

RSA (Rivest-Shamir-Adleman)

RSA was proposed by Rivest, Shamir & Adleman of MIT in 1977. It is a widely used public-key scheme. The algorithm uses number theory and modular arithmetic along with large integers. It is a very secure approach due to cost of factoring large numbers.

RSA Encryption/ Decryption

- For encrypting a message M the sender:
 - obtains public key of recipient $PU = \{e, n\}$
 - computes: $C = M^e \bmod n$, where $0 \leq M < n$
- To decrypt the ciphertext C the owner:
 - uses their private key $PR = \{d, n\}$
 - computes: $M = C^d \bmod n$
- It should be noted that the message M must be smaller than the modulus n

RSA Key Set up

Each user generates a public/private key pair by

- selecting two large primes at random: p, q
- computing their system modulus $n = p \times q$
 - note, $\phi(n) = (p-1) \times (q-1)$
- selecting at random the encryption key e
 - where $1 < e < \phi(n)$, $\gcd(e, \phi(n)) = 1$
- solving the following equation to find decryption key d
 - $e \times d = 1 \bmod \phi(n)$ and $0 \leq d \leq n$
 - By Euler's theorem, $e \times d = 1 + k \times \phi(n)$ for some k

The public encryption key is $PU = \{e, n\}$ which is published and the private decryption key is $PR = \{d, n\}$ which is kept secret.

RSA Example

- **Key Setup**
 - At first, p and q are selected; $p=17$ & $q=11$
 - Then, n is calculated. $n = p \times q = 17 \times 11 = 187$
 - $\phi(n)$ is calculate. $\phi(n) = (p-1) \times (q-1) = 16 \times 10 = 160$
 - e is selected so that $\gcd(e, 160) = 1$; $e=7$
 - d is determined. $d \times e = 1 \bmod 160$ & $0 \leq d \leq 187$. $d=23$ as, $23 \times 7 = 161 = 1 \times 160 + 1$
 - Public key $PU = \{7, 187\}$ is published.
 - Private key $PR = \{23, 187\}$ is kept secret.
- **Encryption/Decryption**

Sample RSA encryption/decryption is:

 - ❖ Given message $M = 88$ (nb. $88 < 187$)
 - ❖ Encryption: $C = 88^7 \bmod 187 = 11$
 - Exploiting the properties of modular arithmetic, following can be done:
 $88^7 \bmod 187 = [(88^4 \bmod 187) \times (88^2 \bmod 187) \times (88^1 \bmod 187)] \bmod 187$
 - ❖ Decryption: $M = 11^{23} \bmod 187 = 88$

RSA Example with text

$$p = 73, q = 151$$

$$n = 11023$$

$$\phi(n) = 10800$$

$$e = 11$$

$$d = 5891$$

Text = How are you?

H=	33	a=	00	y=	24
o=	14	r=	17	o=	14
w=	22	e=	04	u=	20
_	62	_	62	?=	66

$$M_1 = 3314 \quad M_2 = 2262 \quad M_3 = 0017$$

$$M_4 = 0462 \quad M_5 = 2414 \quad M_6 = 2066$$

$$C_1 = 3314^{11} \bmod 11023 = 10260$$

$$C_2 = 2262^{11} \bmod 11023 = 9489$$

$$C_3 = 17^{11} \bmod 11023 = 1782$$

$$C_4 = 462^{11} \bmod 11023 = 727$$

$$C_5 = 2414^{11} \bmod 11023 = 10032$$

$$C_6 = 2006^{11} \bmod 11023 = 2253$$

$$M_1 = 10260^{5891} \bmod 11023 = 3314$$

$$M_2 = 9489^{5891} \bmod 11023 = 2262$$

$$M_3 = 1782^{5891} \bmod 11023 = 0017$$

$$M_4 = 727^{5891} \bmod 11023 = 0462$$

$$M_5 = 10032^{5891} \bmod 11023 = 2414$$

$$M_6 = 2253^{5891} \bmod 11023 = 2006$$