

Feature Engineering for Machine learning in Graphs

Prof. O-Joun Lee

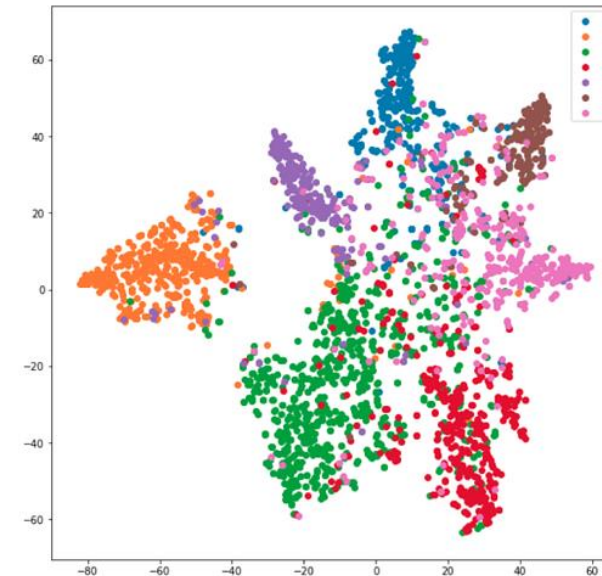
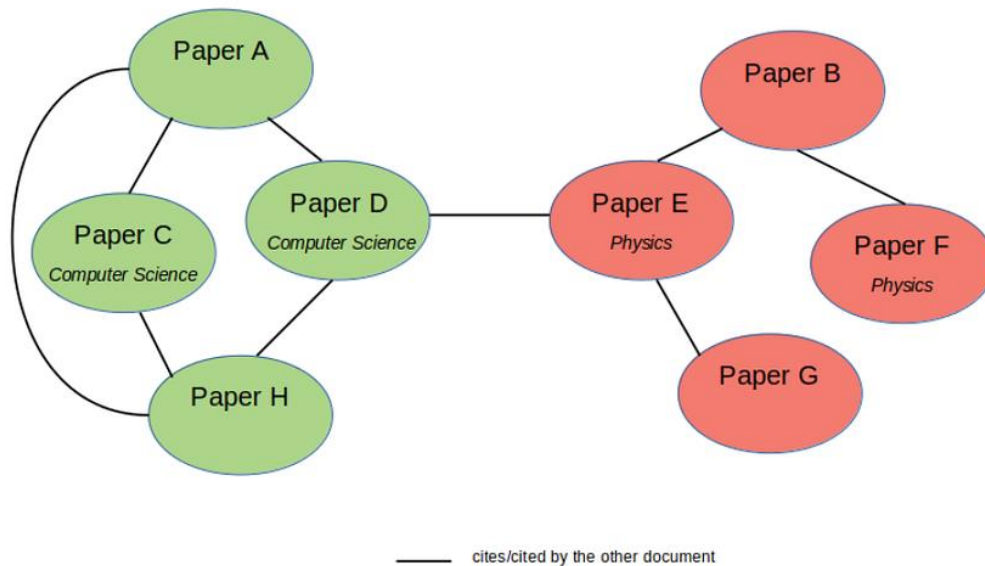
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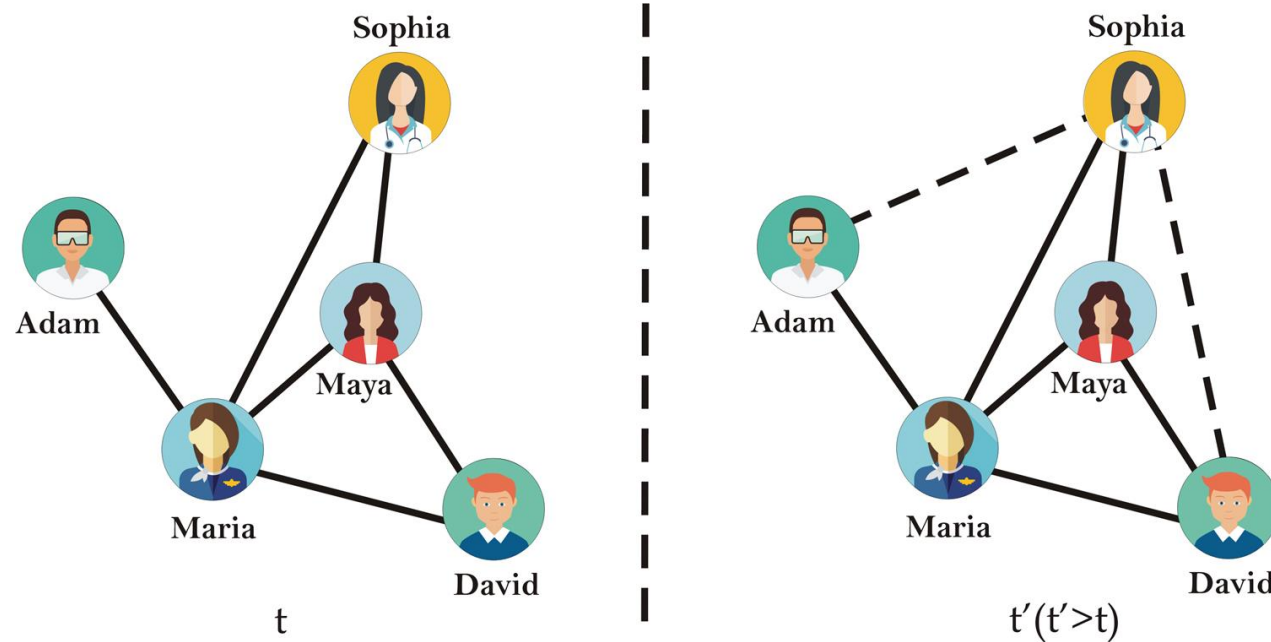


