



1º teste SLT
2021/22

Universidade do Minho

(a) _____ Unidade Curricular _____

Curso _____ Docente _____

ALUNO (b) _____ em ____/____/____

QUESTÃO 1

$P \equiv \{x:=5; z:=x \times 2\}; \text{ if } x \leq y \text{ then } \{x:=x+y; z:=y+x\} \text{ else } x=0$
 $sy=10$

①

$\langle P, s \rangle \xrightarrow{①} \langle z:=x \times 2; \text{ if } \dots, s[x \mapsto 5] \rangle$

$\xrightarrow{②} \langle \text{if } \dots, s[x \mapsto 5, z \mapsto 10] \rangle$

$\xrightarrow{③} \langle x:=x+y; z:=y+x, s[x \mapsto 5, z \mapsto 10] \rangle$

$\Rightarrow \langle z:=y+x, s[x \mapsto 15, z \mapsto 10] \rangle$

$\Rightarrow s[x \mapsto 15, z \mapsto 25]$

(ass_{3,x})

(comp₂)

$\langle x:=5, s \rangle \Rightarrow s[x \mapsto 5]$

(comp₁)

$\langle x:=5; z:=x \times 2, s \rangle \Rightarrow \langle z:=x \times 2, s[x \mapsto 5] \rangle$

... $\xrightarrow{①}$...

(ass₂)

(comp₂)

$\langle z:=x \times 2, s[x \mapsto 5] \rangle \Rightarrow s[x \mapsto 5, z \mapsto 10]$

... $\xrightarrow{②}$...

(if_{in}^{tt})

... $\xrightarrow{③}$...

pp. $\{x \leq y\} s[x \mapsto 5, z \mapsto 10] = 5 \leq 10 = \text{tt}$

② $\langle P, s_0 \rangle \rightarrow s_0[x \mapsto 0, z \mapsto 10]$

(comp_{pre})

1. $\langle x:=5; z:=x \times 2, s_0 \rangle \rightarrow s_0[x \mapsto 5, z \mapsto 10]$

(comp_{pre})

1. $\langle x:=5, s_0 \rangle \rightarrow s_0[x \mapsto 5]$

(ass₂)

2. $\langle z:=x \times 2, s_0[x \mapsto 5] \rangle \rightarrow s_0[x \mapsto 5, z \mapsto 10]$

(ass₂)

2. $\langle \text{if } x \leq y \text{ then } \dots \text{ else } x=0, s_0[x \mapsto 5, z \mapsto 10] \rangle \rightarrow s_0[x \mapsto 0, z \mapsto 10]$

(if_{in}^{tt})

1. $\langle x:=0, s_0[x \mapsto 5, z \mapsto 10] \rangle \rightarrow s_0[x \mapsto 0, z \mapsto 10]$

③ $\{y < 5 \wedge p \wedge z = 0\}$

1. $\{y < 5 \wedge x = 5; z := x + 2 \wedge y < 5 \wedge x = 5\}$

1. $\{y < 5 \wedge x = 5 \wedge y < 5 \wedge x = 5\}$

1. $\{y < 5 \wedge x = 5 \wedge y < 5 \wedge x = 5\}$

2. $\{y < 5 \wedge x = 5 \wedge z := x + 2 \wedge y < 5 \wedge x = 5\}$

2. $\{y < 5 \wedge x = 5 \wedge \text{if } x \leq y \text{ then } \dots \text{ else } x = 0 \wedge x = 0\}$

1. $\{y < 5 \wedge x = 5 \wedge \neg(x \leq y) \wedge x = 0 \wedge x = 0\}$

pf. $y < 5 \wedge x = 5 \wedge x > y \rightarrow 0 = 0$

1. $\{0 = 0 \wedge x = 0 \wedge x = 0\}$

2. $\{y < 5 \wedge x = 5 \wedge x \leq y \wedge z := x + y; z := y + x \wedge x = 0\}$

1. $\{y < 5 \wedge x = 5 \wedge x \leq y \wedge z := x + y \wedge x = 0\}$

pf. $y < 5 \wedge x = 5 \wedge x \leq 5 \rightarrow x + y = 0$ pois Falso $\rightarrow x + y = 0$

2. $\{x = 0 \wedge z := y + x \wedge x = 0\}$

(comp)

(comp)

pf. $y < 5 \rightarrow y < 5 \wedge 5 = 5$ (cons)

(con)

(con)

(if)

(cons)

pf. $0 = 0$ e True

(con)

(comp)

(cons)

(con)

(con)

④ $S_h[P]s = S_h[\text{if } \dots] (S_h[x := 5; z := x + 2]s)$

$= \text{Cond}([x \leq y], S_h[z := x + y] \circ S_h[x := x + y], S_h[x := 0]) (S_h'(z := x + 2) s [z \rightarrow 5])$

$= \text{Cond}(\dots, \dots, \dots) s [x \rightarrow 5, z \rightarrow 10]$

$= \begin{cases} ([z := x + y] \circ [x := x + y]) s [x \rightarrow 5, z \rightarrow 10] & \text{se } [x \leq y] s [x \rightarrow 5, z \rightarrow 10] = \text{tt} \\ [x := 0] s [x \rightarrow 5, z \rightarrow 10] & \text{se } [x > y] s [x \rightarrow 5, z \rightarrow 10] = \text{ff} \end{cases}$

$= \begin{cases} [z := x + y] s [x \rightarrow 5 + 5y, z \rightarrow 10] & \text{se } 5 \leq 5y \\ s [x \rightarrow 0, z \rightarrow 10] & \text{se } 5 > 5y \end{cases}$

$= \begin{cases} s [x \rightarrow 5 + 5y, z \rightarrow 5 + 2 * 5y] & \text{se } 5 \leq 5y \\ s [x \rightarrow 0, z \rightarrow 10] & \text{se } 5 > 5y \end{cases}$

QUESTÃO

①

$[iff^1_{ms}]$

$[iff^2_{ms}]$

$[iff^3_{ms}]$

②

se

1.

