# Image Processing and Pattern Recognition (IPPR) Chapter5:Image Compression

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#### Contents

- Need of Compression
- Coding Redundancy: Huffman Coding
- Interpixel Redundancy: run length coding
- Psychovisual Redundancy: bit plane
- Image Compression Models
- Lossy & Lossless compression: Transform coding and predictive coding.

## Why we need image compression?

One 90 minutes color movie, each second plays 24 frames. When we digitize it, each frame has 512×512 pixels, each pixel has three components R, G, B each one occupies 8 bits respectively, the total byte number is: 90×60×24×3×512×512 = 97,200MB

• A CD may save 600 megabytes data, the movie needs 160 CDs to save.

#### Image Compression

Objective: To reduce the amount of data required to represent an image.

Important in data storage and transmission

- Progressive transmission of images (internet, www)
- Video coding (HDTV, Teleconferencing)
- Digital libraries and image databases
  - Medical imaging
  - Satellite images

#### Lossy Vs. Lossless Compression

#### **Compression techniques**

Information preserving

(loss-less)

Images can be compressed and restored without any loss of information.

Application: Medical images, GIS



Perfect recovery is not possible but provides a large data compression. Example: TV signals, teleconferencing

#### Image Compression

- •Data compression refers to the process of reducing the amount of data required to represent a given quantity of information
  - data are the means where information is conveyed
  - relative data redundancy  $R_D$ :

$$R_D = 1 - \frac{1}{C_R}$$

where  $C_R$  is the compression ratio:

$$C_R = \frac{n_1}{n_2},$$

No of bits used to represent image before compression

$$C_R > = 1$$

No of bits used to represent image after compression

#### Image Compression

- Three basic data redundancy can be identified and exploited
  - coding redundancy
    - to reduce the amount of data representing the same information
  - interpixel redundancy
    - including spatial redundancy, geometric redundancy and interframe redundancy
  - psychovisual redundancy
    - certain information has less relative importance in normal visual processing and can be eliminated

• Assume a discrete random variable  $r_k$  in the interval [0,1] represents the grey levels of an image and that each  $r_k$  occurs with probability  $p_r(r_k)$ :

$$p_r(r_k) = \frac{n_k}{n}$$
  $k = 0,1,...,L-1$ 

and

$$L_{avg} = \sum_{k=0}^{L-1} l(r_k) p_r(r_k)$$

is the average number of bits used to represent each pixel, where  $l(r_k)$  is the number of bits required for each value  $r_k$ 

| r <sub>k</sub> | $p_r(r_k)$ | Code 1 | $l_1(r_k)$ | Code 2 | $l_2(r_k)$ |
|----------------|------------|--------|------------|--------|------------|
| $r_0 = 0$      | 0.19       | 000    | 3          | 11     | 2          |
| $r_1 = 1/7$    | 0.25       | 001    | 3          | 01     | 2          |
| $r_2 = 2/7$    | 0.21       | 010    | 3          | 10     | 2          |
| $r_3 = 3/7$    | 0.16       | 011    | 3          | 001    | 3          |
| $r_4 = 4/7$    | 0.08       | 100    | 3          | 0001   | 4          |
| $r_5 = 5/7$    | 0.06       | 101    | 3          | 00001  | 5          |
| $r_6 = 6/7$    | 0.03       | 110    | 3          | 000001 | 6          |
| $r_7 = 1$      | 0.02       | 111    | 3          | 000000 | 6          |
|                |            |        |            |        |            |

#### TABLE 8.1

Example of variable-length coding.

$$L_{avg} = \sum_{k=0}^{7} l_2(r_k) p_r(r_k)$$

$$= 2(0.19) + 2(0.25) + 2(0.21) + 3(0.16) + 4(0.08)$$

$$+5(0.06)+6(0.03)+6(0.02)$$

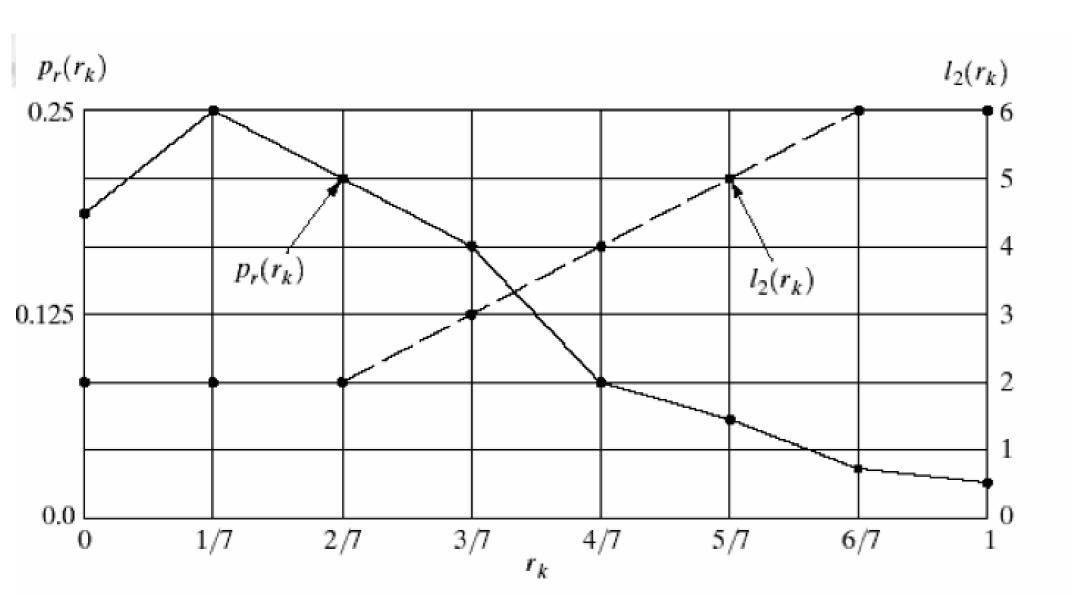
$$= 2.7$$
 bits,

$$R_D = 1 - \frac{1}{1.11} = 0.099$$

| $r_k$       | $p_r(r_k)$ | Code 1 | $l_1(r_k)$ | Code 2 | $l_2(r_k)$ |
|-------------|------------|--------|------------|--------|------------|
| $r_0 = 0$   | 0.19       | 000    | 3          | 11     | 2          |
| $r_1 = 1/7$ | 0.25       | 001    | 3          | 01     | 2          |
| $r_2 = 2/7$ | 0.21       | 010    | 3          | 10     | 2          |
| $r_3 = 3/7$ | 0.16       | 011    | 3          | 001    | 3          |
| $r_4 = 4/7$ | 0.08       | 100    | 3          | 0001   | 4          |
| $r_5 = 5/7$ | 0.06       | 101    | 3          | 00001  | 5          |
| $r_6 = 6/7$ | 0.03       | 110    | 3          | 000001 | 6          |
| $r_7 = 1$   | 0.02       | 111    | 3          | 000000 | 6          |
| •           |            |        |            |        |            |

L<sub>avg</sub> 3 bits/symbol

L<sub>avg</sub> 2.7 bits/symbol



#### FIGURE 8.1

Graphic representation of the fundamental basis of data compression through variable-length coding.

Concept: assign the longest code word to the symbol with the least probability of occurrence.

## Huffman coding

| Huffn              | nan c | ode:  | Consid | er a 6 | symbol                | source |
|--------------------|-------|-------|--------|--------|-----------------------|--------|
|                    | $a_1$ | $a_2$ | $a_3$  | $a_4$  | <b>a</b> <sub>5</sub> | $a_6$  |
| p(a <sub>i</sub> ) | 0.1   | 0.4   | 0.06   | 0.1    | 0.04                  | 0.3    |

#### Huffman coding

 Huffman coding: give the smallest possible number of code symbols per source symbols.

Step 1: Source reduction

| Original source                                 |  | Source reduction                  |                          |                |                     |  |
|---|--|-----------------------------------|--------------------------|----------------|---------------------|--|
| Symbol  | Probability                              | 1                                 | 2                        | 3              | 4                   |  |
| $a_{2}$ $a_{6}$ $a_{1}$ $a_{4}$ $a_{3}$ $a_{5}$ | 0.4<br>0.3<br>0.1<br>0.1<br>0.06<br>0.04 | 0.4<br>0.3<br>0.1<br>0.1<br>• 0.1 | 0.4<br>0.3<br>0.2<br>0.1 | 0.4<br>0.3<br> | <b>►</b> 0.6<br>0.4 |  |

## Huffman coding

#### Step 2: Code assignment procedure

| Original source |       |       | Source reduction |       |            |       |     |      |      |   |
|-----------------|-------|-------|------------------|-------|------------|-------|-----|------|------|---|
| Sym.            | Prob. | Code  | 1                | L     | 2          | 2     | 2   | 3    | 4    | 4 |
| $a_2$           | 0.4   | 1     | 0.4              | 1     | 0.4        | 1     | 0.4 | 1 _  | -0.6 | 0 |
| $a_6$           | 0.3   | 00    | 0.3              | 00    | 0.3        | 00    | 0.3 | 00 - | 0.4  | 1 |
| $a_1$           | 0.1   | 011   | 0.1              | 011   | -0.2       | 010   | 0.3 | 01 🕶 |      |   |
| $a_4$           | 0.1   | 0100  | 0.1              | 0100- | 0.1        | 011 🔫 |     |      |      |   |
| $a_3$           | 0.06  | 01010 | 0.1              | 0101  | <b>4</b> ∫ |       |     |      |      |   |
| $a_5$           | 0.04  | 01011 |                  |       |            |       |     |      |      |   |

