

关系数据库存储树形结构数据的理想实践

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开篇思考

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1 开篇思考

树形结构数据

企业组织

行政区域

商品类目

等级体系

文件分类

资料归档

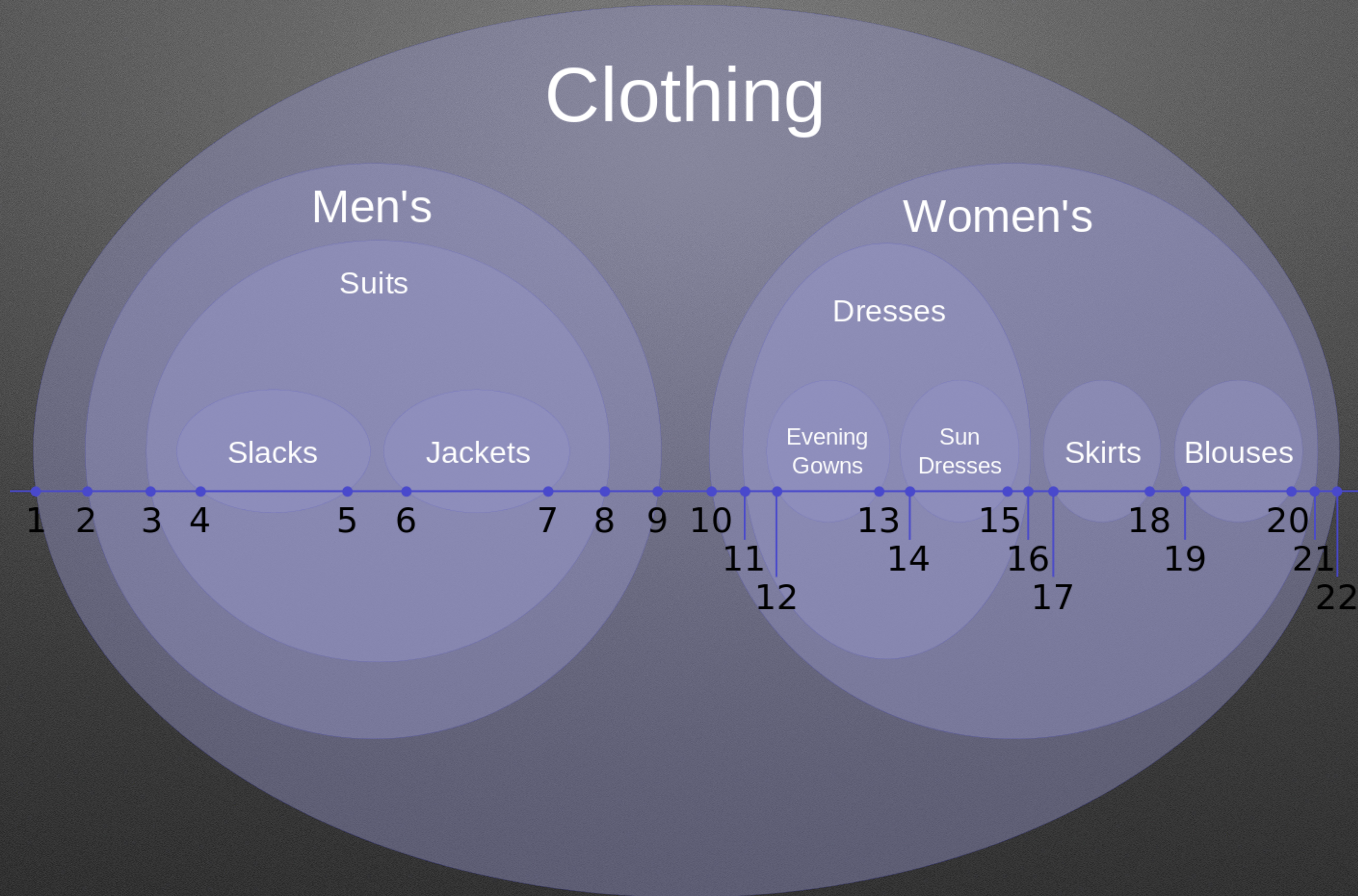
授权体系

.....

