

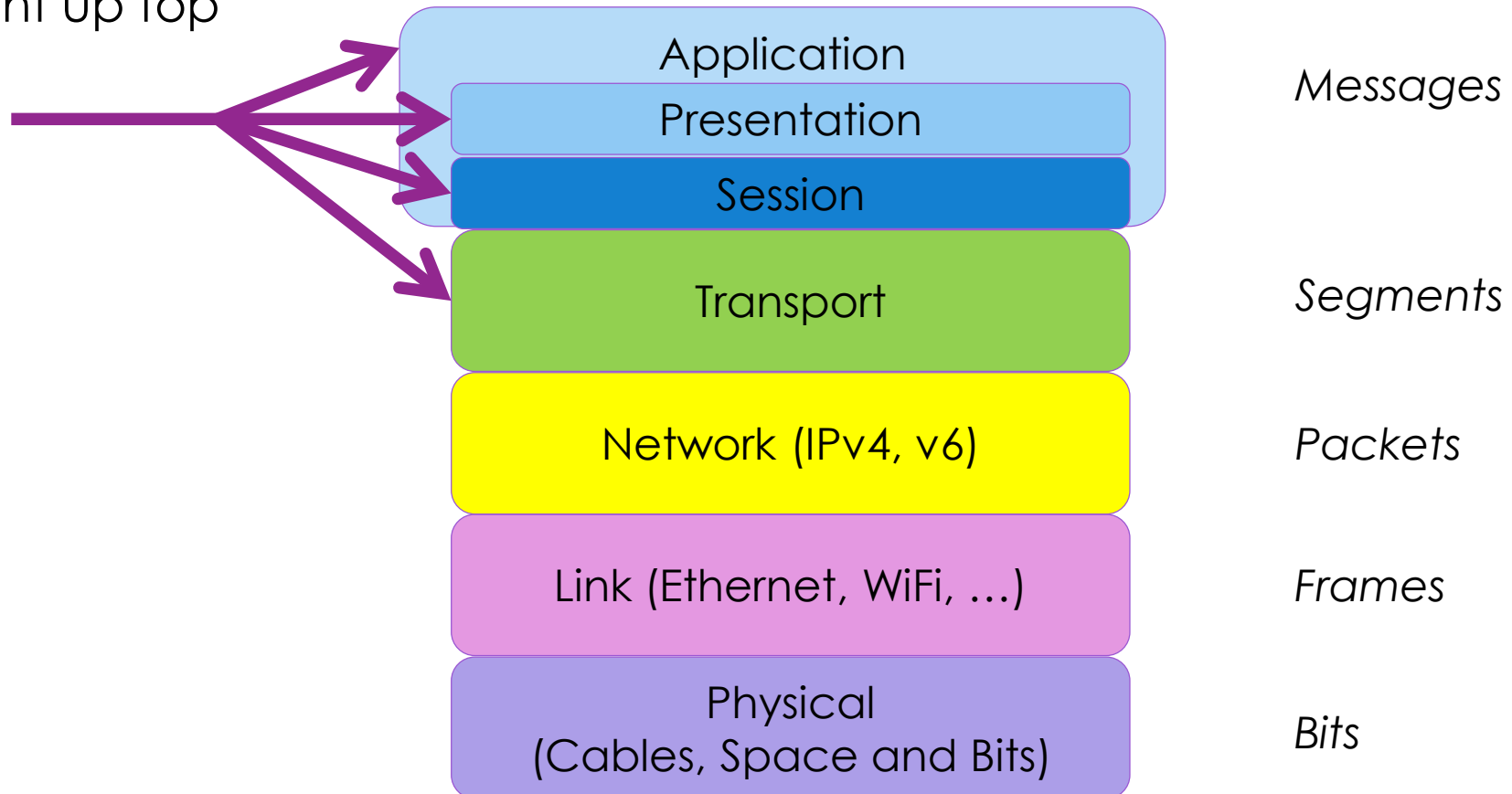
COMP3310/6331 – #16

Realtime communications

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Where are we?

- Right up top



Applications choose their transport

- UDP-based applications:
 - Short messages
 - Low(er) delays
 - Simple request/response transactions
 - Light server touch
 - ARQ suffices
- TCP-based applications:
 - Larger content transfers
 - Longer, and more complex, sessions
 - Reliability matters
 - Packaging and presentation becomes important – TCP is a bytestream

What about real-time traffic?

- Audio/Video applications are obvious
 - And exponentially hard $(N \text{ sites})^2$
- But what about robotics?
 - What about tele-medicine??
- What about broader system control
 - (open a gate, close a valve, ...)?
- When do you trade off timeliness (delays) for reliability (loss)?

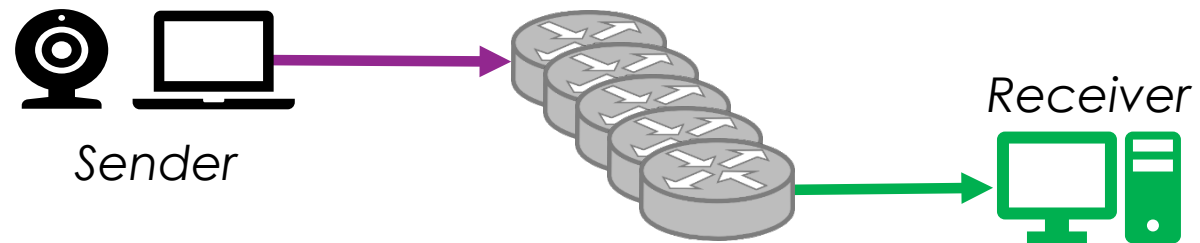


How **real** is real-time?

