
Nonlinear Programming

to accompany
Operations Research: Applications and Algorithms
4th edition, Chapter 11
by Wayne L. Winston

Review of Differential Calculus

- The equation: $\lim_{x \rightarrow a} f(x) = c$
means that as x gets closer to a (but not equal to a), the value of $f(x)$ gets arbitrarily close to c .
- A function $f(x)$ is **continuous** at point a if $\lim_{x \rightarrow a} f(x) = f(a)$.

If $f(x)$ is not continuous at $x=a$, we say that $f(x)$ is **discontinuous** (or has a discontinuity) at a .

- The **derivative** of a function $f(x)$ at $x = a$ (written $f'(a)$) is

defined to be $\lim_{\Delta x \rightarrow 0} \frac{f(a+\Delta x) - f(a)}{\Delta x}$

$f'(a)$: the slope of $f(x)$ at $x=a$.

If $f'(a) > 0$, then $f(x)$ is increasing at $x=a$.

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- **n th-order derivatives, n th-order Taylor series expansion:**
for $0 \leq h \leq b - a$ and some number p between a and $a+h$

$$f(a + h) = f(a) + \sum_{i=1}^n \frac{f^{(i)}(a)}{i!} h^i + \frac{f^{(n+1)}(p)}{(n+1)!} h^{n+1}$$

given that $f^{(n+1)}(x)$ exists for every point on interval $[a, b]$

- The partial derivative of $f(x_1, x_2, \dots, x_n)$ with respect to the variable x_i is written

$$\frac{\partial f}{\partial x_i} = \lim_{\Delta x_i \rightarrow 0} \frac{f(x_1, \dots, x_i + \Delta x_i, \dots, x_n) - f(x_1, \dots, x_i, \dots, x_n)}{\Delta x_i}$$

- The second order partial derivatives: $\frac{\partial^2 f}{\partial x_i \partial x_j} = \frac{\partial^2 f}{\partial x_j \partial x_i}$

1. Introduction to NLP

- A general **nonlinear programming problem** (NLP) can be expressed as follows:

Find the values of decision variables x_1, x_2, \dots, x_n that

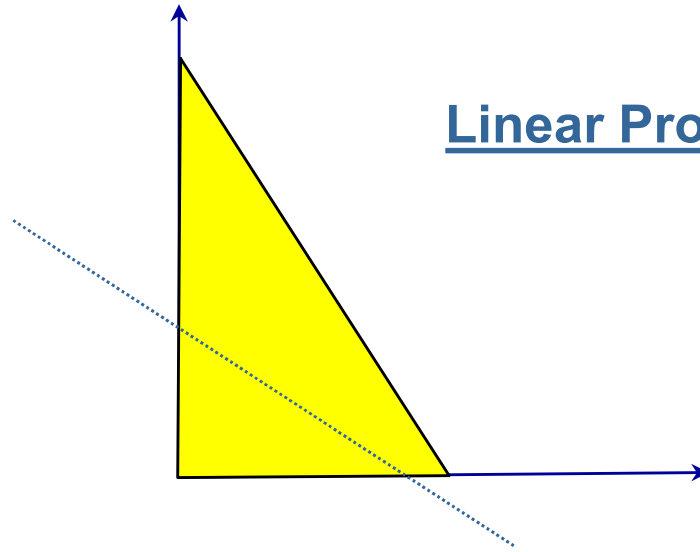
$$\begin{array}{ll} \max & (\text{or min}) \ z = f(x_1, x_2, \dots, x_n) \\ \text{s.t.} & g_1(x_1, x_2, \dots, x_n) (\leq, =, \text{ or } \geq) b_1 \\ & g_2(x_1, x_2, \dots, x_n) (\leq, =, \text{ or } \geq) b_2 \\ & \vdots \\ & g_m(x_1, x_2, \dots, x_n) (\leq, =, \text{ or } \geq) b_m \end{array}$$

x_j : continuous variable

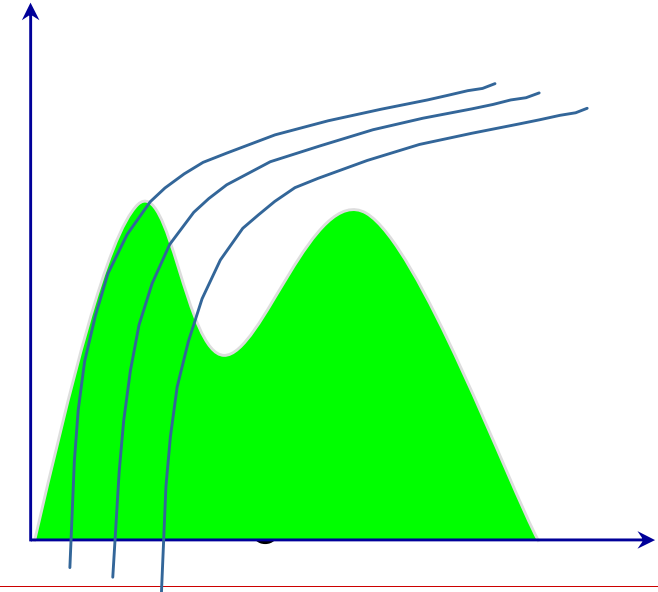
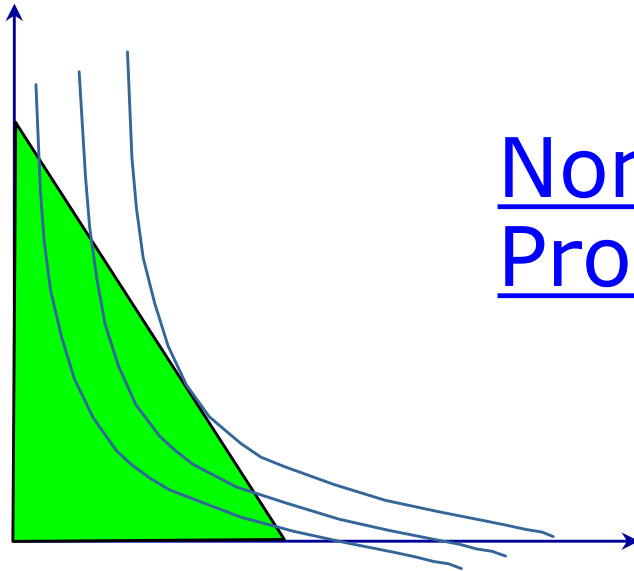
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- As in linear programming, $f(x_1, x_2, \dots, x_n)$ is the NLP's **objective function**, and $g_1(x_1, x_2, \dots, x_n) (\leq, =, \text{ or } \geq) b_1, \dots, g_m(x_1, x_2, \dots, x_n) (\leq, =, \text{ or } \geq) b_m$ are the NLP's **constraints**.
 - An NLP with no constraints is an **unconstrained NLP**.
 - Unconstrained LP?
 - The **feasible region** for NLP above is the set of points (x_1, x_2, \dots, x_n) that satisfy the m constraints in the NLP. A point in the feasible region is a *feasible point*, and a point that is not in the feasible region is an *infeasible point*.

Difficulties of NLP Models

Linear Program:



Nonlinear Programs:



Example: Profit Maximization considering price-demand relation

- It costs $\$c/\text{unit}$ to produce a product.
- Demand (denoted as D) is often modeled as a function of price (denoted as p). For example, $D(p) = 1 - bp$ (linear relation) for a parameter b . The function and its parameters can be determined by statistics regression.
- \rightarrow nonlinear in profit: $(p - c) \cdot D(p)$
- To maximize profit, what is the price and how much should be produced (to satisfy the demand).
- NLP:
$$\begin{aligned} \max \quad & z = (p - c)D(p) \\ \text{s.t.} \quad & 0 \leq p \leq 10 \end{aligned}$$
- Excel file: some examples

Example: Tire Production

- Firerock produces rubber used for tires by combining three ingredients: rubber, oil, and carbon black.
- Costs (cents/pound): rubber (4), oil (1), carbon black (7).
- The rubber used in automobile tires must have
 - ☐ a hardness of between 25 and 35
 - ☐ an elasticity of at least 16
 - ☐ a tensile strength of at least 12
- To manufacture a set of four automobile tires, 100 pounds of product is needed.
- The rubber to make a set of tires must contain between 25 and 60 pounds of rubber and at least 50 pounds of carbon black.

Example (cont'd)

- Define decision variables:

R = pounds of rubber in mixture used to produce four tires

O = pounds of oil in mixture used to produce four tires

C = pounds of carbon black used to produce four tires

- *Statistical analysis* has shown that the hardness, elasticity, and tensile strength of a 100-pound mixture of rubber, oil, and carbon black is

$$\text{Tensile Strength } (TS) = 12.5 - .10(O) - .001(O)^2$$

$$\text{Elasticity } (E) = 17 + .35R - .04(O) - .002(O)^2$$

$$\text{Hardness } (H) = 34 + .10R + .06(O) - .3(C) + .001(R)(O) + .005(O)^2 + .001C^2$$

- Formulate the NLP whose solution will tell Firerock how to minimize the cost of producing the rubber product needed to manufacture a set of automobile tires.

Example: continued

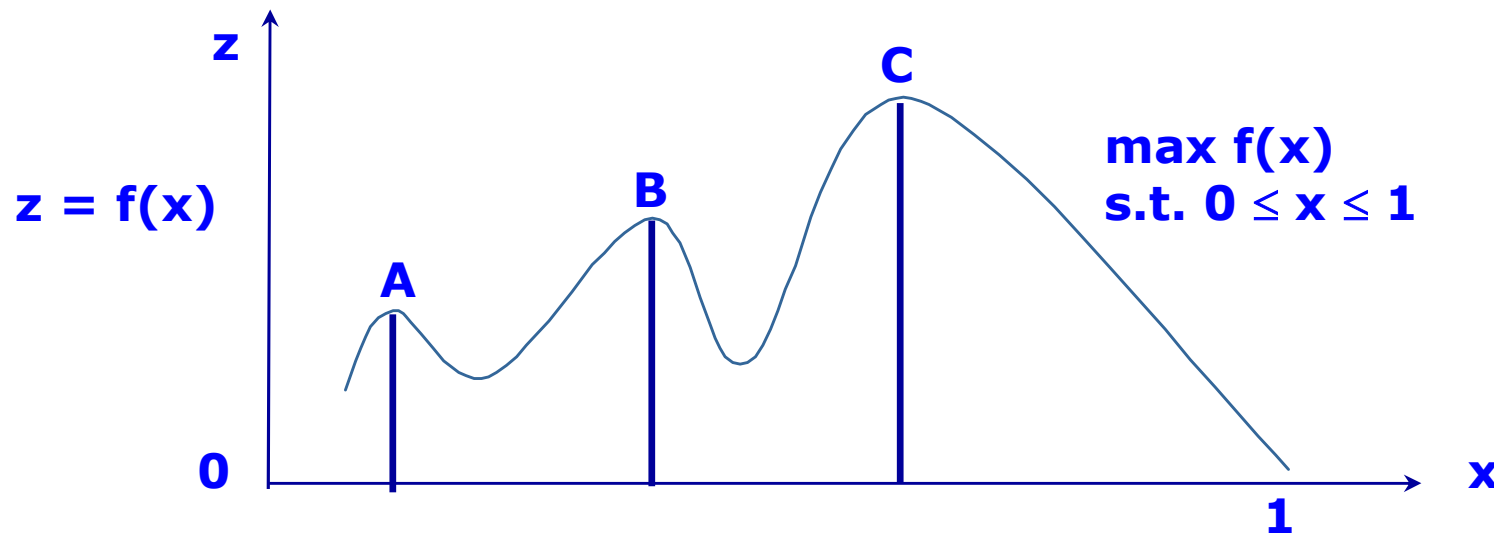
■ Min $4R + O + 7C$
s.t., $TS = 12.5 - .10(O) - .001(O)^2 \geq 12$
 $E = 17 + .35R - .04(O) - .002(O)^2 \geq 16$
 $H = 34 + .10R + .06(O) - .3(C) + .001(R)(O) + .005(O)^2 + .001C^2 \geq 25$
 $H = 34 + .10R + .06(O) - .3(C) + .001(R)(O) + .005(O)^2 + .001C^2 \leq 35$

 $R + O + C = 100$
 $R \leq 60$
 $R \geq 25$
 $C \geq 50$
 $O \geq 0$

Local vs. Global Optima

Definition: Let x be a feasible solution, then

- x is a **global max** if $f(x) \geq f(y)$ for every feasible y .
- x is a **local max** if $f(x) \geq f(y)$ for every feasible y sufficiently close to x (i.e., $x_j - \varepsilon \geq y_j \geq x_j + \varepsilon$ for all j and some small ε).