- 1. Node-level prediction
 - Node centrality
 - Structure-based features (Node degree, Graphlets, ...)
- 2. Link-level prediction
 - Distance-based features
 - Local & Global neighborhood overlap
- 3. Graph-level prediction
 - Graphlet Kernel
 - Weisfeiler-Lehman Kernel

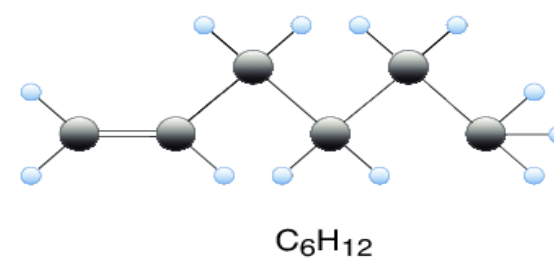
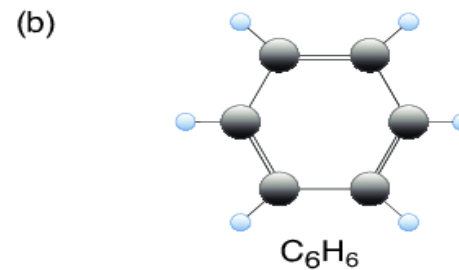
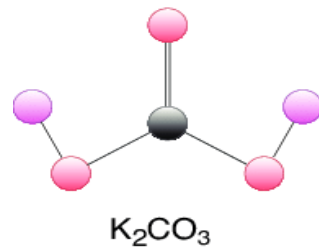
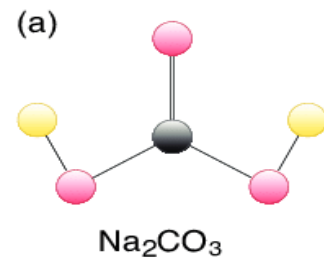
- Predicting the classes or labels of nodes.
- For example, detecting the paper field in a citation network.



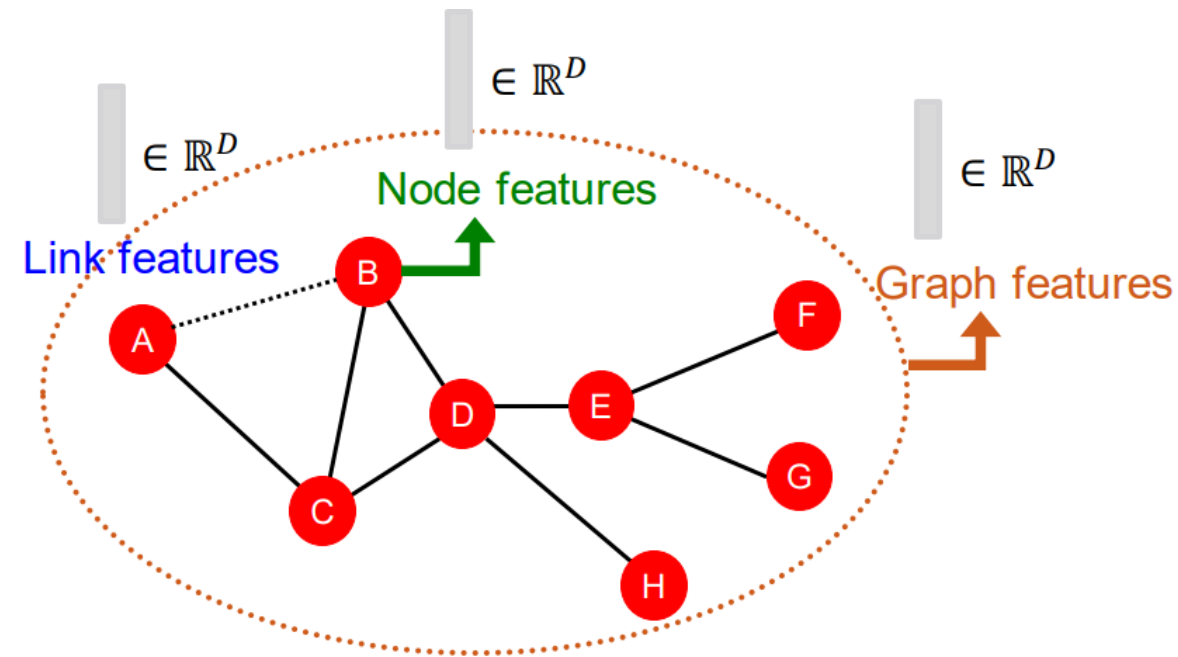
- Predicting the classes or labels of nodes.
- For example, a social networking service suggests possible friend connections based on network data.



