

Test of Distributed Systems

Lecture 11: Linear Temporal Logic

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What do we expect?

Axioms

Proof

Exploiting duality

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Should these statements be true?

- ▶ $\Box p \leftrightarrow \neg \Diamond \neg p$ (duality)
- ▶ $\Box(p \rightarrow q) \rightarrow (\Box p \rightarrow \Box q)$ (distributivity)
- ▶ $\Box p \rightarrow p$ (reflexivity)
- ▶ $p \rightarrow \Diamond p$ (reflexivity)
- ▶ $\Box p \rightarrow \Box \Box p$ (transitivity)
- ▶ $\Diamond p \rightarrow \Diamond \Diamond p$ (transitivity)
- ▶ $\Box p \rightarrow \bigcirc p$ (step)
- ▶ $\bigcirc p \rightarrow \Diamond p$ (step)
- ▶ $\bigcirc p \leftrightarrow \neg \bigcirc \neg p$ (linearity)

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Axioms of LTL

1. Predicate Calculus
2. $\Box(p \rightarrow q) \rightarrow (\Box p \rightarrow \Box q)$ (distributivity of \Box over \rightarrow)
3. $\bigcirc(p \rightarrow q) \rightarrow (\bigcirc p \rightarrow \bigcirc q)$ (distributivity of \bigcirc over \rightarrow)
4. $\Box p \rightarrow (p \wedge \bigcirc p \wedge \bigcirc \Box p)$ (expansion of \Box)
5. $\Box(p \rightarrow \bigcirc p) \rightarrow (p \rightarrow \Box p)$ (induction)
6. $\bigcirc p \leftrightarrow \neg \bigcirc \neg p$ (linearity)

Inference rules:

$$\frac{p, p \rightarrow q}{q}$$

modus ponens

$$\frac{p}{\Box p}$$

$$\frac{p \rightarrow q}{\Box p \rightarrow \Box q}$$

generalization

$$\frac{p \rightarrow q}{\bigcirc p \rightarrow \bigcirc q}$$

$$\frac{p \rightarrow \bigcirc p}{p \rightarrow \Box p}$$

induction

Define:

- $\Diamond p \leftrightarrow \neg \Box \neg p$ (duality)

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\Box is transitive and reflexive

Theorem: $\Box p \leftrightarrow \Box \Box p$.

Proof:

$$\Box p \leftrightarrow \Box \Box p$$

{PC}

- $\Box p \rightarrow \Box \Box p$

{induction}

$$\Box p \rightarrow \bigcirc \Box p$$

{expansion}

- $\Box \Box p \rightarrow \Box p$

{expansion}

\Box distributes over \wedge

Theorem: $\Box(p \wedge q) \rightarrow (\Box p \wedge \Box q)$.

Proof:

$$\Box(p \wedge q) \rightarrow (\Box p \wedge \Box q)$$

{PC}

- $\Box(p \wedge q) \rightarrow \Box p$
{generalization}

$$p \wedge q \rightarrow p$$

{PC}

- $\Box(p \wedge q) \rightarrow \Box q$
{generalization}

$$p \wedge q \rightarrow q$$

{PC}

\Box contraction

Theorem: $p \wedge \bigcirc \Box p \rightarrow \Box p$.

Proof:

$$p \wedge \bigcirc \Box p \rightarrow \Box p$$

{PC}

$$p \wedge \bigcirc \Box p \rightarrow \Box p \wedge \Box \bigcirc \Box p$$

{distribution over \wedge }

$$p \wedge \bigcirc \Box p \rightarrow \Box(p \wedge \bigcirc \Box p)$$

{induction}

$$p \wedge \bigcirc \Box p \rightarrow \bigcirc(p \wedge \bigcirc \Box p)$$

{PC}

$$\bigcirc \Box p \rightarrow \bigcirc(p \wedge \bigcirc \Box p)$$

{generalization}

$$\Box p \rightarrow p \wedge \bigcirc \Box p$$

{expansion}

\Diamond current

Theorem: $p \rightarrow \Diamond p$.

Proof:

$$p \rightarrow \Diamond p$$

{definition}

$$p \rightarrow \neg \Box \neg p$$

{PC}

$$\neg \neg p \rightarrow \neg \Box \neg p$$

{PC}

$$\Box \neg p \rightarrow \neg p$$

{expansion}

◇ next

Theorem: $\bigcirc p \rightarrow \Diamond p$.

Proof:

$$\bigcirc p \rightarrow \Diamond p$$

{definition}

$$\bigcirc p \rightarrow \neg \Box \neg p$$

{linearity}

$$\neg \bigcirc \neg p \rightarrow \neg \Box \neg p$$

{PC}

$$\Box \neg p \rightarrow \bigcirc \neg p$$

{expansion}

\Box implies \Diamond

Theorem: $\Box p \rightarrow \Diamond p$.

Proof:

$$\Box p \rightarrow \Diamond p$$

{PC}

- $\Box p \rightarrow \bigcirc p$

{expansion}

- $\bigcirc p \rightarrow \Diamond p$

{ \Diamond next}

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Distribution

Law: $\Box(p \wedge q) \leftrightarrow (\Box p \wedge \Box q)$.

Theorem: $\Diamond(p \vee q) \leftrightarrow (\Diamond p \vee \Diamond q)$.

Proof:

$$\Diamond(p \vee q) \leftrightarrow (\Diamond p \vee \Diamond q)$$

{definition}

$$\neg\Box\neg(p \vee q) \leftrightarrow (\neg\Box\neg p \vee \neg\Box\neg q)$$

{PC}

$$\neg\Box(\neg p \wedge \neg q) \leftrightarrow \neg(\Box\neg p \wedge \Box\neg q)$$

{PC}

$$\Box(\neg p \wedge \neg q) \leftrightarrow (\Box\neg p \wedge \Box\neg q)$$

{the law above}

More laws

- ▶ $\Diamond p \leftrightarrow p \vee \bigcirc \Diamond p$
- ▶ $\Box(p \rightarrow q) \rightarrow (\Diamond p \rightarrow \Diamond q)$ { generalization }
- ▶ $\Diamond(p \rightarrow q) \rightarrow (\Diamond p \rightarrow \Diamond q)$
- ▶ $\Box \bigcirc p \leftrightarrow \bigcirc \Box p$ { exchange }
- ▶ $\Diamond \bigcirc p \leftrightarrow \bigcirc \Diamond p$ { exchange }
- ▶ $\Diamond \Diamond p \leftrightarrow \Diamond p$

Contraction

Theorem: $\Box\Diamond\Box p \leftrightarrow \Diamond\Box p$.

Proof:

$$\Box\Diamond\Box p \leftrightarrow \Diamond\Box p$$

{PC}

- $\Box\Diamond\Box p \rightarrow \Diamond\Box p$

{expansion}

- $\Diamond\Box p \rightarrow \Box\Diamond\Box p$

{induction}

$$\Diamond\Box p \rightarrow \bigcirc\Diamond\Box p$$

{exchange}

$$\Diamond\Box p \rightarrow \Diamond\bigcirc\Box p$$

{generalization}

$$\Box p \rightarrow \bigcirc\Box p$$

{expansion}

More laws

Theorem (Contraction): $\Diamond\Box\Diamond p \leftrightarrow \Box\Diamond p$.

Proof: exercise.

Theorem: $\Diamond\Box p \rightarrow \Box\Diamond p$.

Proof: exercise.

Theorem: $\Box\Diamond p \rightarrow \Diamond\Box p$.

Proof?

Counterexample!