CMPT 225 D100 Assignment 3 Reports

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The Priority Queue has only one class PQ. It has the following data members:

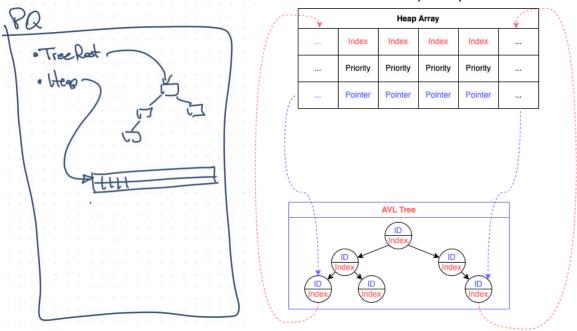
- 1. AVL Tree
 - a. AVL nodes used to construct embedded AVL tree. A definition of an AVL node is:
 - i. ID-typed Task ID is stored in *element*
 - ii. Each AVL node also contains an integer *index*, corresponding to the position of that task in the min-heap array
 - iii. As usual AVL nodes would contain, there are *left, *right, and height contained in each node for tree construction
 - b. A root pointer for AVL tree

- 2. Binary Min-Heap (Array-Implemented)
 - a. A size counter *currentSize* tracking the number of pairs in the array
 - b. An array (vector) of pairs where each pair contains:
 - i. a priority integer for a task t
 - ii. a pointer pointing to an AVL tree node that corresponds to the task linked with that priority
 - c. A temporary ID-typed container *recentlyDeleted* storing a copy of recently deleted task ID. This is used as a convenient container to store returned values in function *deleteMin()*. May be of some other creative usages...

<pre>// Heap // Pair vector stores priority and for t int currentSize; vector<pair<int, *="" avlnode="">> array; ID recentlyDeleted; (for extract_min)</pair<int,></pre>	for task t and a pointer to tree node // Number of elements in heap // The heap array // Keep a copy of recently deleted
3 array of pairs. The 23	

pair = priority

A visualization of the PQ class with embedded AVL tree and Heap Array:



The PQ class has the following accessing methods:

- 1. All AVL tree nodes must be accessible by using the pointers stored in Heap Array.
- 2. Heap elements can be accessed using the indexes store in AVL tree nodes.
- 3. A task ID in an AVL node must have an index associated to a heap element that contains (1) the task's designated priority number, and (2) a pointer pointing to the AVL node itself.

Besides the public operations defined in PQdeclare.h, there are 1 added public function and 19 added private functions.

For debugging purposes, we added one public function called *print_all()*. It prints, if not empty:

- (1) The current size of the heap array
- (2) The task ID with the smallest priority
- (3) All nodes in the AVL tree and their task IDs and heap indexes (performed by printTree())
- (4) The visualization of the tree with the task IDs, the heap indexes, and the addresses of all nodes (performed by displayLinks())
- (5) All pairs in the binary min-heap array with their indexes, priorities, and pointers of AVL nodes

We also added the following 17+2 internal methods for AVL tree constructions and operations and Binary Heap constructions and operations (mostly from part of the textbook author's heap or AVL tree classes):

- (1) Two tree insertions with slight modifications
 - a. Now takes 3 inputs: Task ID, Subtree Root, and Priority
- (2) One tree node removal with slight modifications
 - a. Now swap both the Task ID and the Priority (with its successor)
- (3) One tree balance, four tree rotations, and two supplementary functions *height()* and *max()*
- (4) Three node searches (i.e., find min, find max, find which contains a given Task ID)
 - a. The third function called findNode() is new. It searches and returns the address of a node containing the Task ID intended to search. It returns a null pointer if not found.
- (5) One tree makeEmpty()
- (6) Two tree displaying with displayLinks() from the previous assignment
 - a. Now additionally prints indexes of nodes
- (7) One tree traversal called *traverseTreeIndex()* is newly designed for *deleteMin()* to update pointer addresses in all pairs in the heap array:
 - a. For each node traversed, get its index
 - b. For each index, go to the pair at that index in the heap array
 - c. For each pair reached, set the pointer at the second element to point to the node just traversed
- (8) One heap building
- (9) One heap node percolating down

Using the above accessing methods and functions, we can create, update, and remove a task:

1. Create:

- a. Constructor (with given sets of tasks and priorities)
 - i. Insert all task IDs using AVL nodes while leaving the indexes blank
 - ii. Use findNode() to find node addresses for all task IDs
 - iii. Pair up these addresses (pointers) with their corresponding priorities
 - iv. Insert these pairs into a heap array and build heap
 - v. For each pair in the heap, use the pointer to update the index for each AVL node

b. Insertion

- i. Insert the task ID using an AVL node while leaving the index blank
- ii. Find the address for the inserted task ID after the tree's balanced
- iii. Pair the address (pointer) with the priority and insert to the heap
- iv. After percolating up the heap, use the pointer to update the index for the corresponding AVL node

2. Update Priority

- a. Find the address of the AVL node containing the task to be updated
- b. Use the index stored in the node to access its heap element
- c. Make a copy of the old priority in the heap element
- d. Change the priority to the new one
- e. Percolate up/down if needed
- f. Use the pointer stored in each heap element to update the index for the corresponding AVL node

3. Delete min

- a. Store the old min into the private container recentlyDeleted
- b. Overwrite the top heap element with the last
- c. Reduce current size of the heap array and percolate down
- d. Use the pointer stored in each heap element to update the index for the corresponding AVL node
- e. Remove the AVL node containing the task to be deleted and balance the tree
- f. Use the index in each AVL node to update the pointer address in each heap element (using traverseTreeIndex())
- g. Return the old min stored at the beginning