

SetlX-Quickreference

asdf Element-Placeholder asdf exemplary Implementation
 ... continual Element-List
asdf optional Element

variables	a ;	→ Name has to start with a lowercase letter. Apart from that it can contain any letters, numbers and _ → They can't get statically typed
assignment	a := b ;	→ if the number contains . it will be automatically recognized as a real number
strings	"asdf"	→ can be concatenated via + → the modifiers for lists (see below) are also useable for them
literal strings	'asdf'	→ turns off all processing, e.g. '\n' will be saved as \n in characters, rather than getting processed into a newline-character
undefined Ω	om	
placeholder	—	→ use it if you have to provide a variable for a call because of its syntax but actually don't need this variable
comments	// asdf /* multiple-line asdf */	
output	print(asdf , asdf , ...);	→ you can insert expressions like \$3+2\$ which will be evaluated when printing the output-string
input	a := read("asdf");	→ Prints the argument into the prompt and returns the user-input

rational numbers:

→ they work without overflows and in theory indefinitely accurate because they are stored as fractions
 → 1/3 + 1/2 would return 5/6

different types of functions:

procedure	asdf := procedure(v1 , v2 , ...) { ... return r ; };	
cached / memorized procedure	asdf := cachedProcedure(v1 , ...) { ... return r ; };	→ speeds up computation by saving results of the function in-memory in a lookup-table → only allowed for <i>pure functions</i> a pure functions always returns the same output if it is called with the same input
closure	asdf := closure(v1 , v2 , ...) { ... r := extVar * 2; return r ; };	→ works like a procedure → additionally you are able to access variables which are defined outside the function
mathematical function	f := x -> def ; f := x -> 1.0/(1+x); a := f (2); a = 1/3	→ equals to $f: x \rightarrow def$ which equals to $f(x) = expr$ → useable via a := f (n);
lambda closure- definition	f := in => expression ; f := [in1 , in2] => expression ;	→ equivalent to f := closure(in) { return expression ; } → equivalent to f := closure(in1 , in2) { return expression ; }
call	f (arg);	

control structures:

if-branching	if (test1) { body1 ; } else if (test2) { body2 ; } else { body3 ; }	→ the brackets are always necessary!
switch-branching	switch { case test1 : body1 ; case test2 : body2 ;	→ only one body gets executed

	<pre> ... default : body3; } </pre>	
--	-------------------------------------	--

while-loop	<pre> while (test) { body; } </pre>	
for-loop	<pre> for (i in m) { body; } </pre>	→ iterates through the elements of the set/list like <code>m[i]</code>
abort one iteration	<code>continue;</code>	
abort the loop completely	<code>break;</code>	

predefined real ("reelle") **functions:**

trigonometric	<code>sin(x)</code>	
	<code>asin(x)</code>	equals to $\sin^{-1}(x)$
	<code>sinh(x)</code>	sinus hyperbolises
	<code>cos(x)</code>	
	<code>acos(x)</code>	equals to $\cos^{-1}(x)$
	<code>cosh(x)</code>	cosine hyperbolises
	<code>tan(x)</code>	
	<code>atan(x)</code>	equals to $\tan^{-1}(x)$
	<code>tanh(x)</code>	tangent hyperbolises
exponential	<code>exp(x)</code>	equals to e^x
	<code>x ** a</code>	equals to x^a
logarithmic	<code>log(x)</code>	equals to $\ln(x)$ (natural logarithmic)
	<code>log10(x)</code>	equals to $\log_{10}(x)$
absolute value	<code>abs(x)</code>	equals to $ x $
sign	<code>signum(x)</code>	
square root	<code>sqrt(x)</code>	
3rd-root	<code>cbert(x)</code>	
round up	<code>ceil(x)</code>	rounds up to the next integral number
round down	<code>floor(x)</code>	rounds down to the next integral number
round to nearest	<code>round(x)</code>	also known in german as "kaufmännisches Runden"

sets:

definition by enumeration	<code>{start .. stop}</code>	→ equals to $\{x \in \mathbb{Z} \mid start \leq x \wedge x \leq stop\}$ → any element is only contained once and elements are ordered by their value
definition by step-enumeration	<code>{start, second .. stop}</code>	→ equals to $\{start + n * step \mid n \in \mathbb{N}_0 \wedge start + n * step \leq stop\}$ with $step = second - start$
definition by iterators	<code>{definition : ranges}</code> <code>{n * m : n in {2..10}, m in {2..10}};</code>	→ the set then contains the non-trivial Solutions for the condition which meet the ranges for their elements → equals to $\{n * m \mid n \in \mathbb{N} \wedge 2 \leq n \wedge n \leq 10 \wedge 2 \leq m \wedge m \leq 10\} = \{4, 6, 8, 9, 10, 12, 14, 15, 16, 18, 20, 21, 24, 25, 27, 28, 30, 32, 35, 36, 40, 42, 45, 48, 49, 50, 54, 56, 60, 63, 64, 70, 72, 80, 81, 90, 100\}$
additionally: selection	<code>{definition condition}</code>	→ only elements which fulfill the additional condition are added to the set
summation	<code>+ / m</code>	→ returns the sum of all elements in the set M

product	\ast / m	→ returns the product of all elements in the set M
element-count	$\#(m)$	→ returns the number of elements contained in the set
union $a \cup b$	$a + b$	
intersection $a \cap b$	$a \ast b$	
difference a / b	$a - b$	
power 2^a	$2 \ast \ast a$	
Cartesian product $A \times B$	$a >< b$	
powerset	$\text{pow}(m)$	→ returns the set which contains all possible subsets of m

is a subset $a \subseteq b$	$a <= b$	
is an element $a \in M$	$a \text{ in } m$	
get the element with the highest value	$\text{max}(m)$	
get the element with the lowest value	$\text{min}(m)$	
take a (not pre-defined) element	$\text{from}(m)$	→ Returns a kind of random element from the set: At first, you don't know which one it will be. But if you run the program again, the order of the returned elements is exactly the same. → Removes the element from the set!
get a (not pre-defined) element	$\text{arb}(m)$	→ works like from, but doesn't remove the element from the set
get a (pseudo-) random element	$\text{rnd}(m)$ $\text{rnd}(5)$	→ bad for debugging → computes a random natural number less or equal then 5, via the implicated call $\text{rnd}([1..5])$

general tuples / lists:

- They can be defined and used just like sets.
- $\{\}$ in the definition then become $[\]$
they are definable through enumeration, iterators and selection
→ e.g. a pair $\langle x, y \rangle$ is definable through $[x, y]$
- differences to sets: elements are not ordered and can be contained multiple times

reverse it	$\text{reverse}(l)$	
sort it	$\text{sort}(l)$	→ sorts the elements in the list in ascending order of their values
element-reference	$m[i]$	→ returns the i^{th} element out of the set (ordered ascending by value) $m[-1]$ returns the last, $m[-2]$ the pre-last element and so on → the counting of elements starts at 1!
Subset-reference	$m[a..b];$	→ returns the sub-set of m starting at index a and ending on index b → one of the limits can be omitted
append it to itself	$n \ast l$	
concatenate them as a string	$\text{join}(l, s)$ $\text{join}([1,2,3, \ast \ast]) \quad \Rightarrow "1 \ast 2 \ast 3"$	→ converts the elements of l into strings and concatenates them using the string s as a separator

relations:

definition of relations	$\{[pair-Def] : Condition\}$ $\{[n, n \ast \ast 2] : n \text{ in } \{1..10\}\};$	→ equals to the Function $x \rightarrow x^2$ on the set $\{n \in \mathbb{N} \mid 1 \leq n \wedge n \leq 10\}$ $\Rightarrow \{[1, 1], [2, 4], [3, 9], [4, 16], [5, 25], [6, 36], [7, 49], [8, 64], [9, 81], [10, 100]\}$
domain	$\text{domain}(m)$	
range	$\text{range}(m)$	

logical expressions:

boolean test-operators	$==$	
	$!=$	

	<div><</div> <div><=</div> <div>></div> <div>>=</div> <div>in</div>	<div>also checks \subseteq at sets</div>	
test-junctures	<div>!</div> <div>&&</div> <div> </div>	<div>equals to \neg strongest bind</div> <div>equals to \wedge</div> <div>equals to \vee weakest bind</div>	
all-quantifier	<code>forall(x in m cond)</code>	\rightarrow equal to $\forall x \in m : cond$	
exists-quantifier	<code>exists(x in m cond)</code>	\rightarrow equal to $\exists x \in m : cond$	
implication	<code>a => b</code>		
equivalence	<code>a <==> b</code>		
antivalence	<code>a <!=> b</code>		
convert strings	<code>eval(expr)</code>	\rightarrow the string <code>expr</code> has to be a string which can be parsed as a SetIX-Expression \rightarrow the result of the evaluation of the represented expressions is them returned	