There may be several locally optimal solutions.

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- If the NLP is a maximization problem, then any point $\bar{\mathbf{x}} = (\bar{x}_1, \bar{x}_2, \dots, \bar{x}_n)$ in the feasible region for which $f(\bar{\mathbf{x}}) \geq f(\mathbf{x})$ holds true for all points \mathbf{x} in the feasible region is an **optimal solution** to the NLP.
 - For any NLP (maximization), a feasible point $\mathbf{x}' = (x'_1, x'_2, \dots, x'_n)$ is a **local maximum** if for sufficiently small ε , any feasible point $\mathbf{x} = (x_1, x_2, \dots, x_n)$ having $|x_i - x'_i| < \varepsilon$ for $i = 1, 2, \dots, n$ satisfies $f(\mathbf{x}') \geq f(\mathbf{x})$.

Relations of local and global optima

- For NLPs having multiple local optimal solutions, the Solver may fail to find the optimal solution because it may pick a local optima that is not a global optima.
- NLPs can be solved with LINGO or Excel Solver. However, in general, there is no guarantee that the solution found by them is optimal.

$$\max z = (x-1)(x-2)(x-3)(x-4)(x-5)$$

$$\text{s.t. } x \geq 1$$

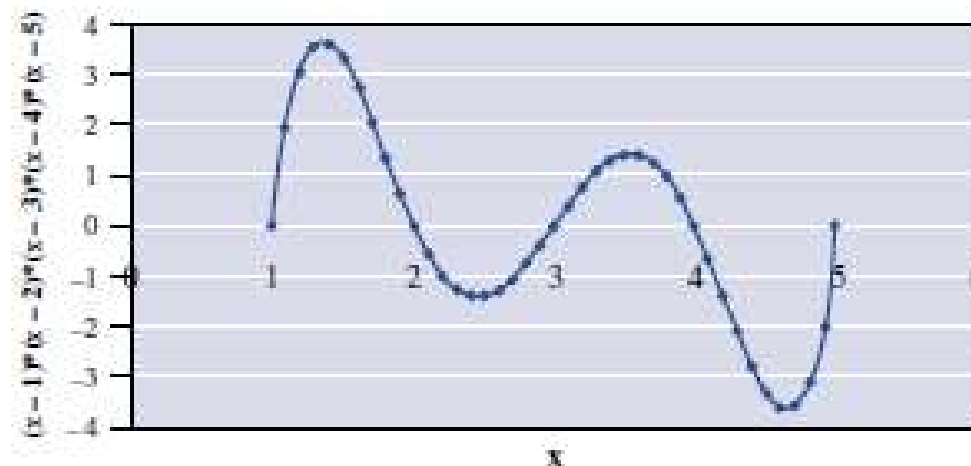
$$x \leq 5$$

Different initial values \Rightarrow ??

In LINGO: INIT:

$x=2$;

ENDINIT



- **When is a locally optimal solution also globally optimal?** Then, LINGO will find the optimal solution to an NLP. ---- Convexity

2 Convex and Concave Functions

- **Theorem**: Consider a general NLP. Suppose **the feasible region S for NLP is a convex set**. If $f(\mathbf{x})$ is **concave (convex)** on S , then any local maximum (minimum) for the NLP is an optimal solution (global optima) to the NLP.

Convex set and convex and concave functions

- Convex set: see (7-convexity-local and global optima.ppt) on convex set
- A function $f(x_1, x_2, \dots, x_n)$ is a **convex function** on a convex set S if for any $\mathbf{x}' \in S$ and $\mathbf{x}'' \in S$
$$f[c\mathbf{x}' + (1-c)\mathbf{x}''] \leq cf(\mathbf{x}') + (1-c)f(\mathbf{x}'')$$
holds for $0 \leq c \leq 1$.
- A function $f(x_1, x_2, \dots, x_n)$ is a **concave function** on a convex set S if for any $\mathbf{x}' \in S$ and $\mathbf{x}'' \in S$
$$f[c\mathbf{x}' + (1-c)\mathbf{x}''] \geq cf(\mathbf{x}') + (1-c)f(\mathbf{x}'')$$
holds for $0 \leq c \leq 1$.

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- A function $f(x_1, x_2, \dots, x_n)$ is a convex function iff - $f(x_1, x_2, \dots, x_n)$ is a concave function, and conversely.
 - The sum of two convex functions is convex and the sum of two concave functions is concave.
 - A linear function is both convex and concave.
 - more ...

 - Suppose that $f(x)$ is a function of a single variable and $f''(x)$ exists for all x in a convex set S . Then $f(x)$ is a convex (concave) function of S if and only if $f''(x) \geq 0$ ($f''(x) \leq 0$) for all x in S . (single variable)

Suppose $f(x_1, x_2, \dots, x_n)$ has continuous second-order partial derivatives for each point $\mathbf{x}=(x_1, x_2, \dots, x_n)$ in a convex set S .

- $f(x_1, x_2, \dots, x_n)$ is a **convex function** on S if and only if for each $x \in S$, all principal minors of H are non-negative.
- $f(x_1, x_2, \dots, x_n)$ is a **concave function** on S if and only if for each $\mathbf{x} \in S$ and $k=1, 2, \dots, n$, all nonzero principal minors have the same sign as $(-1)^k$.
- ✓ The Hessian of $f(x_1, x_2, \dots, x_n)$ is the $n \times n$ matrix whose ij th entry is $\frac{\partial^2 f}{\partial x_i \partial x_j}$, denoted as $H(x_1, x_2, \dots, x_n)$
- ✓ An **i th principal minor** of an $n \times n$ matrix is the determinant of any $i \times i$ matrix obtained by deleting $n - i$ rows and the corresponding $n - i$ columns of the matrix.
- ✓ The **k th leading principal minor** of an $n \times n$ matrix is the determinant of the $k \times k$ matrix obtained by deleting the last $n - k$ rows and columns of the matrix. $H_k(x_1, x_2, \dots, x_n)$ is the k th leading principal minor of the hessian matrix evaluated at the point (x_1, x_2, \dots, x_n) .

Example 1

- $f(x_1, x_2) = x_1^3 + 2x_1x_2 + x_2^2$, then
 - $H(x_1x_2) = \begin{bmatrix} 6x_1 & 2 \\ 2 & 2 \end{bmatrix}$
 - Principal minors ($i=1, 2$): the first principle minors are $6x_1$ and 2 , the second principle minor is the determinant of $H(x_1x_2)$, which is $12x_1 - 4$.
 - Leading principal minors ($k=1, 2$): $H_1(x_1x_2) = 6x_1$ and $H_2(x_1x_2) = 12x_1 - 4$.

Example 2

- Show that $f(x_1, x_2) = -x_1^2 - x_1x_2 - 2x_2^2$ is a concave function on R^2 .
- We have $H(x_1, x_2) = \begin{bmatrix} -2 & -1 \\ -1 & -4 \end{bmatrix}$
- Principal minors ($i=1, 2$): the first principle minors are -2 and -4 . These are both nonpositive. The second principle minor is $7 > 0$. Thus, $f(x_1, x_2)$ is a concave function on R^2 .

3 Unconstrained NLPs with Several Variables

- Consider this unconstrained NLP

$$\begin{array}{ll} \max \text{ (or min)} & f(x_1, x_2, \dots, x_n) \\ \text{s.t.} & (x_1, x_2, \dots, x_n) \in R^n \end{array}$$

- Assume that the first and second partial derivatives of $f(\mathbf{x})$ exist and are continuous at all points.
- A point $\bar{\mathbf{x}}$ having $\frac{\partial f(\bar{\mathbf{x}})}{\partial x_i} = 0$ for $i = 1, 2, \dots, n$ is called a **stationary point** of f .

Single variable: stationary points

THEOREM 4

If $f'(x_0) = 0$ and $f''(x_0) < 0$, then x_0 is a local maximum. If $f'(x_0) = 0$ and $f''(x_0) > 0$, then x_0 is a local minimum.

THEOREM 5

If $f'(x_0) = 0$, and

- 1 If the first nonvanishing (nonzero) derivative at x_0 is an odd-order derivative [$f^{(3)}(x_0)$, $f^{(5)}(x_0)$, and so on], then x_0 is not a local maximum or a local minimum.
- 2 If the first nonvanishing derivative at x_0 is positive and an even-order derivative, then x_0 is a local minimum.
- 3 If the first nonvanishing derivative at x_0 is negative and an even-order derivative, then x_0 is a local maximum.

Multiple variables

- These theorems provide the basics of unconstrained NLP.
 - Necessary condition: If $\bar{\mathbf{x}}$ is a local optima, then $\frac{\partial f(\bar{\mathbf{x}})}{\partial x_i} = 0$ for $i = 1, 2, \dots, n$.
 - If $H_k(\bar{\mathbf{x}}) > 0, k=1, 2, \dots, n$, then a stationary point $\bar{\mathbf{x}}$ is a local minimum.
 - If, for $k=1, 2, \dots, n$, $H_k(\bar{\mathbf{x}})$ is nonzero and has the same sign as $(-1)^k$, then a stationary point $\bar{\mathbf{x}}$ is a local maximum.
 - If $H_n(\bar{\mathbf{x}}) \neq 0$ and the conditions of the previous two theorems do not hold, then a stationary point $\bar{\mathbf{x}}$ is not a local optima.
- If a stationary point \mathbf{x} is not a local extremum, then it is called a **saddle point**.
- If $H_n(\mathbf{x})=0$ for a stationary point \mathbf{x} , then \mathbf{x} may be a local minimum, a local maximum, or a saddle point, and the preceding tests are inconclusive.