A hierarchy

types of clothing

<https://en.wikipedia.org/wiki/File:NestedSetModel.svg>



2 主流方案

主流的存储方法

— 基于关系数据库 (MySQL)

邻接列表模型 (The Adjacency List Model)

邻接列表模型

上述数据模型在关系数据库
MySQL的表结构数据存储

通常如右图所示：

id	node	pid
1	Clothing	0
2	Men's	1
3	Women's	1
4	Suits	2
5	Dresses	3
6	Skirts	3
7	Blouses	3
8	Slacks	4
9	Jackets	4
10	Evening Gowns	5
11	Sun Dresses	5

缺陷



查找子树？查找祖谱？查找深度？完整数据？

3 改进方案

改进的前序遍历树模型

— 基于关系数据库 (MySQL)

The Nested Set Model

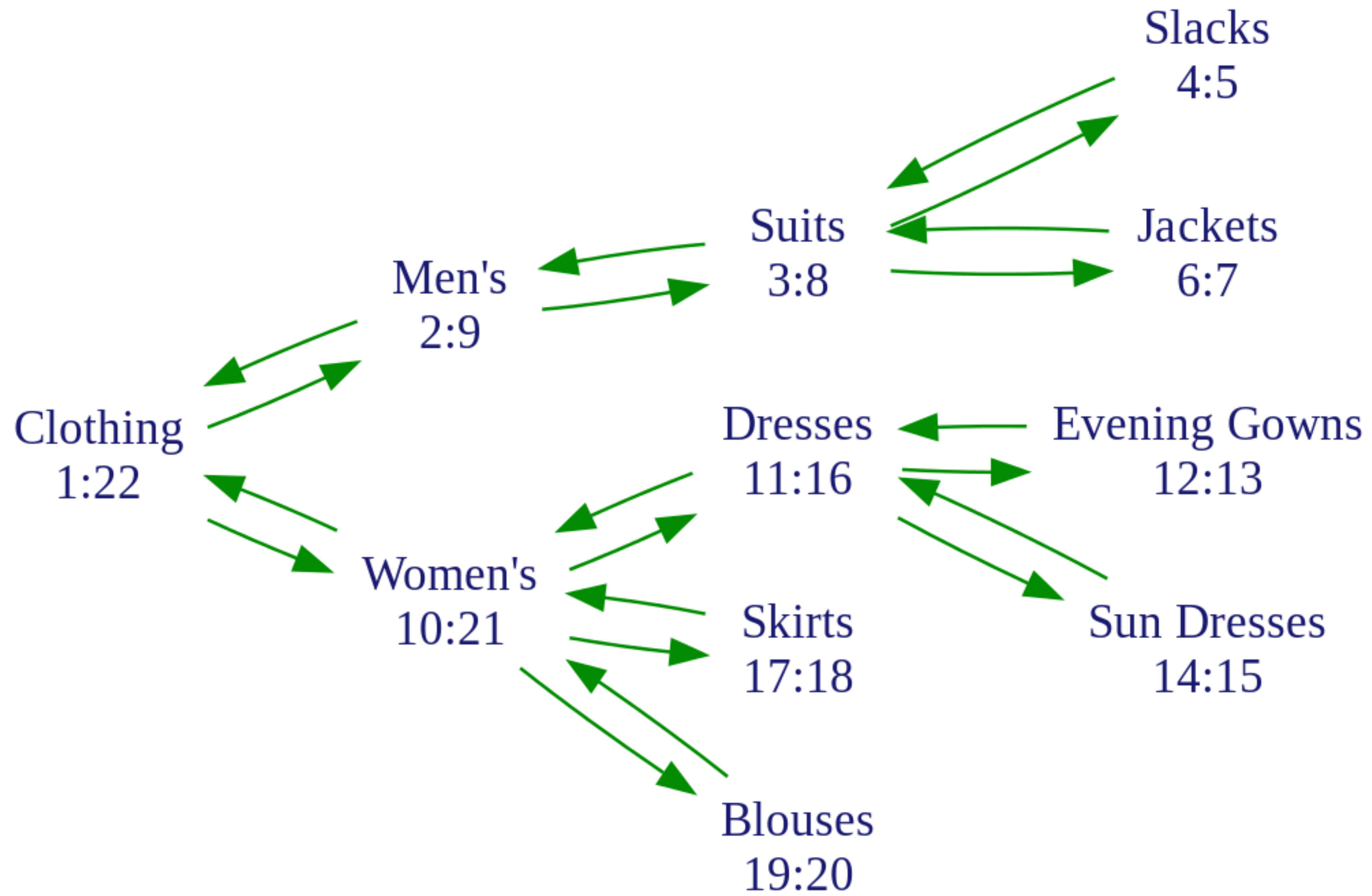
https://en.wikipedia.org/wiki/Nested_set_model