- Classifying a graph itself into different categories.
- Inputs: A collection of graphs.
- For example, determining if a chemical compound is toxic or non-toxic by looking at its graph structure.



- Design features for nodes/links/graphs
- Obtain features for all training data



➤ **Feature extraction:**

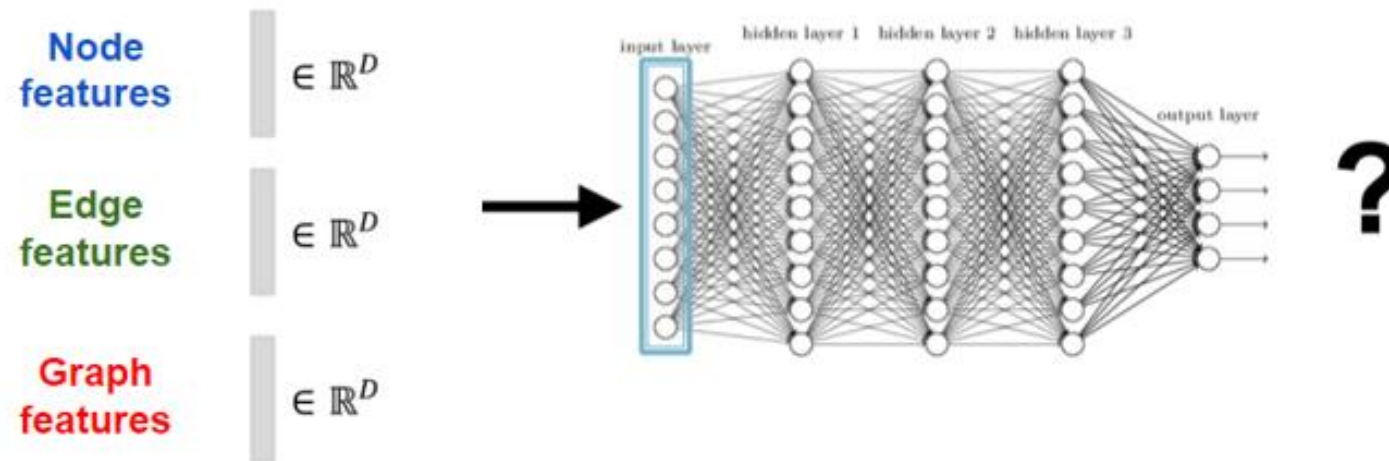
- Node features
- Edge features
- Graph features

➤ **Train a ML model:**

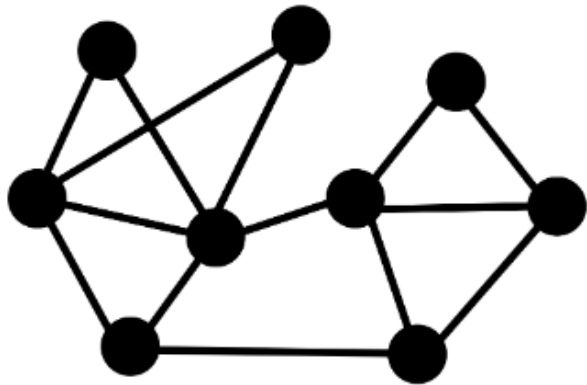
- Logistic Regression
- Random forest
- Neural network, etc.

➤ **Apply the model:**

- Given a new node/link/graph, obtain its features and make a prediction



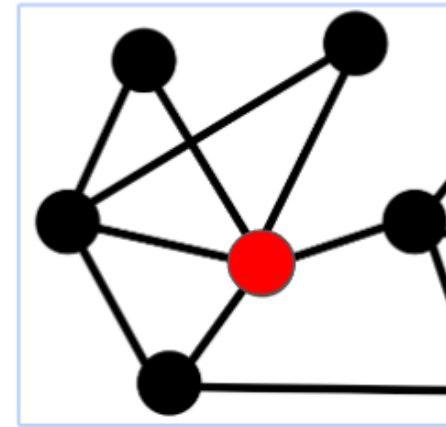
- Using effective features over graphs is the key to achieving good model performance.
- Besides the original node/edge/graph features, can we embed topology structure into node/edge/graph features?



Graph



Node prediction



With neighboring topology

- Goal:

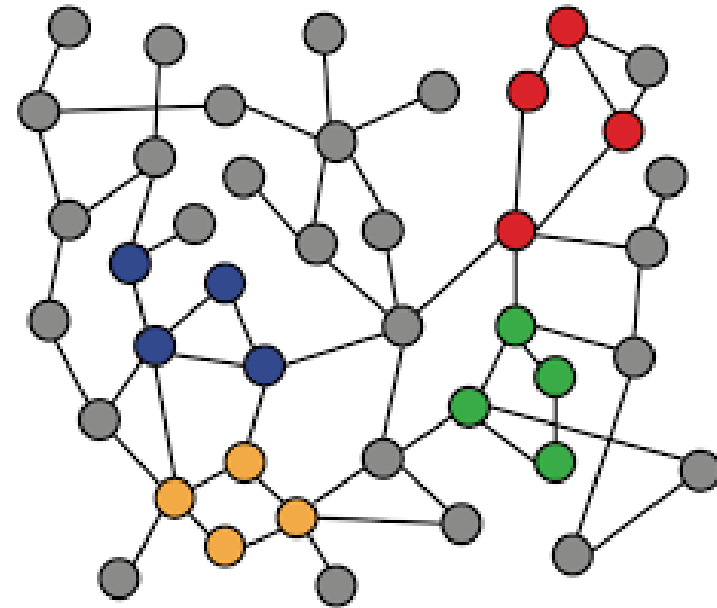
Characterize the structure and position of a node in the network:

- Importance-based features

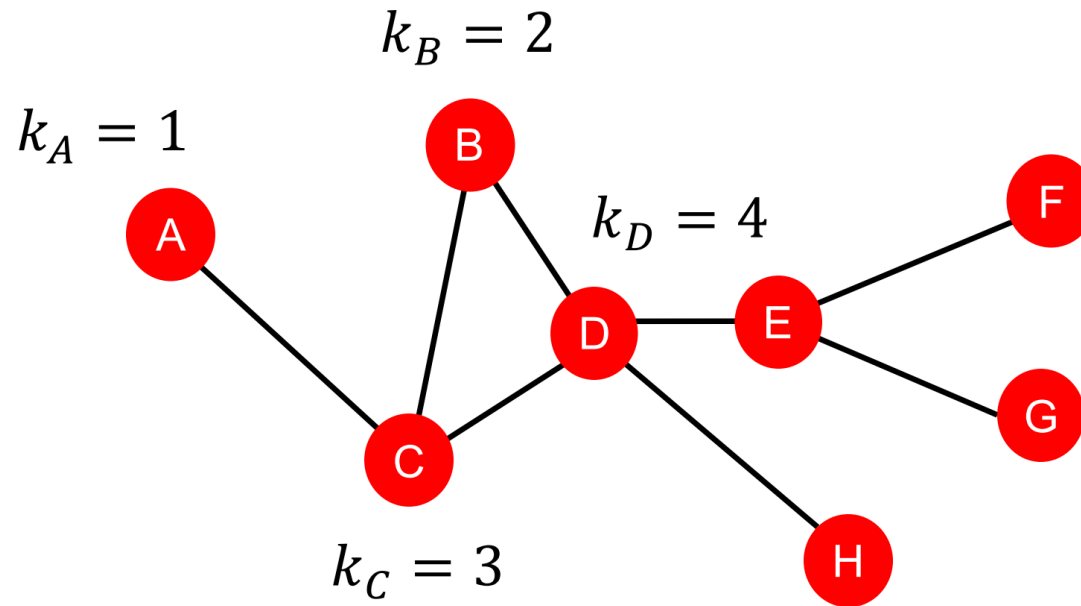
- Node degree
- Node centrality

- Structure-based features

- Node degree
- Clustering coefficient
- Graphlets



- The degree k_v of node v is the number of edges (neighboring nodes) the node has.
- Treats all neighboring nodes equally.



-
- A graph with 6 nodes and 10 edges. One node is highlighted in red.

➤ Eigenvector Centrality

- A node v is important if surrounded by important neighboring nodes $u \in N(v)$.
- We model the centrality of node v as the sum of the centrality of neighboring nodes recursively:

$$c_v = \frac{1}{\lambda} \sum_{u \in N(v)} c_u \quad \longleftrightarrow \quad \lambda \mathbf{c} = \mathbf{A} \mathbf{c}$$

- \mathbf{A} : Adjacency matrix
 $A_{uv} = 1$ if $u \in N(v)$
- \mathbf{c} : Centrality vector
- λ : Eigenvalue

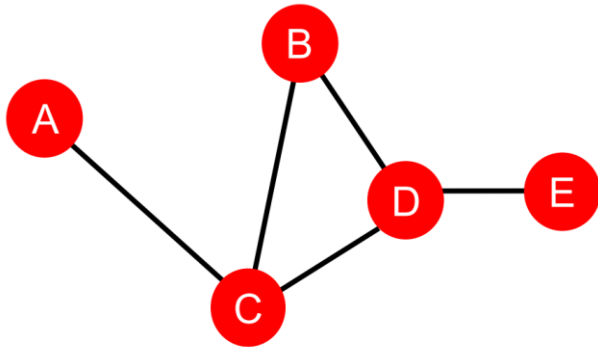
- We can see that centrality \mathbf{c} is the eigenvector for \mathbf{A}

➤ Betweenness Centrality

- A node is important if it **lies on many shortest paths** between other nodes.

$$c_v = \sum_{s \neq v \neq t} \frac{\#(\text{shortest paths between } s \text{ and } t \text{ that contain } v)}{\#(\text{shortest paths between } s \text{ and } t)}$$

- For example:



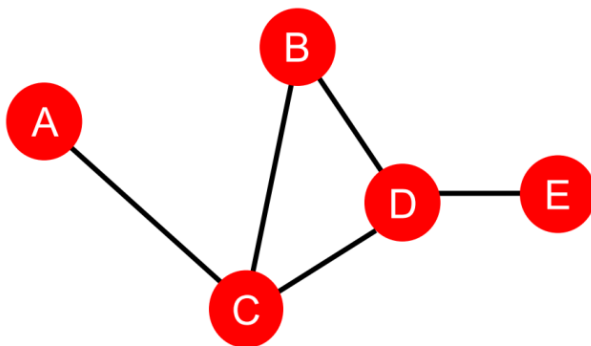
- $C_A = C_B = C_E = 0$
- $C_C = 3$ (A-C-B, A-C-D, A-C-D-E)
- $C_D = 3$ (A-C-D-E, B-D-E, C-D-E)

