

# Quiz (27 Marks)

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Points:

19/27

1

Only OS independent compiler is  
(0.5/0.5 Points)

- ☐ Java compiler
- ☐ Visual basic compiler
- ☐ Turbo C compiler
- ☐ Pascal compiler

2

Question

(2/2 Points)

Consider the translation scheme shown below

$S \rightarrow TR$

$R \rightarrow +T \{ \text{print}(' + '); \} R | \epsilon$

$T \rightarrow \text{num} \{ \text{print}(\text{num.val}); \}$

Here **num** is a token that represents an integer and **num.val** represents the corresponding integer value. For an input string '9 + 5 + 2', this translation scheme will print

- ☐ 9+52+
- ☐ 95+2+
- ☐ 95++2
- ☐ 952++

3

Which languages necessarily need heap allocation in the runtime environment?  
(1/1 Point)

- ☐ Those that support recursion

- ☐ Those that not allow dynamic data structure
- ☐ Those that use dynamic scoping
- ☐ none of these

4

Which of the following is true for line number 3 of given C code.

```
int main ( ) {          /* Line 1 */

int l, N;              /* Line 2 */

for (l = 0, l < N, l++); /* Line 3 */

}
```

(1/1 Point)

- ☐ Only syntactic errors
- ☐ No compilation error
- ☐ Both lexical and syntactic errors
- ☐ Only lexical errors

5

In a bottom-up evaluation of a syntax directed definition, inherited attributes can

(1/1 Point)

- ☐ Be evaluated only if the definition has synthesized attributes
- ☐ Never be evaluated
- ☐ none of these
- ☐ Be evaluated only if the definition is S-attributed

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In the following grammar:

$A \rightarrow B \oplus A/B$

$B \rightarrow C * B/Z$

$Z \rightarrow id$

Which of the following is true?

(0/2 Points)

- ☐ Both ' $\oplus$ ' and '\*' are left associative

- ☐  $\oplus$  is right associative while  $*$  is left associative
- ☐ Both  $\oplus$  and  $*$  are right associative
- ☐  $\oplus$  is left associative while  $*$  is right associative

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Question

(0/1 Point)

Consider the grammar shown below

$$S \rightarrow iEtSS'|a$$

$$S' \rightarrow eS|\epsilon$$

$$E \rightarrow b$$

In the predictive parse table M of this grammar, the entries  $M[S',e]$  and  $M[S',\$]$  respectively are

- ☐  $\{S' \rightarrow eS\}$  and  $\{\}$
- ☐  $\{S' \rightarrow eS, S' \rightarrow \epsilon\}$  and  $\{\}$
- ☐  $\{S' \rightarrow eS\}$  and  $\{S' \rightarrow \epsilon\}$
- ☐  $\{S' \rightarrow \epsilon\}$  and  $\{S' \rightarrow \epsilon\}$

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For the following SDTS, the output printed by a bottom-up parser for the input "aab" is:

$$S \rightarrow aA \text{ \{print 1\}}$$

$$S \rightarrow a \text{ \{print 3\}}$$

$$A \rightarrow Sb \text{ \{print 2\}}$$

(0/2 Points)

- ☐ 1 2 3
- ☐ 2 2 3
- ☐ 2 3 1
- ☐ 3 2 1

Consider that for a grammar G; SLR and LALR parsers have X and Y states, respectively. Then the relationship between X and Y is

(1/1 Point)

- ☐ X is surely less than Y
- ☐ X is surely greater than Y
- ☐ X is surely equal to Y
- ☐ No relationship between X and Y

The 3-address code sequence for the '**A := B + C**' is

(1/1 Point)

<b>S</b> → <b>id</b> := E	{ <i>gen</i> ( <b>id</b> .place = E.place;;)}
E → E <sub>1</sub> + E <sub>2</sub>	{ <i>t</i> = <i>newtemp</i> ( );
	<i>gen</i> ( <i>t</i> = E <sub>1</sub> .place + E <sub>2</sub> .place;;)
	E.place = <i>t</i> ;} {E.place = <b>id</b> .place;}
E → id	

Here, *gen* is a function that generates the output code, and *newtemp* is a function that returns the name of a new temporary variable on every call. Assume that *ti*'s are the temporary variable names generated by *newtemp*.

- ☐ t1:= B; t2 := C; A:= t1+t2
- ☐ t1 := B + C; A: = t1
- ☐ t1 := B + C;
- ☐ A := B + C

Question

(3/3 Points)

For an input '5 \* 6 + 7' the following translation scheme prints

$$S \rightarrow AB$$
$$B \rightarrow *A \{ \text{print}('*'); \} B | \epsilon$$
$$A \rightarrow C+A \{ \text{print}('+'); \} | C$$
$$C \rightarrow S | \text{id} \{ \text{print}(\text{id.value}); \}$$

- ☐ 5 \* + 6 7
- ☐ 5 6 7+\*
- ☐ 5 6\* + 7
- ☐ 5 \* 6 + 7

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A bottom-up parser generates  
(0.5/0.5 Points)

- ☐ Right-most derivation in reverse
- ☐ Left-most derivation in reverse
- ☐ Left-most derivation
- ☐ Right –most derivation

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Question

(3/3 Points)

What is the translation of **xxxxyz** using the following SDT scheme

$$S \rightarrow xxW \{ \text{print}'1' \}$$
$$S \rightarrow y \{ \text{print}'2' \}$$
$$W \rightarrow Sz \{ \text{print}'3' \}$$

- ☐ 1 1 2 3 3
- ☐ 2 3 1 3 1
- ☐ 1 1 2 3 1
- ☐ 2 3 1 1 3

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In context of SLR(1) and LALR (1) tables which is/are false?  
(0/2 Points)

- ☐ The error entries in the tables may be same
- ☐ Shift entries are not identical in both the tables
- ☐ The reduce entries may be same
- ☐ Goto part of both tables may be same

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Match the followings:

### Group-I

- (a) Pointer data type
- (b) Activation Record
- (c) Repeat -Until
- (d) Coercion

### Group-II

- (p) Type Conversion
- (q) Dynamic Data Structure
- (r) Nondeterministic loop
- (s) Recursion

(1/1 Point)

- ☐ a-p, b-r, c-s, d-q
- ☐ a-q, b-s, c-r, d-p
- ☐ a-r, b-q, c-s, d-p
- ☐ a-q, b-r, c-s, d-p

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Compiler translates the source code to

(2/2 Points)

- ☐ Machine code
- ☐ Executable code
- ☐ None of these
- ☐ Binary code

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The Semantic Analysis is/are responsible for ....

(0/1 Point)

- ☐ Check semantics
- ☐ Type checking

☐ Static checking

☐ All of these

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Question

(1/1 Point)

The grammar  $A \rightarrow AA \mid (A) \mid \epsilon$  is not suitable for predictive-parsing because the grammar is

**(A)** ambiguous

**(B)** left-recursive

**(C)** right-recursive

**(D)** an operator-grammar

☐ Only (A)

☐ Only (A), (B), and (C)

☐ Only (B) and (C)

☐ All

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Which phase/phases check the Grammar of the programming.

(1/1 Point)

☐ Code generation

☐ Syntax analysis

☐ Semantic analysis

☐ Code optimization