



Theory of Computations

Formal Modeling Concepts

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Modeling

- What is a Model?
 - A Model is a purposeful abstract representation of reality
 - Purpose
 - You may have more than one model for the same reality for different purpose
 - Abstract
 - Does not represent instances
 - Schematic description

Modeling notation

- Textual
 - Pseudocode, structured English, programs,
- Visual
 - Flow charts, UML, ERD
- Formal
 - Mathematics, logic, graphs

Models

- Enhance better understanding
- Higher level of thinking
- Important in design
- Better communication
- Formal: a medium for proof
 - The existence a required property
 - The absence of a property

Formal Models

- Set Theory
- Graph Theory
- Formal Logic

Sets

- A set is a collection (group) of elements
 - Notation $\{ \}$, $\{a, x, \text{book}\}$ $\{s \mid \text{student}(s)\}$
 - Ordering
 - Repetition
 - Explicit/ Implicit
 - Cardinality (Function)
 - Φ : Empty Set
- Tuple (List)
 - Notation()
 - Order
 - Repetition

Relationships

- Logical (Boolean)
 - Answers a question
 - Membership:
 - *(Input: An item & a Set, Output: Boolean)*
 - Subset:
 - *(Input: Two Sets, Output: Boolean)*
 - Equivalence:
 - *(Input: Two Sets, Output: Boolean)*
 - Note: Subset & Equivalence

Operations

- Produce results
 - Union: $R \cup S$
 - *(Input: Two Sets, Output: a Set)*
 - *A set contains elements in either of input sets*
 - *Commutative, i.e., $R \cup S = S \cup R$*
 - Intersection
 - *(Input: Two Sets, Output: a Set)*
 - *A set contains elements in both input sets*
 - *Commutative*
 - Difference
 - *(Input: Two Sets, Output: a Set)*
 - *A set contains elements in one but not in the second*
 - *Not Commutative*

Operations

- Cartesian Product
- Two sets S and T
 - $S \times T = \{(s, t) \mid s \in S \ \& \ t \in T\}$
 - Cardinality
 - Order

Relation

- A relation, R , between two sets S & T
 - Meaningful association between an element in S with an element in T
 - A set of pairs, one element from S and one element from T
 - $R = \{(s, t) \mid s \in S \text{ \& } t \in T\}$
 - A subset from Cartesian Product
 - Cardinality = number of elements: $|R|$
 - Arity: binary (2), 3, 4, ...number of sets in the relation

Functions

- Two sets, X and Y
- A function $F: X \rightarrow Y$ or $y = f(x)$.
 - X is the Domain
 - Y is the Co-domain
 - y is the Image of x under f
 - Range, the set of images
 - A function is a relation with unique images
 - Cardinality is a function

Sets of Numbers

- N: the set of *Natural Numbers*
- Z: the set of *Integers*
 - N with 0 and negative integers
- Q: the set of *Rational Numbers*
 - Fractions: two numbers from N or Z
- R: The set of *Real Numbers*
 - *Non rational numbers, roots*
 - The set of *Imaginary numbers*

Graph theory

Graphs

- A graph G is a construct of two finite sets, $V=\{v_1, v_2, \dots, v_n\}$ a set of vertices and a set of edges $E=\{e_1, e_2, \dots, e_m\}$.
- Each edge is represented as a pair of vertices from V : $e_i=(v_j, v_k)$, where v_j and $v_k \in V$.
 - Usually, vertices represent objects, or elements, and edges represent relationship
 - Vertices and Edges may be called Nodes and Links

Graphs

- Edges may be directed or undirected
 - Direction may be represented, visually, by an arrow
- A *directed* graph is a type of graph in which each edge (v_i, v_j) is directed, i.e., from v_i to v_j but not vice versa, otherwise it is *undirected*.

Graphs

- A *path*, from v_1 to v_k in a graph $G(V, E)$ is a sequence of vertices $(v_1, v_2, \dots, v_{k-1}, v_k)$ such that $\{(v_1, v_2), (v_2, v_3), \dots, (v_{k-1}, v_k)\}$ exist in E .
- Length of a path is the number of edges in the path

Graphs

- A *cycle* in a graph G is a path (v_i, \dots, v_i)
- A *cyclic graph* is a graph in which there is any cycle, otherwise it is *acyclic*.

Graphs

- *Connected* graph is a graph in which there is a path from every vertex to all other vertices, otherwise it is disconnected.

Trees

- A *tree* is a connected acyclic graph.
- *Root* : a vertex with out-edges, no in-edges
- *Leaves*: Vertices with in-edges, no out-edges
- *Height* of a tree: length of the path from root to a leaf.
 - Sometimes called depth of a tree.

Graphs in CS

- A network is a graph
 - Network of computers
 - Network of people (social network)
- Trees are commonly used in CS
 - Indices in data and database
 - Search tree
- An ontology is a graph or a tree of concepts

Logic

Logic

- What is logic?

