

Fast Fourier Transform (FFT)

Algorithms and Applications

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What is it?

- Efficient implementation of the Discrete Fourier Transform (DFT) or its inverse
- Many types: Split radix, prime factor, Winograd Fourier transform, vector radix
- Focus will be on the Radix-2 Decimation in time FFT Algorithm

What's the Big Deal?

Included in the Top 10 Algorithms of the 20th Century by the IEEE journal: Computing in Science & Engineering

- EKG and EEG signal processing
- Forensic Science
- Image quality measures
- Interpolation and decimation
- Magnetic Resonance Imaging (MRI)
- Noise filtering
- Optical signal processing
- Pattern Recognition
- Video/image compression
- Numerical solution of differential equations

The Discrete Fourier Transform

Discrete Fourier Transform

$$X^F(k) = \sum_{n=0}^{N-1} x(n) e^{(\frac{-j2\pi}{N})kn}, \quad k = 0, 1, \dots, N-1$$

(Forward) Discrete Fourier Transform (DFT)

$$X(n) = \frac{1}{N} \sum_{k=0}^{N-1} X^F(k) e^{(\frac{j2\pi}{N})kn}, \quad n = 0, 1, \dots, N-1$$

Inverse Discrete Fourier Transform (IDFT)

Discrete Fourier Transform

$$X^F(k) = \sum_{n=0}^{N-1} x(n) W_N^{kn}, \quad k = 0, 1, \dots, N-1$$
$$W_N^{kn} = \exp\left[\left(\frac{-j2\pi}{N}\right)kn\right]$$

$O(N^2)$	<i>Computations for direct DFT</i>	N complex multiplications	} For each k
		$N-1$ complex additions	
$O(N \log_2 N)$	<i>Computations using FFT</i>		

Time Comparison

N	1000	10^6	10^9
N^2	10^6	10^{12}	10^{18}
$N \log_2 N$	10^4	20×10^6	30×10^9

Assuming 1 nanosecond per operation...

10^{18} nanoseconds \approx 31.2 years

30×10^9 nanoseconds \approx 30 seconds

The Fast and the Fourier



Jean Joseph Baptiste Fourier



Vinnie "Fast" Fourier

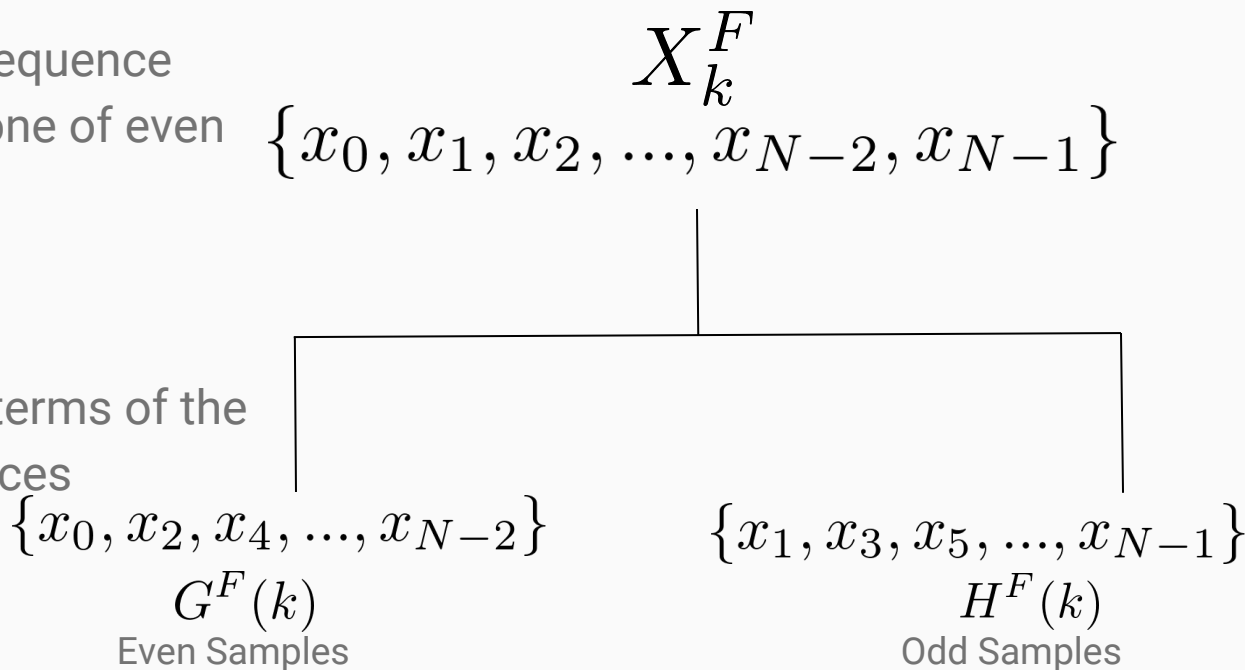
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Decimation in Time (DIT)

Decomposes an N-point sequence into two N/2 sequences (one of even and one of odd samples)

Assume $N = 2^l$, $l = \text{integer}$

Obtain the N-point DFT in terms of the DFTs of these two sequences



$$X^F(k) = \sum_{n=0}^{N-1} x(n) W_N^{kn}, \quad k = 0, 1, \dots, N-1$$

$$= \sum_{\substack{n \text{ even} \\ \text{integer}}} x(n) W_N^{kn} + \sum_{\substack{n \text{ odd} \\ \text{integer}}} x(n) W_N^{kn}$$

