Relational Query Processing	
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Spring 2020	
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How Queries Get Processed	
Start from an SQL query:	
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Start from an SQL query: User Interface	
Select * from A, B, C Where A al = B.b1 and B.b2 = C.c2 and A a0 = "123" (command Line) (command Line)	
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- Start from an SQL query:
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- Transform this query to its equivalent relational algebra expression:

User Interface (Command Line)

Query Parser and

Query Rewrite

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(Command Line) Query Parser and Query Rewrite

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User Interface (Command Line) Query Parser and Query Rewrite Query Optimizer

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 - σ_{A.a0 = "123"}(A _{A.a1 = B.b1} B _{B.b2 = C.c2} C)
- Optimize this expression
 - (σ_{A.a0 = "123"}(A)_{A.a1 = B.b1} B) _{B.b2 = C.c2} C)



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 - (σ_{A.a0} = "123"(A)_{A.a1} = B.b1 B) B.b2 = C.c2 C)
- Generate an equivalent query evaluation

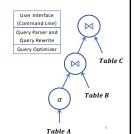
(Command Line) Query Rewrite Query Optimizer

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How Queries Get Processed

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- Generate an equivalent query evaluation plan:



How Queries Get Processed	
• Start from an SQL query: • Select * from A, B, C Where • Aa1 = Bb1 and B.b2 = C.c2 and • Aa0 = "123"	User Interface (Command Line) Query Parser and Query Rewrite
 Transform this query to an equivalent relational algebra expression: σ_{A,a0="123"}(A_{Aa1=Bb1} B_{B,b2=C,c2} C) 	
• Question: • Is the above transformation always	

Conceptual Evaluation of SQL Queries into Relational Algebra Expressions

- Given an SQL query:

 - Select attr1, attr2, . . ., attrn

 From Table1, Table2, . . ., Tablem

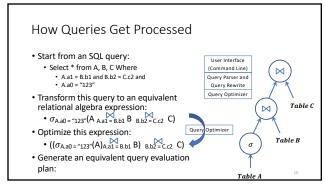
 Where Pred1 and Pred2 and . . . and Pred8
- The equivalent Relational Algebra expression:
 - π attr1, attr2, ..., attrn (σ (Pred1 \land Pred2 \land ... \land Predk)(Table1 X Table2 X ... X Tablem))

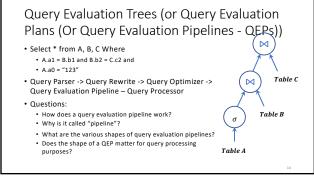
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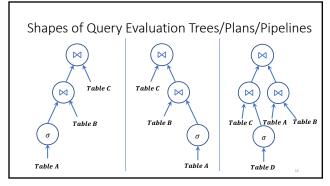
How Queries Get Processed • Start from an SQL query: User Interface (Command Line) Select * from A, B, C Where A.a1 = B.b1 and B.b2 = C.c2 and A.a0 = "123" Query Parser and Query Rewrite Query Optimizer • Transform this query to an equivalent relational algebra expression: • $\sigma_{\text{A-a0}} = \text{``123''}(\text{A}_{\text{A-a1}} = \text{B,b1} \text{B}_{\text{B,b2}} = \text{C,c2} \text{C})$ Table C • Optimize this expression: Table B • (σ_{A.a0} = "123"(A)_{A.a1} = B.b1 B) _{B.b2} = C.c2 C) • Generate an equivalent query evaluation

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plan:





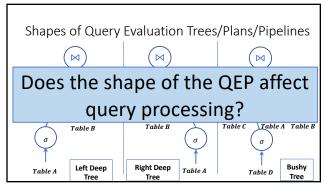


Shapes	of Query	Evaluation	n Trees	/Plans/Pipelines
	M			M
N N	Table C	Table C		M
σ	Table B	Table B	σ	Table C Table A Table B
Table A	Left Deep Tree	Right Deep Tree	Table A	Table D Bushy

assume right-deep tree has a temporary table and temporary table has p tuples. Then the cost of IO is

c*p(c joins with temp table)+ p(save temporary table to disk) + ab(iterate all a,b tuples to construct temporary table)

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The Iterator Model to Evaluate Relational Algebra Operators

- The iterator model:
 - Process one tuple at a time
 - Follow the lazy evaluation scheme
 - Is a pull-based model
 - Does not produce a tuple until it is asked
 - In contrast to "push-based model"
 - Produce tuples without being asked
- Each relational algebra operator is implemented using the iterator model

How to Code an Operator in the Iterator Model?