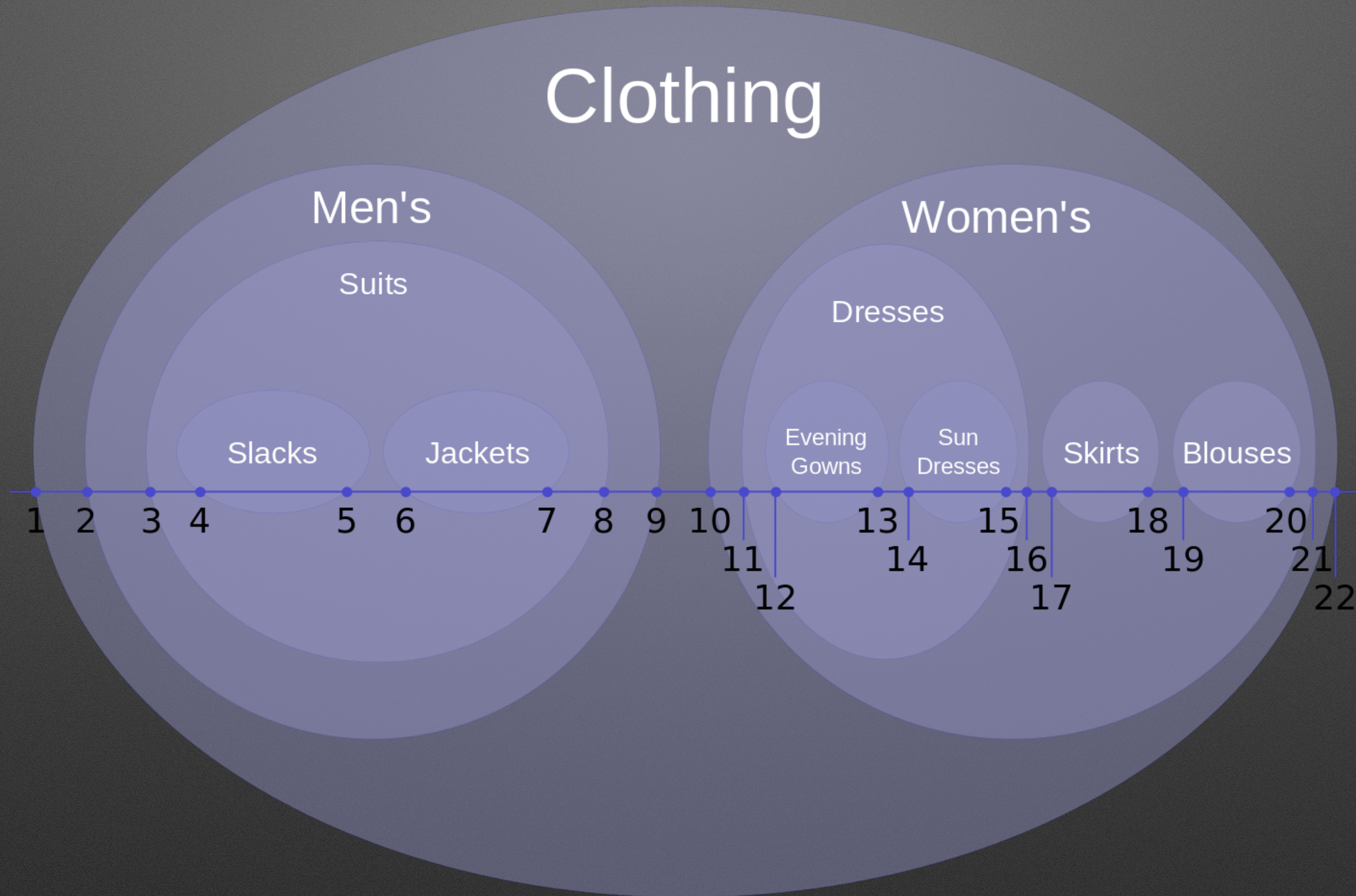
A hierarchy

types of clothing

<https://en.wikipedia.org/wiki/File:Clothing-hierarchy-traversal-2.svg>

The numbering assigned by tree traversal





id	node	pid	depth	lft	rgt
1	Clothing	0	0	1	22
2	Men's	1	1	2	9
3	Women's	1	1	10	21
4	Suits	2	2	3	8
5	Dresses	3	2	11	16
6	Skirts	3	2	17	18
7	Blouses	3	2	19	20
8	Slacks	4	3	4	5
9	Jackets	4	3	6	7
10	Evening Gowns	5	3	12	13
11	Sun Dresses	5	3	14	15