➤ Closeness Centrality

- A node is important if it has **small shortest path lengths** to all other nodes.

$$c_v = \frac{1}{\sum_{u \neq v} \text{shortest path length between } u \text{ and } v}$$

- For example:



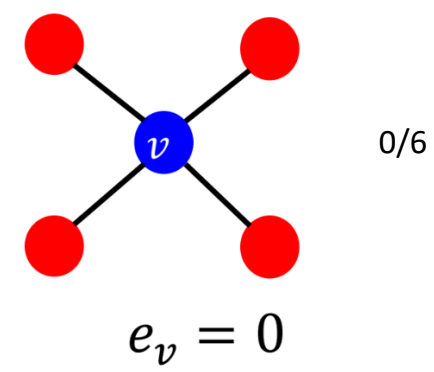
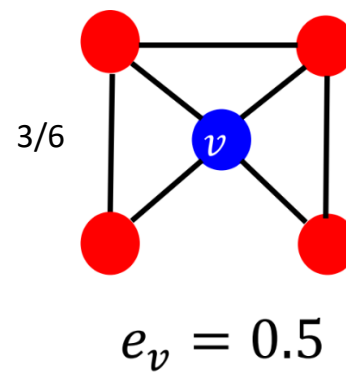
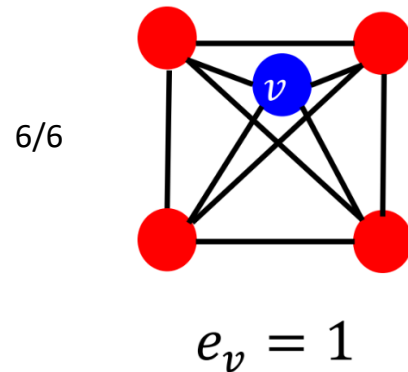
- $C_A = 1/(2 + 1 + 2 + 3) = 1/8$
(A-C-B, A-C, A-C-D, A-C-D-E)
- $C_D = 1/(2 + 1 + 1 + 1) = 1/5$
(D-C-A, D-B, D-C, D-E)

- Measures **how connected v's neighboring nodes** are:

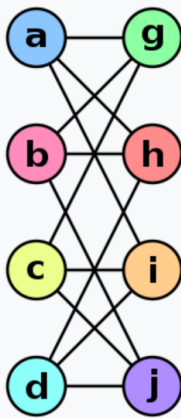
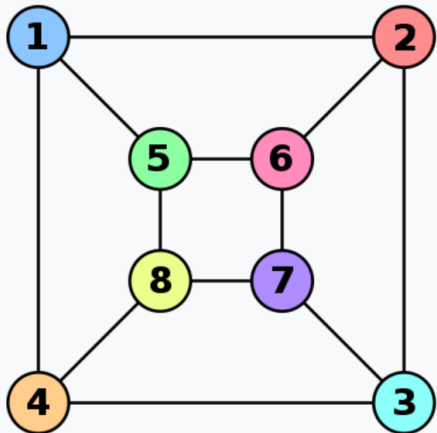
$$e_v = \frac{\#(\text{edges among neighboring nodes})}{\binom{k_v}{2}} \in [0,1]$$

#(node pairs among k_v neighboring nodes)
In our examples below the denominator is 6 (4 choose 2).

- For example:

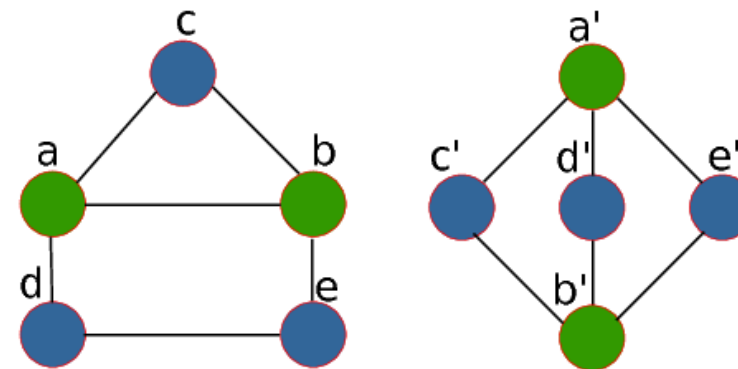


- **Isomorphic graphs** having the same number of vertices, edges, and also the same edge connectivity.