- Each operator has the following interface methods:
 Open():
 Open the operator
 Initialize the operator and its internal state variables

 - GetNext():
 Start from the current state of the operator
 Execute the operator from its current state until is able to retrieve the next tuple
 Report the next tuple

 - Advance and store the current state for the next execution
 If no tuples are available for return (e.g., reached end-of-table), report done/null/end-of-table

 - Close():
 Close the operator
 - Other auxiliary operators:
 HasNext(), Reset() or ReOpen(), among others

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Physical Operators and their Realizations as Iterators

- The Table Scan Operator:
 - Open():
 - Open the table
 - Set the cursor at the beginning of the table
 GetNext():
 - - . Get the next tuple from the table

 - Advance the cursor to the next tuple
 If no tuple exists, then return null or end-of-table
 Else return the tuple

 - Close():
 Close the underlying table

 - Reset():
 Reset the cursor to point to the beginning of the table

Open GetNext GetNext GetNext Table Scan Table A

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Physical Operators and their Realizations as Iterators: The Select Operator

- The Select Operator. $\sigma_{P}(\ldots)$:
 - Upon construction at compile time, Predicate P of the select is passed to the operator to be embedded in it (or passed as a Hint. More on this later).
 • Open():
 - - Issue an Open() operator to the operator under it (could be a table scan operator or any other operator underneath

1. Open 4. Done σ 3. Done 2. Open Here can be any other

operator below or a table-scan

Physical Operators and their Realizations as Iterators: The Select Operator The Select Operator: GetNext(): Try to getNext() from the operator below the select -->(t) Check P on t If true, return t and exit Es repeat the above until a tuple is found or operator below reports Done. Close(): Call Close() for the operator below the select Return done Here can be any other

operator below or a table-scan

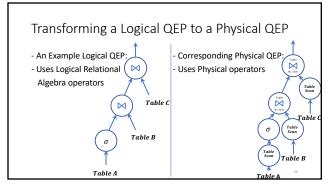
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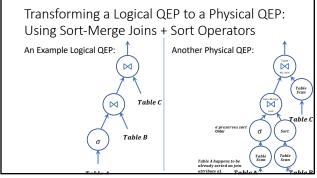
Physical Operators and their Realizations as Iterators: The Join Operator The Join Operator: The Join Operator: The join predicate will be plugged into the inside of the operator at compile time or will be passed as a hint (more on this issue later). State = left_tuple and right_tuple Open(): Open(): State = left_tuple and right_tuple to be null Gentlett(): If left_tuple and right_tuple to be null Gentlett(): If left_tuple and right_tuple to be null Gentlett(): If left_tuple and, then Get the neat tuple from the left operator underneath Whelleth operator did not reach end, of_table do If gentlett(): If left_tuple is null, then Get the neat tuple from the left operator underneath Whelleth operator did not reach end, of_table do If gentlett(): If left_tuple is null, then Get the neat tuple from the left operator underneath Whelleth operator did not reach end, of_table do If gentlett(): If left_tuple is null, then Get the neat tuple from the left operator underneath Whelleth operator did not reach end, of_table do If gentlett(): If left_tuple and right_tuple is null, then Get the neat tuple from the left operator underneath Whelleth operator did not reach end, of_table do If gentlett(): If left_tuple and right_tuple is null, then Get the neat tuple from the left operator underneath Whelleth operator did not reach end, of_table do If the tuple and right_tuple and right operators below the join Between the complete is not the neat tuple from the left operators below or table-scains

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Putting the Operators Together to Construct a Query Evaluation Pipeline

- Connect the operators together to form a QEP.
- Each operator does NOT know the operator underneath it
 And it does not need to know.
- Each operator follows the iterator model protocol (open, getnext, close).
- Thus, operators can communicate successfully and move tuples around to answer the query.