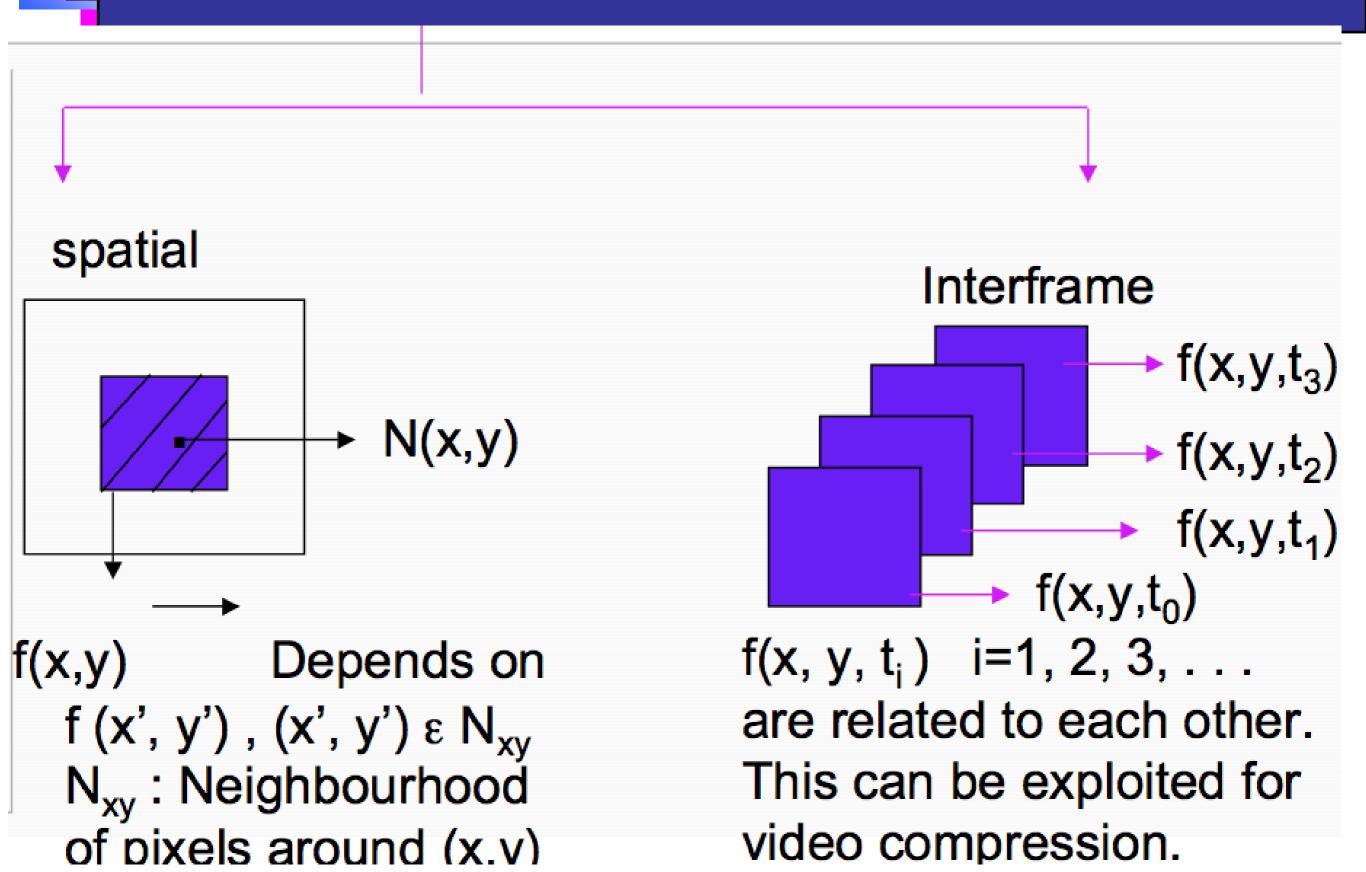
The code is instantaneous uniquely decodable without referencing succeeding symbols.

#### Average length:

$$(0.4)(1) + 0.3(2) + 0.1 \times 3 + 0.1 \times 4 + (0.06 + 0.04)(1)$$

$$0.04) 5 = 2.2 \text{ bits/symbol}$$

#### Interpixel / Interframe Redundancy



#### Interpixel / Interframe Redundancy



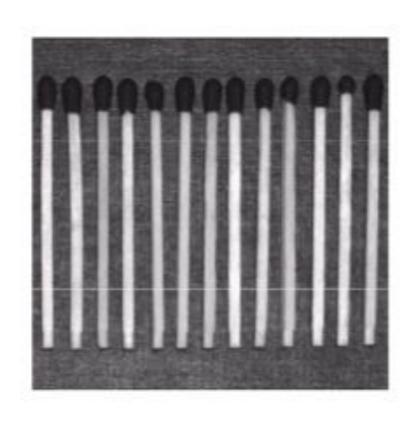
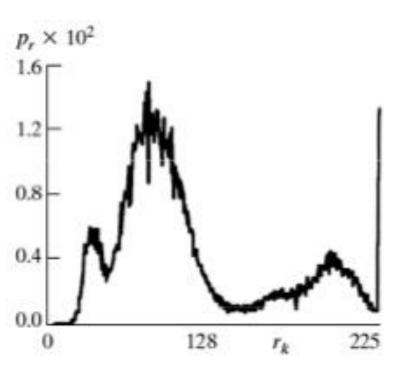


Figure 8.2 Two images and their grey-level pixel histograms and normalized autocorrelation coefficients along one line.



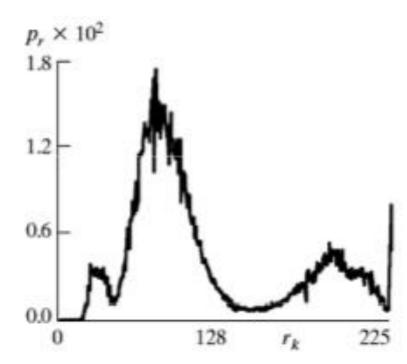
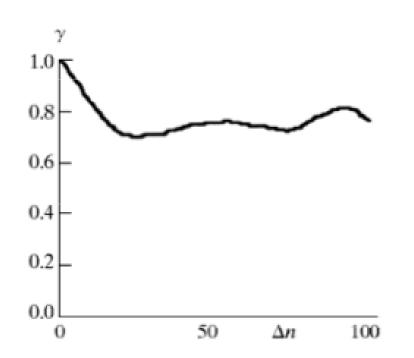




Image Processing and Pattern Recognition (IPPR)

#### Interpixel / Interframe Redundancy



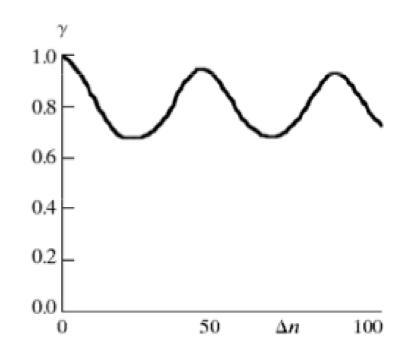


Figure 8.2 Two images and their grey-level pixel histograms and normalized autocorrelation coefficients along one line.



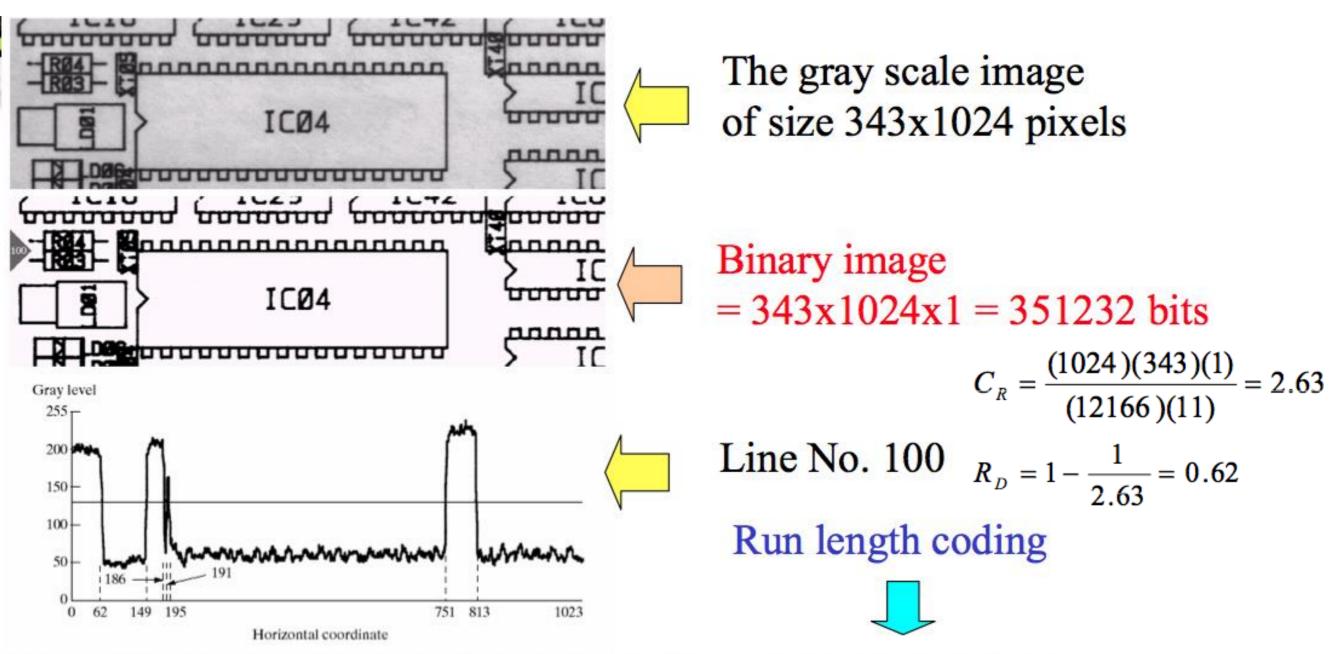
$$r(\Delta n) = \frac{A(\Delta n)}{A(0)},$$