terms:

symbolic Programs = programs/procedures which take functions (contained in strings) and manipulate them
a program which takes strings like "x*3" and finds the derivate of them

function-symbols	A	\rightarrow the name has to start with a uppercase letter. Apart from that it can contain any letters, numbers and _
	<code>^Asdf</code>	\rightarrow used internally to define operators like +
	<code>@asdf</code>	\rightarrow used to case (lowercase) built-in functions into a function-symbol
terms	<code>funcSymbol(value1, value2, ...)</code> <code>Adresse("Musterstr 1", 23456, "Musterstadt")</code>	\rightarrow Terms are never evaluated! They are only used to store data.
undefined	<code>Nil();</code>	
get the function-symbol	<code>fct(Asdf(value))</code>	
get the values/ argument-list	<code>args(Asdf(value))</code>	

matching:

match-branches	<pre> match (Term0) { case pattern1 : body1; case pattern2 : body2; ... default : body3; } match (Pl(3, 5)) { case Pl(t1, t2) : return"\$t1+t2\$ +"; case Mi(t1, t2) : return"\$t1-t2\$ -"; } </pre> \rightarrow "8 via +"	\rightarrow Instead of a Term, a String or a List can also be matched \rightarrow The patterns have to contain placeholders for variables \rightarrow At the evaluation, SetIX tries to insert the values of <code>Term0</code> into these placeholders to create a (new) term which is equal to <code>Term0</code> . If this succeeds, the corresponding body gets executed (with the specified variables filled accordingly).. If not, the default-body gets executed.
----------------	---	--

vectors:

definition	<code>v1 := la_vector([1, 1/2, 1/3]);</code> <code>v2 := <<1 1/2 1/3>>;</code> $\rightarrow v1 == v2 == <<1.0 0.5 0.333333333333>>$ note: 1/3 is only printed rounded	\rightarrow only real valued vectors, but with any dimension, are supported \rightarrow all vectors are column-vectors, via concept
accessor	<code>v[i]</code>	\rightarrow gives the <i>i</i> -th element of the vector back
addition / subtraction	<code>v := v1 + v2;</code> <code>v := v1 - v2;</code> <code>v += v1;</code> <code>v -= v2;</code>	
scalar multiplication	<code>v := <<1 1/2 1/3>> * (1/2);</code> <code>v *= (1/3);</code>	\rightarrow * is commutative
scalar product	<code>v := v1 * v2;</code>	
cross product	<code>v := v1 >< v2;</code>	\rightarrow only defined for three-dimensional products

matrices:

definition	<pre>m1 := la_matrix([[1,2],[3,4]]); m2 := << <<1 2>> <<3 4>> >>; → m1 == m2 == << <<1.0 2.0>> <<3.0 4.0>> >></pre>	→ only real valued matrices are supported
transforming vectors	<pre>v := <<1 2 3>>; m1 := la_matrix(v);</pre>	→ returns an $n \times 1$ -matrix → the column-vector gets transformed into a one-row-matrix
addition / subtraction	<pre>m := m1 + m2; m := m1 - m2; m += m1; m -= m2;</pre>	
scalar multiplication	<pre>m := << <<1 2>> <<3 4>> >> * (3); m *= (1/3);</pre>	→ * is commutative
matrix multiplication	<pre>a * b; a * v;</pre>	→ only possible if a is a $m \times n$ -matrix and b is a $n \times k$ -matrix (returning a $n \times k$ -matrix) → if v is a n-dimensional vector, it automatically is interpreted as an nx1-matrix and the result will be converted to an m-dimensional vector
exponentiation	<pre>a ** 2;</pre>	→ only possible for square matrices
inverse	<pre>a ** -1;</pre>	→ only possible for non-singular matrices
transposing	<pre>a!;</pre>	
Dimension m	<pre>#a;</pre>	
Dimension n	<pre>#a[1];</pre>	
Determinant	<pre>la_det(a);</pre>	→ the result might be a small non-zero value, even if the matrix is really singular (due to rounding errors)