- Very application specific
 - How much can you tolerate a single lost packet?
 - How much can you tolerate a delayed packet?
 - How much can you notice a delay?
 - E.g. streaming vs videoconf
 - And different audio/video delays (synch) are worse
- TCP for streaming (one-way), with CDNs
 - Delays less crucial, win on reliability
- UDP for (two-way) interactive, real-time
 - low-delay, low-overheads

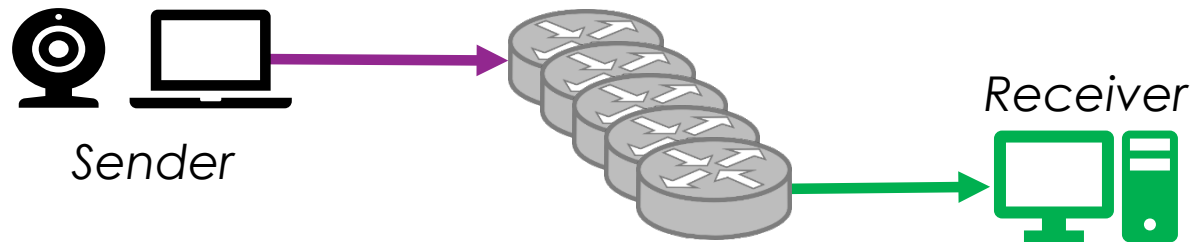
Why is it hard?

- Real-time media – e.g. videoconference
- Internet is ‘best-effort’
 - Unless you have circuits
- Any **delay** is a problem – **variable** delay is another, “**loss**” is another
 - Various kinds of loss