where 
$$A(\Delta n) = \frac{1}{N - \Delta n} \sum_{y=0}^{N-1-\Delta n} f(x, y) f(x, y + \Delta n)$$

Interpixel redundancy:
Parts of an image are
highly correlated.

In other words, we can predict a given pixel from its neighbor.

#### Run Length Coding



Line 100: (1,63) (0,87) (1,37) (0,5) (1,4) (0,556) (1,62) (0,210)

Total 12166 runs, each run use 11 bits → Total = 133826 Bits

#### Psychovisual Redundancy

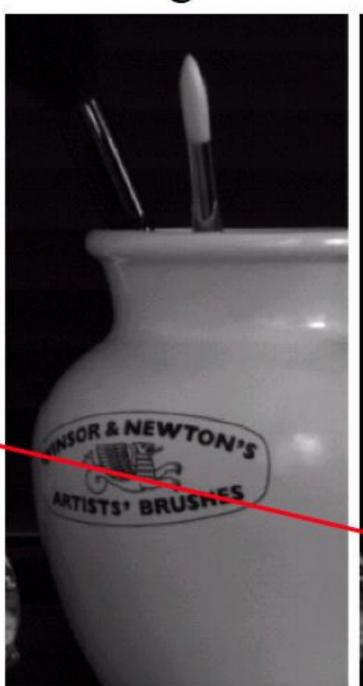
- Certain information simply has less relative importance than other information in normal visual processing
- This information is said to be psychovisually redundant
  - can be eliminated without significantly impairing the quality of image perception

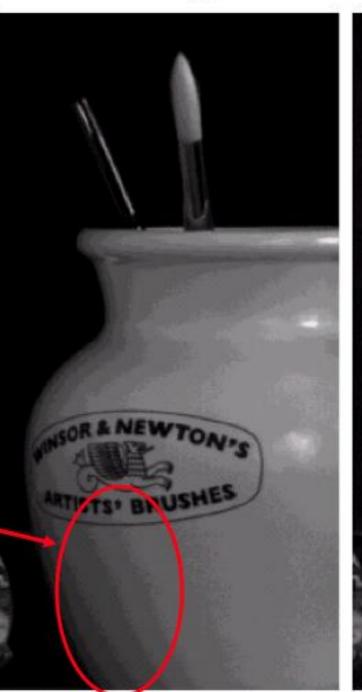
#### Psychovisual Redundancy

8-bit gray scale image

4-bit gray scale image

4-bit IGS image







False contours

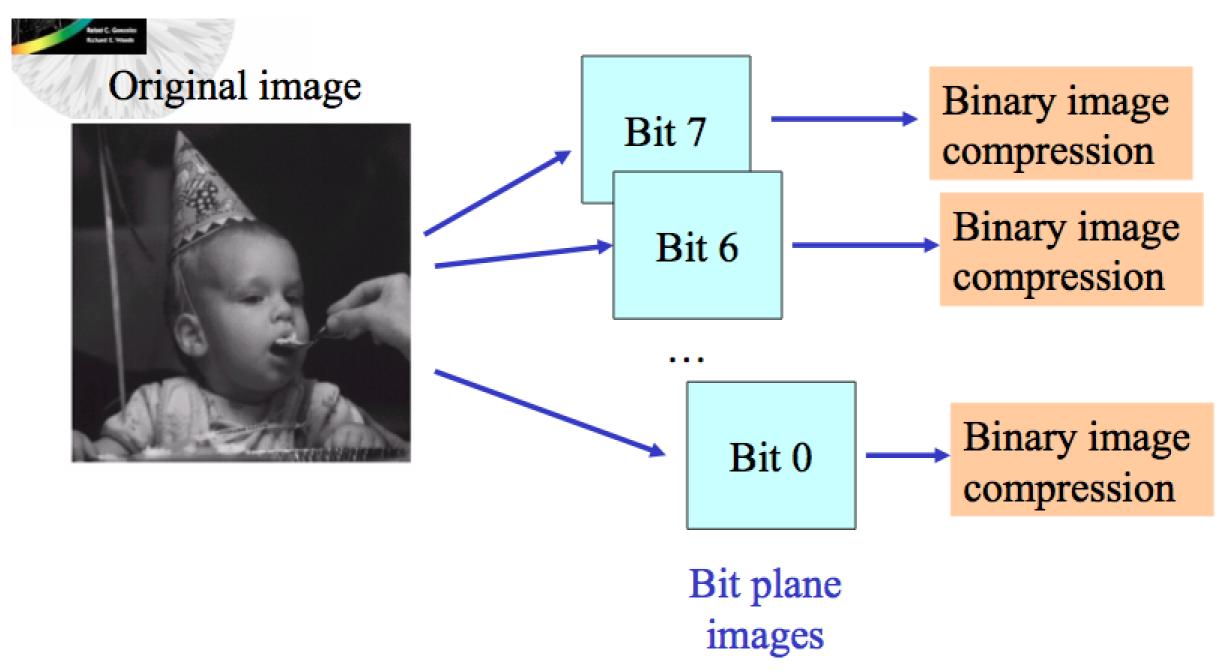
#### Psychovisual Redundancy

| Pixel | Gray Level | Sum       | IGS Code |
|-------|------------|-----------|----------|
| i - 1 | N/A        | 0000 0000 | N/A      |
| i     | 01101100   | 01101100  | 0110     |
| i + 1 | 1000 1011  | 1001 0111 | 1001     |
| i + 2 | 1000 0111  | 1000 1110 | 1000     |
| i + 3 | 1111 0100  | 1111 0100 | 1111     |

TABLE 8.2 IGS quantization procedure.

A sum-initially to zero-is first formed from the current 8-bit grey-level value and the four least significant bits of a previously generated sum. If the 4 most significant bits of the current value are 1111<sub>2</sub>, however, 0000<sub>2</sub> is added instead. The 4 most significant bits of the resulting sum are used as the coded pixel value.

