 (carry), V_1 (sum).

• $\overline{C_0}$ is connected to Q_0 . • $\overline{G_1}$ is connected to R_0 . • G_0 is connected to Q_1 . • G_1 is connected to R_1 .

• Logic 1 is connected to both Q_2 and R_2 . • Logic 0 is connected to both Q_3 and R_3 .

5 After $S_1 = 0, S_0 = 0, C_0 = 0$, the output V will equal to:

☐ A $F + G$ ☐ B $F - G$ ☐ C $F + G + 1$ ☒ D $F - G - 1$ ☐ E $F - G + 1$

6 After $S_1 = 0, S_0 = 0, C_0 = 1$, the output V will equal to:

☐ A $F + G$ ☒ B $F - G$ ☐ C $F + G + 1$ ☐ D $F - G - 1$ ☐ E $F - G + 1$

7 After $S_1 = 0, S_0 = 1, C_0 = 1$, the output V will equal to:

☐ A $F + G$ ☐ B $F - G$ ☒ C $F + G + 1$ ☐ D $F - G - 1$ ☐ E $F - G + 1$

8 After $S_1 = 1, S_0 = 0, C_0 = 0$, the output V will equal to:

☐ A $F + G$ ☐ B $F - G$ ☐ C $F + 1$ ☒ D $F - 1$ ☐ E F

9 After $S_1 = 1, S_0 = 0, C_0 = 1$, the output V will equal to:

☐ A $F + G$ ☐ B $F - G$ ☐ C $F + 1$ ☐ D $F - 1$ ☒ E F

10 After $S_1 = 1, S_0 = 1, C_0 = 1$, the output V will equal to:

☐ A $F + G$ ☐ B $F - G$ ☒ C $F + 1$ ☐ D $F - 1$ ☐ E F

11 The outputs of 16 registers are connected to the inputs of a bus. Each register stores 32 bits. The bus selects the outputs of one of the input registers according to the selection inputs of the bus. The bus must have the following number of selection inputs:

- ☐ A 3 ☒ B 4 ☐ C 5 ☐ D 16 ☐ E 32

12 The outputs of 16 registers are connected to the inputs of a bus. Each register stores 32 bits. The bus selects the outputs of one of the input registers according to the selection inputs of the bus. The bus must have the following number of multiplexers:

- ☐ A 3 ☐ B 4 ☐ C 5 ☐ D 16 ☒ E 32

13 A memory unit contains 2048 words. Each word is 32 bits. The number of input address lines to this memory unit is:

- ☐ A 5 ☐ B 10 ☒ C 11 ☐ D 16 ☐ E 32

14 A memory unit contains 2048 words. Each word is 32 bits. The number of data output lines out of this memory unit is:

- ☐ A 5 ☐ B 10 ☐ C 11 ☐ D 16 ☒ E 32

15 After applying arithmetic shift right to a 6-bit register containing the binary number 101100 it becomes:

- ☐ A 010110 ☒ B 110110 ☐ C 011001 ☐ D 011000 ☐ E 111100

16 Comparing two numbers to know whether they are equal or not equal, can be done by:

- ☐ A bitwise and ☐ B bitwise or ☒ C bitwise xor ☐ D bitwise complement ☐ E arithmetic shift

17 Consider the ADD instruction which increases the value of the accumulator register by the value located at the effective address. If the instruction uses direct addressing and it contains the address 10 in its address field. The memory word at address 10 contains the value 30, and the memory word at address 30 contains the value 45. After the instruction is executed, the accumulator register increases by:

- ☐ A 10 ☒ B 30 ☐ C 40 ☐ D 45 ☐ E 75

18 Consider the ADD instruction which increases the value of the accumulator register by the value located at the effective address. If the instruction uses indirect addressing and it contains the address 10 in its address field. The memory word at address 10 contains the value 30, and the memory word at address 30 contains the value 45. After the instruction is executed, the accumulator register increases by:

- ☐ A 10 ☐ B 30 ☐ C 40 ☒ D 45 ☐ E 75

19 The following register contains the address of the instruction that is going to be executed:

- ☒ A Program counter ☐ B Accumulator ☐ C Instruction register ☐ D Input register ☐ E Data register

20 Assume that 2 clock cycles are required on average to execute any instruction. The time required to execute a program containing 10,000 instructions on a computer with a clock rate of 5,000,000 clock cycles per second is:

- ☐ A 0.001 seconds ☐ B 0.002 seconds ☒ C 0.004 seconds ☐ D 500 seconds ☐ E 1000 seconds

Practice Sheet



Qa \Rightarrow A 4×1 multiplexer with selection inputs S_1 and S_0 selects input 0 when $S_1S_0 = 00$, selects input 1 when $S_1S_0 = 01$, selects input 2 when $S_1S_0 = 10$, and selects input 3 when $S_1S_0 = 11$.

For questions 1 to 6, consider an electronic circuit which has:

- Inputs: S_1, S_0, C_0 , the binary number $A = A_1A_0$, and the binary number $B = B_1B_0$.
- output: the binary number $D = D_1D_0$.

The circuit consists of the following items:

- A 4×1 multiplexer with selection inputs S_1, S_0 , and 4 inputs: Q_0, Q_1, Q_2, Q_3 , and output: M_0 .
- Another 4×1 multiplexer with selection inputs S_1, S_0 , and 4 inputs: R_0, R_1, R_2, R_3 , and output: M_1 .
- A full-adder with inputs: A_0, M_0, C_0 and outputs: C_1 (carry), D_0 (sum).
- Another full-adder with inputs: A_1, M_1, C_1 and outputs: C_2 (carry), D_1 (sum).
- B_0 is connected to Q_0 . • B_1 is connected to R_0 . • $\overline{B_0}$ is connected to Q_1 . • $\overline{B_1}$ is connected to R_1 .
- Logic 0 is connected to both Q_2 and R_2 . • Logic 1 is connected to both Q_3 and R_3 .

1 When $S_1 = 0, S_0 = 0, C_0 = 0$:

- ☒ A $D = A + B$ ☐ B $D = A - B$ ☐ C $D = A + 1$ ☐ D $D = A - 1$ ☐ E $D = A$

2 When $S_1 = 0, S_0 = 1, C_0 = 1$:

- ☐ A $D = A + B$ ☒ B $D = A - B$ ☐ C $D = A + 1$ ☐ D $D = A - 1$ ☐ E $D = A$

3 When $S_1 = 1, S_0 = 0, C_0 = 0$:

- ☐ A $D = A + B$ ☐ B $D = A - B$ ☐ C $D = A + 1$ ☐ D $D = A - 1$ ☒ E $D = A$

4 When $S_1 = 1, S_0 = 0, C_0 = 1$:

- ☐ A $D = A + B$ ☐ B $D = A - B$ ☒ C $D = A + 1$ ☐ D $D = A - 1$ ☐ E $D = A$

5 When $S_1 = 1, S_0 = 1, C_0 = 0$:

- ☐ A $D = A + B$ ☐ B $D = A - B$ ☐ C $D = A + 1$ ☒ D $D = A - 1$ ☐ E $D = A$

6 When $S_1 = 1, S_0 = 1, C_0 = 1$:

- ☐ A $D = A + B$ ☐ B $D = A - B$ ☐ C $D = A + 1$ ☐ D $D = A - 1$ ☒ E $D = A$