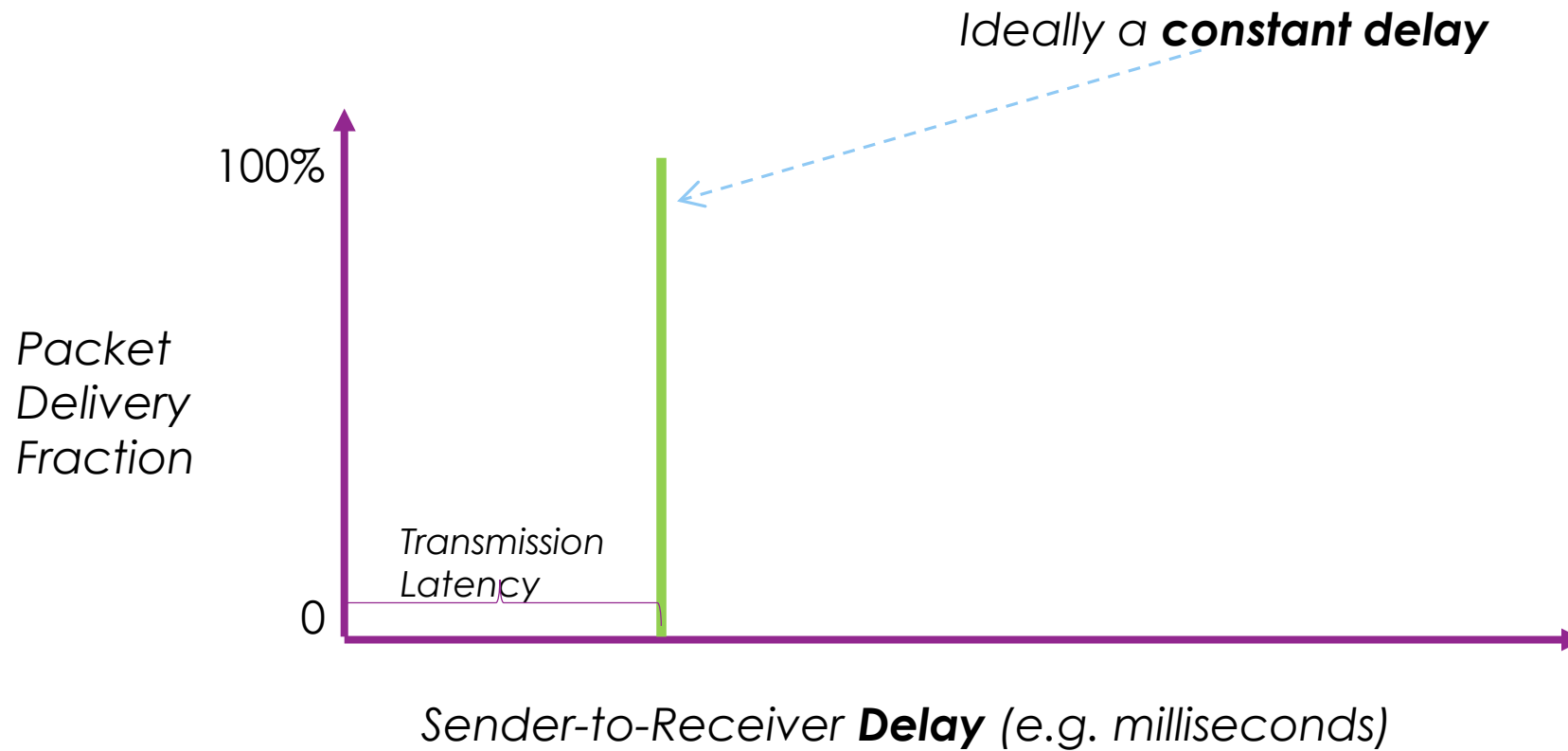


Network delays

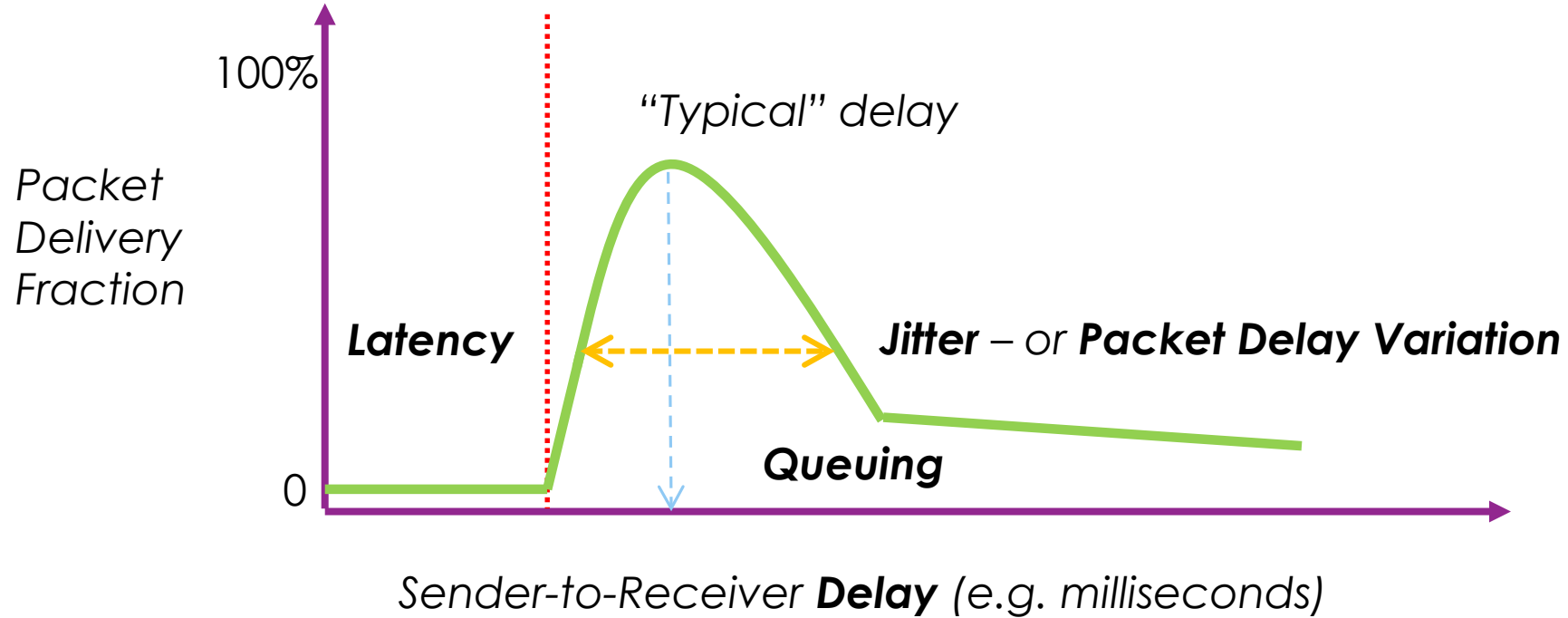
- Sender is sending a constant stream of audio/video samples
 - Bandwidth may vary depending on codec
 - Receiver expects to receive a constant stream!
- But all those routers don't belong to us... neither do the paths
- We need the path(s), the bandwidth, and capacity on the routers



