

SetlX-Quickreference

asdf Element-Placeholder asdf exemplary Implementation
... continual Element-List
asdf optional Element

variables	a ;	→ Name has to start with a lowercase letter. Apart from that it can contain any letters, numbers and _ → They can't get statically typed
assignment	a := b ;	→ if the number contains . it will be automatically recognized as a real number
strings	"asdf"	→ can be concatenated via + → the modifiers for lists (see below) are also useable for them
literal strings	'asdf'	→ turns off all processing, e.g. '\n' will be saved as \n in characters, rather than getting processed into a newline-character
undefined Ω	om	
placeholder	—	→ use it if you have to provide a variable for a call because of its syntax but actually don't need this variable
comments	// asdf /* multiple-line asdf */	
output	print(asdf , asdf , ...);	→ you can insert expressions between two \$, which will be evaluated when printing the output-string → print("The answer is \$6*7\$!");
input	a := read("asdf");	→ Prints the argument into the prompt and returns the user-input

rational numbers:

→ they work without overflows and in theory indefinitely accurate because they are stored as fractions
 → 1/3 + 1/2 would return 5/6

different types of functions:

procedure	asdf := procedure(v1 , v2 , ...) { ... return r ; };	
cached / memorized procedure	asdf := cachedProcedure(v1 , ...) { ... return r ; };	→ speeds up computation by saving results of the function in-memory in a lookup-table → only allowed for <i>pure functions</i> a pure functions always returns the same output if it is called with the same input
closure	asdf := closure(v1 , v2 , ...) { ... r := extVar * 2; return r ; };	→ works like a procedure → additionally you are able to access variables which are defined outside the function
lambda procedure definition	f := x -> definition ; f := x -> 1.0/(1+x); a := f (2); a = 1/3	→ equals to $f: x \rightarrow \text{definition}$ which equals to $f(x) = \text{definition}$ → useable via f(n) ;
lambda closure- definition	f := in > expression ; f := [in1 , in2] > expression ;	→ equivalent to f := closure(in) { return expression ; } → equivalent to f := closure(in1 , in2) { return expression ; }
default argument	closure(a := 2) { ... }	
call	f (arguments);	

control structures:

If-branching	if (test1) { body1 ; } else if (test2) { body2 ; } else { body3 ; }	→ the brackets are always necessary!
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switch-branching	<pre>switch { case test1 : body1; case test2 : body2; ... default : body3; }</pre>	→ only one body gets executed
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while-loop	<pre>while (test) { body; }</pre>	
for-loop	<pre>for (i in m) { body; }</pre>	→ iterates through the elements of the set/list like <code>m[i]</code>
abort one iteration	<code>continue;</code>	
abort the loop completely	<code>break;</code>	

predefined real ("reelle") functions:

trigonometric	<code>sin(x)</code>	
	<code>asin(x)</code>	equals to $\sin^{-1}(x)$
	<code>sinh(x)</code>	sinus hyperbolises
	<code>cos(x)</code>	
	<code>acos(x)</code>	equals to $\cos^{-1}(x)$
	<code>cosh(x)</code>	cosine hyperbolises
	<code>tan(x)</code>	
	<code>atan(x)</code>	equals to $\tan^{-1}(x)$
	<code>tanh(x)</code>	tangent hyperbolises
exponential	<code>exp(a)</code>	equals to e^a
	<code>x ** a</code>	equals to x^a
logarithmic	<code>log(x)</code>	equals to $\ln(x)$ (natural logarithmic)
	<code>log10(x)</code>	equals to $\log_{10}(x)$
absolute value	<code>abs(x)</code>	equals to $ x $
sign	<code>signum(x)</code>	returns -1.0 or 0.0 or 1.0
square root	<code>sqrt(x)</code>	
3rd-root	<code>cbrt(x)</code>	
round up	<code>ceil(x)</code>	rounds up to the next integral number
round down	<code>floor(x)</code>	rounds down to the next integral number
round to nearest	<code>round(x)</code>	also known in German as "kaufmännisches Runden"

sets:

definition by enumeration	<code>{start .. stop}</code>	→ equals to $\{x \in \mathbb{Z} \mid start \leq x \wedge x \leq stop\}$ → any element is only contained once and elements are ordered by their value
definition by step-enumeration	<code>{start, second .. stop}</code>	→ equals to $\{start + n * step \mid n \in \mathbb{N}_0 \wedge start + n * step \leq stop\}$ with $step = second - start$
definition by iterators	<pre>{definition : ranges} {n * m : n in {2..10}, m in {2..10}};</pre>	→ the set then contains the non-trivial Solutions for the condition which meet the ranges for their elements → equals to $\{n * m \mid n \in \mathbb{N} \wedge 2 \leq n \wedge n \leq 10 \wedge 2 \leq m \wedge m \leq 10\} = \{4, 6, 8, 9, 10, 12, 14, 15, 16, 18, 20, 21, 24, 25, 27, 28, 30, 32, 35, 36, 40, 42, 45, 48, 49, 50, 54, 56, 60, 63, 64, 70, 72, 80, 81, 90, 100\}$
additionally: selection	<code>{definition condition}</code>	→ only elements which fulfil the additional condition are added to the set