Example 28

- Find all local maxima, local minima, and saddle points for

$$f(x_1, x_2) = x_1^2 x_2 + x_1 x_2^3 - x_1 x_2$$

4 The Method of Steepest Ascent

- The method of steepest ascent can be used to approximate a function's **stationary point** having $\nabla f(\mathbf{x}) = 0$ (candidates for optimal solutions).
- Given a vector $\mathbf{x} = (x_1, x_2, \dots, x_n) \in R^n$, the length of \mathbf{x} (written $||\mathbf{x}||$) is $||\mathbf{x}|| = (x_1^2 + x_2^2 + \dots + x_n^2)^{\frac{1}{2}}$
- For any vector \mathbf{x} , the unit vector $\mathbf{x}/||\mathbf{x}||$ is called the normalized version of \mathbf{x} .
- A **direction** can be represented by only one normalized vector.

- Consider a function $f(x_1, x_2, \dots, x_n)$, all of whose **partial derivatives exist at every point**.
- A **gradient vector** for $f(x_1, x_2, \dots, x_n)$, written $\nabla f(\mathbf{x})$, is

$$\nabla f(x) = \left[\frac{\partial f(x)}{\partial x_1}, \frac{\partial f(x)}{\partial x_2}, \dots, \frac{\partial f(x)}{\partial x_n} \right]$$

- Suppose we are at a point \mathbf{v} and we move from \mathbf{v} a small distance δ in a direction \mathbf{d} . Then for a given δ , the maximal increase in the value of $f(\mathbf{x})$ will occur if we choose

$$\mathbf{d} = \frac{\nabla f(\mathbf{v})}{\|\nabla f(\mathbf{v})\|}$$

- In other words, if we move a small distance away from \mathbf{v} and we want $f(\mathbf{x})$ to increase as quickly as possible, then we should move in the direction of $\nabla f(\mathbf{v})$.

Procedure of steepest ascent method

- Begin at any point \mathbf{v}_0 , and then move in the direction of $\nabla f(\mathbf{v}_0)$, a maximum rate of increase for f at \mathbf{v}_0 . For some nonnegative value of t_0 , we move to a point $\mathbf{v}_1 = \mathbf{v}_0 + t_0 \nabla f(\mathbf{v}_0)$.
- t_0 solves the following one-dimensional optimization problem:

$$\begin{array}{ll} \max & f(\mathbf{v}_0 + t_0 \nabla f(\mathbf{v}_0)) \\ \text{s. t.,} & t_0 \geq 0 \end{array}$$

This single-variable NLP may be solved by the methods using differentials or, if necessary, by a search procedure such as the Golden Section Search.

- If $\|\nabla f(\mathbf{v}_1)\|$ is sufficiently small (say, less than 0.01) (**termination condition**), we may terminate the algorithm with the knowledge that \mathbf{v}_1 is near a stationary point \mathbf{v}' having $\nabla f(\mathbf{v}') = 0$.

Example 29:

- Use the method of steepest ascent to approximate the solution to

$$\max z = -(x_1 - 3)^2 - (x_2 - 2)^2 = f(x_1, x_2)$$

$$s.t., (x_1, x_2) \in R^2$$

We arbitrarily choose to begin at the point $v_0 = (1, 1)$. Because $\nabla f(x_1, x_2) = (-2(x_1 - 3), -2(x_2 - 2))$, we have $\nabla f(1, 1) = (4, 2)$. Thus, we must choose t_0 to maximize

$$f(t_0) = f[(1, 1) + t_0(4, 2)] = f(1 + 4t_0, 1 + 2t_0) = -(-2 + 4t_0)^2 - (-1 + 2t_0)^2$$

Setting $f'(t_0) = 0$, we obtain

$$-8(-2 + 4t_0) - 4(-1 + 2t_0) = 0$$

$$20 - 40t_0 = 0$$

$$t_0 = 0.5$$

Our new point is $v_1 = (1, 1) + 0.5(4, 2) = (3, 2)$. Now $\nabla f(3, 2) = (0, 0)$, and we terminate the algorithm. Because $f(x_1, x_2)$ is a concave function, we have found the optimal solution to the NLP.

5 Constrained NLP – KKT Conditions

- A general NLP:

$$\begin{array}{ll}\text{max (or min)} & f(x_1, x_2, \dots, x_n) \\ \text{s.t.} & g_i(x_1, x_2, \dots, x_n) \leq 0 \text{ for } i = 1, 2, \dots, q \\ & g_i(x_1, x_2, \dots, x_n) = 0 \text{ for } i = q + 1, q + 2, \dots, m \\ & x_j \geq \mathbf{0}, \leq \mathbf{0}, \text{ or unrestricted for } j = 1, 2, \dots, n\end{array}$$

- Associate multipliers $\lambda_1, \lambda_2, \dots, \lambda_m$ with the constraints
- Construct Lagrangian function as

$$\begin{aligned} & L(x_1, x_2, \dots, x_n, \lambda_1, \lambda_2, \dots, \lambda_m) \\ & = f(x_1, x_2, \dots, x_n) - \sum_{i=1}^m \lambda_i g_i(x_1, x_2, \dots, x_n) \end{aligned}$$

- The **KKT conditions** are **necessary** for a feasible point $\bar{\mathbf{x}} = (\bar{x}_1, \bar{x}_2, \dots, \bar{x}_n)$ to solve the NLP. (Of course, meanwhile satisfying all the original constraints and sign restrictions).

5.1 The (Karush)-Kuhn-Tucker Conditions

■ KKT conditions:

$$\frac{\partial L}{\partial x_j} = \frac{\partial f(\bar{\mathbf{x}})}{\partial x_j} - \sum_{i=1}^m \bar{\lambda}_i \frac{\partial g_i(\bar{\mathbf{x}})}{\partial x_j} (\leq, \geq, \text{ or } =) 0 \quad (j = 1, 2, \dots, n)$$

$$\bar{\lambda}_i \cdot g_i(\bar{\mathbf{x}}) = 0 \quad (i = 1, 2, \dots, q) \quad \text{complementary conditions}$$

$$\bar{x}_j \cdot \frac{\partial L}{\partial x_j} = 0 \quad (j = 1, 2, \dots, n)$$

$$\bar{\lambda}_i (\geq 0 \text{ (max NLP)}, \leq 0 \text{ (min NLP)}) \quad (i = 1, 2, \dots, q)$$

$$\bar{\lambda}_i \text{ unrestricted} \quad (i = q + 1, \dots, m)$$

		Max NLP	Min NLP
Variable	$x_j \geq 0$	$\partial L / \partial x_j \leq 0$	$\partial L / \partial x_j \geq 0$
	$x_j \leq 0$	$\partial L / \partial x_j \geq 0$	$\partial L / \partial x_j \leq 0$
	x_j unrestricted	$\partial L / \partial x_j = 0$	$\partial L / \partial x_j = 0$

5.2 Sufficient conditions

- Consider a maximization (minimization) NLP as in the proceeding page 27. If $f(\mathbf{x})$ is a concave (convex) function and the feasible region formed by all the constraints is convex set, then any feasible point $\bar{\mathbf{x}}$ satisfying the necessary KKT conditions is an optimal solution.
- The feasible region defined by $g_i(\mathbf{x}) \leq 0$ is convex set if $g_i(\mathbf{x})$ is a convex function.
- If all the constraints are defined by convex functions in terms of \leq direction, the feasible region is convex set.
- The feasible region defined by linear constraint is convex set.

5.3 Special NLP – 1

- Simplify the KKT conditions for the following NLPs in which all the constraints are equality constraints and all variables are unrestricted.

$$\begin{array}{ll} \max \text{ (or min)} & f(x) \\ \text{s.t.} & g_1(x_1, x_2, \dots, x_n) = b_1 \\ & g_2(x_1, x_2, \dots, x_n) = b_2 \\ & \vdots \\ & g_m(x_1, x_2, \dots, x_n) = b_m \end{array}$$

- The KKT conditions: $\frac{\partial L}{\partial x_j} = \frac{\partial f(\bar{\mathbf{x}})}{\partial x_j} + \sum_{i=1}^m \bar{\lambda}_i \frac{\partial g_i(\bar{\mathbf{x}})}{\partial x_j} = 0$
- The original constraints: $\frac{\partial L}{\partial \lambda_i} = g_i(\bar{\mathbf{x}}) - b_i = 0$
- A point $(x_1, x_2, \dots, x_n, \lambda_1, \lambda_2, \dots, \lambda_m)$ that maximizes (minimizes) $L(x_1, x_2, \dots, x_n, \lambda_1, \lambda_2, \dots, \lambda_m)$ must satisfy

$$\frac{\partial L}{\partial x_1} = \frac{\partial L}{\partial x_2} = \dots = \frac{\partial L}{\partial x_n} = \frac{\partial L}{\partial \lambda_1} = \frac{\partial L}{\partial \lambda_2} = \dots = \frac{\partial L}{\partial \lambda_m} = 0$$

Example 30

- A company is planning to spend \$10,000 on advertising. It costs \$3,000 per minute to advertise on television and \$1,000 per minute to advertise on radio. If the firm buys x minutes of television advertising and y minutes of radio advertising, then its revenue in thousands of dollars is given by $f(x,y) = -2x^2 - y^2 + xy + 8x + 3y$.
- How can the firm maximize its revenue?

Special NLP – 2

- The (Karush-)Kuhn-Tucker conditions are used to solve NLPs **(26)**:

$$\begin{array}{l} \max(\text{ or } \min) f(x_1, x_2, \dots, x_n) \\ \text{s.t. } g_1(x_1, x_2, \dots, x_n) \leq b_1 \\ \quad g_2(x_1, x_2, \dots, x_n) \leq b_2 \\ \quad \quad \quad \vdots \\ \quad g_m(x_1, x_2, \dots, x_n) \leq b_m \end{array}$$

- Associate multipliers $\lambda_1, \lambda_2, \dots, \lambda_m$ with the constraints
- The **Kuhn-Tucker conditions** are necessary for a point $\bar{\mathbf{x}} = (\bar{x}_1, \bar{x}_2, \dots, \bar{x}_n)$ to solve the NLP.

