

## **RSA (Rivest-Shamir-Adleman)**

RSA was proposed by Rivest, Shamir & Adleman of MIT in 1977. It is a widely used public-key scheme. The algorithm uses number theory and modular arithmetic along with large integers. It is very a very secure approach due to cost of factoring large numbers.

### **RSA Encryption/ Decryption**

- For encrypting a message M the sender:
  - obtains public key of recipient  $PU=\{e,n\}$
  - computes:  $C = M^e \bmod n$ , where  $0 \leq M < n$
- To decrypt the ciphertext C the owner:
  - uses their private key  $PR=\{d,n\}$
  - computes:  $M = C^d \bmod n$
- It should be noted that the message M must be smaller than the modulus n

### **RSA Key Set up**

Each user generates a public/private key pair by

- selecting two large primes at random: p, q
- computing their system modulus  $n = p \cdot q$ 
  - note,  $\phi(n)=(p-1)(q-1)$
- selecting at random the encryption key e
  - where  $1 < e < \phi(n)$ ,  $\gcd(e, \phi(n))=1$
- solving the following equation to find decryption key d
  - $e \cdot d \equiv 1 \bmod \phi(n)$  and  $0 \leq d \leq n$
  - By Euler's theorem,  $e \cdot d = 1 + k \cdot \phi(n)$  for some k

The public encryption key is  $PU=\{e,n\}$  which is published and the private decryption key is  $PR=\{d,n\}$  which is kept secret.

### **RSA Example**

- **Key Setup**
  - At first, p and q are selected;  $p=17$  &  $q=11$
  - Then, n is calculated.  $n = p \times q = 17 \times 11 = 187$
  - $\phi(n)$  is calculate.  $\phi(n)=(p-1) \times (q-1) = 16 \times 10 = 160$
  - e is selected so that  $\gcd(e, 160) = 1$ ;  $e=7$
  - d is determined.  $d \times e \equiv 1 \bmod 160$  &  $0 \leq d \leq 187$ .  $d=23$  as,  $23 \times 7 = 161 = 10 \times 160 + 1$
  - Public key  $PU= \{7,187\}$  is published.
  - Private key  $PR= \{23,187\}$  is kept secret.
- **Encryption/Decryption**

Sample RSA encryption/decryption is:

  - ❖ Given message  $M = 88$  (nb.  $88 < 187$ )
  - ❖ Encryption:  $C = 88^7 \bmod 187 = 11$ 
    - Exploiting the properties of modular arithmetic, following can be done:  
 $88^7 \bmod 187 = [(88^4 \bmod 187) \times (88^2 \bmod 187) \times (88^1 \bmod 187)] \bmod 187$
  - ❖ Decryption:  $M = 11^{23} \bmod 187 = 88$

### RSA Example with text

$$p = 73, q = 151$$

$$n = 11023$$

$$\phi(n) = 10800$$

$$e = 11$$

$$d = 5891$$

Text = How are you?

H=	33	a=	00	y=	24
o=	14	r=	17	o=	14
w=	22	e=	04	u=	20
_	62	_	62	?=	66

$$M_1 = 3314 \quad M_2 = 2262 \quad M_3 = 0017$$

$$M_4 = 0462 \quad M_5 = 2414 \quad M_6 = 2066$$

$$C_1 = 3314^{11} \bmod 11023 = 10260$$

$$C_2 = 2262^{11} \bmod 11023 = 9489$$

$$C_3 = 17^{11} \bmod 11023 = 1782$$

$$C_4 = 462^{11} \bmod 11023 = 727$$

$$C_5 = 2414^{11} \bmod 11023 = 10032$$

$$C_6 = 2006^{11} \bmod 11023 = 2253$$

$$M_1 = 10260^{5891} \bmod 11023 = 3314$$

$$M_2 = 9489^{5891} \bmod 11023 = 2262$$

$$M_3 = 1782^{5891} \bmod 11023 = 0017$$

$$M_4 = 727^{5891} \bmod 11023 = 0462$$

$$M_5 = 10032^{5891} \bmod 11023 = 2414$$

$$M_6 = 2253^{5891} \bmod 11023 = 2006$$