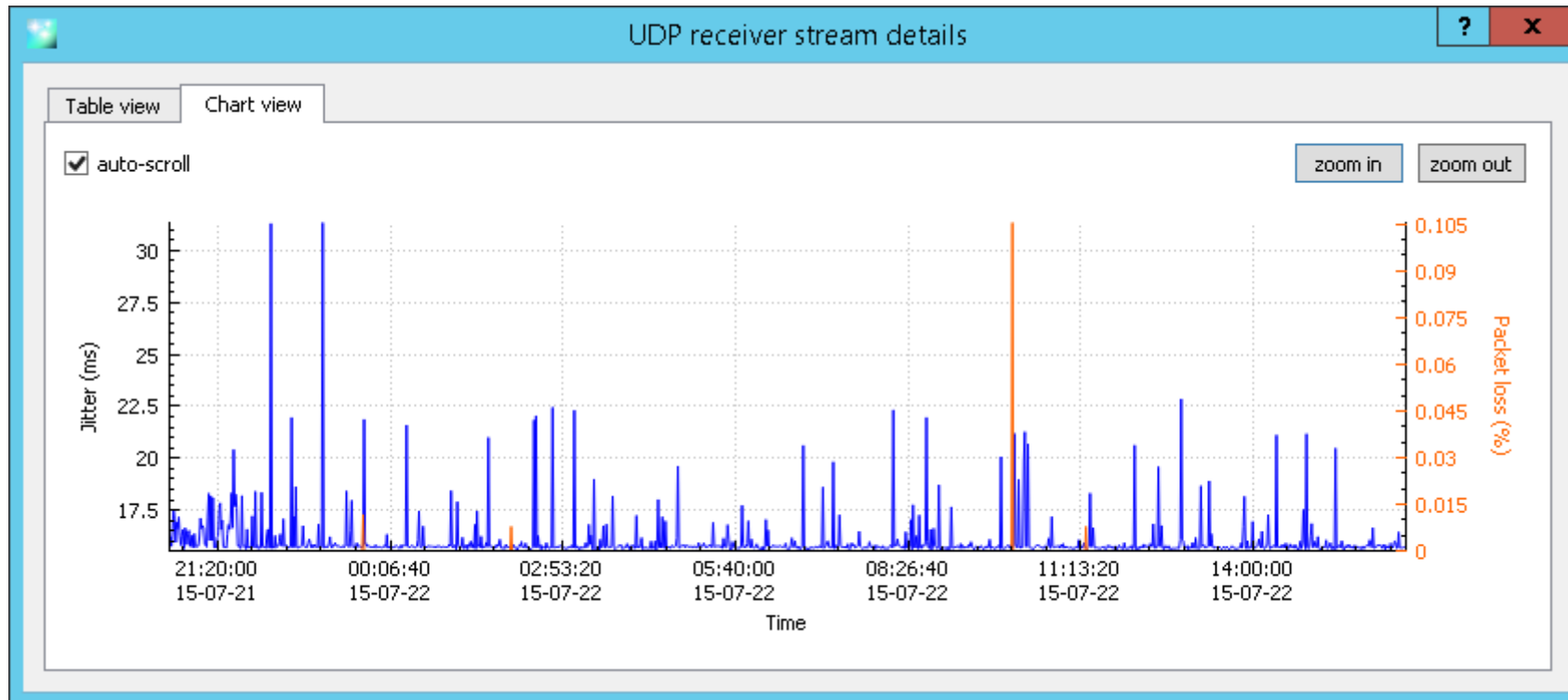
Network delay elements



Network delay elements

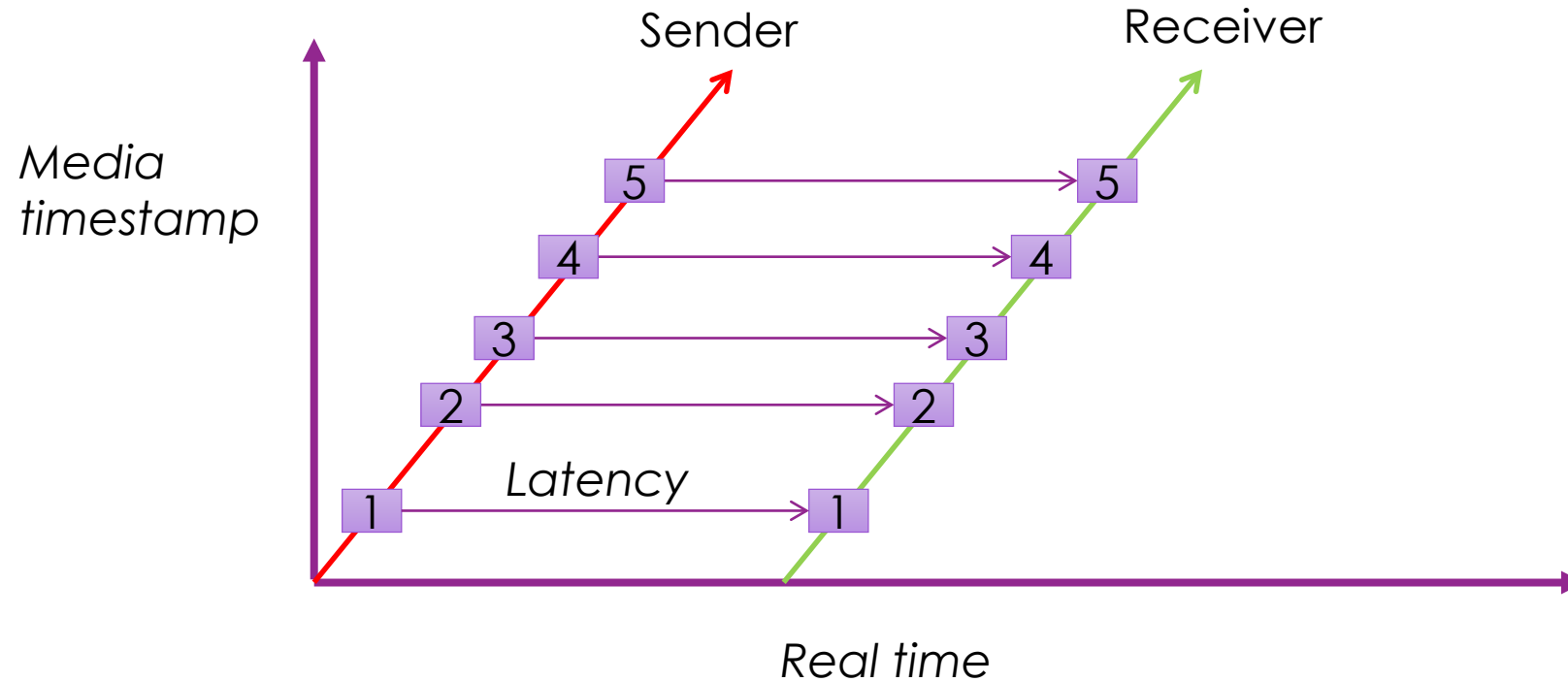


And jitter changes over time



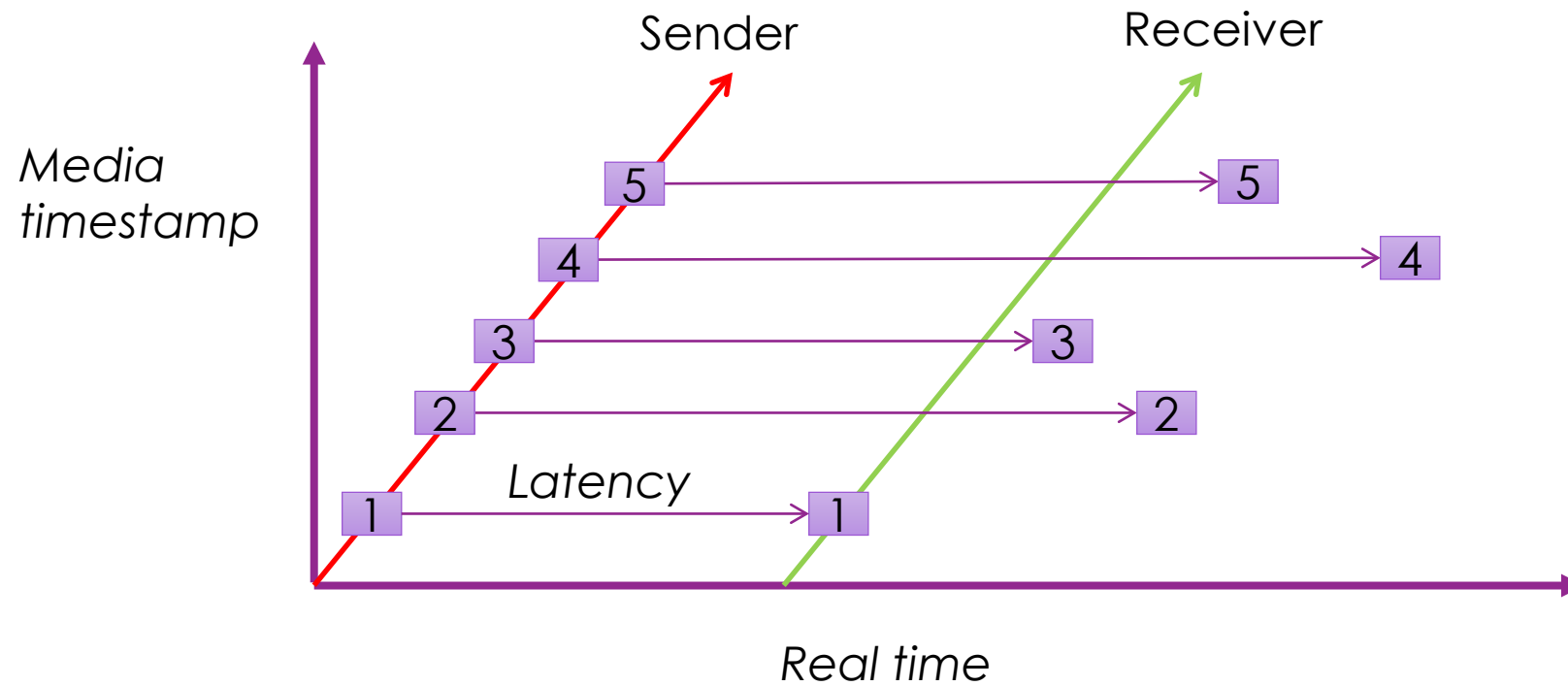
Ideally, sending packets