KKT necessary conditions for NLP (26)

- If $\bar{\mathbf{x}} = (\bar{x}_1, \bar{x}_2, \dots, \bar{x}_n)$ is an optimal solution to NLP, then $\bar{\mathbf{x}}$ must **satisfy the m constraints** in the NLP, and there must exist **multipliers** $\lambda_1, \lambda_2, \dots, \lambda_m$ **satisfying**

□ for the maximization NLP:

$$\begin{aligned}\frac{\partial f(\bar{x})}{\partial x_j} - \sum_{i=1}^{i=m} \bar{\lambda}_i \frac{\partial g_i(\bar{x})}{\partial x_j} &= 0 & (j = 1, 2, \dots, n) \\ \bar{\lambda}_i [b_i - g_i(\bar{x})] &= 0 & (i = 1, 2, \dots, m) \\ \bar{\lambda}_i &\geq 0 & (i = 1, 2, \dots, m)\end{aligned}$$

□ for the minimization NLP:

$$\begin{aligned}\frac{\partial f(\bar{x})}{\partial x_j} + \sum_{i=1}^{i=m} \bar{\lambda}_i \frac{\partial g_i(\bar{x})}{\partial x_j} &= 0 & (j = 1, 2, \dots, n) \\ \bar{\lambda}_i [b_i - g_i(\bar{x})] &= 0 & (i = 1, 2, \dots, m) \\ \bar{\lambda}_i &\geq 0 & (i = 1, 2, \dots, m)\end{aligned}$$

Special NLP – 3

- NLPs (30) in which the variables are nonnegative:

$$\begin{aligned} \max \text{ (or min) } & z = f(x_1, x_2, \dots, x_n) \\ \text{s.t.} & \quad g_1(x_1, x_2, \dots, x_n) \leq b_1 \\ & \quad g_2(x_1, x_2, \dots, x_n) \leq b_2 \\ & \quad \vdots \\ & \quad g_m(x_1, x_2, \dots, x_n) \leq b_m \\ & \quad -x_1 \leq 0 \\ & \quad -x_2 \leq 0 \\ & \quad \vdots \\ & \quad -x_n \leq 0 \end{aligned}$$

- If $\bar{\mathbf{x}} = (\bar{x}_1, \bar{x}_2, \dots, \bar{x}_n)$ is an optimal solution to NLP, then $\bar{\mathbf{x}}$ must **satisfy the m constraints and sign restrictions** in the NLP, and there must exist **multipliers** $\lambda_1, \lambda_2, \dots, \lambda_m$ **satisfying the KKT conditions as below.**

KKT Necessary conditions for NLP (30)

□ for the maximization NLP:

$$\begin{aligned}\frac{\partial f(\bar{x})}{\partial x_j} - \sum_{i=1}^{i=m} \bar{\lambda}_i \frac{\partial g_i(\bar{x})}{\partial x_j} &\leq 0 \quad (j = 1, 2, \dots, n) \\ \bar{\lambda}_i [b_i - g_i(\bar{x})] &= 0 \quad (i = 1, 2, \dots, m) \\ \left[\frac{\partial f(\bar{x})}{\partial x_j} - \sum_{i=1}^{i=m} \bar{\lambda}_i \frac{\partial g_i(\bar{x})}{\partial x_j} \right] \bar{x}_j &= 0 \quad (j = 1, 2, \dots, n) \\ \bar{\lambda}_i &\geq 0 \quad (i = 1, 2, \dots, m)\end{aligned}$$

□ for the minimization NLP:

$$\begin{aligned}\frac{\partial f(\bar{x})}{\partial x_j} + \sum_{i=1}^{i=m} \bar{\lambda}_i \frac{\partial g_i(\bar{x})}{\partial x_j} &\geq 0 \quad (j = 1, 2, \dots, n) \\ \bar{\lambda}_i [b_i - g_i(\bar{x})] &= 0 \quad (i = 1, 2, \dots, m) \\ \left[\frac{\partial f(\bar{x})}{\partial x_j} + \sum_{i=1}^{i=m} \bar{\lambda}_i \frac{\partial g_i(\bar{x})}{\partial x_j} \right] \bar{x}_j &= 0 \quad (j = 1, 2, \dots, n) \\ \bar{\lambda}_i &\geq 0 \quad (i = 1, 2, \dots, m)\end{aligned}$$

KKT Necessary conditions for NLP (30)

□ for the maximization NLP:

$$\frac{\partial f(\bar{x})}{\partial x_j} - \sum_{i=1}^{i=m} \bar{\lambda}_i \frac{\partial g_i(\bar{x})}{\partial x_j} + \bar{\mu}_j = 0 \quad (j = 1, 2, \dots, n)$$

$$\bar{\lambda}_i [b_i - g_i(\bar{x})] = 0 \quad (i = 1, 2, \dots, m)$$

$$\left[\frac{\partial f(\bar{x})}{\partial x_j} - \sum_{i=1}^{i=m} \bar{\lambda}_i \frac{\partial g_i(\bar{x})}{\partial x_j} \right] \bar{x}_j = 0 \quad (j = 1, 2, \dots, n)$$

$$\bar{\lambda}_i \geq 0 \quad (i = 1, 2, \dots, m) \quad \text{Because } \bar{\mu}_j \geq 0, \text{ equivalently}$$

$$\bar{\mu}_j \geq 0 \quad (j = 1, 2, \dots, n)$$

$$\frac{\partial f(\bar{x})}{\partial x_j} - \sum_{i=1}^{i=m} \bar{\lambda}_i \frac{\partial g_i(\bar{x})}{\partial x_j} \leq 0 \quad (j = 1, 2, \dots, n)$$

$$\bar{\lambda}_i [b_i - g_i(\bar{x})] = 0 \quad (i = 1, 2, \dots, m)$$

$$\left[\frac{\partial f(\bar{x})}{\partial x_j} - \sum_{i=1}^{i=m} \bar{\lambda}_i \frac{\partial g_i(\bar{x})}{\partial x_j} \right] \bar{x}_j = 0 \quad (j = 1, 2, \dots, n)$$

$$\bar{\lambda}_i \geq 0 \quad (i = 1, 2, \dots, m)$$

KKT Necessary conditions for NLP (30)

□ for the minimization NLP:

$$\frac{\partial f(\bar{x})}{\partial x_j} + \sum_{i=1}^{i=m} \bar{\lambda}_i \frac{\partial g_i(\bar{x})}{\partial x_j} - \bar{\mu}_j = 0 \quad (j = 1, 2, \dots, n)$$

$$\bar{\lambda}_i [b_i - g_i(\bar{x})] = 0 \quad (i = 1, 2, \dots, m)$$

$$\left[\frac{\partial f(\bar{x})}{\partial x_j} + \sum_{i=1}^{i=m} \bar{\lambda}_i \frac{\partial g_i(\bar{x})}{\partial x_j} \right] \bar{x}_j = 0 \quad (j = 1, 2, \dots, n)$$

$$\bar{\lambda}_i \geq 0 \quad (i = 1, 2, \dots, m) \quad \text{Because } \bar{\mu}_j \geq 0, \text{ equivalently}$$

$$\bar{\mu}_j \geq 0 \quad (j = 1, 2, \dots, n)$$

$$\frac{\partial f(\bar{x})}{\partial x_j} + \sum_{i=1}^{i=m} \bar{\lambda}_i \frac{\partial g_i(\bar{x})}{\partial x_j} \geq 0 \quad (j = 1, 2, \dots, n)$$

$$\bar{\lambda}_i [b_i - g_i(\bar{x})] = 0 \quad (i = 1, 2, \dots, m)$$

$$\left[\frac{\partial f(\bar{x})}{\partial x_j} + \sum_{i=1}^{i=m} \bar{\lambda}_i \frac{\partial g_i(\bar{x})}{\partial x_j} \right] \bar{x}_j = 0 \quad (j = 1, 2, \dots, n)$$

$$\bar{\lambda}_i \geq 0 \quad (i = 1, 2, \dots, m)$$

5.4 Geometrical interpretation of KKT conditions

- Three KKT conditions for (26) hold at a point \bar{x} if and only if ∇f is a linear combination of $\nabla g_1, \nabla g_2, \dots, \nabla g_m$, and the weight multiplying ∇g_i in this linear combination equals 0 if the i th constraint in (26) is nonbinding.

FIGURE 43
Example of
Kuhn-Tucker
Conditions: Both
Constraints Binding

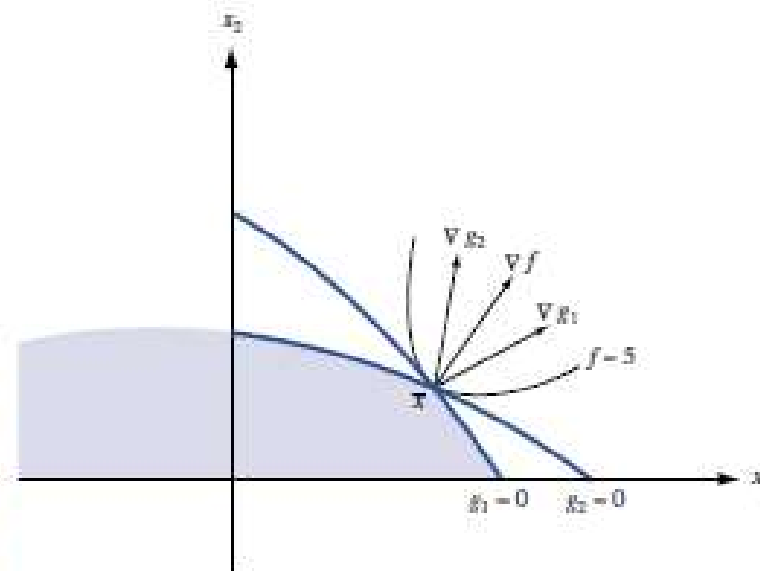
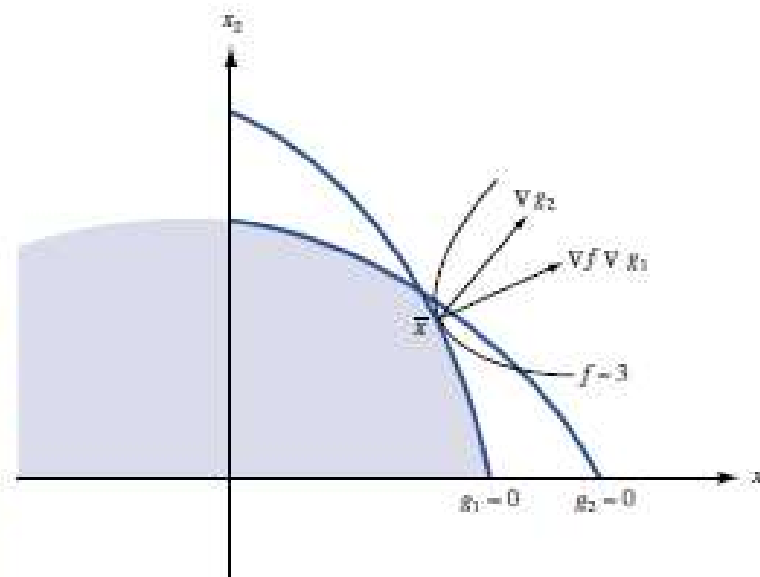


FIGURE 44
Example of
Kuhn-Tucker
Conditions: One
Constraint Binding and
One Constraint
Nonbinding



Example 33

- A monopolist can purchase up to 17.25 oz of a chemical for \$10/oz. At a cost of \$3/oz, the chemical can be processed into an ounce of product 1; or, at a cost of \$5/oz, the chemical can be processed into an ounce of product 2. If x_1 oz of product 1 are produced, it sells for a price of $$(30 - x_1)$ per ounce. If x_2 oz of product 2 are produced, it sells for a price of $$(50 - 2x_2)$ per ounce. Determine how the monopolist can maximize profits.
- Solution. Let
 - x_1 =ounces of product 1 produced
 - x_2 =ounces of product 2 produced
 - x_3 =ounces of chemical processed

5.5 Constraint Qualifications

- For the theorems in this section to hold, the functions g_1, g_2, \dots, g_m must satisfy certain regularity conditions (constraint qualifications).
- Unless a constraint qualification or regularity condition is satisfied at an optimal point $\bar{\mathbf{x}}$, the Kuhn-Tucker conditions may fail to hold at $\bar{\mathbf{x}}$.
- When the constraints are linear, these regularity assumptions are always satisfied.
- One constraint qualification - **linear independence CQ**: If all g_i are continuous, and the gradients of all binding constraints (including any binding nonnegativity constraints) at optimal solution \mathbf{x} form a set of linearly independent vectors, then the KKT conditions must hold at \mathbf{x} .

-
- A set of vectors is said to be **linearly dependent** if one of the vectors in the set can be defined as a linear combination of the others; if no vector in the set can be written in this way, then the vectors are said to be **linearly independent**.
 - Linear Dependent: $a_1V_1 + a_2V_2 + \dots + a_kV_k = 0$ for a_i not all zero.