Even Indices

Odd Indices

}

$n = 2r$

$n = 2r + 1$

$r = 0, 1, \dots, N/2 - 1$

$$= \sum_{r=0}^{(N/2)-1} x(2r) W_N^{2rk} + \sum_{r=0}^{(N/2)-1} x(2r+1) W_N^{(2r+1)k}$$

Math Magic

$$= \sum_{r=0}^{(N/2)-1} x(2r)(W_N^2)^{rk} + W_N^k \sum_{r=0}^{(N/2)-1} x(2r+1)(W_N^2)^{rk}$$

Note that

$$W_N^2 = \left[\frac{-j2(2\pi)}{N} \right] = \left[\frac{-j2\pi}{N/2} \right] = W_{N/2}$$

$$X^F(k) = \sum_{r=0}^{(N/2)-1} x(2r)W_{N/2}^{rk} + W_N^k \sum_{r=0}^{(N/2)-1} x(2r+1)W_{N/2}^{rk}$$

$$= G^F(k) + W_N^k H^F(k), \quad k = 0, 1, \dots, \frac{N}{2} - 1$$

Data-Flow Diagrams



Image from:
https://zibbet.s3.amazonaws.com/uploads/photo/file/7833065/il_fullxfull.130323217.jpg

Butterfly Computations

N = 8

$$X^F(k) = G^F(k) + W_N^k H^F(k), \quad k = 0, 1, \dots, \frac{N}{2} - 1$$

$X^F[0]$

$X^F[2]$

$X^F[4]$

$X^F[6]$

$X^F[1]$

$X^F[3]$

$X^F[5]$

$X^F[7]$

Butterfly Computations

N = 8

$X^F[0] \longrightarrow$

$X^F[2] \longrightarrow$

$X^F[4] \longrightarrow$

$X^F[6] \longrightarrow$

$X^F[1] \longrightarrow$

$X^F[3] \longrightarrow$

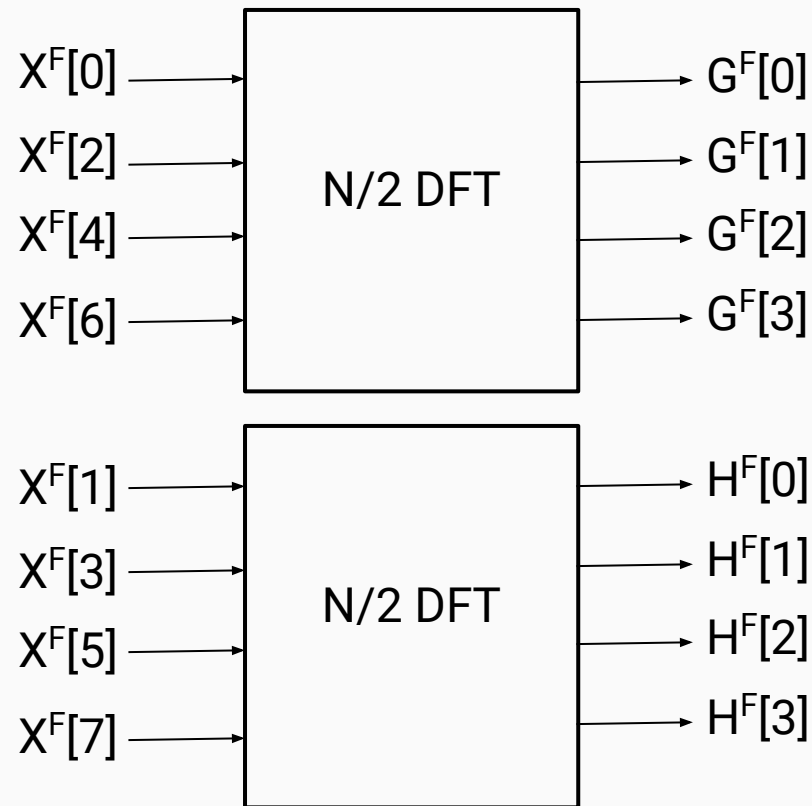