Ideally, network delay is constant (PDV=0) – like a telephone circuit



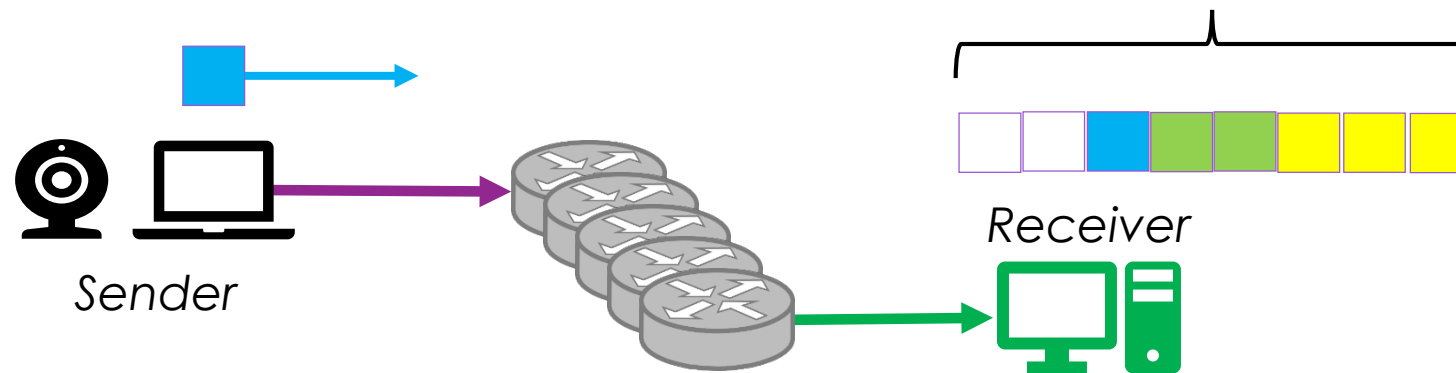
Reality...

Packet delays are random,
and packet order can be messy

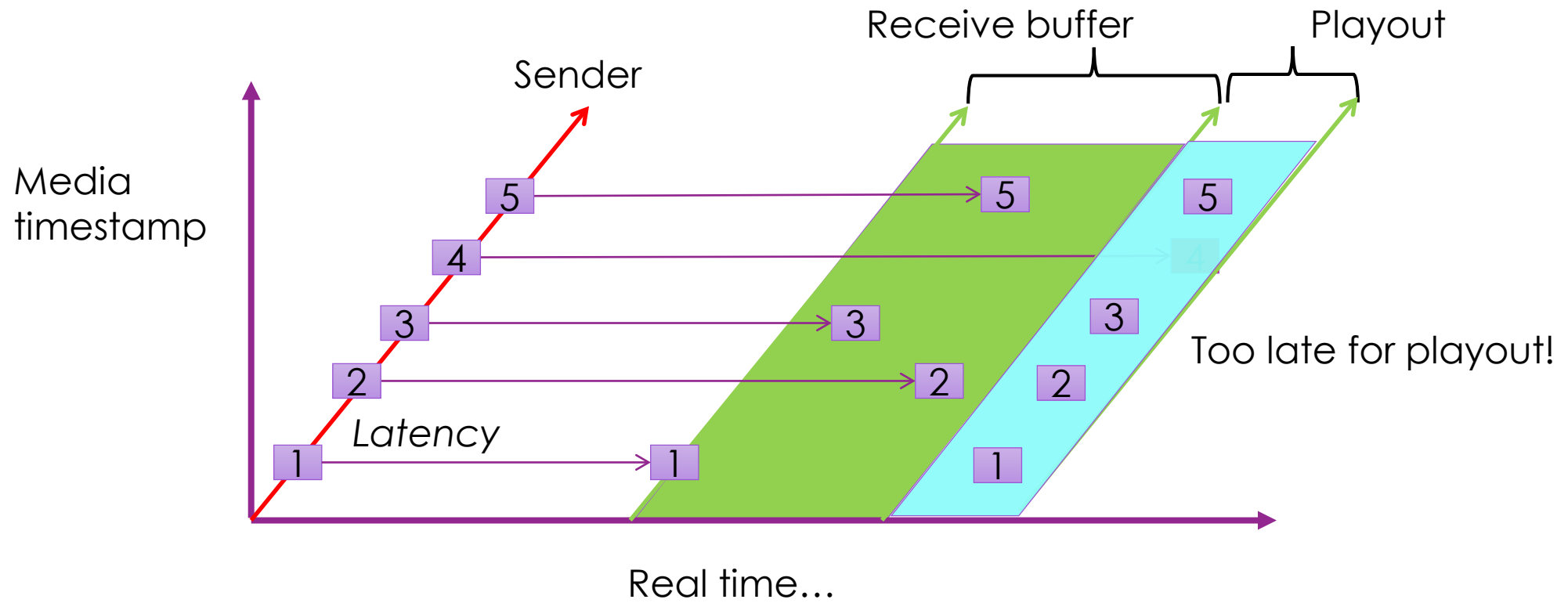


Buffering (TCP-lite)