summation	$+/\mathbf{m}$	→ returns the sum of all elements in the set M
product	$*/\mathbf{m}$	→ returns the product of all elements in the set M
element-count	$\#(\mathbf{m})$	→ returns the number of elements contained in the set
union	$a \cup b$	$\mathbf{a} + \mathbf{b}$
intersection	$a \cap b$	$\mathbf{a} * \mathbf{b}$
difference	a / b	$\mathbf{a} - \mathbf{b}$
power	2^a	$2 ** \mathbf{a}$
Cartesian product	$A \times B$	$\mathbf{a} >< \mathbf{b}$
powerset	$\text{pow}(\mathbf{m})$	→ returns the set which contains all possible subsets of m

is a set	<code>isSet(a);</code>	→ returns true or false
is a subset $a \subseteq b$	<code>a <= b</code>	
is an element $a \in M$	<code>a in m</code>	
get the element with the highest value	<code>max(m)</code>	
get the element with the lowest value	<code>min(m)</code>	
take a (not pre-defined) element	<code>from(m)</code>	→ Returns a kind of random element from the set: At first, you don't know which one it will be. But if you run the program again, the order of the returned elements is exactly the same. → Removes the element from the set!
get a (not pre-defined) element	<code>arb(m)</code>	→ works like from, but doesn't remove the element from the set
get a (pseudo-) random element	<code>rnd(m)</code> <code>rnd(5)</code>	→ bad for debugging → computes a random natural number less or equal then 5, via the implicated call <code>rnd([1..5])</code>

general tuples / lists:

- They can be defined and used just like sets.
- `{ }` in the definition then become `[]`
they are definable through enumeration, iterators and selection
→ e.g. a pair $\langle x, y \rangle$ is definable through `[x, y]`
- differences to sets: elements are not ordered and can be contained multiple times

reverse it	<code>reverse(l)</code>	
sort it	<code>sort(l)</code>	→ sorts the elements in the list in ascending order of their values
check if a variable holds a list	<code>isList(a);</code>	→ returns true or false
element-reference	<code>m[i]</code>	→ returns the i^{th} element out of the set (ordered ascending by value) <code>m[-1]</code> returns the last, <code>m[-2]</code> the pre-last element and so on → the counting of elements starts at 1!
Subset-reference	<code>m[a..b];</code>	→ returns the sub-set of <code>m</code> starting at index <code>a</code> and ending on index <code>b</code> → one of the limits can be omitted
append it to itself	<code>n * l</code>	
concatenate them as a string	<code>join(l, s)</code> <code>join([1,2,3, ""], "1*2*3")</code>	→ converts the elements of <code>l</code> into strings and concatenates them using the string <code>s</code> as a separator

relations:

definition of relations	<code>{[pair-Def] : Condition}</code> <code>{[n, n**2] : n in {1..10}};</code>	→ equals to the Function $x \rightarrow x^2$ on the set $\{n \in \mathbb{N} \mid 1 \leq n \wedge n \leq 10\}$ <code>[[1, 1], [2, 4], [3, 9], [4, 16], [5, 25], [6, 36], [7, 49], [8, 64], [9, 81], [10, 100]]</code>
domain	<code>domain(m)</code>	
range	<code>range(m)</code>	

. / .

logical expressions:

boolean test-operators	==		
	!=		
	<		
	<=	also checks \subseteq at sets	
	>		
	>=		
	in		
	notin		
test-junctures	!	equals to \neg strongest bind	
	&&	equals to \wedge	
		equals to \vee weakest bind	
all-quantifier	forall(x in m condition)	→ equal to $\forall x \in m : condition$	
exists-quantifier	exists(x in m condition)	→ equal to $\exists x \in m : condition$	
implication	a => b		
equivalence	a <==> b		
antivalence	a <!=> b		
convert strings	eval(expr)	→ the string expr has to be a string which can be parsed as a SetIX-Expression → the result of the evaluation of the represented expressions is them returned	

terms:

symbolic Programs = programs/procedures which take functions (contained in strings) and manipulate them
a program which takes strings like "x*3" and finds the derivate of them