$X^F[5] \longrightarrow$

$X^F[7] \longrightarrow$

$$X^F(k) = G^F(k) + W_N^k H^F(k), \quad k = 0, 1, \dots, \frac{N}{2} - 1$$

Butterfly Computations

$N = 8$

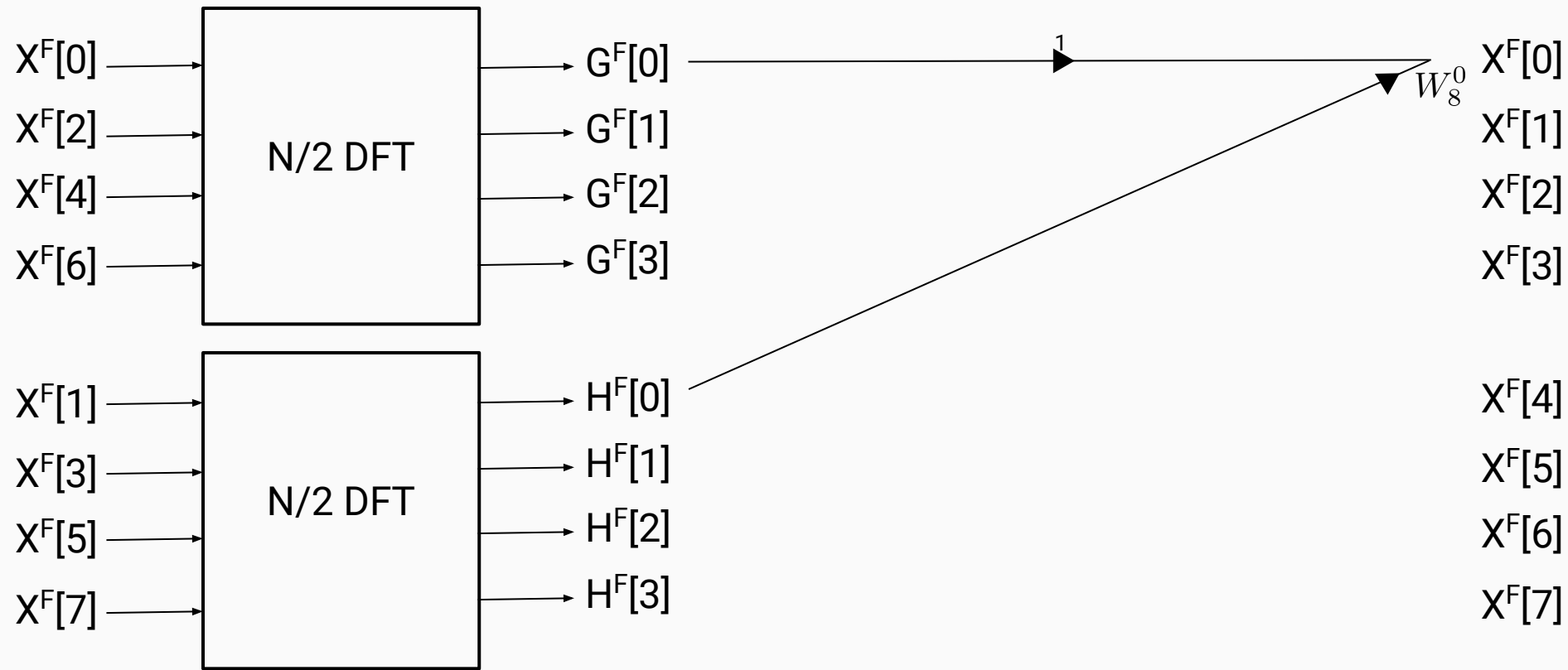


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Butterfly Computations

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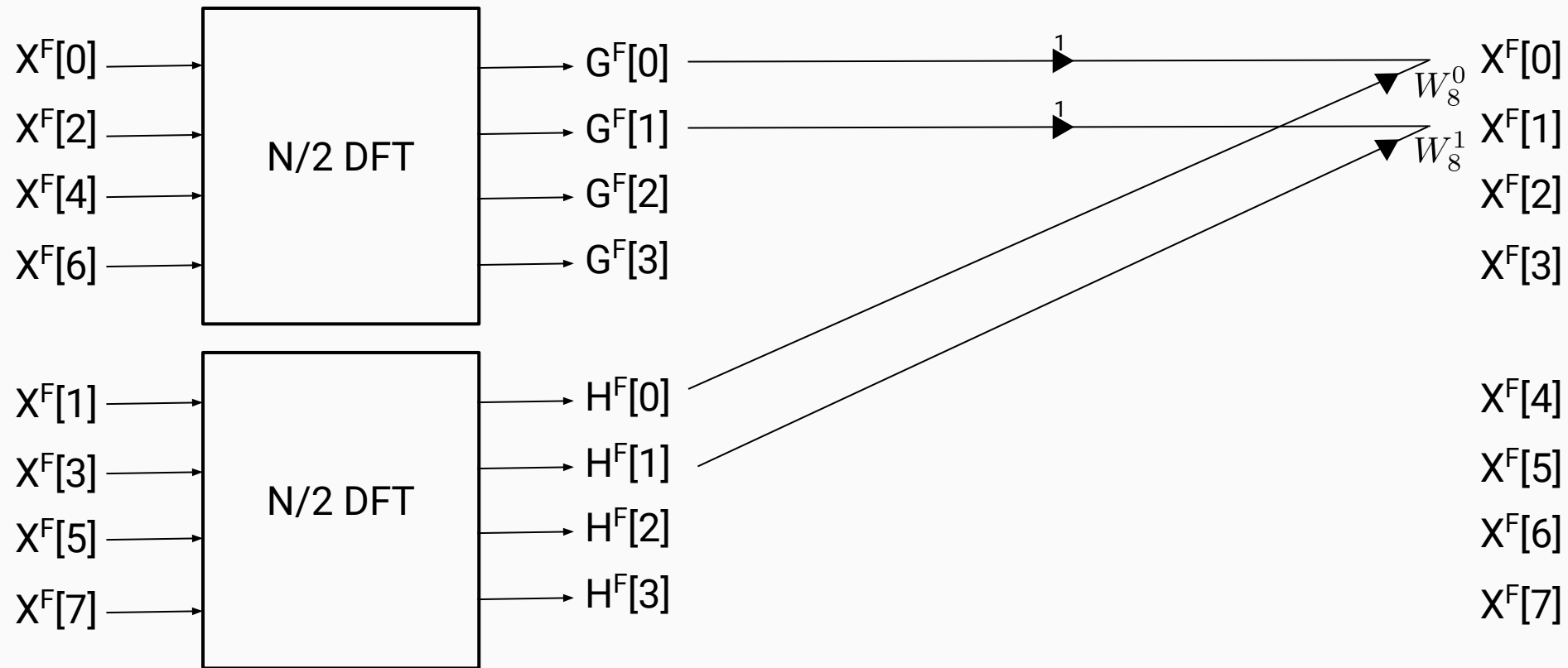
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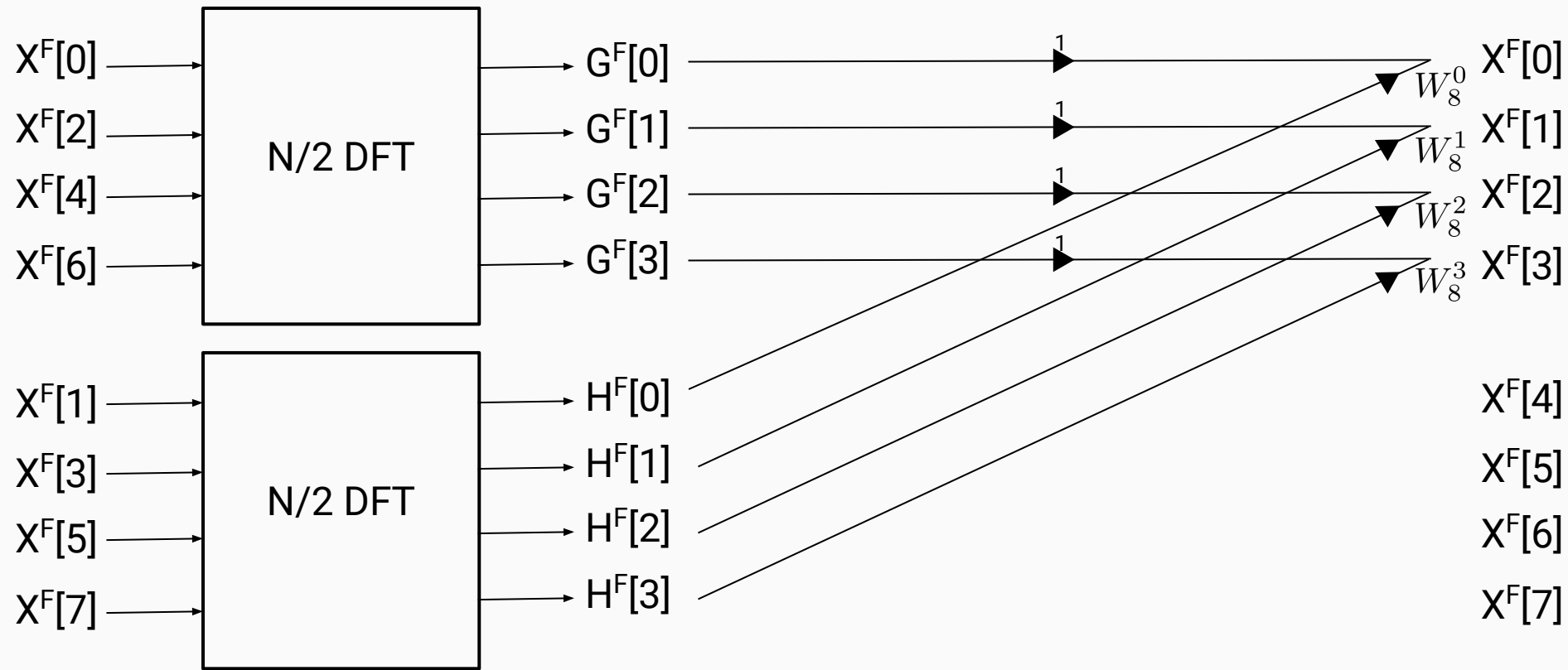
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Butterfly Computations

