variables	a;	 Names have to start with any letter and after that can contain any letters, numbers and _ → They are not statically typed
assignment	a := b;	\rightarrow if the number contains . it will be automatically recognized as a real number
strings	"asdf"	→ can be concatenated via + → the modifiers for lists (see below) are also useable for them
literal strings	'asdf'	→ turns of all processing, e.g. '\n' will be saved as \n in characters, rather than getting processed into a newline-character
undefined Ω	om	
placeholder	_	→ use it if you have to provide a variable for a call because of its syntax but actually don't need this variable
comments	<pre>// asdf /* multiple-line asdf */</pre>	
output	<pre>print(asdf, asdf,);</pre>	 → you can insert expressions between two \$, which will be evaluated when printing the output-string → print("The answer is \$6*7\$!");
input	a := read("asdf");	→ Prints the argument into the prompt and returns the user-input

rational numbers:

 \rightarrow they work without overflows and in theory indefinitely accurate because they are stored as fractions \rightarrow 1/3 + 1/2 would return 5/6

different types of functions:

,, ,,	, ,	
procedure	<pre>asdf := procedure(v1, v2,) { return r; };</pre>	
cached / memorized procedure	<pre>asdf := cachedProcedure(v1,) { return r; };</pre>	 → speeds up computation by saving results of the function in-memory in a lookuptable → only allowed for pure functions a pure functions always returns the same output if it is called with the same input
closure	<pre>asdf := closure(v1, v2,) { r := extVar * 2; return r; };</pre>	 → works like a procedure → additionally you are able to access variables which are defined outside the function
lambda procedure definition	f := x -> def; f := x -> 1.0/(1+x); a := f(2);	\rightarrow equals to $f: x \rightarrow def$ which equals to $f(x) = def$. In the non-mathmatical realm they are identical to $f:=$ procedure(x) { return def; } \rightarrow useable via $f(n)$;
lambda closure- definition	f := in => def; f := [in1, in2] => def;	<pre>→equivalent to f := closure(in) { return def; } →equivalent to f := closure(in1, in2) { return def; }</pre>
default argument	fkt(a := 2) { }	→ fkt = closure / procedure / cachedProcedure
call-by-reference	fkt(rw a) { }	→ fkt = closure / procedure / cachedProcedure → passes a as a reference (a "link" to the original variable) which allows the function to read and write the original variable
call	f(arguments)	

control structures:

```
If-branching
                    if (test1) {
                                                                 \rightarrow the brackets are always necessary!
                    body1;
} else if (test2) {
                       body2;
                    } else {
                     body3;
                    switch {
  case test1 : body1;
  case test2 : body2;
switch-branching
                                                                 \rightarrow only one body gets executed
                        default : body3;
                    }
while-loop
                    while (test) {
                        body;
                    }
                    for (i in m) {
                                                                 → iterates through the elements of the set/list like m[i]
for-loop
                        body;
abort one iteration continue;
abort the loop
                    break;
completely\\
```

predefined real ("reelle") functions:

preaejmea re			
trigonometric	sin(x)		
	asin(x)	equals to $sin^{-1}(x)$	
	sinh(x)	sinus hyperbolises	
	cos(x)		
	acos(x)	equals to $cos^{-1}(x)$	
	cosh(x)	cosine hyperbolises	
	tan(x)		
	atan(x)	equals to $tan^{-1}(x)$	
	tanh(x)	tangent hyperbolises	
exponential	exp(a)		equals to e^a
	x ** a		equals to x^a
logarithmic	log(x)		equals to $ln(x)$ (natural logarithmic)
	log10(x)		equals to $log_{10}(x)$
absolute value	abs(x)		equals to $ x $
sign	signum(x)		returns -1.0 or 0.0 or 1.0
square root	sqrt(x)		
3rd-root	cbrt(x)		
round up	ceil(x)		rounds up to the next integral number
round down	floor(x)		rounds down to the next integral number
round to nearest	round(x)		also known in German as "kaufmännisches Runden"

sets:

3613.		
definition by enumeration	{start stop}	\rightarrow equals to $\{x \in \mathbb{Z} \mid start \le x \land x \le stop\}$ \rightarrow any element is only contained once and elements are ordered by their value
definition by step- enumeration	{start, second stop}	$ \rightarrow \text{equals to } \{start + n*step \mid n \in \mathbb{N}_0 \ \land start + n*step \leq stop \} \text{ with } step = second - start $
definition by iterators	{definition : ranges} {n * m : n in {210}, m in {210}};	→ the set then contains the non-trivial Solutions for the condition which meet the ranges for their elements → equals to $\{n*m \mid n \in \mathbb{N} \land 2 \le n \land n \le 10 \land 2 \le m \land m \le 10\} = \{4, 6, 8, 9, 10, 12, 14, 15, 16, 18, 20, 21, 24, 25, 27, 28, 30, 32, 35, 36, 40, 42, 45, 48, 49, 50, 54, 56, 60, 63, 64, 70, 72, 80, 81, 90, 100\}$
additionally: selection	{definition condition}	→ only elements which fulfil the additional condition are added to the set
summation	+/m	→ returns the sum of all elements in the set M
product	*/m	→ returns the product of all elements in the set M
element-count	#m	→ returns the number of elements contained in the set
union $a \cup b$	a + b	
intersection $a \cap b$	a * b	
difference a/b	a - b	
power 2 ^a	2 ** a	
Cartesian $A \times B$ product	a >< b	
powerset	pow(m)	→ returns the set which contains all possible subsets of m
is a set	isSet(a)	→ returns true or false
is a subset $a \subseteq b$	a <= b	
is an element $a \in M$	a in m	
get the element with the highest value	max(m)	
get the element with the lowest value	min(m)	
take an element from the start / the end of the list	<pre>fromE(m) fromB(m)</pre>	→ Removes the element from the set!
take a (not pre- defined) element	from(m)	 → Returns a kind of random element from the set: At first, you don't know which one it will be. But if you run the program again, the order of the returned elements is exactly the same. → Removes the element from the set!
get a (not pre- defined) element	arb(m)	→ works like from, but doesn't remove the element from the set
get a (pseudo-) random element	<pre>rnd(m) rnd(5) random()</pre>	→ bad for debugging → computes a random non-negative number less or equal then 5, via the implicated call rnd([15]) → computes a random non-negative number froim the range [0, 1]
concatenate as a string	<pre>join(1, s) join([1,2,3, "*")</pre>	\rightarrow converts the elements of 1 into strings and concatenates them using the string s as a separator

general tuples / lists:

- \rightarrow They can be defined and used just like sets.
- \rightarrow { } in the definition then become []

- they are definable through enumeration, Iterators and selection \rightarrow e.g. a pair $\langle x, y \rangle$ is definable through [x, y] \rightarrow differences to sets: elements are not ordered and can be contained multiple times and the following functions are additionally available

reverse it	reverse(1)	
sort it	sort(1)	\rightarrow sorts the elements in the list in ascending order of their values
check if a variable holds a list	isList(a)	→ returns true or false
element-reference	m[i]	→ returns the i th element out of the set (ordered ascending by value) m[-1] returns the last, m[-2] the pre-last element and so on → the counting of elements starts at 1!
Subset-reference	m[ab]	\rightarrow returns the sub-set of m starting at index a and ending on index b \rightarrow one of the limits can be omitted
appends 1 n-times it to itself	1 * n	→ n needs to be a natural number or zero

relations:

	{[pair-Def] : Condition} {[n, n**2] : n in {110}};	→ equals to the Function $x \to x^2$ on the set $\{n \in \mathbb{N} \mid 1 \le n \land n \le 10\}$ $\mathbb{B}\{[1, 1], [2, 4], [3, 9], [4, 16], [5, 25], [6, 36], [7, 49], [8, 64], [9, 81], [10, 100]\}$
domain	<pre>domain(m)</pre>	
range	range(m)	

logical expressions:

logical expres.	310113.		
boolean	==		
test-operators	!=		
	<		
	<=	also checks ⊆ at sets	
	>		
	>=		
	in		
	notin		
test-junctures	!	equals to ¬ strongest bind	
	&&	equals to Λ	
	П	equals to V weakest bind	
all-quantifier	forall((x in m condition)	\rightarrow equal to $\forall x \in m : condition$
exists-quantifier	exists((x in m condition)	$ ightarrow$ equal to $\exists x \in m : condition$
implication	a => b		
equivalence	a <==> b		
antivalence	a = b		
convert strings	eval(ex	xpr)	→ the string expr has to be a string which can be parsed as a SetIX-Expression → the result of the evaluation of the represented expressions is them returned

terms:

 $symbolic\ Programs = programs/procedures\ which\ take\ functions\ (contained\ in\ strings)\ and\ manipulate\ them$ $a\ program\ which\ takes\ strings\ like\ "x*3"\ and\ finds\ the\ derivate\ of\ them$

function symbols	@Asdf	\to the name has to start with @, followed by any letter. Apart from that it can contain any letters, numbers and _
	@@@asdf	→ used internally to define operators like +
terms	<pre>@funcSymbol(value1, value2,) Adresse("Musterstr 1", 23456, "Musterstadt")</pre>	→ Terms are never evaluated! They are only used to store data.
undefined	@Nil()	
get the function- symbol	<pre>fct(Asdf(value))</pre>	
get the values/ argument-list	args(Asdf(value))	

matching:

```
match-branches
                      match (Term0) {
                                                                         \rightarrow Instead of a Term, a String or a List can also be matched
                          case pattern1 : body1;
                                                                         \rightarrow The patterns have to contain placeholders for variables
                           case pattern2 : body2;

ightarrow At the evaluation, SetIX tries to insert the values of Term0 into these placeholders
                                                                           to create a (new) term which is equal to Term0.
                          default : body3;
                                                                           If this succeeds, the corresponding body gets executed (with the specified
                      }
                                                                           variables filled accordingly)..
                                                                           If not, the default-body gets executed.
                      match (P1(3, 5)) {
   case P1(t1, t2) : return"$t1+t2$ +";
   case Mi(t1, t2) : return"$t1-t2$ -";
                      }

→ "8 via +"
```

(end of the lecture "Grundlagen und Logik")

vectors:

definition	v1 := la_vector([1, 1/2, 1/3]); v2 := <<1 1/2 1/3>>; > v1 == v2 == <<1.0 0.5 0.333333333333>> note: 1/3 is only printed rounded	 → only real valued vectors, but with any dimension, are supported → all vectors are column-vectors, via concept → it is possible to add a , between the values, altough this would archive nothing. The componets of a vector are really seperated by the whitspace between them
accessor	v[i]	ightarrow gives the $i-th$ element of the vector back
addition / subtraction	<pre>v := v1 + v2; v := v1 - v2; v += v1; v -= v2;</pre>	
scalar multiplication	v := <<1 1/2 1/3>> * (1/2); v *= (1/3);	→* is commutative
scalar product	v := v1 * v2;	
cross product	v := v1 >< v2;	→ only defined for three-dimensional products

matrices:

definition	m1 := la_matrix([[1,2],[3,4]]); m2 := << <<1 2>> <<3 4>> >>;	→ only real valued matrices are supported
transforming vectors	<pre>v := <<1 2 3>>; m1 := la_matrix(v);</pre>	\rightarrow returns an $n \times 1$ -matrix \rightarrow the column-vector gets transformed into a one-row-matrix
addition / subtraction	m := m1 + m2; m := m1 - m2; m += m1; m -= m2;	
scalar multiplication	m := << <<1 2>> <<3 4>> >> * (3); m *= (1/3);	→* is commutative
matrix multiplication	a * b; a * v;	\rightarrow possible if a is a $m \times n$ -matrix and b is a $n \times k$ -matrix (returns a $n \times k$ -matrix) \rightarrow if v is a n-dimensional vector, it automatically is interpreted as an nx1-matrix and the result will be converted to an m-dimensional vector
exponentiation	a ** 2;	→ only possible for square matrices
inverse	a ** -1;	→ only possible for non-singular matrices
transposing	a!;	
Dimension m	#a;	
Dimension n	#a[1];	
Determinant	la_det(a);	→ the result might be a small non-zero value, even if the matrix is really singular (due to rounding errors)

manual error-handling:

handling exceptions	<pre>try { // normal statements } catch (e) { // error-handling }</pre>	→ also possible: catchUsr and catchLng
throwing exceptions	<pre>throw("message");</pre>	\rightarrow it is strongly advised to then use catchUsr(e) to handle the exception risen by throw, to avoid masking of exceptions thrown by the interpreter

debugging:

ggg.		
tracing	<pre>trace(true); trace(false);</pre>	ightarrow all assignments written in this area, will be "documented" in the console-output
watch variables	<pre>stop("message");</pre>	 → stops the execution, prints the provided message, and waits until you press Enter without an input → if you enter All, the value of all available variables in the current scope will be printed
test assertions	<pre>assert(condition, "message");</pre>	 → if the condition evaluates to true, nothing happens → otherwise, the execution gets terminated and the provided message is printed