#### Bit Plane Coding



Example of binary image compression: Run length coding

#### Bit Plane Coding





Original gray scale image



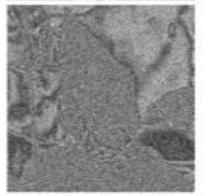
Bit 6



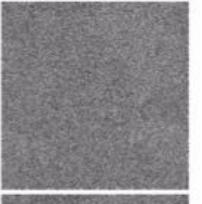
Bit 5



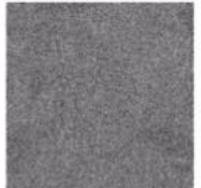
Bit 3



Bit 2



Bit 1

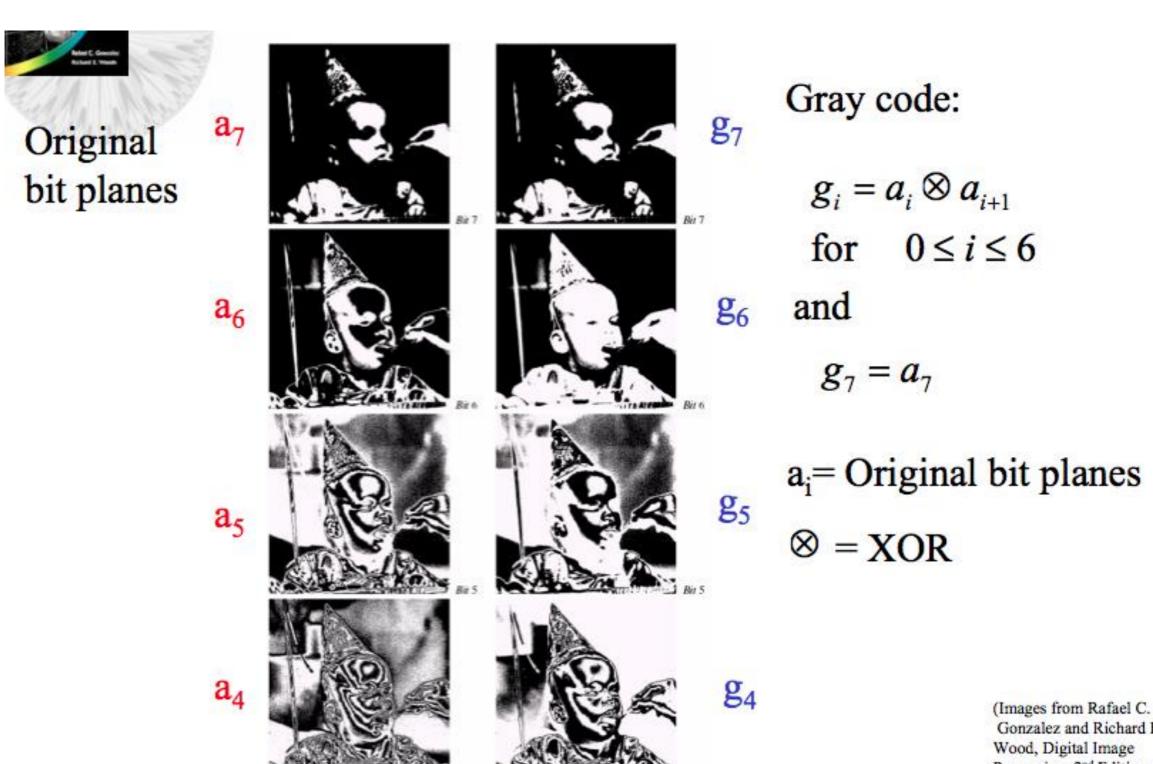


Bit 0

Bit 4

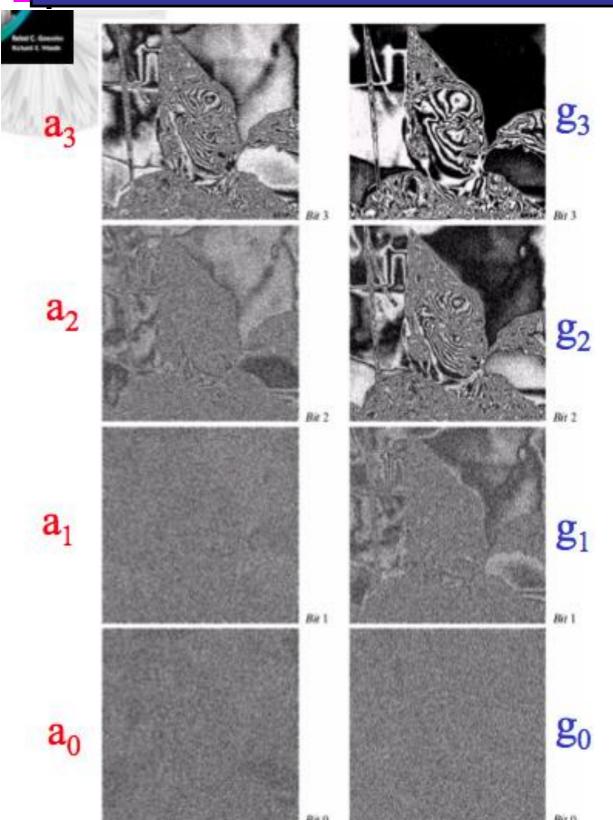
mages from Rafael C. Fonzalez and Richard E. Food, Digital Image rocessing, 2nd Edition.

#### Gray Coded Bit Planes



Gonzalez and Richard E. Processing, 2nd Edition.

#### Gray Coded Bit Planes

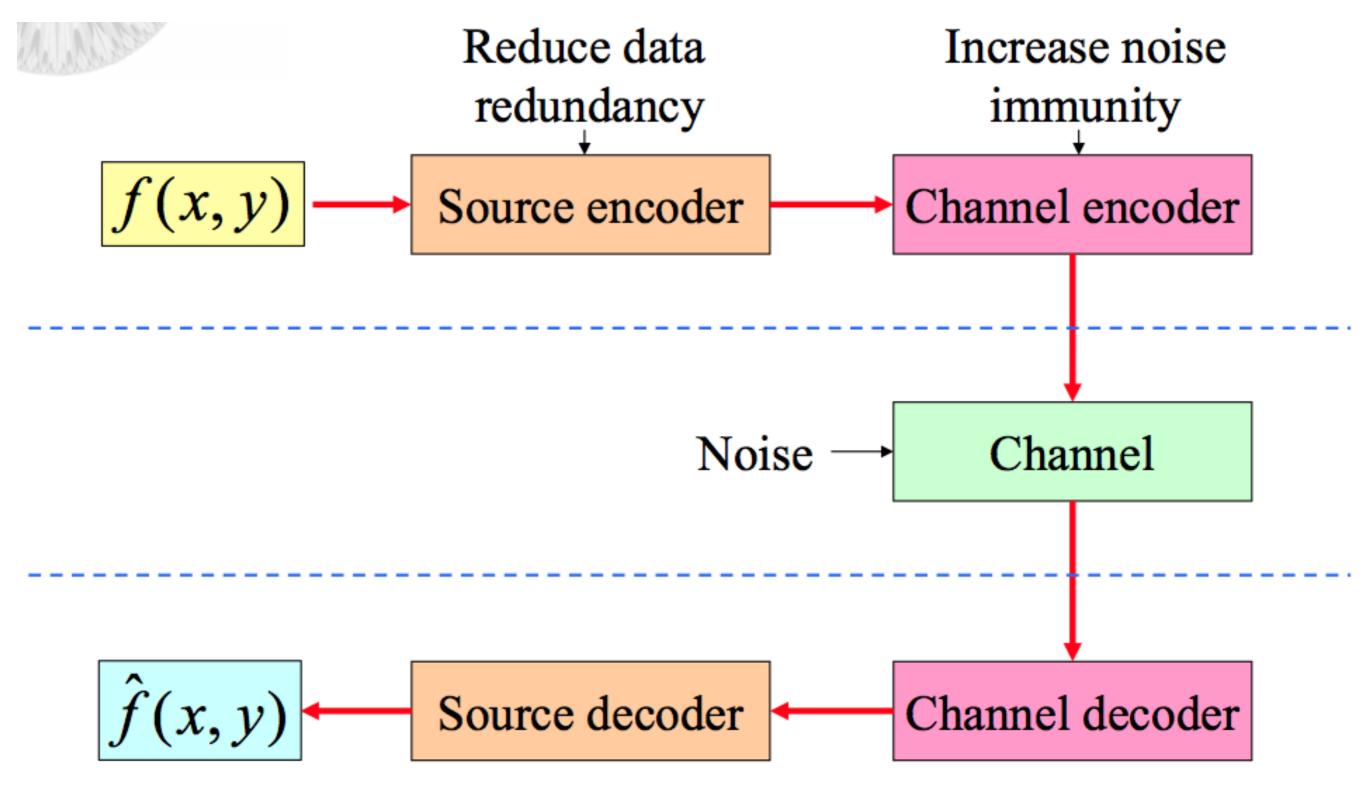


There are less 0-1 and 1-0 transitions in grayed code bit planes.

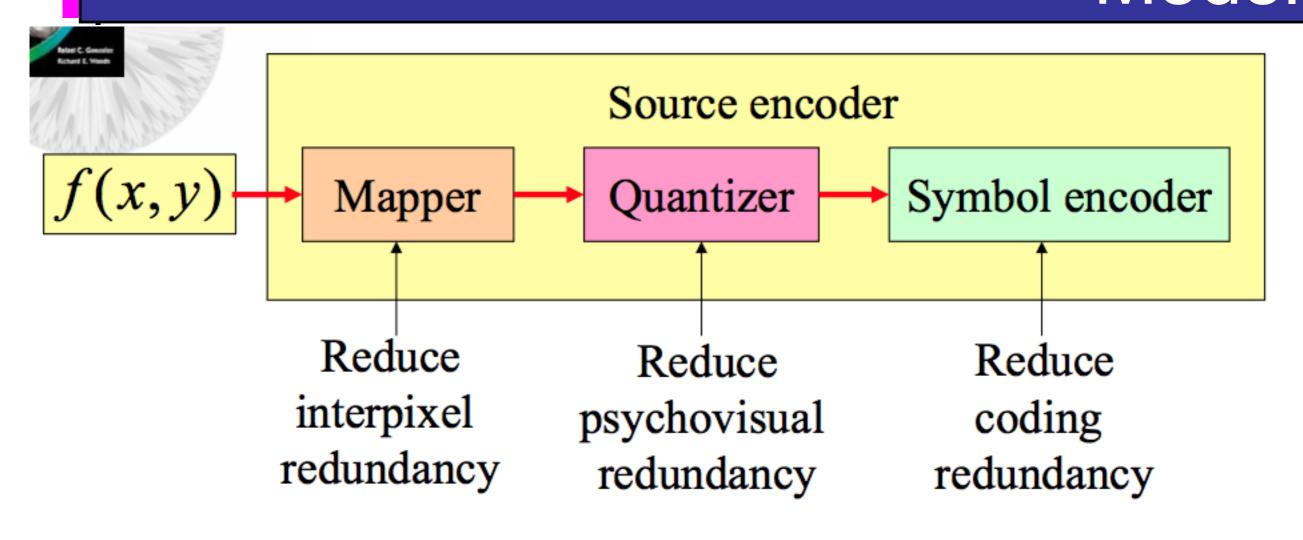
Hence gray coded bit planes are more efficient for coding.