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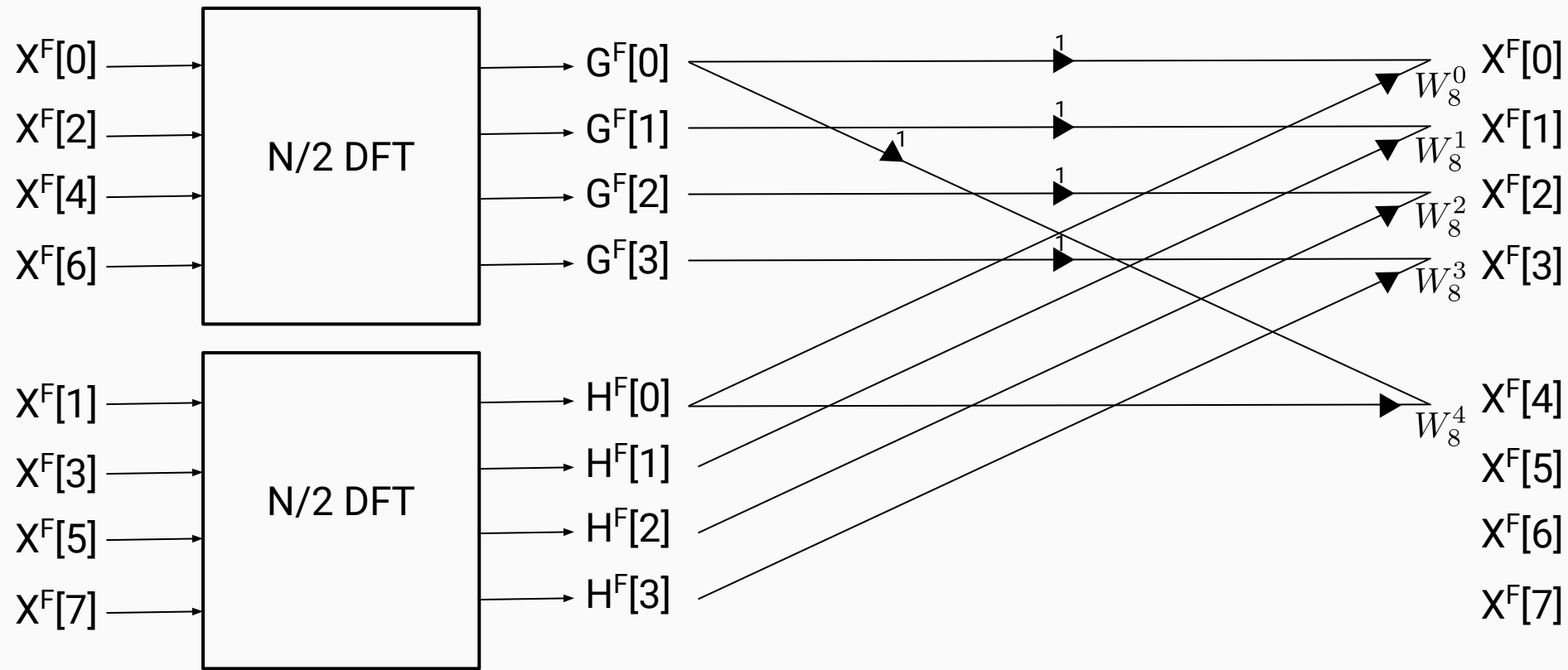
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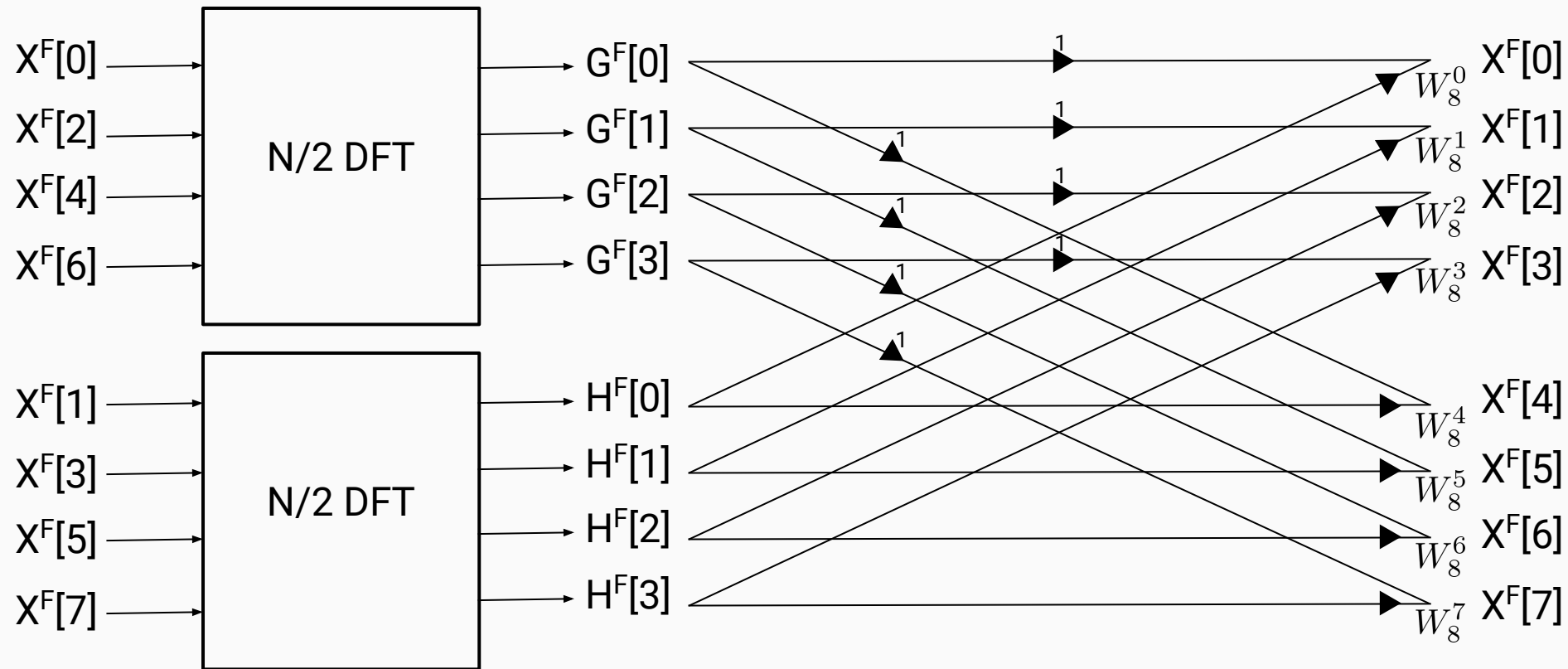
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Butterfly Computations

