

Language Design and Overview of COOL

CS143
Lecture 2

Prof. Aiken CS 143 Lecture 2

1

Grade Weights

- Project 50%
 - I, II 10% each
 - III, IV 15% each
- Midterm 15%
- Final 25%
- Written Assignments 10%
 - 2.5% each

Prof. Aiken CS 143 Lecture 2

2

Lecture Outline

- Today's topic: language design
 - Why are there new languages?
 - Good-language criteria
- History of ideas:
 - Abstraction
 - Types
 - Reuse
- Cool
 - The Course Project

Prof. Aiken CS 143 Lecture 2

3

Programming Language Economics 101

- Languages are adopted to fill a void
 - Enable a previously difficult/impossible application
 - Orthogonal to language design quality (almost)
- Programmer training is the dominant cost
 - Languages with many users are replaced rarely
 - Popular languages become ossified
 - But easy to start in a new niche . . .

Prof. Aiken CS 143 Lecture 2

4

Why So Many Languages?

- Application domains have distinctive and conflicting needs
- Examples:

Prof. Aiken CS 143 Lecture 2

5

Topic: Language Design

- No universally accepted metrics for design
- Claim: “A good language is one people use”

Prof. Aiken CS 143 Lecture 2

6

Language Evaluation Criteria

Characteristic	Criteria		
	Readability	Writeability	Reliability
Data types	*	*	*
Abstraction		*	*
Type checking			*
Exception handling			*

Prof. Aiken CS 143 Lecture 2

7

History of Ideas: Abstraction

- Abstraction = detached from concrete details
- Abstraction necessary to build software systems
 - In fact, any complex system
- Modes of abstraction
 - Via languages/compiler:
 - Higher-level code, few machine dependencies
 - Via subroutines
 - Abstract interface to behavior
 - Via modules
 - Export interfaces; hide implementation
 - Via abstract data types
 - Bundle data with its operations

Prof. Aiken CS 143 Lecture 2

8

History of Ideas: Types

- Originally, few types
 - FORTRAN: scalars, arrays
 - LISP: no static type distinctions
- Realization: Types help
 - Allow the programmer to express abstraction
 - Allow the compiler to check against many frequent errors
 - Sometimes to the point that programs are guaranteed “safe”
- More recently
 - Lots of interest in types
 - Experiments with various forms of parameterization
 - Best developed in functional programming

Prof. Aiken CS 143 Lecture 2

9

History of Ideas: Reuse

- Reuse = exploit common patterns in software systems
 - Goal: mass-produced software components
 - Reuse is difficult
- Two popular approaches
 - Type parameterization (`List(int)`, `List(double)`)
 - Classes and inheritance: C++ derived classes
 - Combined in C++, Java
- Inheritance allows
 - Specialization of existing abstraction
 - Extension, modification, hiding behavior

Prof. Aiken CS 143 Lecture 2

10

Trends

- Language design
 - Many new special-purpose languages
 - Popular languages to stay
- Compilers
 - More needed and more complex
 - Driven by increasing gap between
 - new languages
 - new architectures
 - Venerable and healthy area

Prof. Aiken CS 143 Lecture 2

11

Why Study Languages and Compilers ?

5. Increase capacity of expression
4. Improve understanding of program behavior
3. Increase ability to learn new languages
2. Learn to build a large and reliable system
1. See many basic CS concepts at work

Prof. Aiken CS 143 Lecture 2

12

Cool Overview

- Classroom Object Oriented Language
- Designed to
 - Be implementable in a short time
 - Give a taste of implementation of modern
 - Abstraction
 - Static typing
 - Reuse (inheritance)
 - Memory management
 - And more ...
- But many things are left out

Prof. Aiken CS 143 Lecture 2

13

A Simple Example

```
class Point {  
  x : Int ← 0;  
  y : Int ← 0;  
};
```

- Cool programs are sets of class definitions
 - A special class **Main** with a special method **main**
 - No separate notion of subroutine
- class = a collection of attributes and methods
- Instances of a class are objects

Prof. Aiken CS 143 Lecture 2

14

Cool Objects

```
class Point {  
  x : Int ← 0;  
  y : Int; (* use default value *)  
};
```