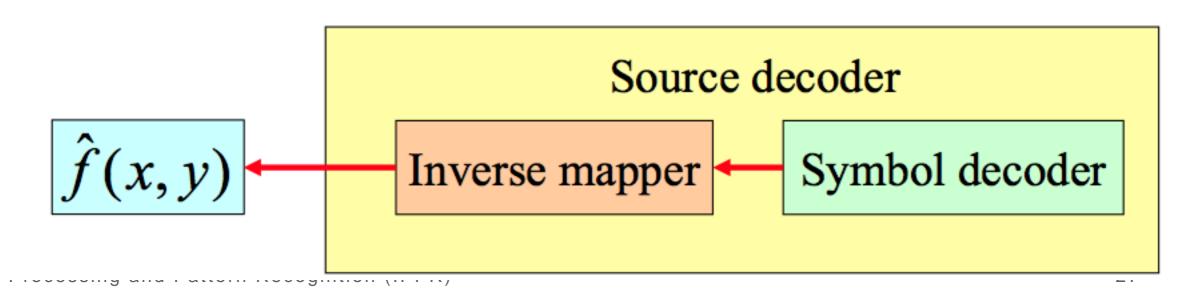
(Images from Rafael C. Gonzalez and Richard E. Wood, Digital Image Processing, 2<sup>nd</sup> Edition.

#### Image Compression Models



## Source Encoder and Decoder Models

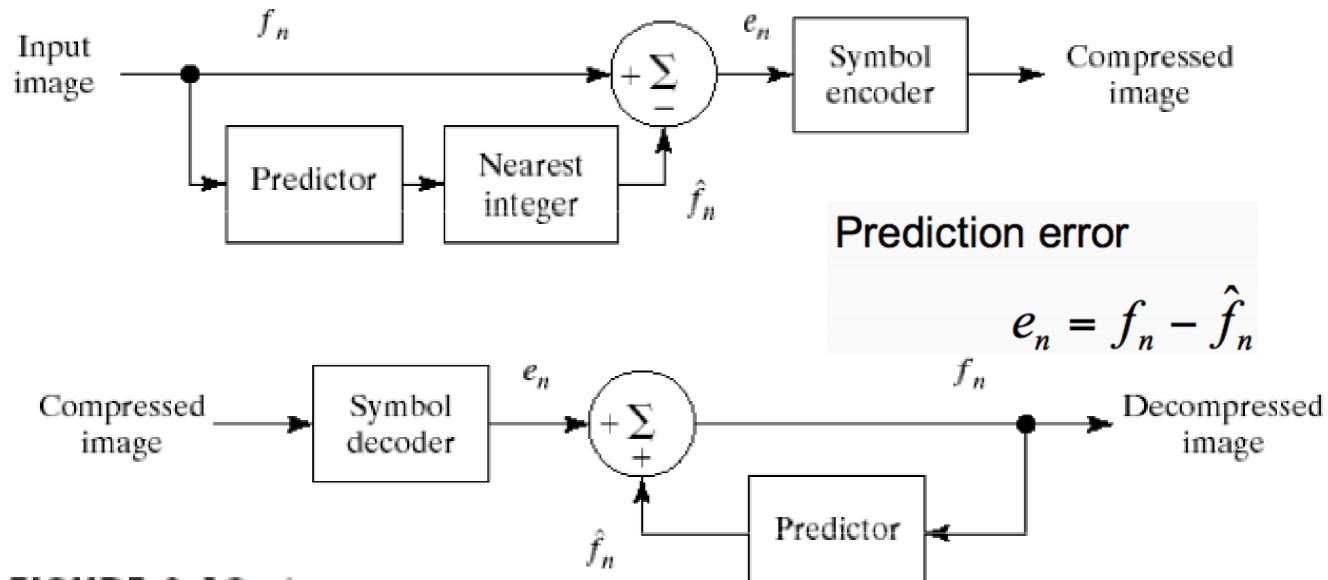




#### Error-Free Compression (Lossless)

- Error-free compression means lossless compression
- In numerous applications error-free compression is the only acceptable means of data reduction.
- They normally provide compression ratios of 2 to 10.

## Lossless Predictive Coding Model



#### FIGURE 8.19 A

lossless predictive coding model:

- (a) encoder;
- (b) decoder.

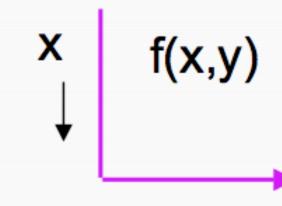
#### Prediction error: $e_n = f_n - f_n$ $e_n$ is coded using a variable length code

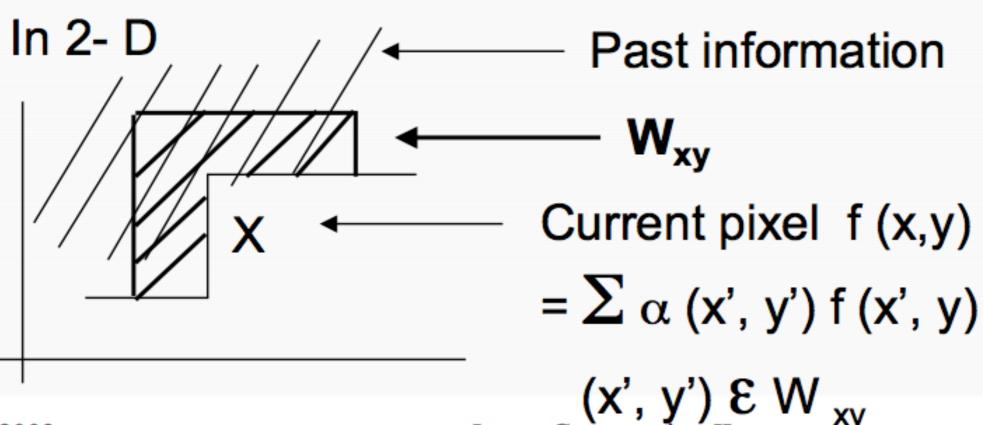
#### Lossless Predictive Coding Model

Example 1: 
$$\hat{\mathbf{f}}_{n} = \operatorname{Int}\left(\sum_{i=1}^{m} \alpha_{i} f_{n-i}\right)$$

 $\rightarrow$  Linear predictor; m = order of predictor

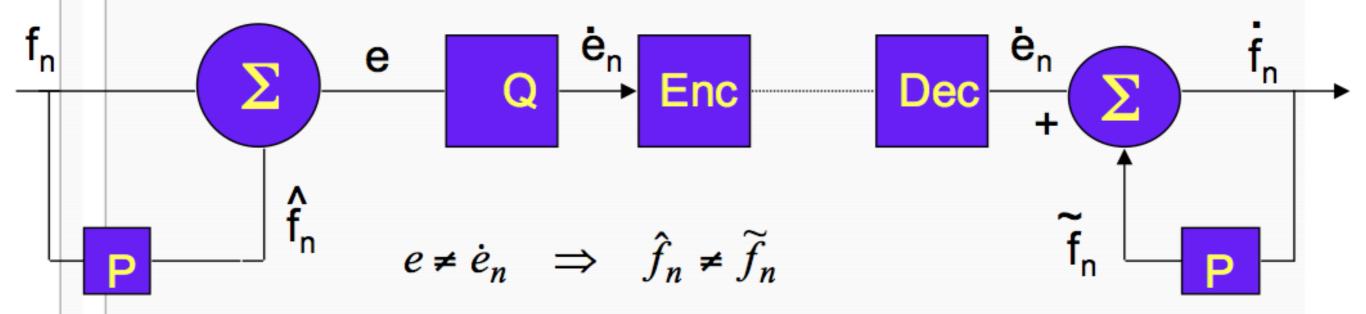
Example 2: 
$$\hat{f}_n(x,y)=Int\left(\sum_{i=1}^m \alpha_i f(x,y-i)\right)$$





 Unlike the error-free approaches, lossy encoding is based on the concept of compromising the accuracy of the reconstructed image in exchange for increased compression.

