$$\frac{\overline{A \times B \vdash A \times B} \stackrel{ax}{=} \times_{e_2}}{\underline{A \times B \vdash A \times B} \stackrel{ax}{=} \times_{e_1}} \times_{e_1}$$

$$\frac{A \times B \vdash B \times A}{\cdot \vdash (A \times B) \to (B \times A)} \times_{i}$$

A category, C, consists of a collection of objects  $A, \ldots, Z \in \mathsf{Obj}(C)$  and a collection of morphisms  $f, \ldots, z \in \mathsf{Mor}(C)$  such that the following hold:

- i. There are two functions  $\operatorname{src}:\operatorname{\mathsf{Mor}}\to\operatorname{\mathsf{Obj}}$  and  $\operatorname{\mathsf{tar}}:\operatorname{\mathsf{Mor}}\to\operatorname{\mathsf{Obj}}$ , such that, we write  $f:A\longrightarrow B$  if  $\operatorname{\mathsf{src}}(f)=A$  and  $\operatorname{\mathsf{tar}}(f)=B$
- ii. For any  $A \in \mathsf{Obj}(\mathcal{C})$ , there is a morphism  $\mathsf{id}_A : A \longrightarrow A$ .
- iii. Given any two morphisms  $f: A \longrightarrow B$  and  $g: B \longrightarrow C$  there must be a morphism  $f; g: A \longrightarrow C$ . In addition,

a. 
$$f$$
;  $id_B = id_A$ ;  $f = f : A \longrightarrow B$ 

b. 
$$f;(g;h)=(f;g);h$$
 for  $f:A\longrightarrow B,$   $g:B\longrightarrow C,$  and  $h:C\longrightarrow D$ 

A cartesian closed category is a category C with the following additional data:

- i. There is a special object  $1 \in \mathsf{Obj}(\mathcal{C})$  such that for any  $A \in \mathsf{Obj}(\mathcal{C})$ , there is a unique morphism  $i : A \longrightarrow 1$ .
- ii. For any  $A, B \in \mathsf{Obj}(\mathcal{C})$ , there is an object  $A \times B \in \mathsf{Obj}(\mathcal{C})$ . In addition,

a. 
$$\pi_1: A \times B \longrightarrow A$$

b. 
$$\pi_2: A \times B \longrightarrow B$$

c. 
$$\langle f, g \rangle : A \longrightarrow B \times C$$
 for any  $f : A \longrightarrow B$  and  $g : A \longrightarrow C$ 

iii. For any  $A, B \in \mathsf{Obj}(\mathcal{C})$ , there is an object  $A \to B \in \mathsf{Obj}(\mathcal{C})$ . In addition,

a. 
$$f^*: G \longrightarrow A \to B$$
 for any  $f: G \times A \longrightarrow B$ 

b. 
$$\mathsf{app}_{A,B}: (A \to B) \times A \longrightarrow B$$

$$A ::= 1 \mid A \times B \mid A \rightarrow B$$

$$\lambda x.x:A\to A$$

$$\lambda x.\lambda y.y:A\to B\to B$$

$$\lambda x.(x,x):A\to (A\times A)$$

$$\lambda f.\lambda g.\lambda x.g(fx):(A\to B)\to (B\to C)\to (A\to C)$$

$$\frac{1}{(\lambda x.t_2) t_1 \rightsquigarrow [t_1/x]t_2} \beta \qquad \frac{1}{\mathsf{fst}(t_1, t_2) \rightsquigarrow t_1} fst$$

$$\frac{}{\operatorname{snd}\left(t_{1},t_{2}\right)\leadsto t_{2}}\ snd$$

A formula A is valid iff there exists a term t and a context  $\Gamma$ , such that,  $\Gamma \vdash t : A$ .

$$\overline{\Gamma, A \vdash A} \stackrel{ax}{ax} \qquad \qquad \Gamma \times A \xrightarrow{\pi_2} A$$

$$\overline{\Gamma \vdash A} \qquad \Gamma \vdash B \qquad \qquad \Gamma \xrightarrow{f} A \qquad \Gamma \xrightarrow{g} B \qquad \qquad \Gamma \xrightarrow{\langle f, g \rangle} A \times B$$

$$\underline{\Gamma \vdash A \times B}_{\Gamma \vdash A} \times e_1 \qquad \qquad \underline{\Gamma} \xrightarrow{f} A \times B \qquad \qquad \Gamma \xrightarrow{\pi_1} A \qquad \qquad \Gamma \xrightarrow{f} A \times B \xrightarrow{\pi_1} A$$

$$\underline{\Gamma \vdash A \times B}_{\Gamma \vdash B} \times e_1 \qquad \qquad \underline{\Gamma} \xrightarrow{f} A \times B \xrightarrow{\pi_2} B$$

$$\underline{\Gamma, A \vdash B}_{\Gamma \vdash A \to B} \xrightarrow{r} A \xrightarrow{g} A \xrightarrow{gpp_{A,B}} B$$

$$\underline{\Gamma \vdash A \to B}_{\Gamma \vdash B} \xrightarrow{\Gamma \vdash A} \rightarrow e \qquad \qquad \underline{\Gamma} \xrightarrow{\langle f, g \rangle} (A \to B) \times A \xrightarrow{gpp_{A,B}} B$$