Graph G	Graph H	An isomorphism between G and H
		$f(a) = 1$ $f(b) = 6$ $f(c) = 8$ $f(d) = 3$ $f(g) = 5$ $f(h) = 2$ $f(i) = 4$ $f(j) = 7$

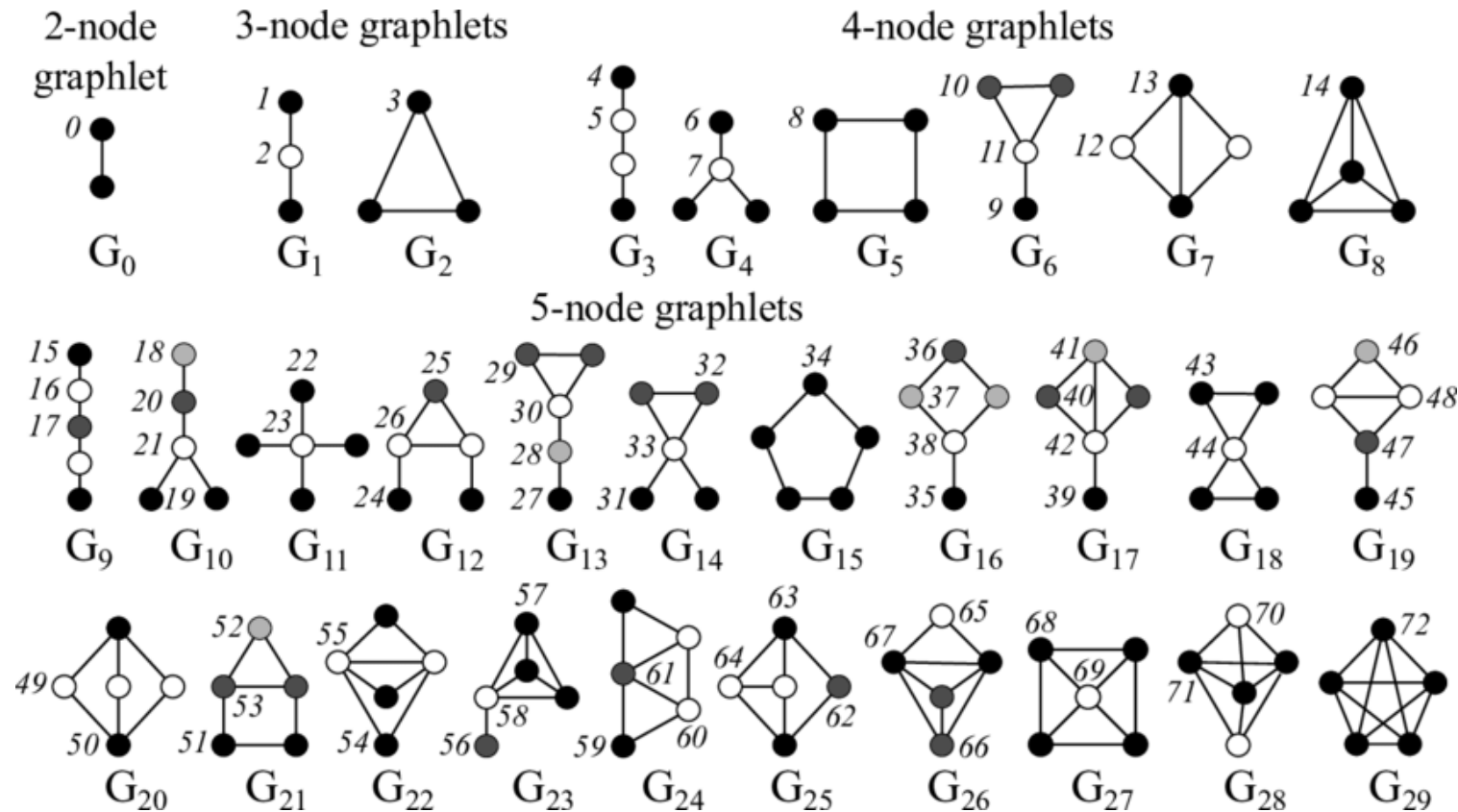


Isomorphic

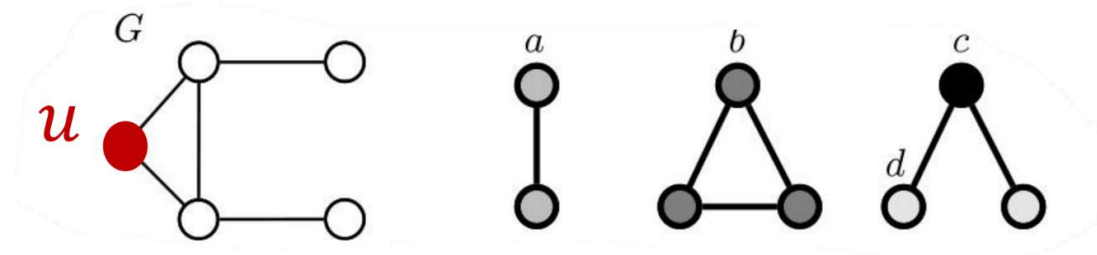


Non-Isomorphic

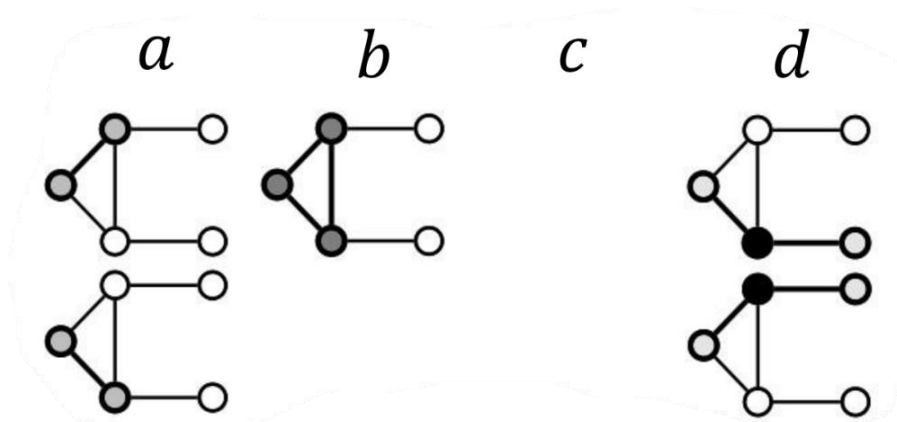
- Graphlets are **induced, non-isomorphic** subgraphs that **describe the structure of node u 's network neighborhood**.