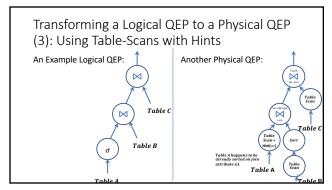


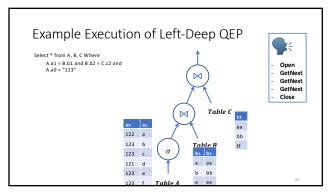


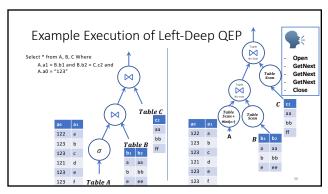
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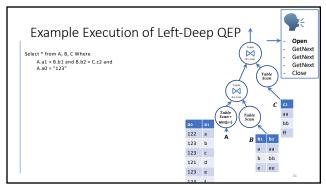
Physical Operators

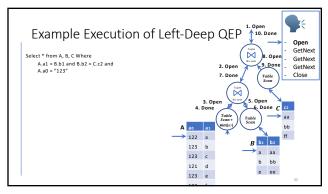
- · Table-Scan with select predicate as a hint
- Index-Scan (Useful for index-only plans)
 Index-Scan with select predicate as a hint
- Temp (Create a temporary table out of the incoming tuples from the lower-level operator)
- Create-index (Create index on the fly)
 Table-Copy (Make copy of table)
- Page-oriented NL-join, Block NL-join, Indexed NL-join, Hash join, Sort-Merge Join
- Bulk-load table from file
- Insert, delete, update
- Collect/update statistics Check some integrity constraint