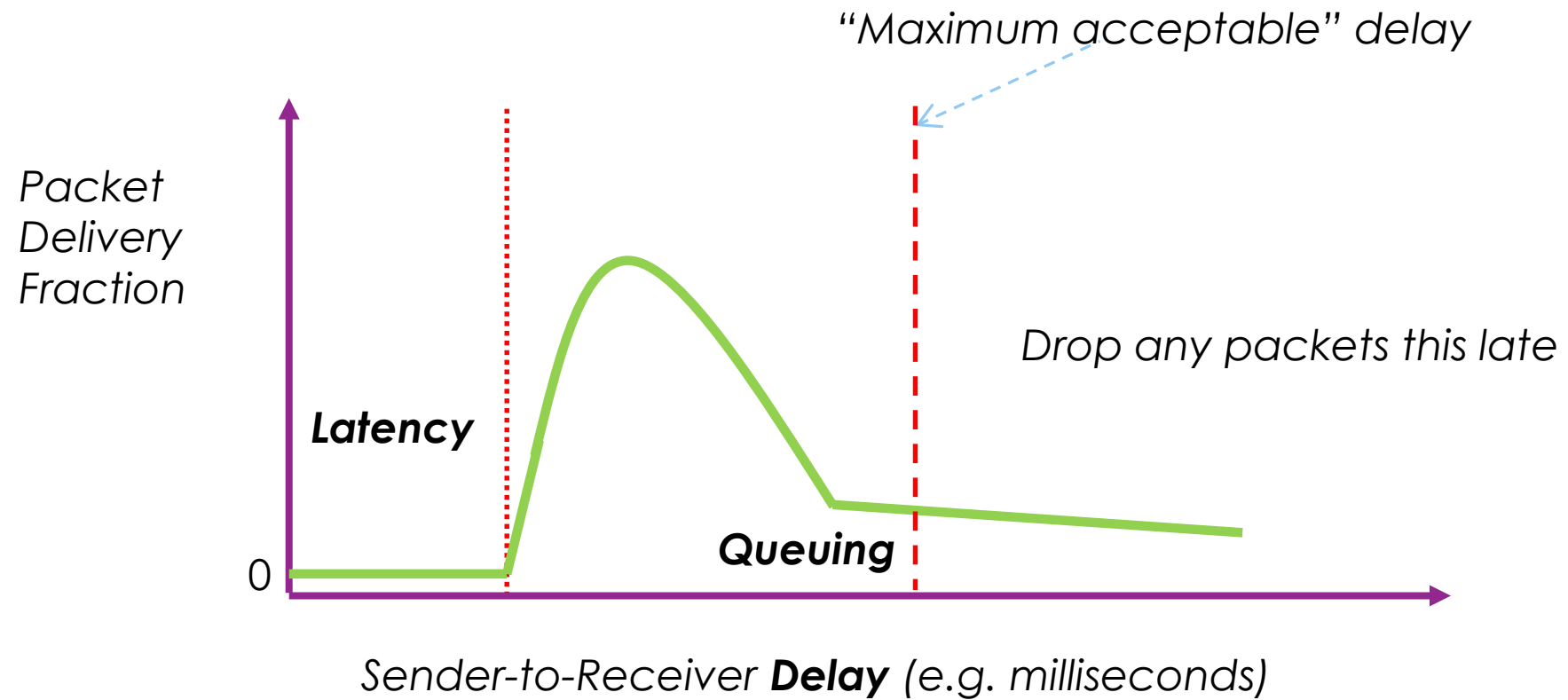
- Sender is sending a constant stream of audio/video samples
 - Receiver expects to receive a constant stream!
- So Receiver makes a playout buffer – smooth out variable delays
 - Measured in bytes, but effectively in time



This solves everything? ...



Network delay elements



It's a trade-off

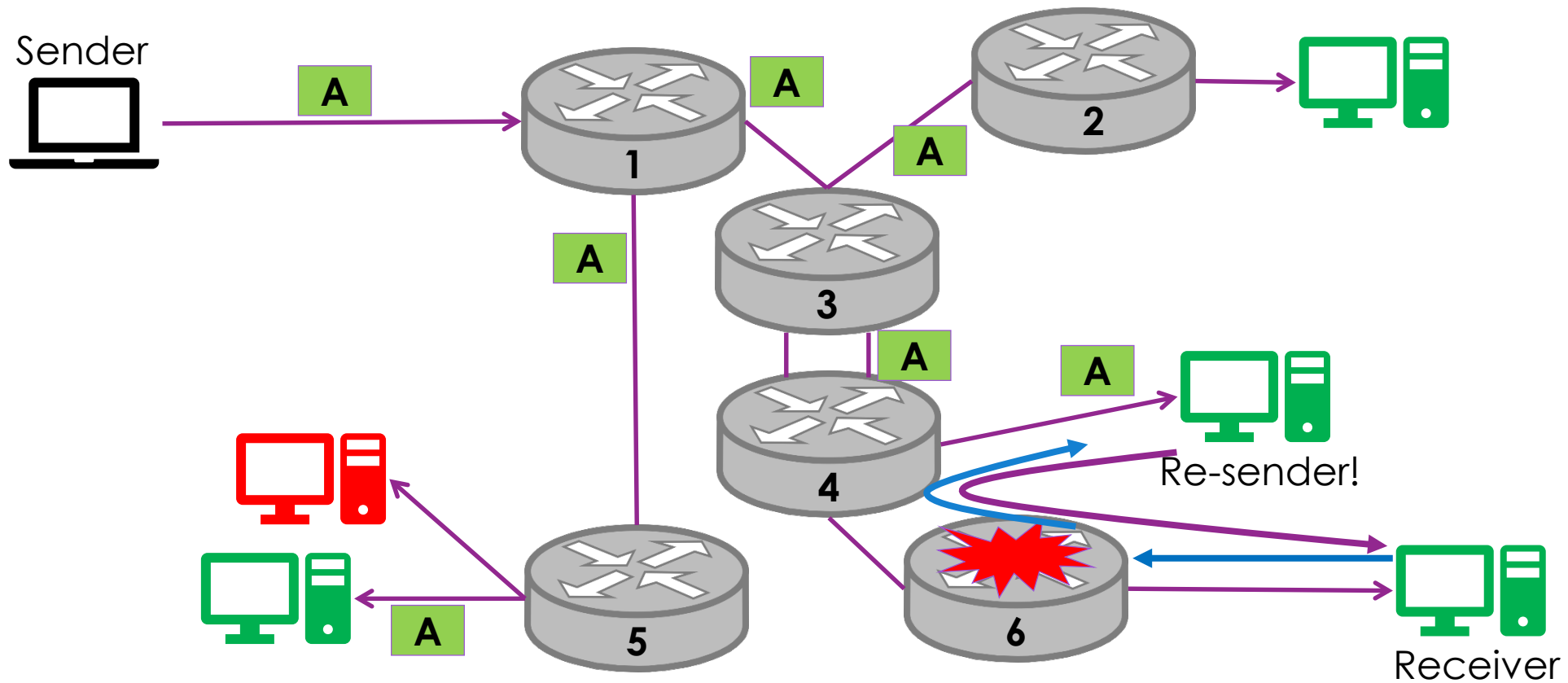
- **Big** buffer:
 - Fewer packets lost to delays = More tolerant of jitter/PDV
 - Greater delay between transmission and playout
- **Small** buffer:
 - More packets may be lost to delays = Less tolerant of jitter/PDV
 - Smaller delay between transmission and playout
- The smaller the (desired) delay, the harder to deal with loss – so you “glitch”
 - Might be ok for video, less so for audio, but ...