$N = 8$

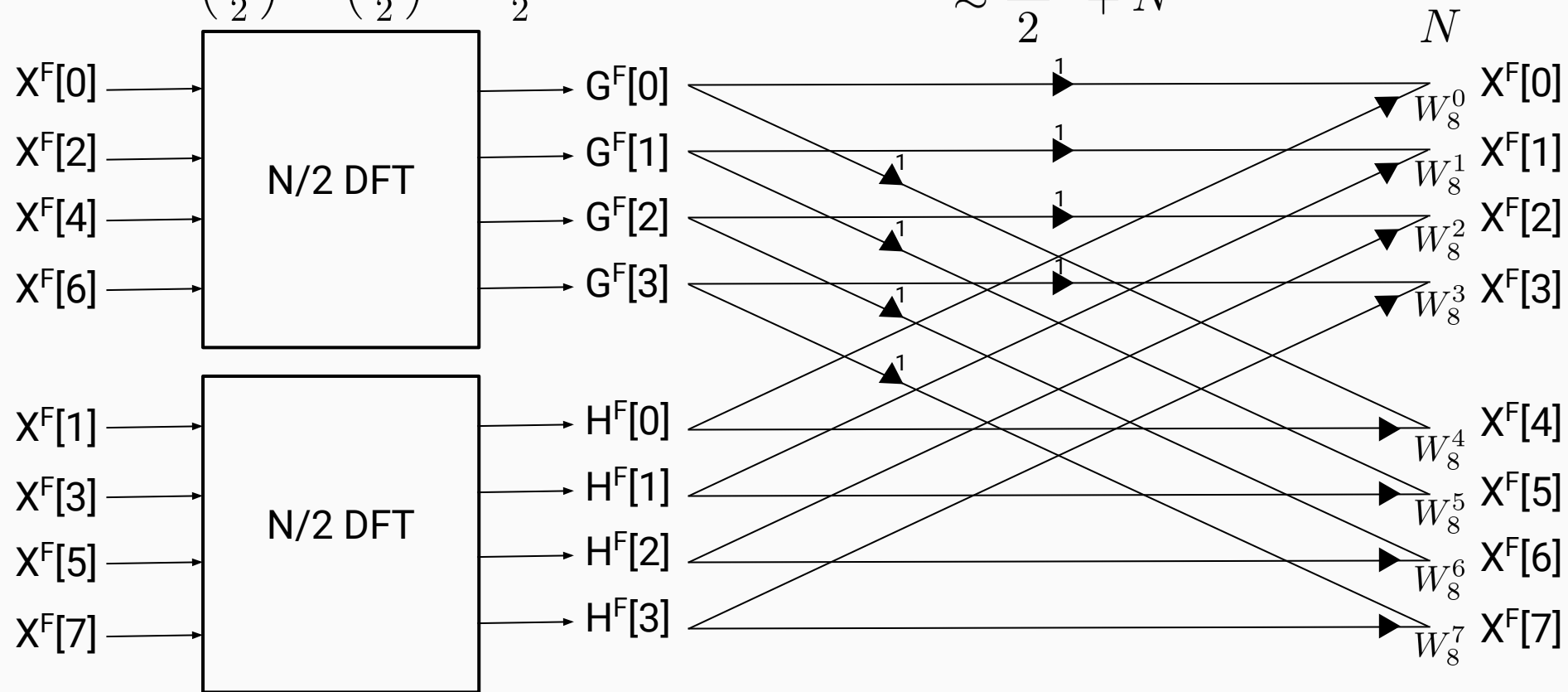
$$X^F(k) = G^F(k) + W_N^k H^F(k), \quad k = 0, 1, \dots, \frac{N}{2} - 1$$



Butterfly Computations

$$N = 8 \quad \left(\frac{N}{2}\right)^2 + \left(\frac{N}{2}\right)^2 = \frac{N^2}{2}$$

$$\approx \frac{N^2}{2} + N$$



Why Stop Here?

Let's keep splitting: *each* $\frac{N}{2}pt \rightarrow 2 \frac{N}{4}pt$ *DFTs*

How many times? $\frac{N}{2}, \frac{N}{4}, \dots, \frac{N}{2^{v-1}}, \frac{N}{2^v} = 1, \quad v = \log_2 N$

Cost?

$$1 : \frac{N}{2} \rightarrow 2\left(\frac{N}{2}\right)^2 + N = \frac{N^2}{2} + N$$

$$2 : \frac{N}{4} \rightarrow 2\left(2\left(\frac{N}{4}\right)^2 + \frac{N}{2}\right) + N = \frac{N^2}{4} + 2N$$

$$3 : \frac{N}{8} \rightarrow 2\left[2\left(2\left(\frac{N}{8}\right)^2 + \frac{N}{4}\right) + \frac{N}{2}\right] + N = \frac{N^2}{8} + 3N$$

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Cost?