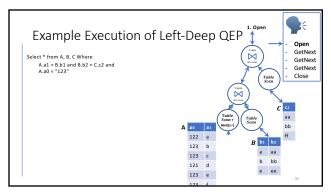


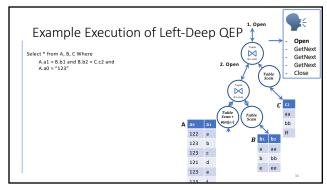


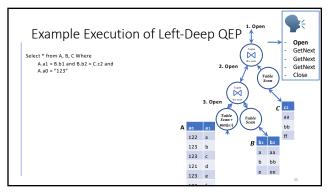


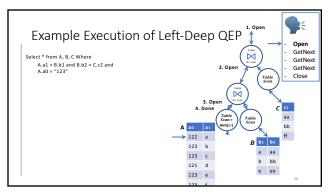


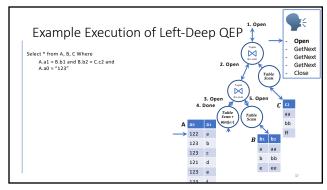


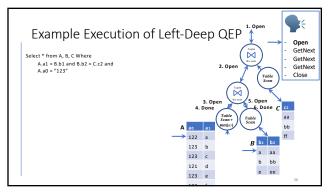


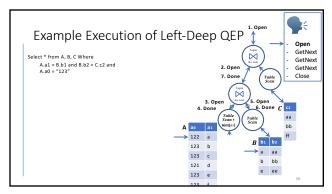


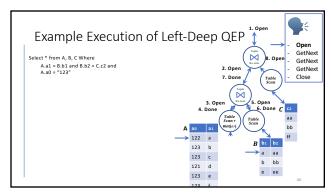


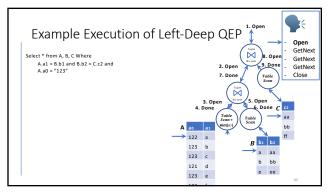


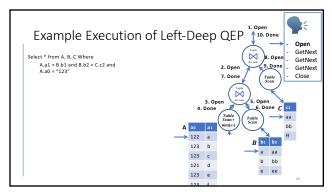


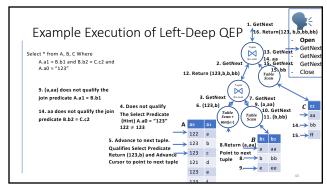


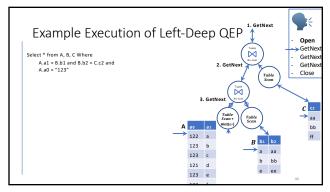


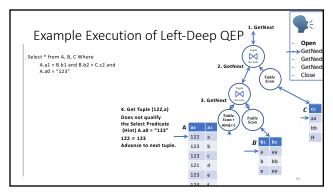


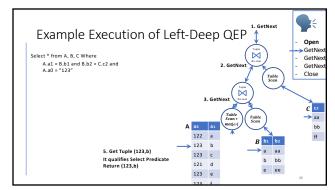


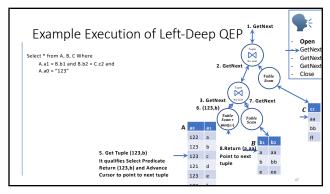


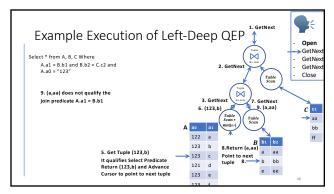


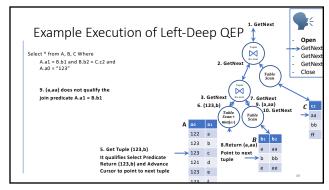


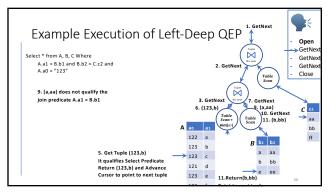


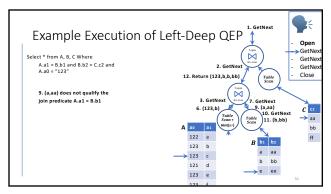


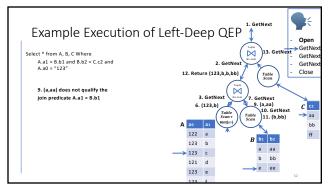


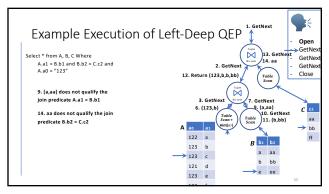


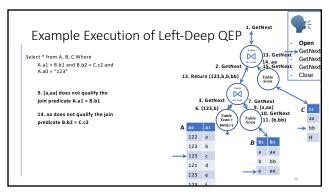


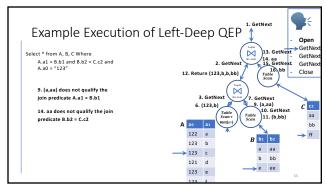


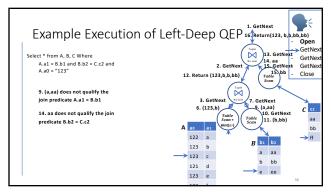


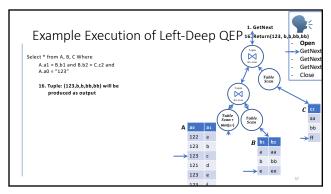


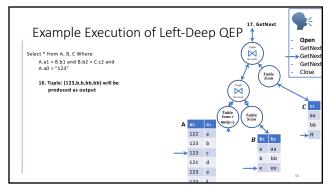


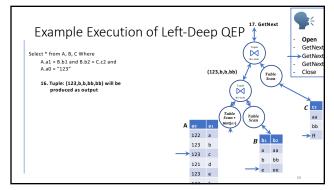


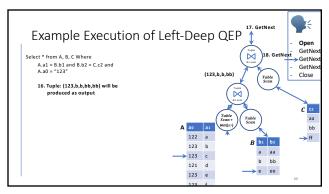


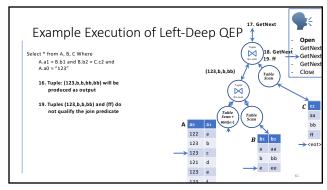


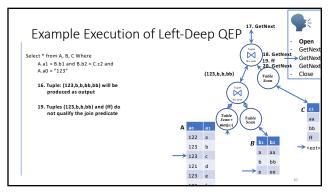


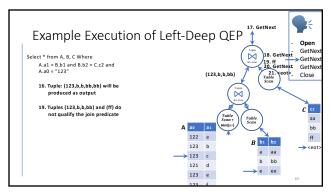


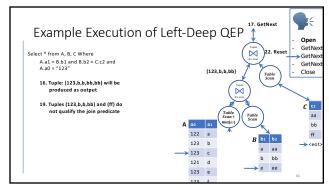