Fixes?

- Retransmission
 - You know something is lost
 - Request a resend of the problem packet
 - ARQ -> TCP-like!
 - Takes *round-trip-time* plus *transmission time* plus *queuing delays*
 - Huge buffer/delays implied
 - Generally not done
 - *In multicast, it may not be the sender who retransmits!*

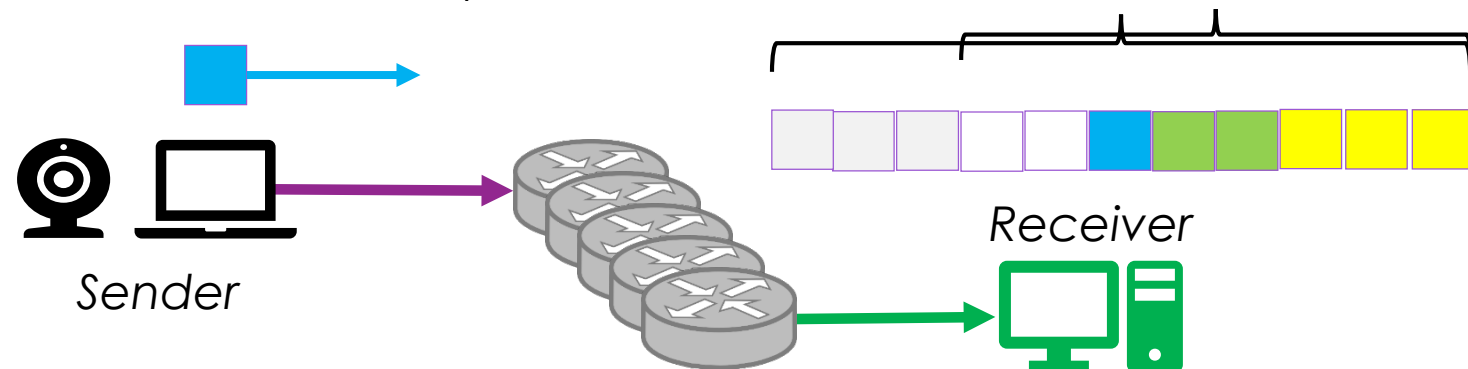
IP Multicast: UDP

Everyone is a listener, and a **sender**



Fixes?

- Elastic buffers
 - When things go bad, you adapt – stretch the buffer
 - And playout *slower*
 - Easy to stretch/squeeze audio and video...
 - Shrink the buffer when things get better
 - Want to keep as close to real-time as possible



Fixes?

- Error-correction

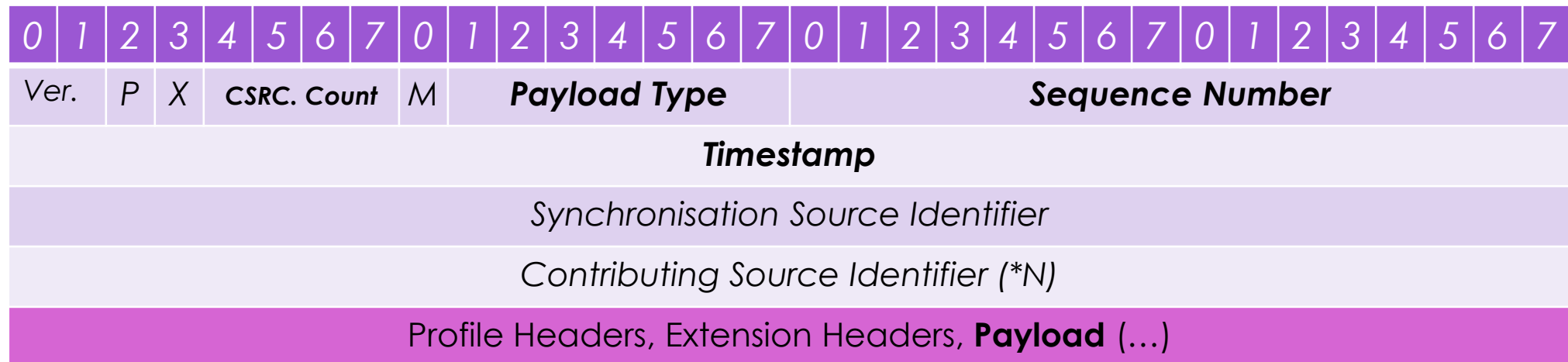
- Encode media for interpolation between packets
 - Anything missing, you have an approximation
 - Similar to compression algorithms, progressive-display images, ...
- (Adaptive) Forward Error Correction
 - Anything broken, you can reconstruct
 - Useful for (control) reliability, or where retransmission is too expensive