- The expression “**new Point**” creates a new object of class **Point**
- An object can be thought of as a record with a slot for each attribute

x	y
0	0

Prof. Aiken CS 143 Lecture 2

15

Methods

- A class can also define methods for manipulating the attributes

```
class Point {  
  x : Int ← 0;  
  y : Int ← 0;  
  movePoint(newx : Int, newy : Int): Point {  
    { x ← newx;  
      y ← newy;  
      self;  
    } -- close block expression  
  }; -- close method  
}; -- close class
```

- Methods can refer to the current object using **self**

Prof. Aiken CS 143 Lecture 2

16

Information Hiding in Cool

- Methods are global
- Attributes are local to a class
 - They can only be accessed by the class's methods

- Example:

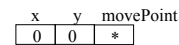
```
class Point {
    . . .
    x () : Int { x };
    setx (newx : Int) : Int { x ← newx };
};
```

Prof. Aiken CS 143 Lecture 2

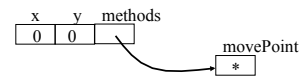
17

Methods

- Each object knows how to access the code of a method
- As if the object contains a slot pointing to the code



- In reality implementations save space by sharing these pointers among instances of the same class



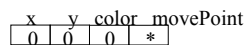
Prof. Aiken CS 143 Lecture 2

18

Inheritance

- We can extend points to colored points using subclassing => class hierarchy

```
class ColorPoint inherits Point {
    color : Int ← 0;
    movePoint(newx : Int, newy : Int): Point {
        { color ← 0;
          x ← newx; y ← newy;
          self;
        }
    };
};
```



Prof. Aiken CS 143 Lecture 2

19

Cool Types

- Every class is a type
- Base classes:
 - Int for integers
 - Bool for boolean values: true, false
 - String for strings
 - Object root of the class hierarchy
- All variables must be declared
 - compiler infers types for expressions

Prof. Aiken CS 143 Lecture 2

20

Cool Type Checking

```
x : A;
x ← new B;
```

- Is well typed if **A** is an ancestor of **B** in the class hierarchy
 - Anywhere an **A** is expected a **B** can be used
- Type safety:
 - A well-typed program cannot result in runtime type errors

Prof. Aiken CS 143 Lecture 2

21

Method Invocation and Inheritance

- Methods are invoked by dispatch
- Understanding dispatch in the presence of inheritance is a subtle aspect of OO languages

```
p : Point;
p ← new ColorPoint;
p.movePoint(1,2);
```

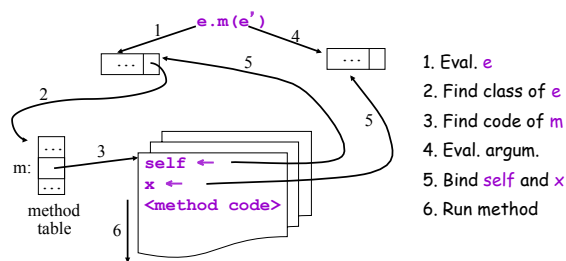
- **p** has static type **Point**
- **p** has dynamic type **ColorPoint**
- **p.movePoint** must invoke the **ColorPoint** version

Prof. Aiken CS 143 Lecture 2

22

Method Invocation

- Example: invoke one-argument method **m**



Prof. Aiken CS 143 Lecture 2

23

Other Expressions

- Expression language
 - every expression has a type and a value
 - Loops: **while E loop E pool**
 - Conditionals: **if E then E else E fi**
 - Case statement: **case E of x : Type ⇒ E; ... esac**
 - Arithmetic, logical operations
 - Assignment: **x ← E**
 - Primitive I/O: **out_string(s), in_string(), ...**
- Missing features:
 - arrays, floating point operations, exceptions, ...

Prof. Aiken CS 143 Lecture 2

24

Cool Memory Management

- Memory is allocated every time **new** is invoked
- Memory is deallocated automatically when an object is no longer reachable
- Done by the garbage collector (GC)
 - There is a Cool GC

Course Project

- A complete compiler
 - Cool ==> MIPS assembly language
 - No optimizations
- Split in 4 programming assignments (PAs)
- There is adequate time to complete assignments
 - But start early and please follow directions
- Individual or team
 - max. 2 students