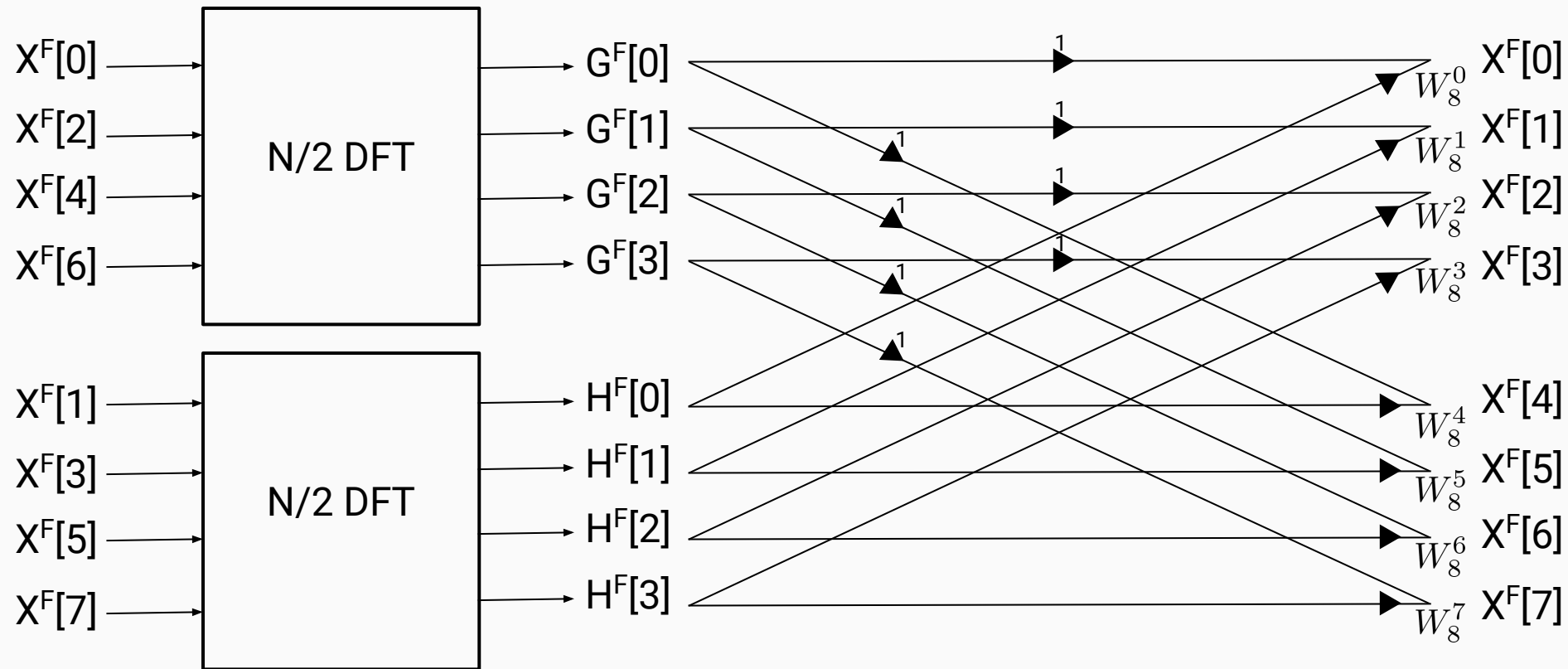
$$v : \frac{N}{2^v} = 1 \rightarrow \frac{N^2}{2^v} + vN = \frac{N^2}{N} + N \log_2 N$$