数据库表设计SQL

```
CREATE DATABASE IF NOT EXISTS `treeDB`;
CREATE TABLE IF NOT EXISTS `treeDB`.`nested` (
  `id` INT NOT NULL COMMENT '节点ID',
  `node` VARCHAR(64) CHARACTER SET 'utf8' NOT NULL COMMENT '节点名称',
  `pid` INT NOT NULL COMMENT '父节点ID',
  `depth` INT NOT NULL COMMENT '深度Level',
  `lft` INT NOT NULL COMMENT '左值',
  `rgt` INT NOT NULL COMMENT '右值',
  PRIMARY KEY (`id`),
  INDEX `depth_index` (`depth` ASC),
  INDEX `lft_index` (`lft` ASC),
  INDEX `rgt_index` (`rgt` ASC))
ENGINE = InnoDB DEFAULT CHARACTER SET = utf8 COMMENT = '前序遍历树模型表';
```


常用操作SQL

#插入根节点 (Adding New Nodes)

```
SET @vid:=1,@vnode:='Clothing',@vpid:=0,@vdepth:=0,@vlft:=1,@vrgt:=2;  
INSERT INTO `treeDB`.`nested`(`id`,`node`,`pid`,`depth`,`lft`,`rgt`)  
VALUES(@vid,@vnode,@vpid,@vdepth,@vlft,@vrgt);
```


常用操作SQL

#插入节点 (Adding New Nodes) — First Child

```
SET @vid:=2,@vnode:='Men\'s',@vpid:=1;
SELECT @vprgt:=rgt,@vpdepth:=depth
FROM `treeDB`.`nested` WHERE id=@vpid;
UPDATE `treeDB`.`nested`
SET rgt=CASE WHEN rgt>=@vprgt THEN rgt+2 END,
    lft=CASE WHEN lft>=@vprgt THEN lft+2 END;
INSERT INTO `treeDB`.`nested`(`id`,`node`,`pid`,`depth`,`lft`,`rgt`)
VALUES(@vid,@vnode,@vpid,@vpdepth+1,@vprgt,@vprgt+1);
```


常用操作SQL

#插入节点 (Adding New Nodes) — Not First Child

```
SET @vid:=3,@vnode:='Women\'s',@vlid:=1;
SELECT @vlpid:=pid,@vlrgt:=rgt,@vldepth:=depth
FROM `treeDB`.`nested` WHERE id=@vlid;
UPDATE `treeDB`.`nested`
SET rgt=CASE WHEN rgt>@vlrgt THEN rgt+2 END,
    lft=CASE WHEN lft>@vlrgt THEN lft+2 END;
INSERT INTO `treeDB`.`nested`(`id`,`node`,`pid`,`depth`,`lft`,`rgt`)
VALUES(@vid,@vnode,@vlpid,@vldepth,@vlrgt+1,@vlrgt+2);
```


常用操作SQL

#返回某结点的子树 (Retrieving a Full/Subordinates Tree)
— Given Parent Node Index

```
SET @vid:=2;  
SELECT Child.id,Child.node,Child.pid,Child.depth,Child.lft,Child.rgt  
FROM `treeDB`.`nested` AS Parent,`treeDB`.`nested` AS Child  
WHERE Child.lft BETWEEN Parent.lft AND Parent.rgt  
AND Parent.id=@vid  
ORDER BY Child.id;
```


常用操作SQL

#返回某结点的祖谱路径 (Retrieving a Single Path)

— Given Child Node Index

```
SET @vid:=11;  
SELECT Parent.id,Parent.node,Parent.pid,Parent.depth,Parent.lft,Parent.rgt  
FROM `treeDB`.`nested` AS Parent,`treeDB`.`nested` AS Child  
WHERE Child.lft BETWEEN Parent.lft AND Parent.rgt  
AND Child.id=@vid  
ORDER BY Child.id;
```


常用操作SQL

#返回所有的叶子节点 (Retrieving all the Leaf)

```
SELECT id,node,pid,depth,lft,rgt  
FROM `treeDB`.`nested` WHERE rgt=lft+1;
```


常用操作SQL

#删除节点 (Deleting Nodes)

```
SET @vid:=10;
SELECT @vleft:=left,@vright:=right,@vwidth:=right-left+1
FROM `treeDB`.`nested` WHERE id=@vid;
DELETE FROM `treeDB`.`nested` WHERE left BETWEEN @vleft AND @vright;
UPDATE `treeDB`.`nested`
SET right=CASE WHEN right>@vright THEN right-@vwidth END,
    left=CASE WHEN left>@vright THEN left-@vwidth END;
```


Thanks

To Be Continued...