2.

3.

③

QUESTÃO 2

①

$$\text{Diff}_1^1 \frac{\langle c_1, s \rangle \rightarrow s' \quad \text{se } \llbracket b_1 \rrbracket_s = \text{tt}}{\langle \text{if } b_1 \text{ then } c_1 \text{ else if } b_2 \text{ then } c_2 \text{ else } c_3, s \rangle \rightarrow s'}$$

$$\text{Diff}_1^2 \frac{\langle c_2, s \rangle \rightarrow s' \quad \text{se } \llbracket b_1 \rrbracket_s = \text{ff} \text{ e } \llbracket b_2 \rrbracket_s = \text{tt}}{\langle \text{if } b_1 \text{ then } c_1 \text{ else if } b_2 \text{ then } c_2 \text{ else } c_3, s \rangle \rightarrow s'}$$

$$\text{Diff}_1^3 \frac{\langle c_3, s \rangle \rightarrow s' \quad \text{se } \llbracket b_1 \rrbracket_s = \text{ff} \text{ e } \llbracket b_2 \rrbracket_s = \text{ff}}{\langle \text{if } \dots, s \rangle \rightarrow s'}$$

②

$$\frac{\frac{\frac{\frac{\phi \wedge b_1}{\psi} \quad \text{①}}{\phi \wedge \neg b_1 \wedge b_2} \quad \text{②}}{\phi \wedge \neg b_1 \wedge \neg b_2} \quad \text{③}}{\phi \wedge (\text{if } b_1 \text{ then } c_1 \text{ else if } b_2 \text{ then } c_2 \text{ else } c_3) \wedge \psi}$$

P

se $\llbracket \phi \rrbracket_s = \text{tt}$ e $\langle P, s \rangle \rightarrow s'$ então pode acontecer um dos seguintes casos

1. $\llbracket b_1 \rrbracket_s = \text{tt}$ e $\langle c_1, s \rangle \rightarrow s'$ então temos $\llbracket \phi \wedge b_1 \rrbracket_s = \text{tt}$, logo por ①, $\llbracket \psi \rrbracket_s = \text{tt}$
2. $\llbracket b_1 \rrbracket_s = \text{ff}$ e $\llbracket b_2 \rrbracket_s = \text{tt}$ e $\langle c_2, s \rangle \rightarrow s'$ então $\llbracket \phi \wedge \neg b_1 \wedge b_2 \rrbracket_s = \text{tt}$ por ②, $\llbracket \psi \rrbracket_s = \text{tt}$
3. $\llbracket b_1 \rrbracket_s = \text{ff}$ e $\llbracket b_2 \rrbracket_s = \text{ff}$ e $\langle c_3, s \rangle \rightarrow s'$ então $\llbracket \phi \wedge \neg b_1 \wedge \neg b_2 \rrbracket_s = \text{tt}$, por ③, $\llbracket \psi \rrbracket_s = \text{tt}$

Portanto $\models \phi \wedge (\text{if } b_1 \text{ then } c_1 \text{ else if } b_2 \text{ then } c_2 \text{ else } c_3) \wedge \psi$

③

$$CS[\text{if } b_1 \text{ then } c_1 \text{ else if } b_2 \text{ then } c_2 \text{ else } c_3] =$$

$$CB[\llbracket b_1 \rrbracket : \text{BRANCH}(CS[c_1], CB[\llbracket b_2 \rrbracket : \text{BRANCH}(CS[c_2], CS[c_3])])]$$

QUESTÃO 3

- ① $\text{Asgn}(\text{skip}) = \emptyset$
 $\text{Asgn}(x := a) = \{x\}$
 $\text{Asgn}(C_1; C_2) = \text{Asgn}(C_1) \cup \text{Asgn}(C_2)$
 $\text{Asgn}(\text{if } b \text{ then } C_1 \text{ else } C_2) = \text{Asgn}(C_1) \cup \text{Asgn}(C_2)$
 $\text{Asgn}(\text{while } b \text{ do } C) = \text{Asgn}(C)$

- ② $\& \langle C, s \rangle \rightarrow s' \text{ e } x \in \text{Asgn}(C) \text{ então } sx = s'x$?

Prova por indução na derivação de $\langle C, s \rangle \rightarrow s'$. caso a última regra da derivação seja

[skip] $\langle \text{skip}, s \rangle \rightarrow s$ logo $sx = sx$ ✓

[ass] $\langle y := a, s \rangle \rightarrow s[y \mapsto a]s$ e $\text{Asgn}(y := a) = \{y\}$

se $x \neq y$ temos $sx = s[y \mapsto a]x$ ✓

[comp] $\& \langle C_1; C_2, s \rangle \xrightarrow{\text{①}} s' \text{ e } x \notin \text{Asgn}(C_1; C_2) \text{ ②}$

De ① temos $\langle C_1, s \rangle \xrightarrow{\text{③}} s'' \text{ e } \langle C_2, s'' \rangle \xrightarrow{\text{④}} s' \text{ por algum } s''$

de ② temos $x \notin \text{Asgn}(C_1) \text{ e } x \notin \text{Asgn}(C_2) \text{ ⑤ ⑥}$

De 3+5, por HI, temos $sx = s''x \rightarrow \text{logo } sx = s'x$ ✓

De 4+5, por HI, temos $s''x = s'x$