Logic

- Logic is the branch of science that studies the relationship between premises (Assumptions) & results (Conclusion).
- Logic is a branch of philosophy.
- In CS, Symbolic Logic
 - Well-Formed-Formulae (WFF)
 - Propositional
 - Predicate

Propositional Logic

- *Atomic Formula* is a proposition.
- A *proposition* is a statement that can be either true or false, but not both.
 - P: It is hot
 - Q: The air conditioner is on
 - R: Lights are on
 - P and Q: it is hot and the air conditioner is on

Propositional Logic

- WFF in Propositional Logic:
 - A proposition is a formula (Atomic Formula)
 - If F is a formula, then $\sim F$ (not F) is also a formula
 - If F & G are Formulae, then
 - F and G (Conjunction) is a formula
 - F or G (Disjunction) is a formula
 - If F then G is a formula
 - F iff G is a formula (if and only if)
 - Nothing Else is a formula
- Note: Algebraic rules & DeMorgan's

Evaluation of a Formula

- A Proposition has a value (True or False)
- An *interpretation* of a formula is an assignment of each of its atoms by a value (true or false).
- A formula has a value (T or F) under each interpretation.

Evaluation of a Formula

- A formula is said to be *valid* if it is true under all its interpretations, otherwise it is *invalid*.
- A formula is said to be *inconsistent* if it is false under all its interpretations, otherwise it is *consistent*.
 - A valid formula is called a “tautology”
 - An inconsistent formula is called a contradiction

Predicate Logic

- A *predicate* is a logic statement that has an argument. Usually the predicate applies to the argument(s).
 - Student (John): John is a student
 - Has-book(John): John has book
 - Likes(John, Mary): John Likes Mary
 - Give(John, Mary, book)
- An argument may be a variable
 - Student(x) has x as a variable
 - When x is substituted with a value, it is called an instance of the predicate.
 - Instantiation is the process.

Predicate Logic

- WFF in Predicate Logic
 - A predicate is a formula (Atomic Formula)
 - If F is a formula, then $\sim F$ (not F) is also a formula
 - If F & G are Formulae, then
 - F and G (Conjunction) is a formula
 - F or G (Disjunction) is a formula
 - If F then G is a formula
 - F iff G is a formula

Predicate Logic

- If F is a formula that contains a variable x defined over a domain D , then
(“***For All***” x) F is evaluated as true only if it is true for each $x \in D$.
- If F is a formula that contains a variable x defined over a domain D , then
(“***There Exists***” x) F is evaluated as true if at least one value of $x \in D$ makes F true.
- Nothing Else is a formula

Notes on Quantifiers

- For All & There Exists are called quantifiers.
 - For all = for each = for every
 - There exist = for some = for at least one
 - \sim for all (x) = there exist (\sim x)
 - \sim there exist(x) = for all (\sim x)
 - \sim for all (x) \neq for all (\sim x)
- A quantifier has a scope of variables.
- If all the variables in a formula are quantified, it is called *bound*, otherwise it is *loose*.

Evaluation of Formula

- The same definitions of Validity and Consistency apply in general to predicates.
- “For all x ” $F(x)$ is true if F is true for all values of $x \in D$
- “There Exists x ” $F(x)$ is true if any value of $x \in D$ makes F true

EXAMPLES

- All basketball players are tall
 - $\forall X \text{ play}(X, \text{basketball}) \rightarrow \text{tall}(X)$
- John like anyone who likes books
 - $\text{like}(X, \text{book}) \rightarrow \text{like}(\text{john}, X)$
- Nobody likes taxes
 - $\neg \exists X \text{ likes}(X, \text{taxes})$
- There is a person who writes computer class
 - $\exists X \text{ write}(X, \text{computer class})$
- John did not study but he is lucky
 - $\neg \text{study}(\text{john}) \wedge \text{luky}(\text{john})$
- All cats and dogs are animals
 - $\forall X \forall Y \text{ cats}(X) \wedge \text{dogs}(Y) \rightarrow \text{animals}(X) \wedge \text{animals}(Y)$

Logical Consequence

- Given a set of formulae, G , we can say F is a logical consequence G , $G \models F$, if we can prove F using a subset of G .
 - $\{F \cup G\}$ is consistent
 - $\{\sim F \cup G\}$ is inconsistent

Proof Procedure

- The path from the assumptions to the result is called the *proof procedure*.

Logic in CS

- In AI there are many useful application of Logic.
 - Knowledge & knowledge base
 - Knowledge is composed of logic statements
 - Inference is the mechanism of using knowledge to infer more statements from asserted ones.
 - Semantic Web
 - RDF a language of semantic web based on logic, each RDF statement is a clause (OAV)
 - OWL a language for semantic web based on logic.
 - Association rules and logic statements