Lossy compression: uses a quantizer to compress further the number of bits required to encode the 'error'. First consider this:



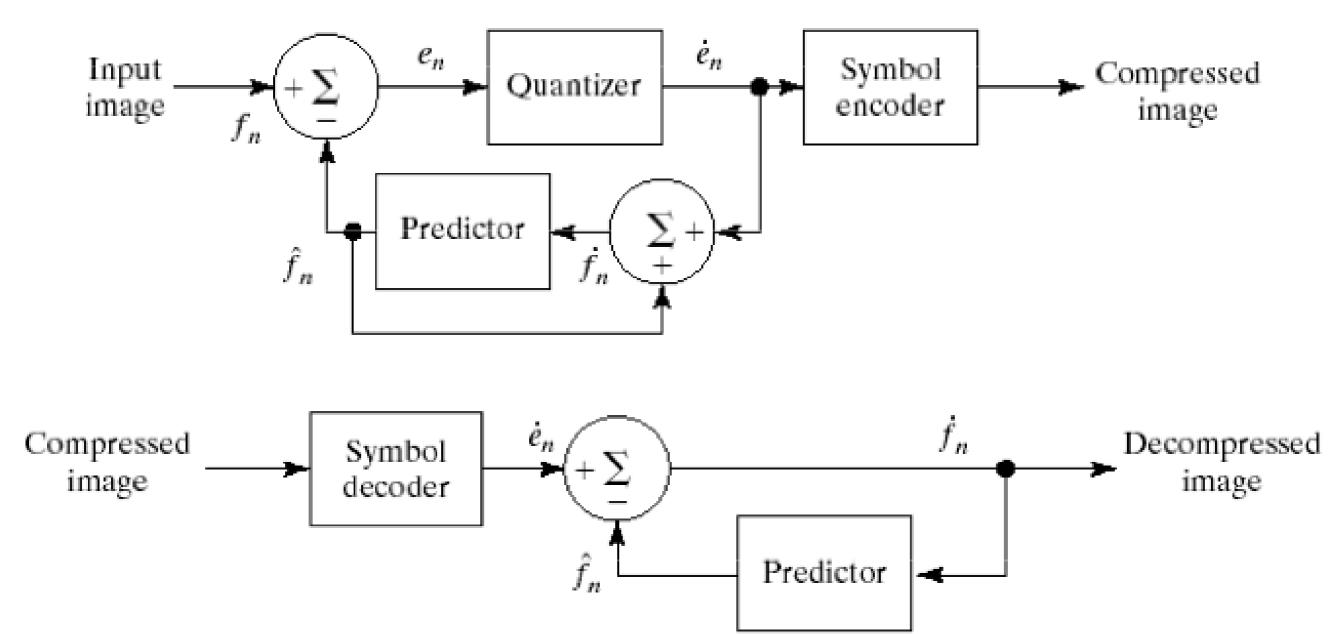
Notice that, unlike in the case of loss-less prediction, in lossy prediction the predictors P "see" different inputs at the encoder and decoder

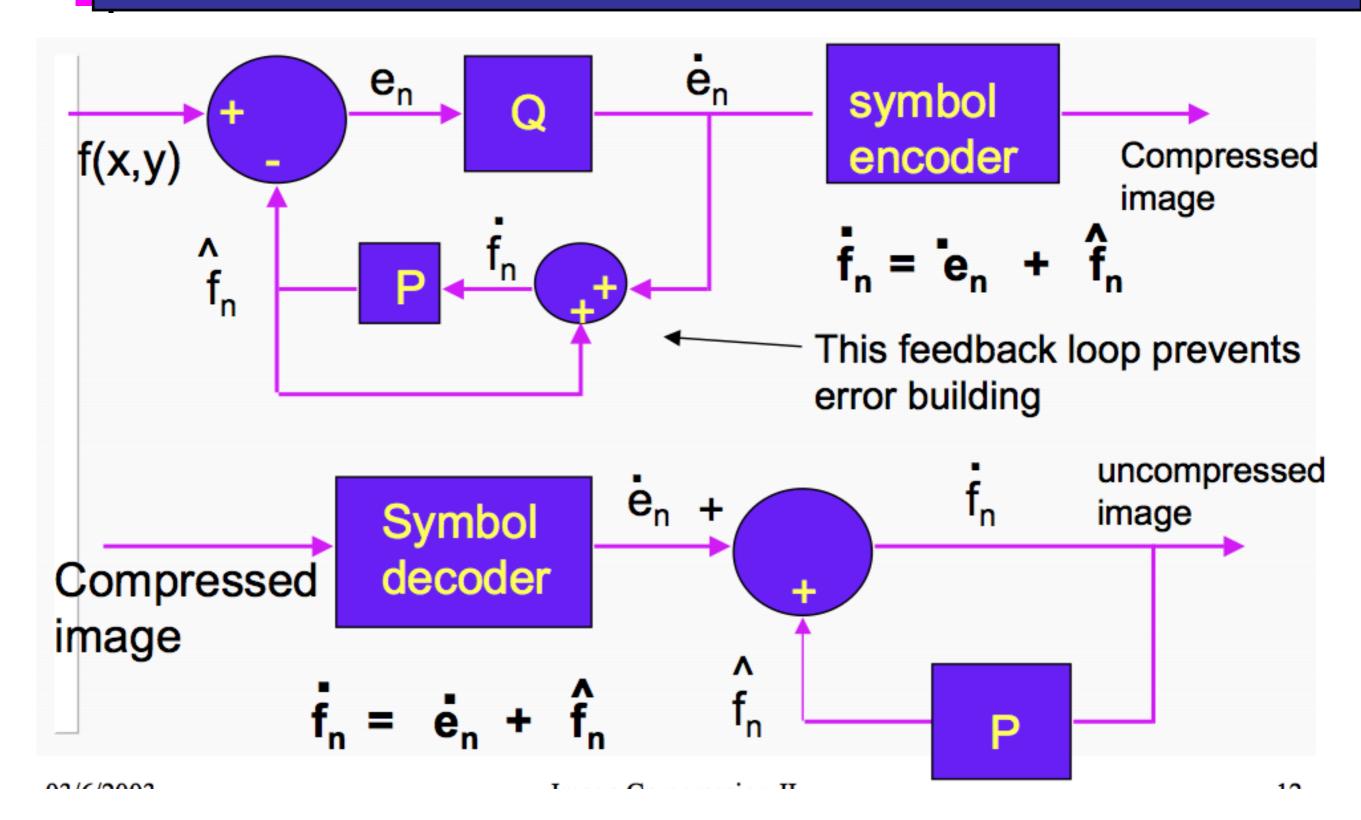
This results in a gradual buildup of error which is due to the quantization error at the encoder site.

In order to minimize this buildup of error due to quantization we should ensure that 'Ps' have the same input in both the cases.

| n | $f_n$ | $\hat{f}_n$ | $e_n$ | $\dot{e}_n$ | $\dot{f}_n$ |
|---|-------|-------------|-------|-------------|-------------|
| 0 | 0     | -           | -     | -           | 0           |
| 1 | 1     | 0           | 1     | 2           | 2           |
| 2 | 2     | 1           | 1     | 2           | 4           |
| 3 | 3     | 2           | 1     | 2           | 6           |
| 4 | 4     | 3           | 1     | 2           | 8           |
| 5 | 5     | 4           | 1     | 2           | 10          |
| 6 | 6     | 5           | 1     | 2           | 12          |

$$\dot{f}_n = \dot{e}_n + \tilde{f} = \dot{e}_n + \dot{f}_{n-1}$$





#### **Example:**

$$\hat{f}_n = \alpha \, \dot{f}_{n-1}$$
and 
$$\dot{e}_n = \begin{cases} +\xi & e_n > 0 \\ -\xi & e_n < 0 \end{cases}$$

 $0 < \alpha < 1$  prediction coefficient

$$\dot{f}_n = \dot{e}_n + \hat{f}_n$$
$$= \dot{e}_n + \alpha \, \dot{f}_{n-1}$$

Note: The quantizer used here is-- floor  $(e_n/2)*2$ . This is different from the one used in the earlier example. Note that this would result in a worse response if used without Feedback (output will be flat at "0").

| n | $f_n$ | $\hat{f}_n$ | $e_n$ | $\dot{e}_n$ | $\dot{f}_n$ |
|---|-------|-------------|-------|-------------|-------------|
| 0 | 0     | -           | -     | -           | 0           |
| 1 | 1     | 0           | 1     | 0           | 0           |
| 2 | 2     | 0           | 2     | 2           | 2           |
| 3 | 3     | 2           | 1     | 0           | 2           |
| 4 | 4     | 2           | 2     | 2           | 4           |

$$\dot{f}_n = \dot{e}_n + \hat{f}_n$$

$$\hat{f}_n = \dot{f}_{n-1}$$
 is used here for computation

{14, 15, 14, 15, 13, 15, 15, 14, 20, 26, 27, 28, 27, 27, 29, 37, 37, 62, 75, 77, 78, 79, 80, 81, 81, 82, 82}