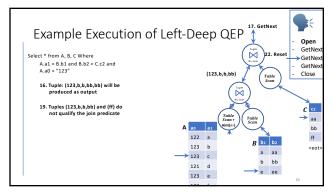


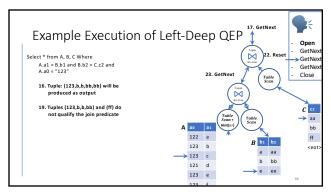


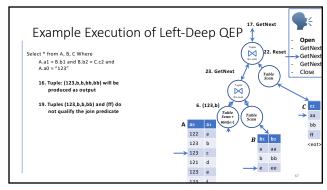


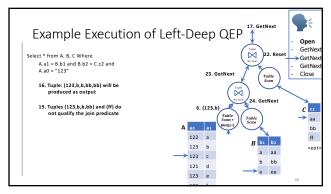


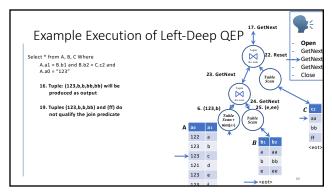


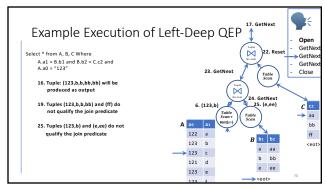


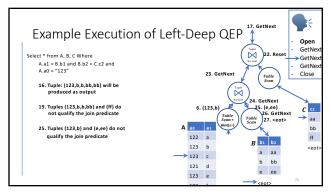


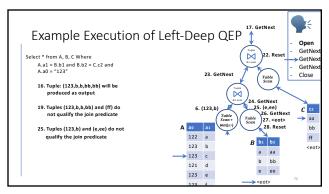


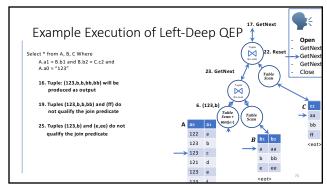


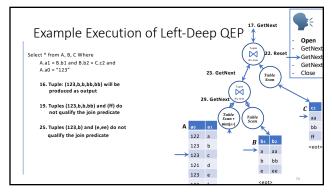


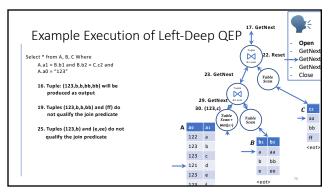


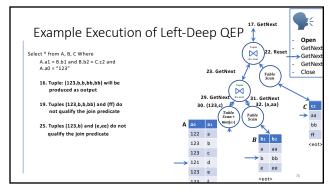


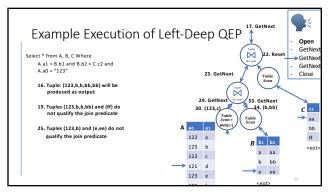


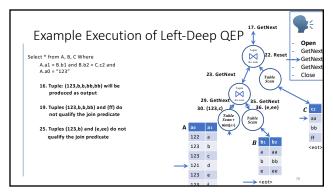


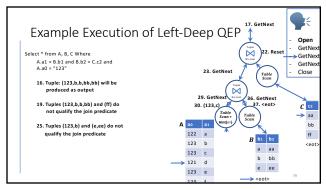


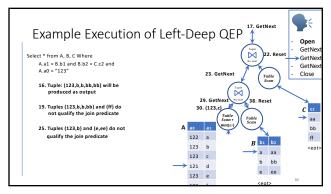


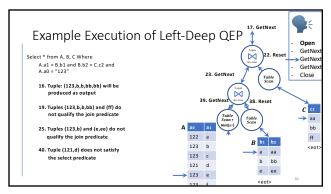


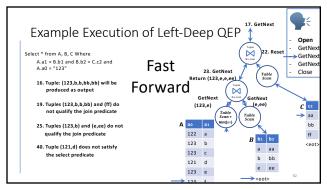


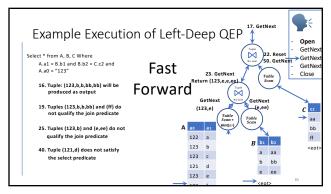


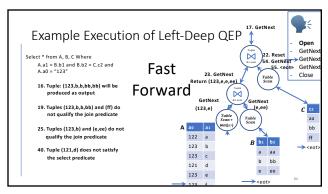


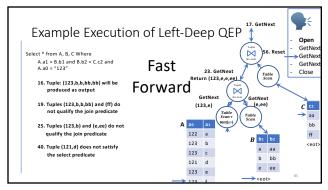


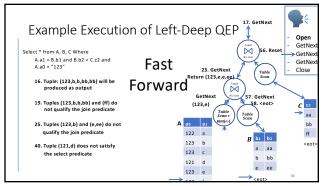


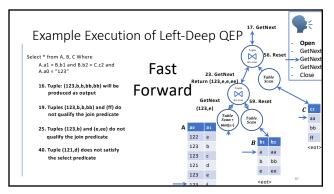


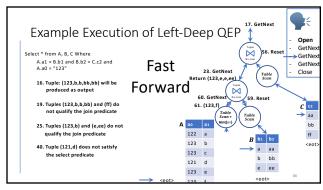


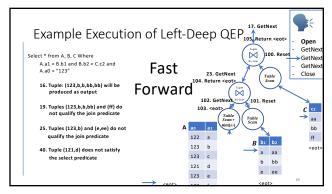


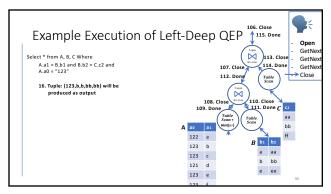


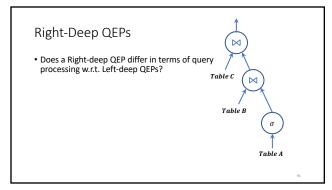


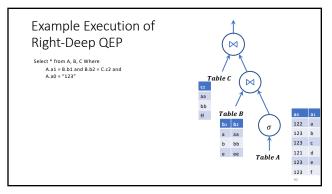


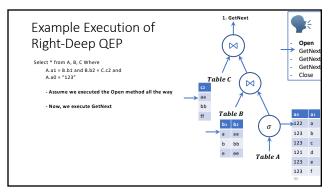


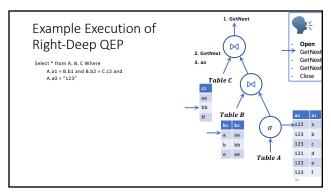


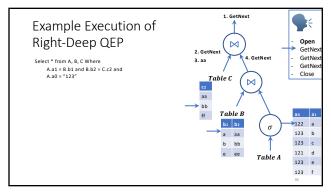