$\approx O(N \log_2 N)$ for N large

Butterfly Computations

$N = 8$

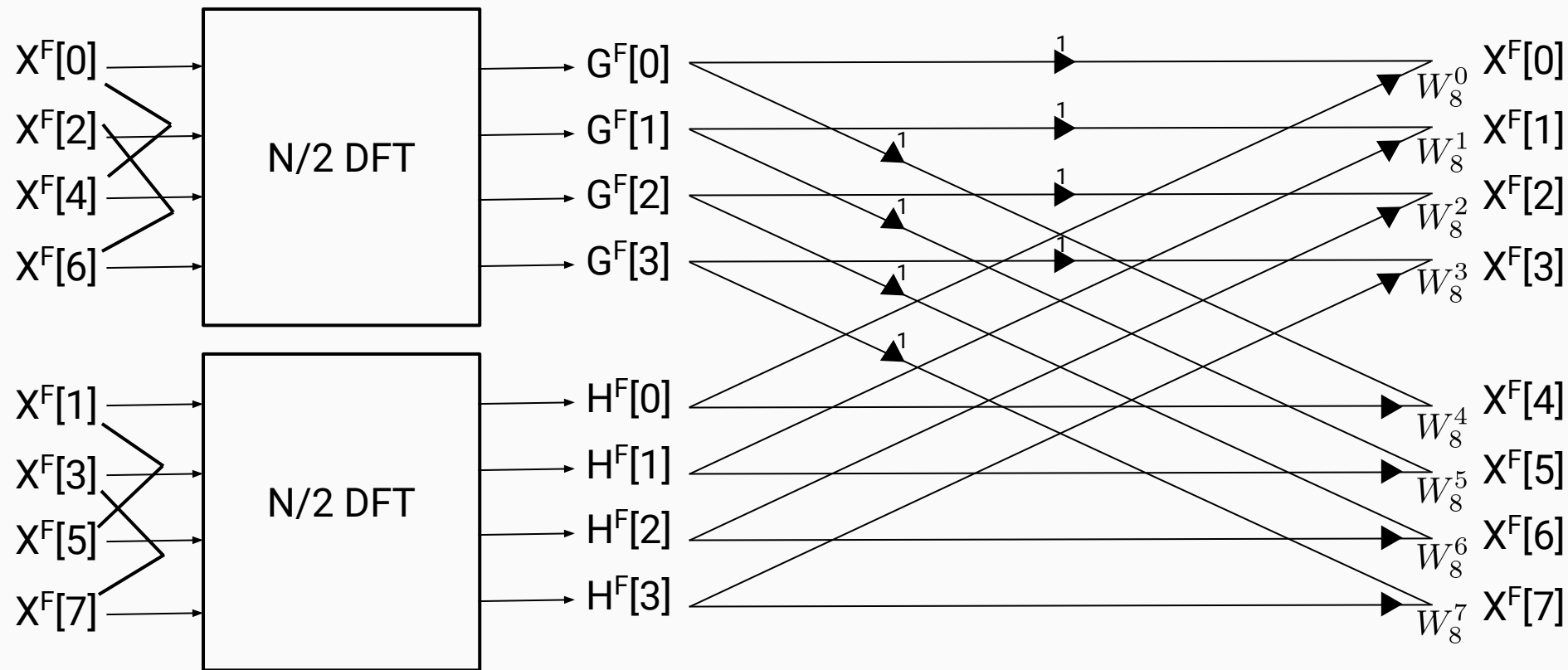
$$X^F(k) = G^F(k) + W_N^k H^F(k), \quad k = 0, 1, \dots, \frac{N}{2} - 1$$



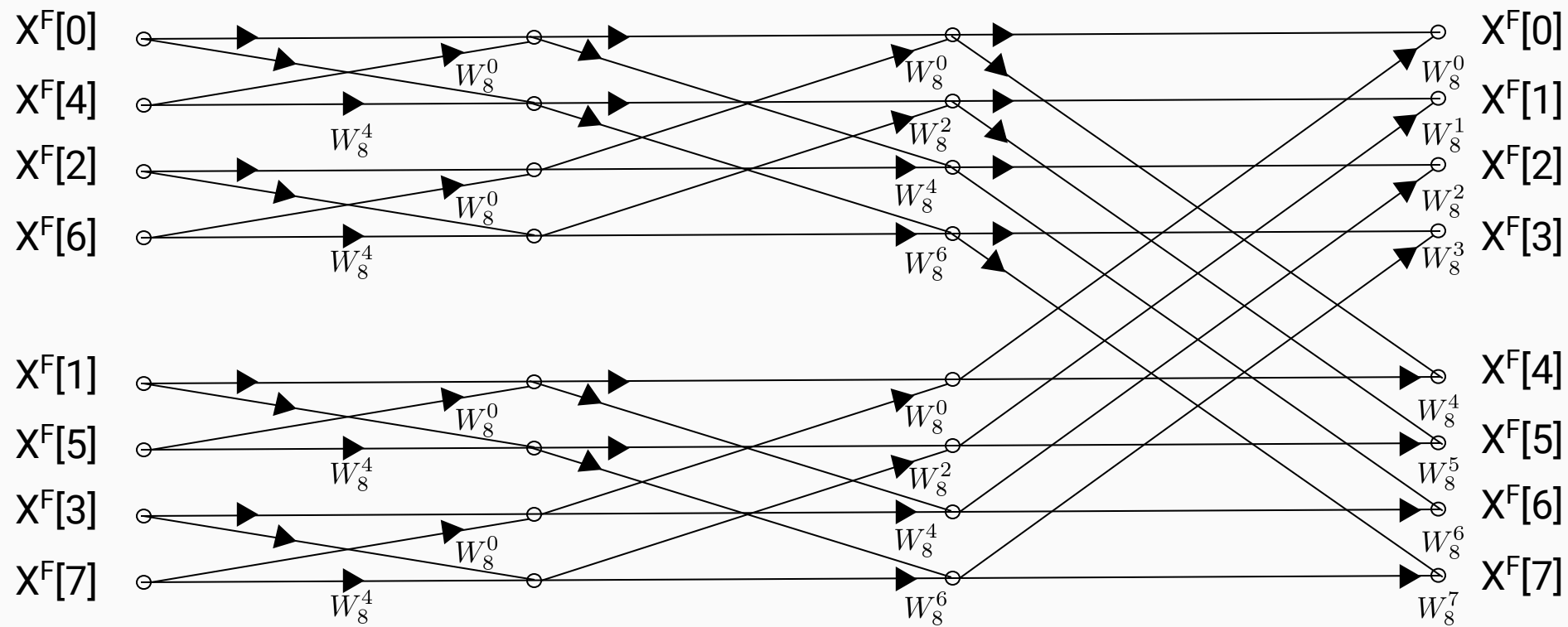
Butterfly Computations

$N = 8$

$$X^F(k) = G^F(k) + W_N^k H^F(k), \quad k = 0, 1, \dots, \frac{N}{2} - 1$$



Butterfly Computations



Thank You

References

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