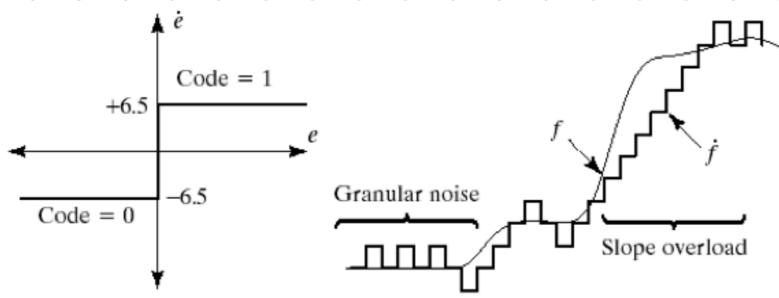
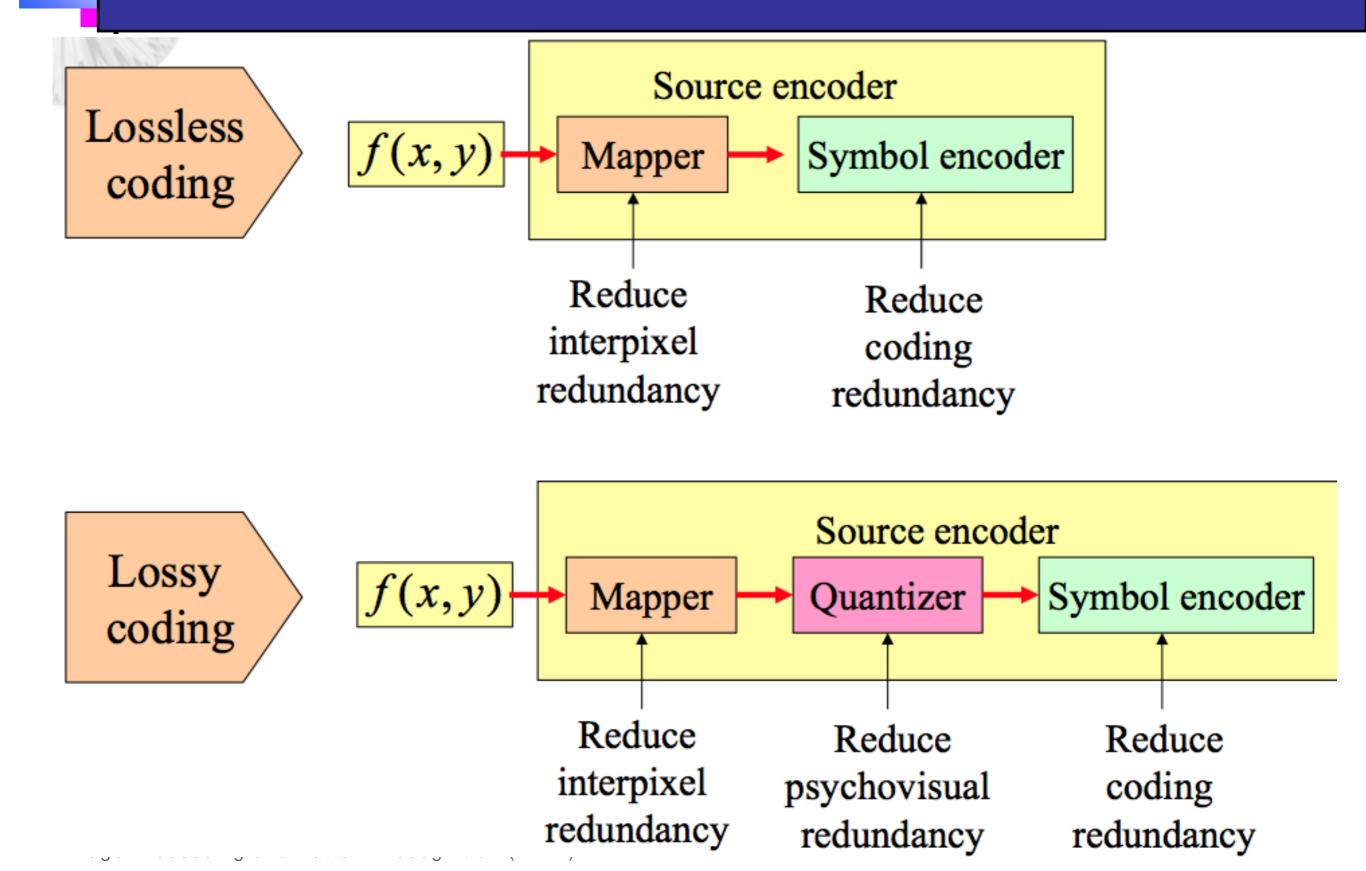


FIGURE 8.22 An example of delta modulation.

| Input            |  |  | Encoder   |   |  | Decoder                     | Error         |
|------------------|--|--|---|---|--|-----------------------------|---------------|
| n                | f  | $\hat{f}$  | e   | ė   | f  | $\hat{f}$ $\dot{f}$         | $[f-\dot{f}]$ |
| 0<br>1<br>2<br>3 | 14<br>15<br>14<br>15<br>29<br>37<br>47<br>62<br>75<br>77 | 20.5<br>14.0<br>20.5<br>14.0<br>20.5<br>27.0<br>33.5<br>40.0<br>46.5<br>53.0 | 1.0<br>-6.5<br>1.0<br>8.5<br>10.0<br>13.5<br>22.0<br>28.5<br>24.0 | 6.5<br>-6.5<br>6.5<br>6.5<br>6.5<br>6.5<br>6.5<br>6.5 | 14.0<br>20.5<br>14.0<br>20.5<br>27.0<br>33.5<br>40.0<br>46.5<br>53.0<br>59.6 | — 14.0 20. 20.5 14. 14.0 20 | 5             |
| ÷                | :  |  | :   | ·   | :  |                             |               |

## Lossless Vs. Lossy Coding

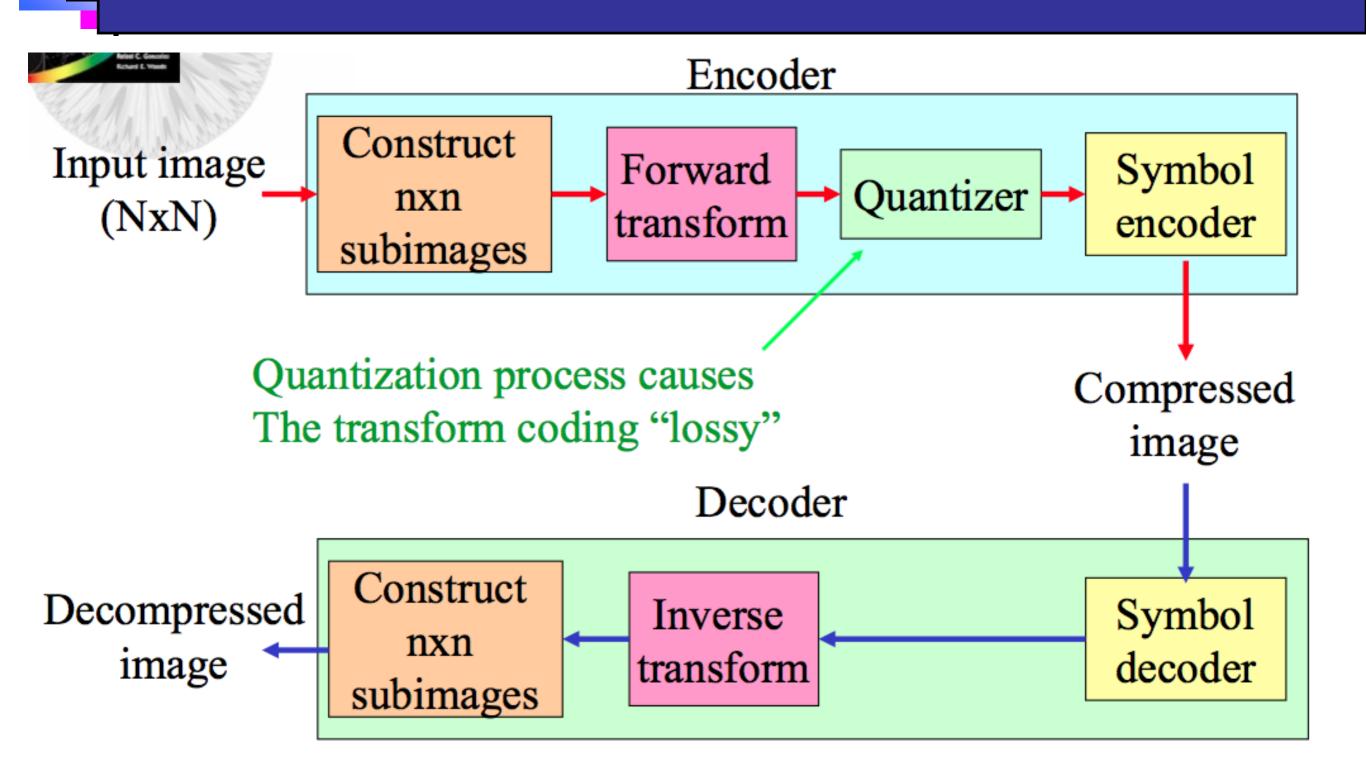


# Lossless Vs. Lossy Coding

|            | DATA COMPRESSION                            |  |  |  |  |
|------------|---|--|--|--|--|
| FACTORS    | LOSSLESS COMPRESSION                        | LOSSY COMPRESSION  |  |  |  |
| Definition | Lossless compression is a class of data     | Lossy compression is the class of data encoding methods      |  |  |  |
|            | compression algorithms that allow the       | that uses inexact approximations to represent the content.   |  |  |  |
|            | original data to be perfectly reconstructed | These techniques are used to reduce the data size for        |  |  |  |
|            | from the compressed data <sup>[7]</sup> .   | storage, handling, and transmitting content <sup>[8]</sup> . |  |  |  |
| Algorithm  | RLW, LZW, Arithmetic encoding, Huffman      | Transform coding, DCT, DWT, Fractal compression,             |  |  |  |
|            | coding, Shannon-Fano coding                 | Rectangle Segmentation and Sparse Matrix Storage             |  |  |  |
|            |   | (RSSMS).   |  |  |  |
| Uses       | Text or programs, images and sound          | Images, audio and video.                                     |  |  |  |
| Images     | RAW, BMP, and PNG are all Lossless          | JPEG and GUI are lossy image formats.                        |  |  |  |
|            | formats.                                    |  |  |  |  |
| Audio      | WAV, FLAC, and ALAC are all Lossless        | MP3, MP4, and OGG are lossy audio formats.                   |  |  |  |
|            | formats.                                    |  |  |  |  |
| Video      | Few lossless video formats are in common    | Common formats like H-264, MKV, and WMV are all              |  |  |  |
|            | consumer use, they would result in video    | lossy. H-264 can provides smaller files with higher          |  |  |  |
|            | files taking up a huge amount of space.     | qualities than previous generation of video codec because    |  |  |  |
|            |   | it has a "smaller" algorithm that's better at choosing the   |  |  |  |
|            |   | data to throw out.   |  |  |  |
| Advantages | It maintains quality.                       | It can make a multimedia file much smaller than its          |  |  |  |
|            | Conversion in any other format possible     | original size.It can reduce file sizes much more than        |  |  |  |
|            | without loss of audio information.          | lossless compression.  |  |  |  |
| Disadvanta | It doesn't reduce the file size as much as  | Conversion to another format only with loss of audio         |  |  |  |
| ges        | lossy compression. Lossless encoding        | information.It cannot be used in all types of files because  |  |  |  |
|            | technique cannot achieve high levels of     | it works by removing data. Text and data cannot be           |  |  |  |
|            | compression.                                | compressed because they do not have redundant                |  |  |  |
| Р          |   | information.   |  |  |  |