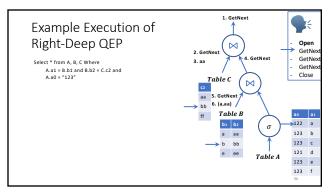


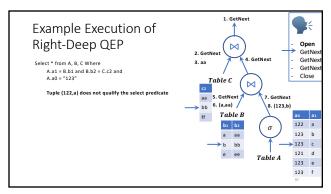


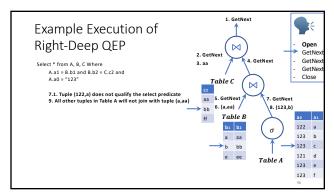


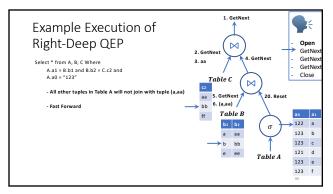


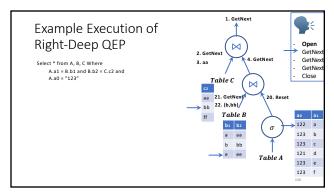


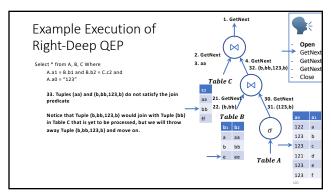


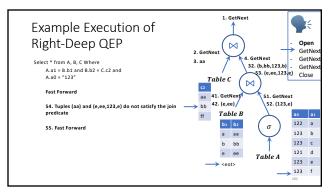


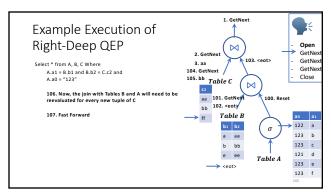


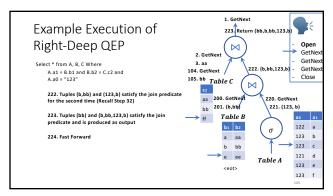


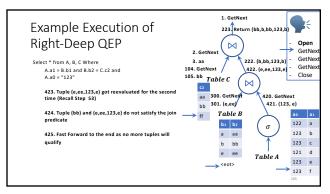


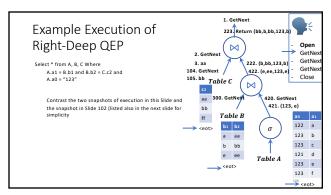


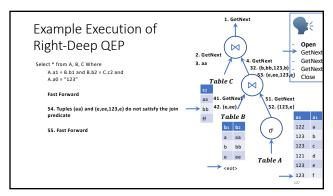


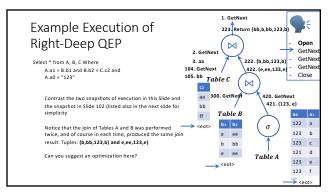










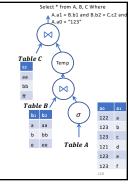


Optimizing Right-deep QEPs • Will add a temp operator on top of each join operator: Table C • Perform the once and save its results Next time, instead of re-evaluating it, access the temp table Table B ff Select * from A, B, C Where bı bı 122 a A.a1 = B.b1 and B.b2 = C.c2 and A.a0 = "123" 123 b a aa 123 с bb 121 d 123 e 123 f

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Optimizing Right-deep QEPs

- Note the addition of the Temp operator above the join of B and A
- Temp will create a temporary table that is the result of the subtree below it
 Temp = B B.D1=Al1 (OA.B0="123"(A))
- The temporary table is created the first time the Temp operator is invoked

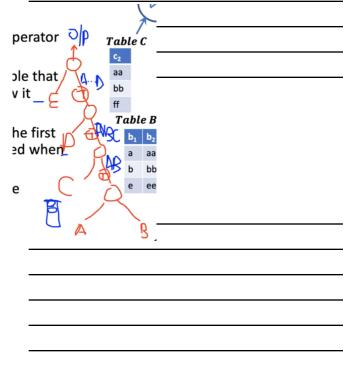


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How to optimize right deep QEP

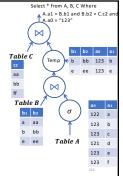
By adding a temporary table(result of subtree below it) above a join operator. So we don't need to iterate to get every join permutation.

i.e. If there is a query that E joins D joins C joins A and B.
We-add a temporary table above each join operator



Optimizing Right-deep QEPs

- Note the addition of the Temp operator above the join of B and A
- Temp will create a temporary table that is the result of the subtree below it
 Temp = B _{8.b1} = 1, a1 (σ_{A.80} = "123"(A))
- The temporary table is created the first time the Temp operator is invoked when processing Tuple (aa) of Table C
- For the remaining tuples in C, the temporary table is accessed



Comparison between Left- and Right-deep QEPs

- Left-deep Trees are the more favorable
- Allow for non-blocking operators
- Do not need to write temporary tables on disk
- Do not have to read the temporary tables from disk over and over for every tuple of the outer tables
- Make the query optimization job easier as we do not have to search over all possible trees that can evaluate the query
- \bullet But still exponential as we need to know the order of the joins (which join to perform first).

рас	e to store temporary results
•	

However, temporary table is blocking (get all results from subtree until done) and need