Example 34: Necessity of Constraint Qualification

Show that the Kuhn–Tucker conditions fail to hold at the optimal solution to the following NLP:

$$\begin{aligned} \max z &= x_1 \\ \text{s.t.} \quad &x_2 - (1 - x_1)^3 \leq 0 \\ &x_1 \geq 0, x_2 \geq 0 \end{aligned} \tag{56}$$

Solution If $x_1 > 1$, then the first constraint in (56) implies that $x_2 < 0$. Thus, the optimal z -value for (56) cannot exceed 1. Because $x_1 = 1$ and $x_2 = 0$ is feasible and yields $z = 1$, $(1, 0)$ must be the optimal solution to NLP (56).

From Theorem 10, the following are two of the Kuhn–Tucker conditions for (56).

$$1 + 3\lambda_1(1 - x_1)^2 = -\mu_1 \tag{57}$$

$$\mu_1 \geq 0 \tag{58}$$

At the optimal solution $(1, 0)$, (57) implies $\mu_1 = -1$, which contradicts (58). Thus, the Kuhn–Tucker conditions are not satisfied at $(1, 0)$. We now show that at the point $(1, 0)$ the Linear Independence Constraint Qualification is violated. At $(1, 0)$ the constraints $x_2 - (1 - x_1)^3 \leq 0$ and $x_2 \geq 0$ are binding. Then

$$\nabla(x_2 - (1 - x_1)^3) = [0, 1]$$

$$\nabla(-x_2) = [0, -1]$$

Because $[0, 1] + [0, -1] = [0, 0]$, these gradients are linearly dependent. Thus, at $(1, 0)$ the gradients of the binding constraints are linearly dependent, and the constraint qualification is not satisfied.

LINGO

- If LINGO displays the message DUAL CONDITIONS:SATISFIED then you know it has found the point satisfying the Kuhn-Tucker conditions. Unless satisfying the sufficient conditions, LINGO might return a solution that is not optimal. Use different initial solutions to test the optimality.

6 The Method of Feasible Directions

- This method, a modification of the steepest ascent method, can be used to solve the NLP with linear constraints.

$$\begin{array}{ll}\max & z = f(\mathbf{x}) \\ \text{s.t.} & A\mathbf{x} \leq \mathbf{b} \\ & \mathbf{x} \geq \mathbf{0}\end{array}$$

- To solve, begin with a feasible solution \mathbf{x}^0 (perhaps by using the two-phase simplex algorithm).
- Next, find a direction to move away from \mathbf{x}^0 , which makes the new solution **remain feasible** and increase the value of z .
- Let \mathbf{d}^0 be a solution to the following LP:

$$\begin{array}{ll}\max & z = \nabla f(\mathbf{x}^0) \cdot \mathbf{d} \\ \text{s.t.} & A\mathbf{d} \leq \mathbf{b} \\ & \mathbf{d} \geq \mathbf{0}\end{array}$$

-
- Choose our new point \mathbf{x}^1 to be $\mathbf{x}^1 = \mathbf{x}^0 + t^0(\mathbf{d}^0 - \mathbf{x}^0)$, where t^0 solves

$$\begin{aligned} \max \quad & f(\mathbf{x}^0 + t^0(\mathbf{d}^0 - \mathbf{x}^0)) \\ \text{s.t.}, \quad & 0 \leq t^0 \leq 1 \end{aligned}$$

It is an NLP with a single variable.

- Continue in this fashion and generate directions of movement $\mathbf{d}^1, \mathbf{d}^2, \dots, \mathbf{d}^{n-1}$ and new points $\mathbf{x}^2, \mathbf{x}^3, \dots, \mathbf{x}^n$.
- We terminate the algorithm if $\mathbf{x}^k = \mathbf{x}^{k-1}$ or successive points are sufficiently close together. Return \mathbf{x}^{k-1} as the solution to NLP.

Example 37

Perform two iterations of the feasible directions method on the following NLP:

$$\begin{aligned} \max z &= f(x, y) = 2xy + 4x + 6y - 2x^2 - 2y^2 \\ \text{s.t.} \quad &x + y \leq 2 \\ &x, y \geq 0 \end{aligned}$$

Begin at the point (0,0).

7 Quadratic Programming

- A **quadratic programming** (QP) is an NLP in which each term in the objective function is of degree 2, 1, or 0 (quadratic function) and all constraints are linear.
- LINGO, Excel and **Wolfe's method** (a modified version of Phase I of the two-phase simplex to find a point satisfying the KKT conditions) may be used to solve QP problems.
- Wolfe's method is guaranteed to obtain the optimal solution to a QP if all leading principal minors of the objective function's Hessian are positive (positive definite).
- In practice, the method of complementary pivoting is most often used to solve QPs (Shapiro, 1979).

Example 35: Portfolio Optimization

- I have \$1,000 to invest in three stocks. Let S_i be the random variable representing the annual return on \$1 invested in stock i . Thus, if $S_i = 0.12$, \$1 invested in stock i at the beginning of a year was worth \$1.12 at the end of the year. We are given the following information:
- $E(S_1) = 0.14$, $E(S_2) = 0.11$, $E(S_3) = 0.10$, $\text{var}(S_1) = 0.20$, $\text{var}(S_2) = 0.08$, $\text{var}(S_3) = 0.18$, $\text{cov}(S_1, S_2) = 0.05$, $\text{cov}(S_1, S_3) = 0.02$, $\text{cov}(S_2, S_3) = 0.03$.
- Formulate a QP that can be used to find the portfolio that attains an expected annual return of at least 12% and minimizes the variance of the annual dollar return on the portfolio.

8 Unconstrained NLPs with One Variable

$$\max (or \min) f(x)$$

$$\text{s.t.} \quad x \in [a, b]$$

- There are three types of points for which the NLP can have a local maximum or minimum (these points are often called *extremum candidates*).
 - Points where $a < x < b$, $f'(x) = 0$ (called a **stationary point** of $f(x)$)
 - Points where $f'(x)$ does not exist
 - Boundary points (endpoints) a and b of the interval $[a, b]$
- To find the optimal solution for the NLP, find all the local optima. The optimal solution is the local maximum (or minimum) having the largest (or smallest) value of $f(x)$.

Case 1: stationary points

THEOREM 4

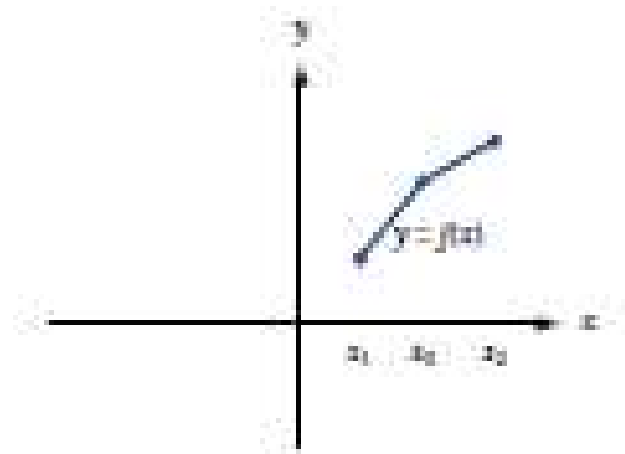
If $f'(x_0) = 0$ and $f''(x_0) < 0$, then x_0 is a local maximum. If $f'(x_0) = 0$ and $f''(x_0) > 0$, then x_0 is a local minimum.

THEOREM 5

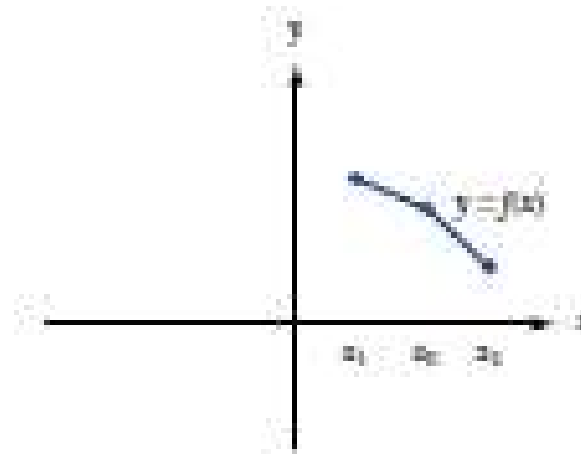
If $f'(x_0) = 0$, and

- 1 If the first nonvanishing (nonzero) derivative at x_0 is an odd-order derivative [$f^{(3)}(x_0)$, $f^{(5)}(x_0)$, and so on], then x_0 is not a local maximum or a local minimum.
- 2 If the first nonvanishing derivative at x_0 is positive and an even-order derivative, then x_0 is a local minimum.
- 3 If the first nonvanishing derivative at x_0 is negative and an even-order derivative, then x_0 is a local maximum.

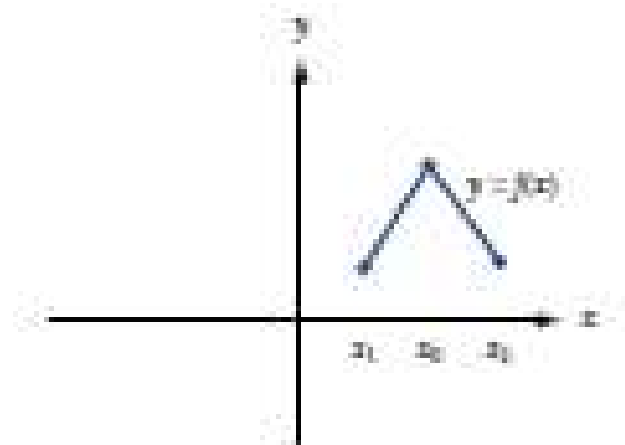
Case 2: non-differentiable points



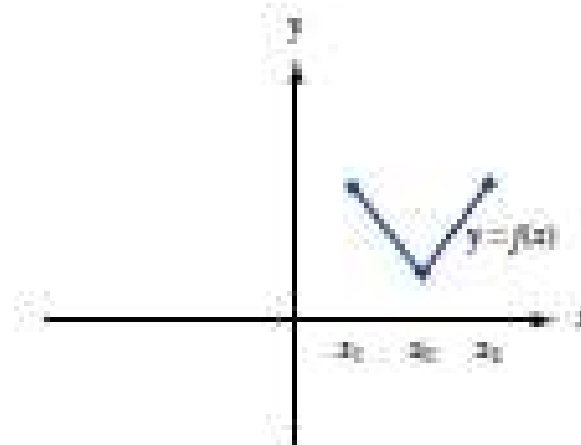
■ x_0 not a local extremum



■ x_0 not a local extremum



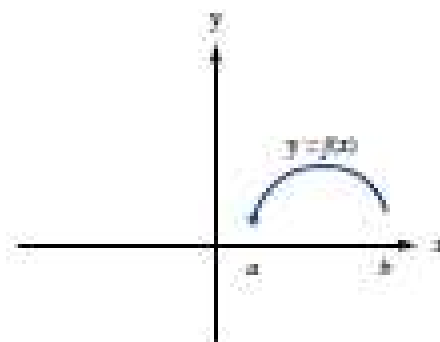
■ x_0 is a local maximum



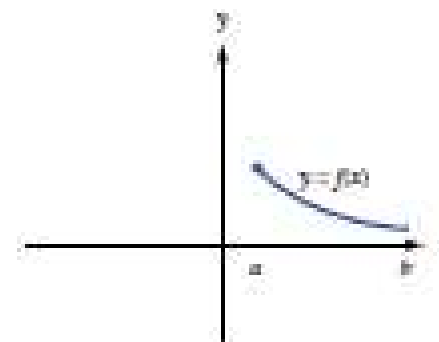
■ x_0 is a local minimum

Case 3: endpoints

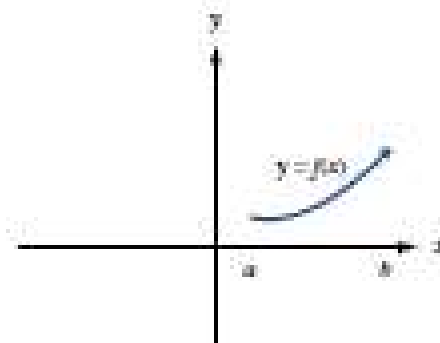
- If $f'(a) > 0$, then a is a local minimum.
- If $f'(a) < 0$, then a is a local maximum.
- If $f'(b) > 0$, then b is a local maximum.
- If $f'(b) < 0$, then b is a local minimum.
- If $f'(a) = 0$ or $f'(b) = 0$, evaluate $f(x)$ at some point $a < x < b$ sufficient close to a or b .