Image F

### Transform Coding



Examples of transformations used for image compression: DFT and DCT

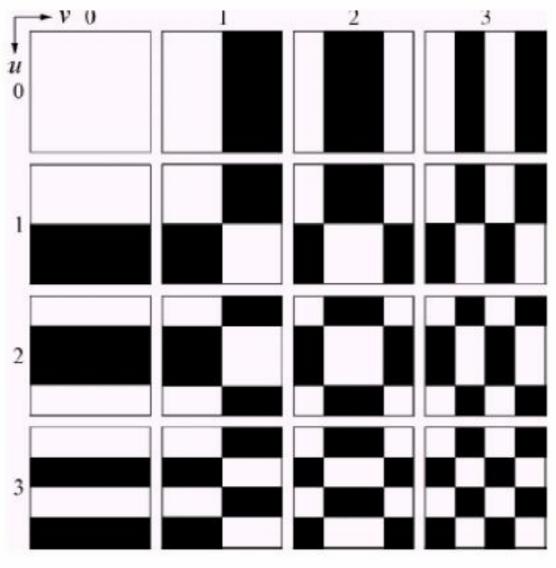
#### Transform Coding

- Parameters that effect transform coding performance:
  - 1. Type of transformation
  - 2. Size of subimage
  - 3. Quantization algorithm

#### Hadamard Transform



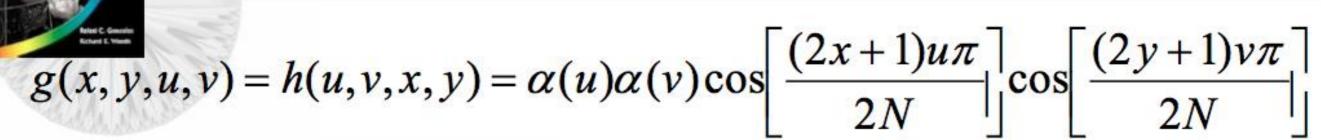
$$g(x, y, u, v) = h(u, v, x, y) = \frac{1}{N} (-1)^{\sum_{i=0}^{m-1} \lfloor b_i(x) p_i(u) + b_i(y) p_i(v) \rfloor}$$

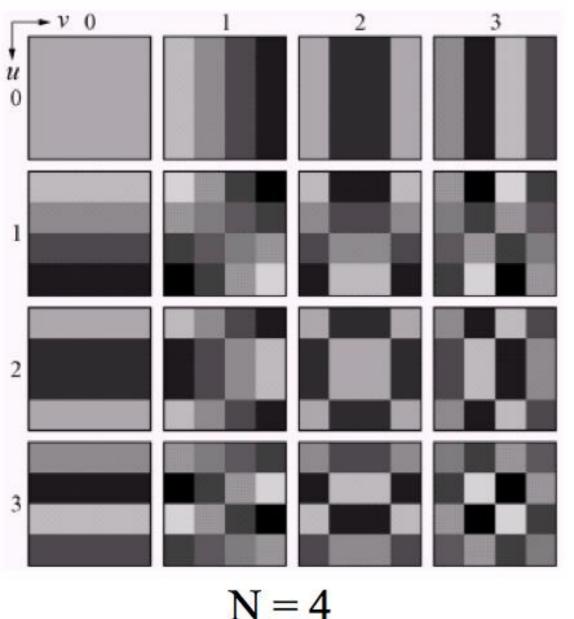


$$N = 4$$

Advantage: simple, easy to implement Disadvantage: not good packing ability

#### Discrete Cosine Transform





$$\alpha(u) = \begin{cases} \sqrt{\frac{1}{N}} & \text{for } u = 0 \\ \sqrt{\frac{2}{N}} & \text{for } u = 1, \dots, N-1 \end{cases}$$

DCT is one of the most frequently used transform for image compression. For example, DCT is used in JPG files.

Advantage: good packing ability, modulate computational complexity

#### Transform Coding Example



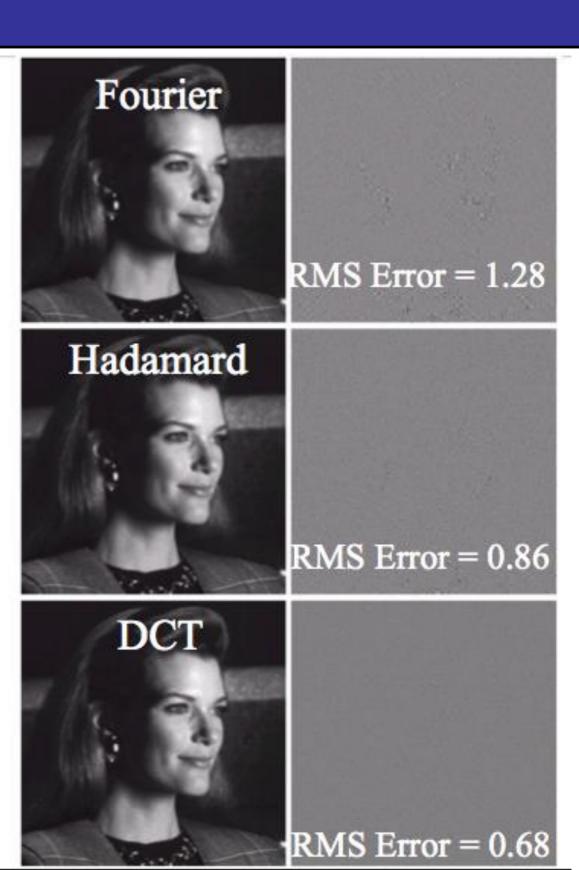


Original image 512x512 pixels

Subimage size: 8x8 pixels = 64 pixels

Quatization by truncating 50% of coefficients (only 32 max cofficients are kept.)

(Images from Rafael C. Gonzalez and Richard E. Wood, Digital Image Processing, 2<sup>nd</sup> Edition.



### Transform Coding Example

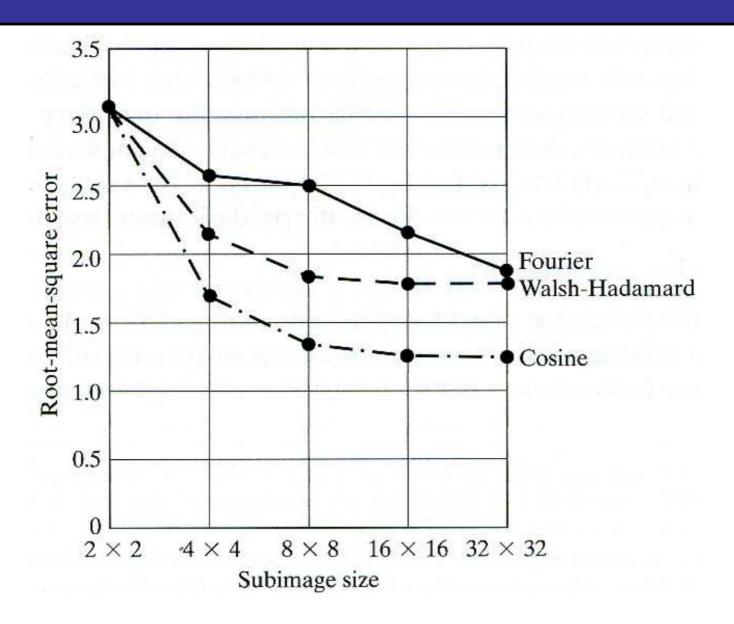
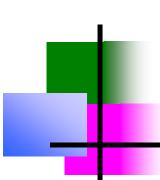


Figure 8.33 illustrates graphically the impact of subimage size on transform coding reconstruction error. The data plotted were obtained by dividing the monochrome image of Fig. 8.23 into subimages of size  $n \times n$ , for n = 2, 4, 8, 16, and 32, computing the transform of each subimage, truncating 75% of the resulting coefficients, and taking the inverse transform of the truncated arrays.



### Thank you !!!