- Parallel-transmission

- Send several copies at the same time
 - At (multiple) different (lower) qualities
 - Low quality audio/video still better than no audio/video!
 - And also useful for control reliability



Realtime Transport Protocol (RTP)



- UDP payload (or TCP)
- Allows for any media (stream) encodings
- Allows for multiple sources to be merged, identified, and synchronised
 - Lip-synch audio to video

RTP Feedback

- As sender, would be useful to know:
 - How well is my stream getting through?
 - Should I adjust rate, encoding, error-correction, retransmission, ...
 - How many endpoints are receiving it? (multicast)
 - And identify them
- **RTP Control Protocol (RTCP)**
 - Bidirectional, *out-of-band* signalling
 - But need to limit bandwidth usage!
 - Sender Reports, Receiver Reports
 - Statistics: Packets received/lost, delay, delay variation,
 - Source Description, Hello/Goodbye, ...
 - Also provides heartbeat, distance-measure, retransmission-request channel, ...

All together: a videoconference

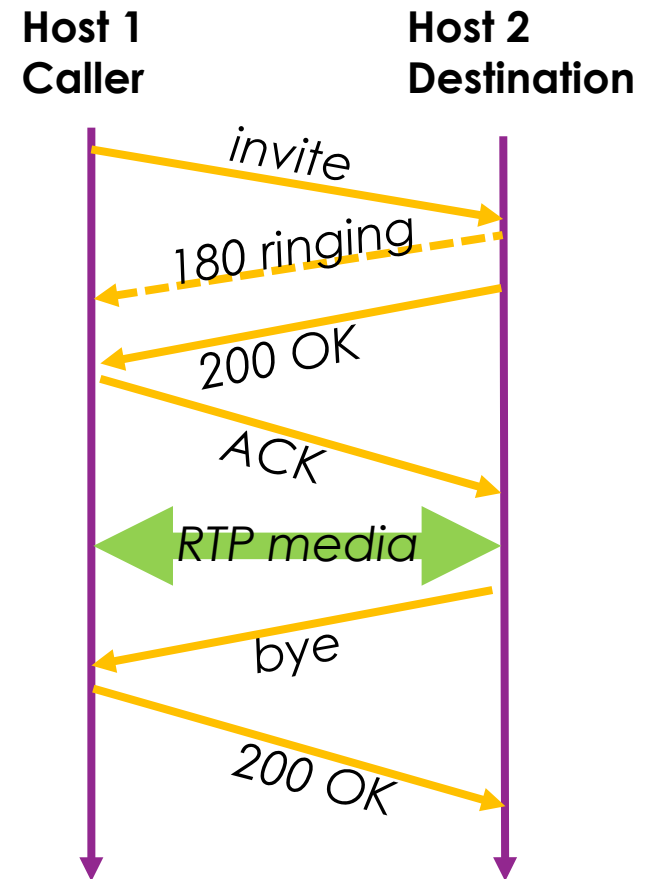
- Multiple sites, multiple media streams
- Need to:
 - Establish a call: *Session Initiation Protocol (SIP)*
 - Negotiate the details: *Session Description Protocol (SDP)*
 - Deliver the media: *RTP and RTCP*
 - Playout the media, as reliably and quickly as you can: buffers etc

Session Initiation Protocol (SIP)

- Open (IETF) protocol to establish and tear-down calls (rfc3261+++)
 - Doesn't care how you transport the media
 - Is not used by Skype, Messenger, WhatsApp, Facetime, Zoom, Lync, ...
 - Competes with ITU H.323
- Widely used for Voice-over-IP (VoIP) = Internet Telephony
- Includes proxies, registrars, redirectors, border controllers, and gateways
 - Useful for “mobile” users, large directories, NATs, PSTN connections, etc.

SIP signalling

- Looks a lot like a phone call
- Looks a lot like HTTP
 - Commands, Responses, Code-classes
 - Runs over TCP and UDP
 - (and not – no XYZ codes?)



Non-realtime realtime?

- One-way media transmission: Streaming
 - Less interactive
 - Play out from a file, not necessarily a capture-device
 - Less sensitive to delays (“almost live”)
 - Still have issues with bandwidth and jitter
- Now manage playout buffers through sliding windows
 - Receiver ‘pulls’ content
 - Fills its buffers as content is played out
 - (Progress bars on video clients)
 - If bandwidth is not sufficient, client pauses – or other magic happens

Real Time Streaming Protocol (RTSP)

- Establishes a streaming session and negotiates media transport
 - rtp/rtcp
 - http, ...
- Looks a lot like HTTP. And SIP.
 - OPTIONS (what can you do?)
 - DESCRIBE (what can this file give me?)
 - SETUP (get ready to send one or more streams, over protocol P)
 - PLAY (play, from time T1 to time T2)
 - ...and a dozen more...
- Over RTP/RTCP can adapt bandwidth, encoding, ...

Or *sigh* use HTTP

- **Because it gets through firewalls**
- Because it has so many extensions
 - It's an application and a transport protocol
- Use HEAD requests to learn about media and options
- Use GET with Range requests
 - Download pieces for the playout buffer
 - Can ask server to encode adaptively
- Note:
 - No server state required!
 - Inherit HTTP proxies, caches and CDNs!
- HTML5 has video player built-in (to kill Flash)

In closing

- Perfect, low-delay, real-time over a best-effort network is really hard
 - Very application-specific.
 - Brains are adaptive, robots less so.
- All of the smarts is in the end-points
 - Unless you build circuits (RSVP, QoS, ...)
- Audio/video works 'ok'
 - In your context...
- Increasing use of real-time traffic for device control
 - But maybe the Internet isn't the best choice
 - But maybe it's the only choice