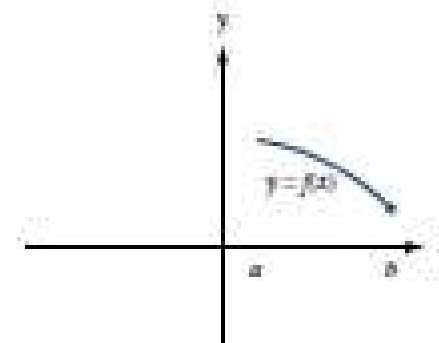
■ $f'(a) > 0$
 a is a local minimum



■ $f'(a) < 0$
 a is a local maximum



■ $f'(b) > 0$
 b is a local maximum



■ $f'(b) < 0$
 b is a local minimum

Example: Profit Maximization by Monopolist

- It costs a monopolist \$5/unit to produce a product.
- If he produces x units of the product, then each can be sold for $10-x$ dollars.
- To maximize profit, how much should the monopolist produce.
- Note that: Demand (denoted as d) is often modeled as a function of price (denoted as p). For example, $d = 1 - bp$ (linear relation) for a parameter b . \rightarrow nonlinear in revenue and profit.

Example: Solution

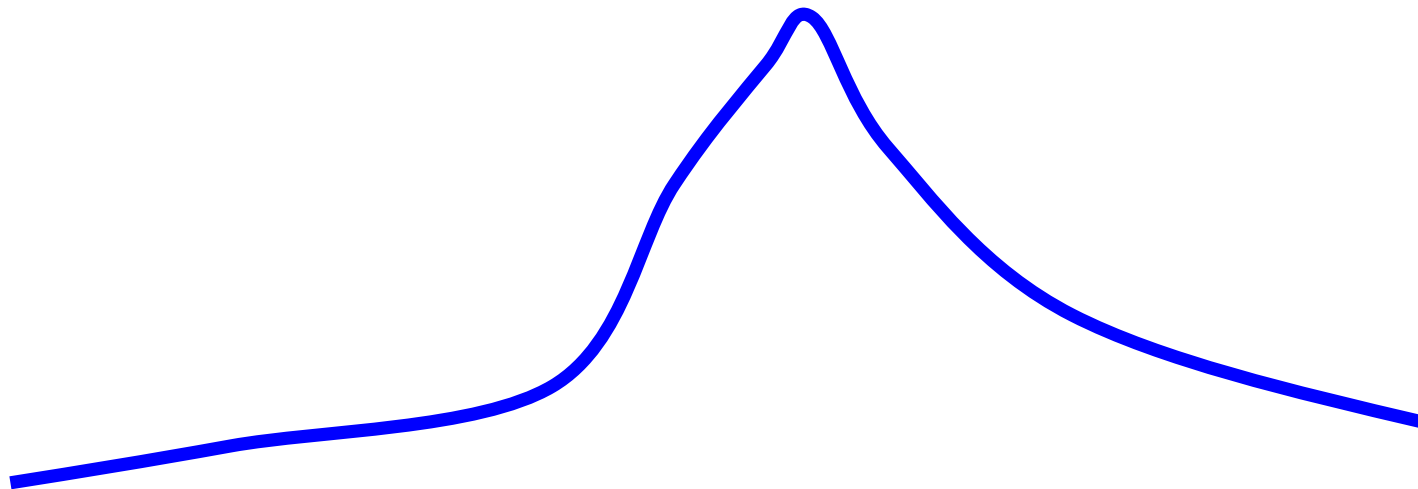
- The monopolist wants to solve the NLP

$$\begin{aligned} \max P(x) &= x(10 - x) - 5x = 5x - x^2 \\ \text{s.t. } 0 &\leq x \leq 10 \end{aligned}$$

- The extremum candidates can be classified as
 - Case 1 check tells us $x=2.5$ is a local maximum yielding a profit $P(2.5)=6.25$.
 - $P'(x)$ exists for all points in $[0,10]$, so there are no Case 2 candidates.
 - $a = 0$ has $P'(0) = 5 > 0$ so $a = 0$ is a local minimum; $b=10$ has $P'(10)=-15<0$, so $b = 10$ is a local minimum

8.1 Local Search

- The some numerical methods can be used if the function is a unimodal function.
- A function $f(x)$ is **unimodal** on $[a,b]$ if for some point \bar{x} on $[a,b]$, $f(x)$ is strictly increasing on $[a, \bar{x}]$ and strictly decreasing on $[\bar{x}, b]$.
- Not necessary concave or even $f'(x)$ may not exist
- A single variable function is **unimodal** if there is at most one local maximum (or at most one local minimum) .



-
- The optimal solution of the NLP is some point on the interval $[a,b]$. By evaluating $f(x)$ at two points x_1 and x_2 on $[a,b]$, we may reduce the size of the interval in which the solution to the NLP must lie.
 - After evaluating $f(x_1)$ and $f(x_2)$, one of these cases must occur. It can be shown in each case that the optimal solution will lie in a subset of $[a,b]$.
 - Case 1: $f(x_1) < f(x_2)$ and $\bar{x} \in (x_1, b]$
 - Case 2: $f(x_1) = f(x_2)$ and $\bar{x} \in (a, x_2]$
 - Case 3: $f(x_1) > f(x_2)$ and $\bar{x} \in (a, x_2]$
 - The interval in which \bar{x} must lie – either $[a, x_2)$ or $(x_1, b]$ – is called the **interval of uncertainty**.

The three cases:

FIGURE 33
If $f(x_1) < f(x_2)$,
 $x \in (x_1, b]$

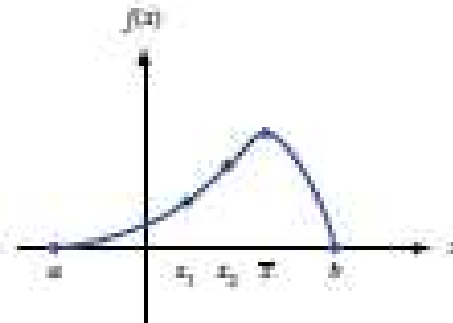
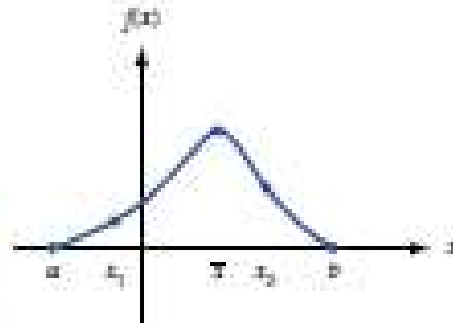


FIGURE 34
If $f(x_1) = f(x_2)$,
 $x \in [a, x_2]$

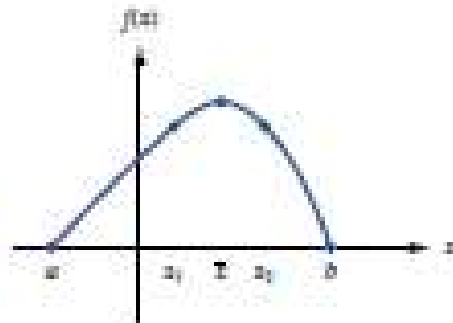
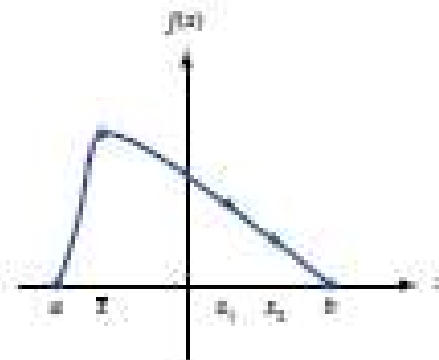
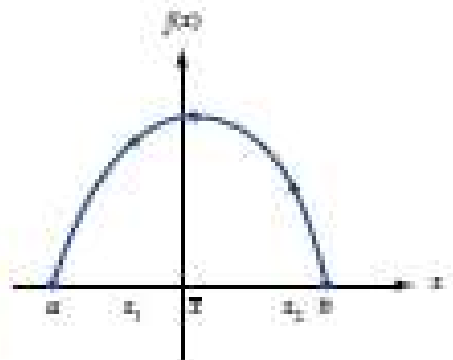


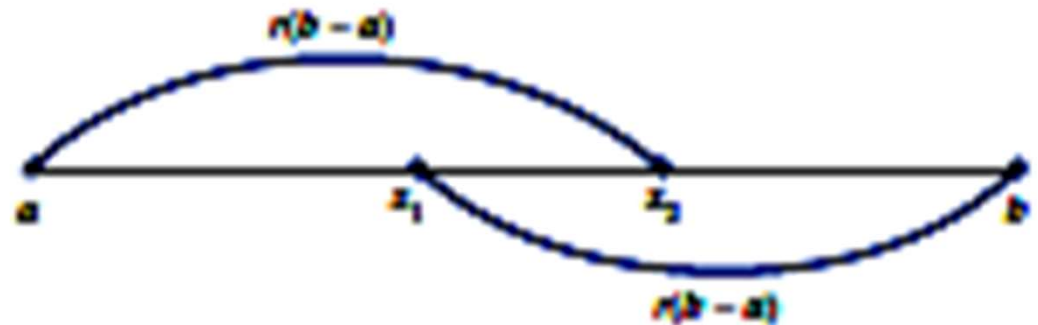
FIGURE 35
If $f(x_1) > f(x_2)$,
 $x \in [a, x_2]$



-
- Many search algorithms use these ideas to reduce the interval of uncertainty. Most of these algorithms proceed as follows:
 - Begin with the interval of uncertainty for x being $[a,b]$. Evaluate $f(x)$ at two judiciously chosen points x_1 and x_2 .
 - Determine which of cases 1-3 holds, and find a reduced interval of uncertainty.
 - Evaluate $f(x)$ at two new points (the algorithm specifies how the two new points are chosen). Return to step 2 unless the length of the interval of uncertainty is sufficiently small.
 - Such search algorithms: golden section search, bisection search, Fibonacci search

Golden section search

- How to choose the points to evaluate:
 - $x_1 = b - r(b-a)$ and $x_2 = a + r(b-a)$
- After each iteration, the interval of uncertainty is reduced by r times (check the 3 cases).
 - $b - x_1 = r(b-a)$ and $x_2 - a = r(b-a)$
- After k iterations, the interval of uncertainty $= r^k(b-a)$
- Determine r :
 - $L/rL = rL/(1-r)L$
 - $x_1 = b - r(b-a) = a + r[r(b-a)]$
→ $r = (5^{1/2} - 1)/2 = 0.618$



Example

Use Golden Section Search to find

$$\max -x^2 - 1$$

$$\text{s.t.} \quad -1 \leq x \leq 0.75$$

with the final interval of uncertainty having a length less than $\frac{1}{4}$.

