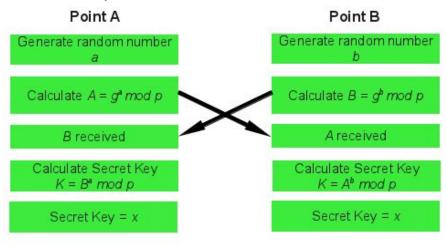
## **Experiment 7**

**Aim:** Write a program to implement key exchange using Diffie-Hallman key exchange.

**Theory:** Diffie-Hellman key exchange, also called exponential key exchange, is a method of digital encryption that uses numbers raised to specific powers to produce decryption keys on the basis of components that are never directly transmitted, making the task of a would-be code breaker mathematically overwhelming. To implement Diffie-Hellman, the two end users Alice and Bob, while communicating over a channel they know to be private, mutually agree on positive whole numbers p and q, such that p is a prime number and q is a generator of p. The generator q is a number that, when raised to positive whole-number powers less than p, never produces the same result for any two such whole numbers.



p = prime number, g = generator

The value of p may be large but the value of q is usually small. Once Alice and Bob have agreed on p and q in private, they choose positive whole-number personal keys a and b, both less than the prime-number modulus p. Neither user divulges their personal key to anyone; ideally they memorize these numbers and do not write them down or store them anywhere. Next, Alice and Bob compute public keys a\* and b\* based on their personal keys according to the formulas: -

$$a^* = [q^a] \mod p$$
  
 $b^* = [q^b] \mod p$ 

The two users can share their public keys a\* and b\* over a communications medium assumed to be insecure, such as the Internet or a corporate wide area network (WAN). From these public keys, x and y can be generated by either user on the basis of their own personal keys.

Alice computes x using the formula

$$x = [(b^*)^a] \mod p$$

Bob computes x using the formula

$$y = [(a^*)^b] \mod p$$

The value of x and y turns out to be the same according to either of the above two formulas. The two users can therefore, in theory, communicate privately over a public medium with an encryption method of their choice using the decryption key x.

## **Source Code:**

```
#include <iostream>
#include<math.h>
using namespace std;
long long int calc(long long int a, long long int b, long long int N)
  if(b == 1)
    return a;
  else
     return (((long long int)pow(a, b)) % N);
}
int main() {
  long long int N, G, x, a, y, b, ka, kb;
  cout << "\nEnter the values for N and G\n";
  cin>>N>>G;
  cout << "\nEnter the private key for A ";
  cin>>a;
  cout << "Enter the private key for B";
  cin>>b;
  cout << "\nThe private key of A: " << a;
  cout << "\nThe private key of B: " << b;
  x = calc(G, a, N);
  y = calc(G, b, N);
  cout << "\n\nAfter exchange of keys";
  cout << "\nkey recieved by A is (y):" << y;
  cout <<"\nkey recieved by B is (x):"<<x;
  ka = calc(y, a, N);
```

## **Output:**

```
kunal@DESKTOP-AITAEP7:/mnt/c/Users/Admin/Desktop/college/7th Semester/Information and network kunal@DESKTOP-AITAEP7:/mnt/c/Users/Admin/Desktop/college/7th Semester/Information and network Enter the values for N and G

Enter the private key for A 3
Enter the private key for B 6

The private key of A: 3
The private key of B: 6

After exchange of keys key recieved by A is (y):13
key recieved by B is (x):6

Actual key for the A is : 12
Actual Key for the B is : 12

Both users have matching keys, thus successful
```

## **Learning Outcomes:**

The personal keys a and b, which are critical in the calculation of x, have not been transmitted over a public medium. Because it is a large and apparently random number, a potential hacker has almost no chance of correctly guessing x, even with the help of a powerful computer to conduct millions of trials.