function-symbols	A	→ the name has to start with a uppercase letter. Apart from that it can contain any letters, numbers and _
	^Asdf	→ used internally to define operators like +
	@asdf	→ used to case (lowercase) built-in functions into a function-symbol
terms	funcSymbol(value1, value2, ...) Adresse("Musterstr 1", 23456, "Musterstadt")	→ Terms are never evaluated! They are only used to store data.
undefined	Nil();	
get the function-symbol	fct(Asdf(value))	
get the values/ argument-list	args(Asdf(value))	

matching:

match-branches	<pre>match (Term0) { case pattern1 : body1; case pattern2 : body2; ... default : body3; } match (P1(3, 5)) { case P1(t1, t2) : return "\$t1+t2\$ +"; case Mi(t1, t2) : return "\$t1-t2\$ -"; } → "8 via +"</pre>	<p>→ Instead of a Term, a String or a List can also be matched</p> <p>→ The patterns have to contain placeholders for variables</p> <p>→ At the evaluation, SetlX tries to insert the values of Term0 into these placeholders to create a (new) term which is equal to Term0.</p> <p>If this succeeds, the corresponding body gets executed (with the specified variables filled accordingly)..</p> <p>If not, the default-body gets executed.</p>
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(end of the lecture "Grundlagen und Logik")

vectors:

definition	<pre>v1 := la_vector([1, 1/2, 1/3]); v2 := <<1 1/2 1/3>>; → v1 == v2 == <<1.0 0.5 0.333333333333>> note: 1/3 is only printed rounded</pre>	<p>→ only real valued vectors, but with any dimension, are supported</p> <p>→ all vectors are column-vectors, via concept</p>
accessor	<pre>v[i]</pre>	→ gives the <i>i-th</i> element of the vector back
addition / subtraction	<pre>v := v1 + v2; v := v1 - v2; v += v1; v -= v2;</pre>	
scalar multiplication	<pre>v := <<1 1/2 1/3>> * (1/2); v *= (1/3);</pre>	→ * is commutative
scalar product	<pre>v := v1 * v2;</pre>	
cross product	<pre>v := v1 >> v2;</pre>	→ only defined for three-dimensional products

matrices:

definition	<pre>m1 := la_matrix([[1,2],[3,4]]); m2 := << <<1 2>> <<3 4>> >>; → m1 == m2 == << <<1.0 2.0>> <<3.0 4.0>> >></pre>	→ only real valued matrices are supported
transforming vectors	<pre>v := <<1 2 3>>; m1 := la_matrix(v);</pre>	<p>→ returns an $n \times 1$-matrix</p> <p>→ the column-vector gets transformed into a one-row-matrix</p>
addition / subtraction	<pre>m := m1 + m2; m := m1 - m2; m += m1; m -= m2;</pre>	
scalar multiplication	<pre>m := << <<1 2>> <<3 4>> >> * (3); m *= (1/3);</pre>	→ * is commutative
matrix multiplication	<pre>a * b; a * v;</pre>	<p>→ only possible if a is a $m \times n$-matrix and b is a $n \times k$-matrix (returning a $n \times k$-matrix)</p> <p>→ if v is a n-dimensional vector, it automatically is interpreted as an nx1-matrix and the result will be converted to an m-dimensional vector</p>
exponentiation	<pre>a ** 2;</pre>	→ only possible for square matrices
inverse	<pre>a ** -1;</pre>	→ only possible for non-singular matrices
transposing	<pre>a!;</pre>	
Dimension <i>m</i>	<pre>#a;</pre>	
Dimension <i>n</i>	<pre>#a[1];</pre>	
Determinant	<pre>la_det(a);</pre>	→ the result might be a small non-zero value, even if the matrix is really singular (due to rounding errors)

manual error-handling:

handling exceptions	<pre>try { ... // normal statements } catch (e) { ... // error-handling }</pre>	→ also possible: catchUsr and catchLng
throwing exceptions	<pre>throw(e); throw("Left boundary a has to be less than right boundary b!");</pre>	→ it is strongly advised to then use catchUsr(e) to handle the exception risen by throw, to avoid masking of exceptions thrown by the interpreter

debugging:

tracing	<pre>trace(true); ... trace(false);</pre>	→ all assignments written in this area, will be "documented" in the console-output
watch variables	<pre>stop("message");</pre>	→ stops the execution, prints the provided message, and waits until you press Enter without an input → if you enter the name of a variable, its current value is printed into the console → if you enter All, the value of all available variables in the current scope will be printed
test assertions	<pre>assert(condition, "message");</pre>	→ if the condition evaluates to true, nothing happens → otherwise, the execution gets terminated and the provided message is printed