Bisection Search

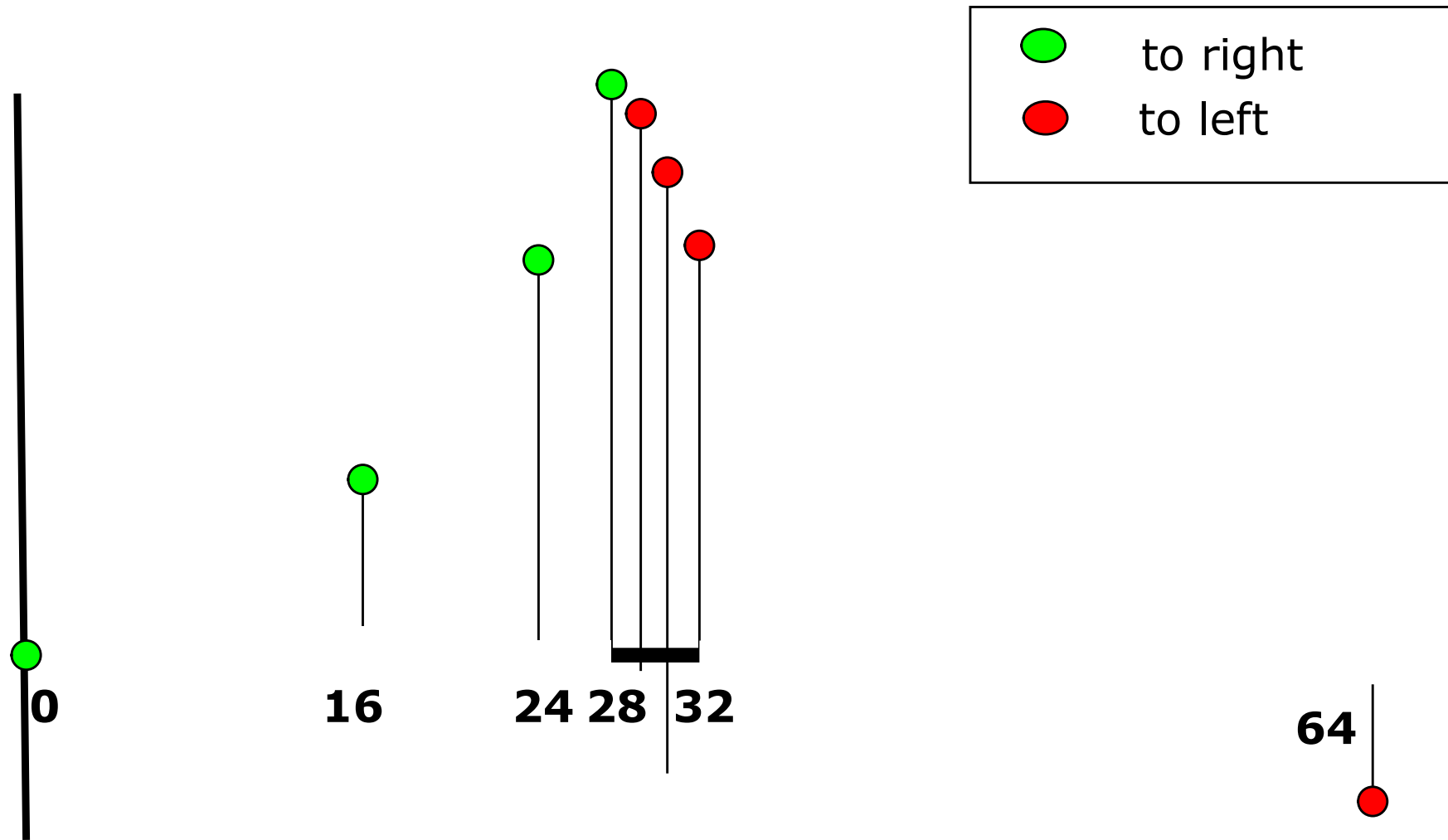
If $f(x)$ is concave (or simply unimodal) and differentiable

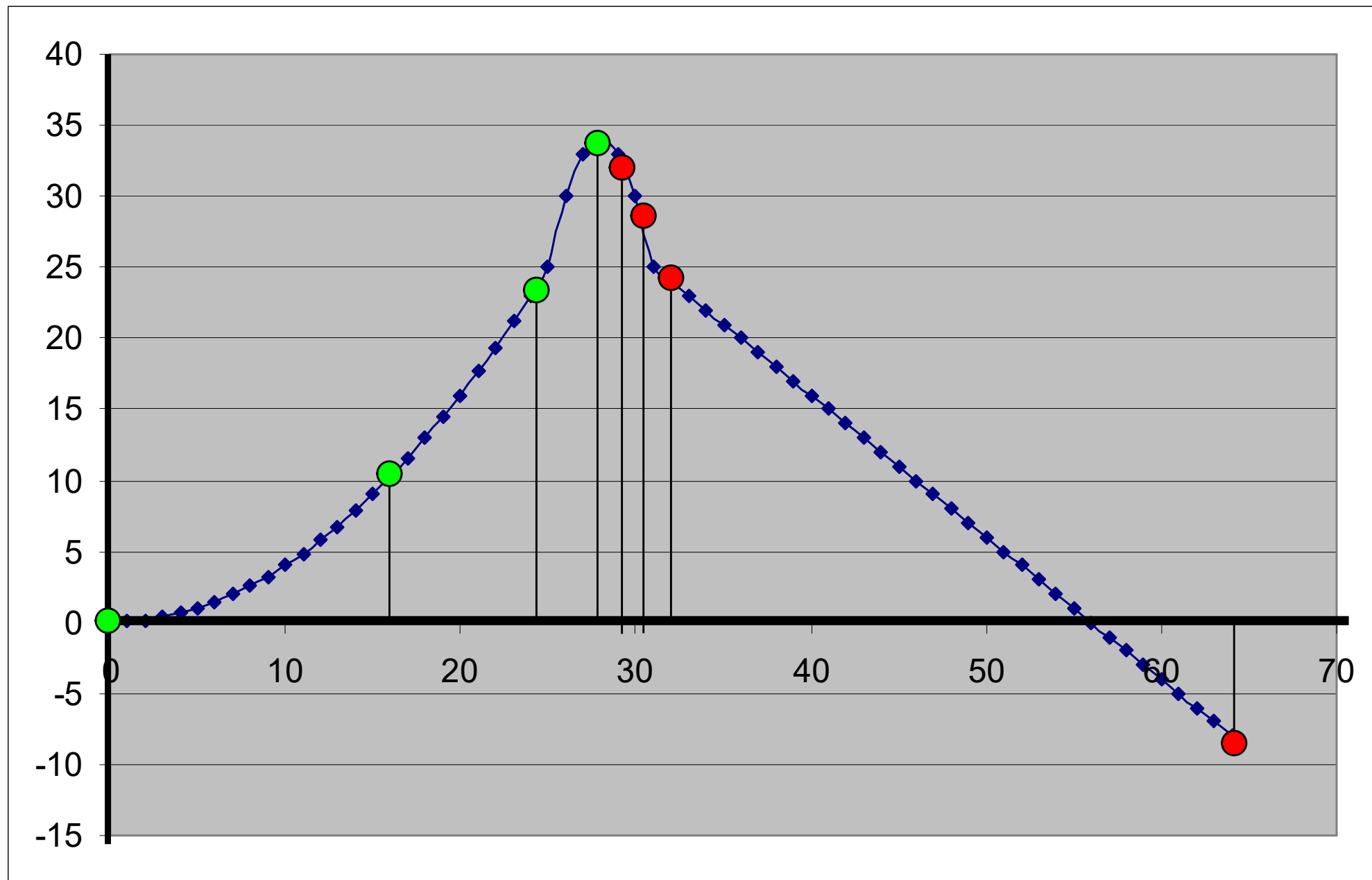
$$\begin{aligned} \max \quad & f(x) \\ \text{s.t.} \quad & a \leq x \leq b \end{aligned}$$

Bisection (or Bolzano) Search:

- Step 1. Begin with the region of uncertainty for x as $[a, b]$. Evaluate $f'(x)$ at the midpoint $x_M = (a+b)/2$.
- Step 2. If $f'(x_M) > 0$, then eliminate the interval up to x_M . If $f'(x_M) < 0$, then eliminate the interval beyond x_M .
- Step 3. Evaluate $f'(x)$ at the midpoint of the new interval. Return to Step 2 until the interval of uncertainty is sufficiently small.

Determine by taking a derivative if a local maximum is to the right or left.





Fibonacci Search

- Instead of taking derivatives (which may be computationally intensive), use two function evaluations to determine updated interval.
- Fibonacci Search
- Step 1. Begin with the region of uncertainty for q as $[a, b]$. Evaluate $f(q_1)$ and $f(q_2)$ for 2 **symmetric** points $q_1 < q_2$.
- Step 2. If $f(q_1) \leq f(q_2)$, then eliminate the interval up to q_1 . If $f(q_1) > f(q_2)$, then eliminate the interval beyond q_2 .
- Step 3. Select a second point symmetric to the point already in the new interval, rename these points q_1 and q_2 such that $q_1 < q_2$ and evaluate $f(q_1)$ and $f(q_2)$. Return to Step 2 until the interval is sufficiently small.

On Fibonacci search

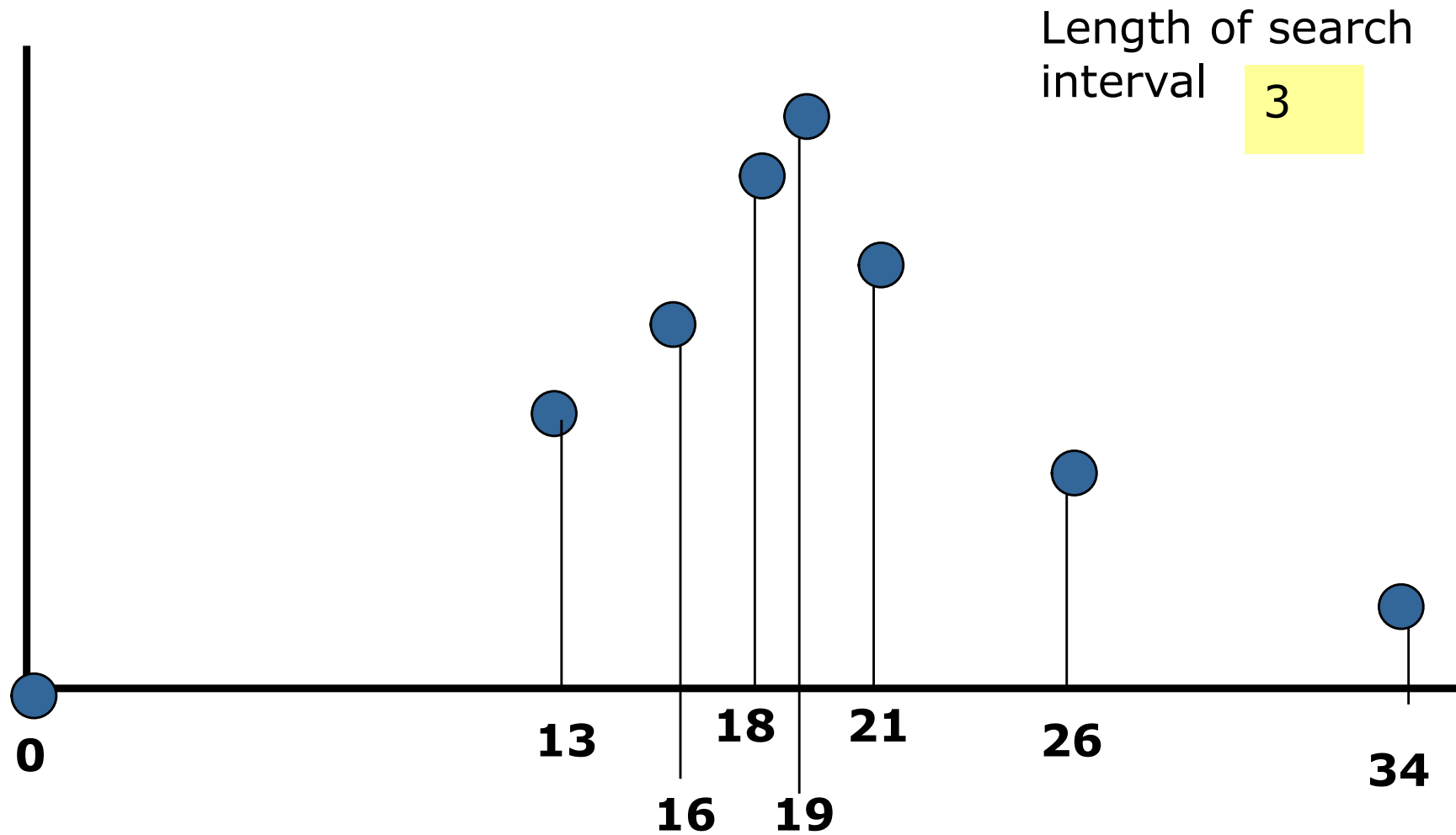
Fibonacci numbers: 1, 1, 2, 3, 5, 8, 13, 21, 34

At iteration 1, the length of the search interval is the k th Fibonacci number for some k

At iteration j , the length of the search interval is the $k-j+1$ Fibonacci number.

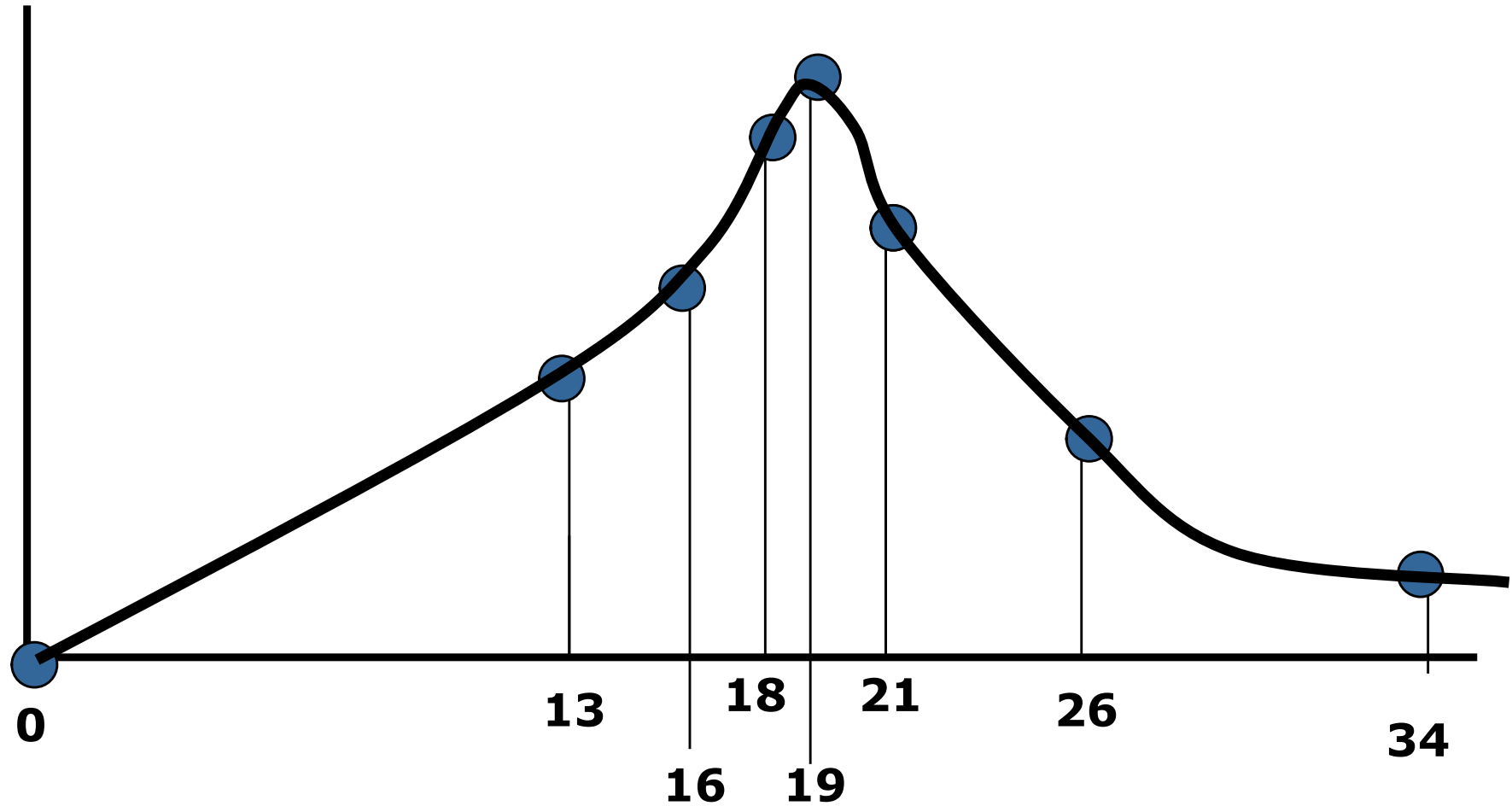
The technique converges to the optimal when the function is unimodal.

Finding a local maximum using Fibonacci Search

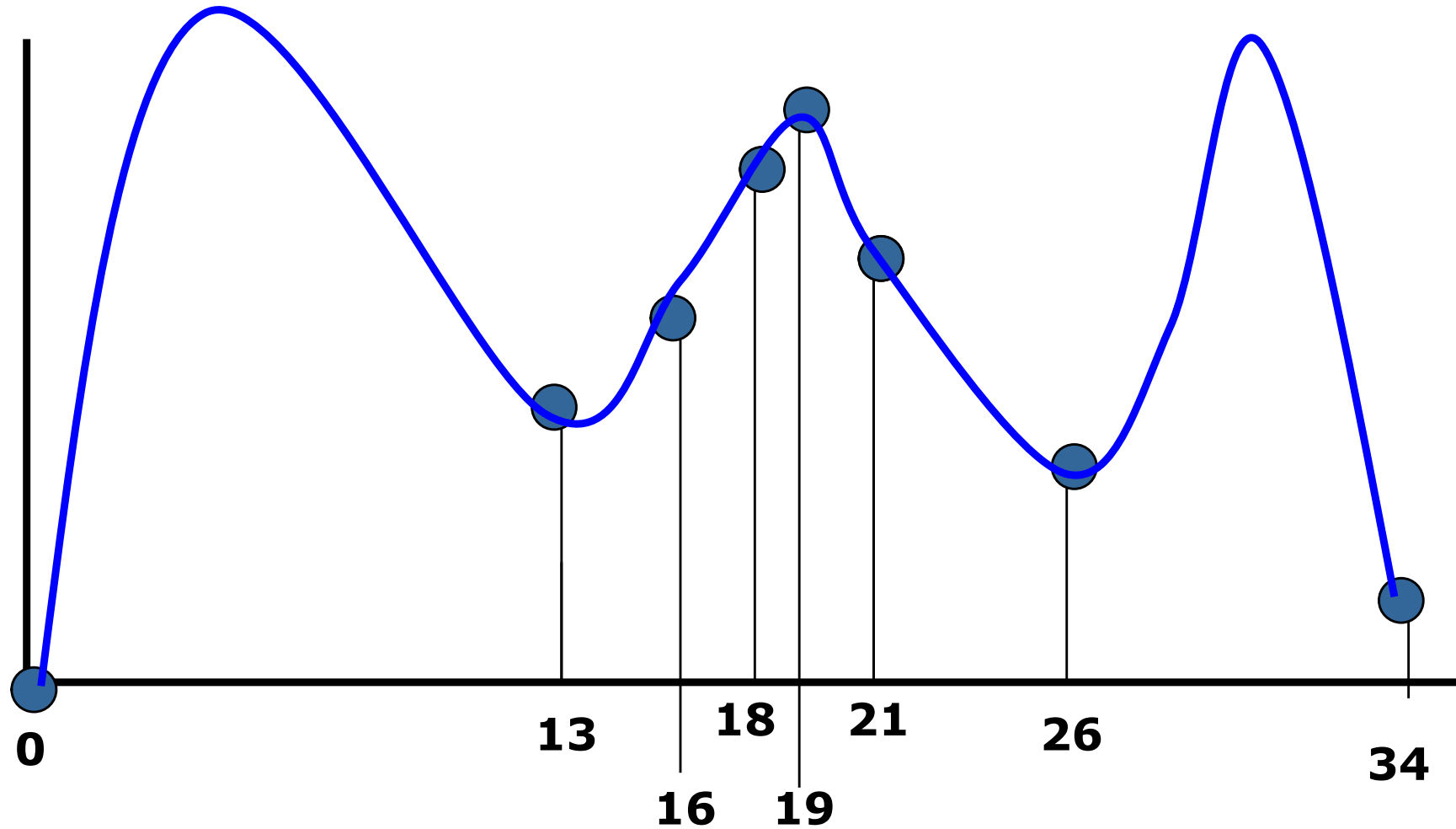


Where the maximum may be

The search finds a local maximum, but not necessarily a global maximum.



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Number of function evaluations in Fibonacci Search

- As new point is chosen symmetrically, the length l_k of successive search intervals is given by: $l_k = l_{k+1} + l_{k+2}$.
- Solving for these lengths given a final interval length of 1, $l_n = 1$, gives the Fibonacci numbers: 1, 2, 3, 5, 8, 13, 21, 34,...
- Thus, if the initial interval has length 34, it takes 8 function calculations to reduce the interval length to 1.
- Remark: If the function is convex or unimodal, then Fibonacci search converges to the global maximum.