- Graphlet Degree Vector (GDV): A count vector of graphlets rooted at a given node.
- Graphlet degree vector provides a measure of a node's local network topology (more detail than node degrees or clustering coefficient.)

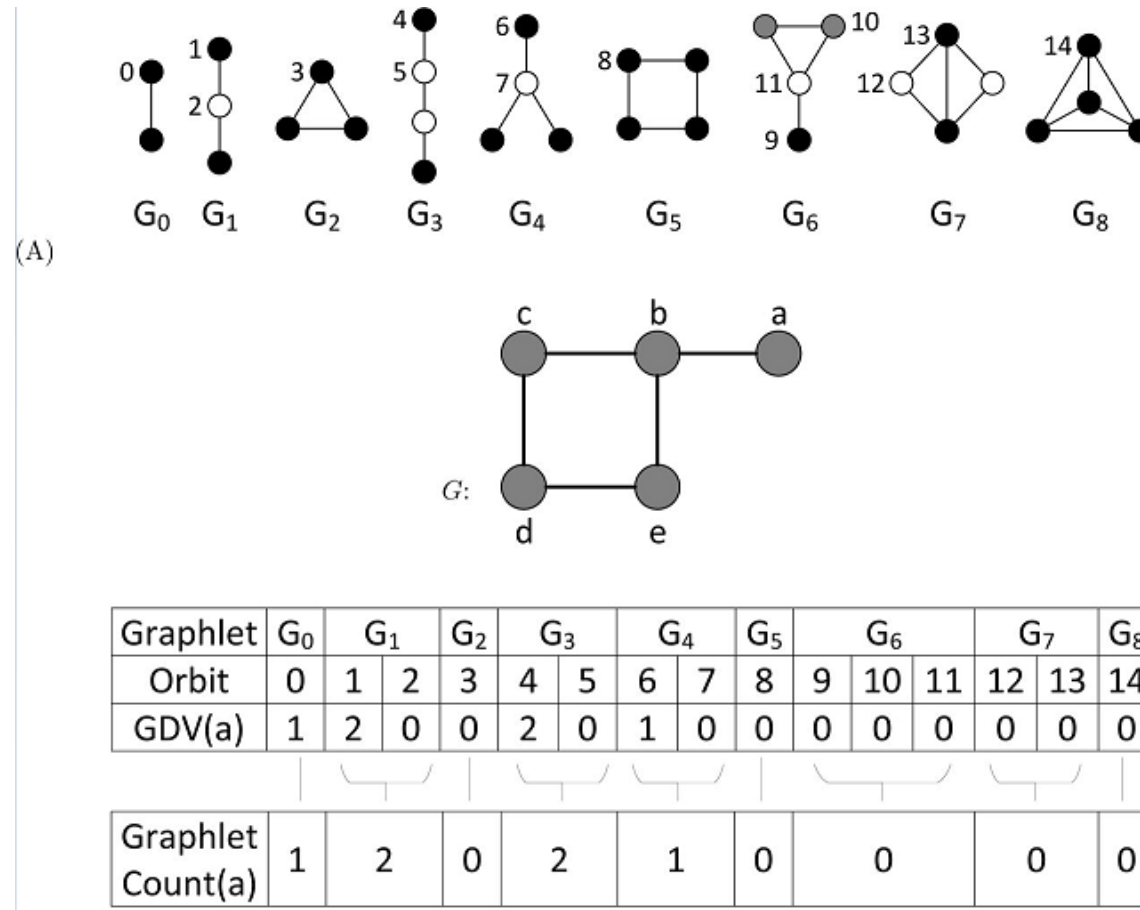


Select graphlets up to 3 nodes



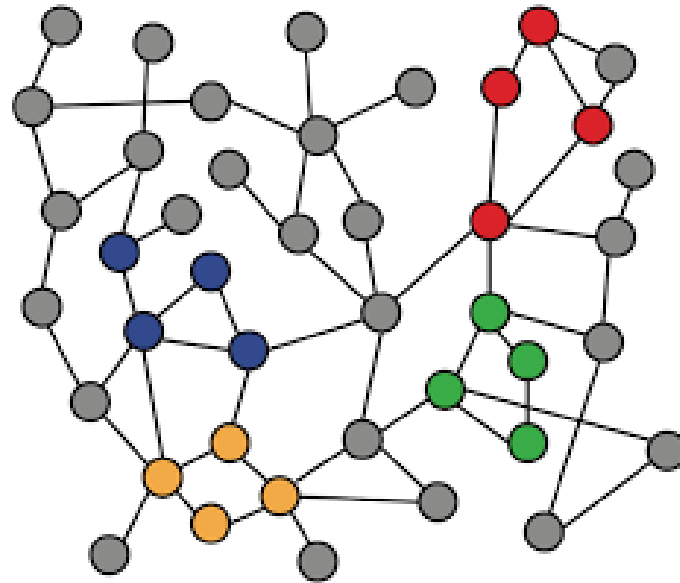
Graphlet instances for node u

- Graphlets are induced, non-isomorphic subgraphs that describe the structure of node u 's network neighborhood.
- Example:

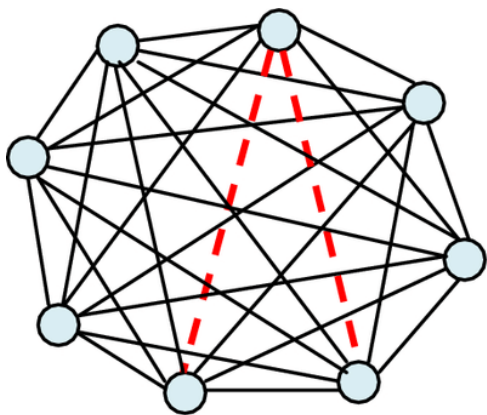


- Importance-based features: capture the importance of a node in a graph, useful for predicting influential nodes in a graph
 - Node degree
 - Node centrality (Eigenvector, Betweenness, Closeness Centrality)
- Structure-based features: capture topological properties of local neighborhood around a node, useful for predicting a particular role a node plays in a graph
 - Node degree
 - Clustering coefficient
 - Graphlets

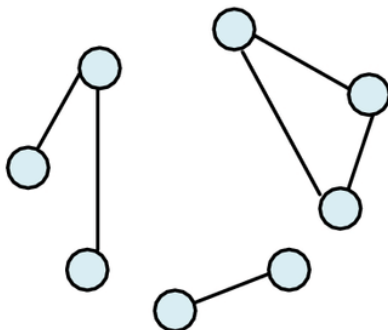
- Distance-based Features
- Local Neighborhood Overlap Feature
- Global Neighborhood Overlap Feature



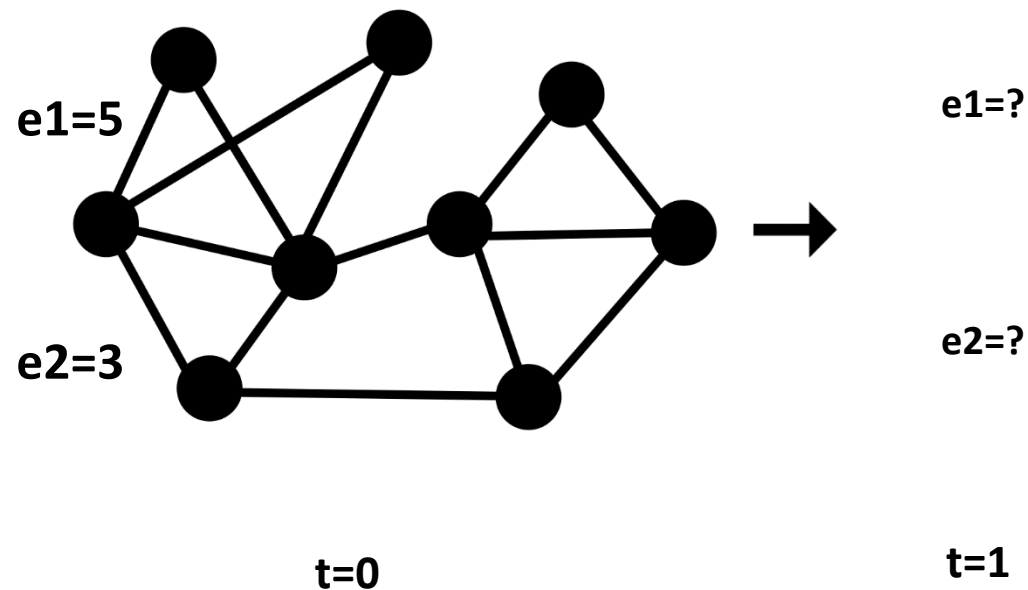
- Predict new links based on existing links.
- Predict edge-level features/labels.



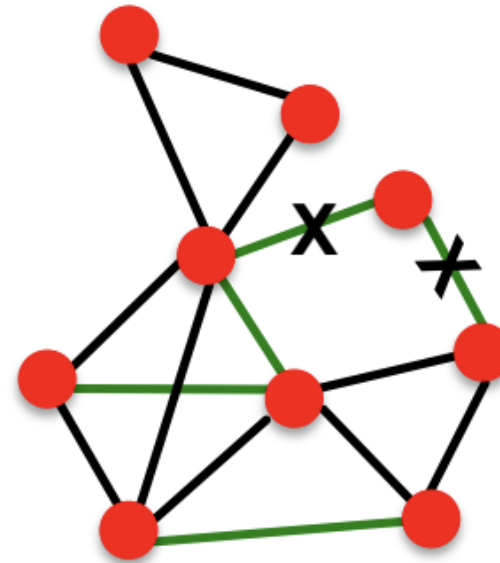
(a)
Training



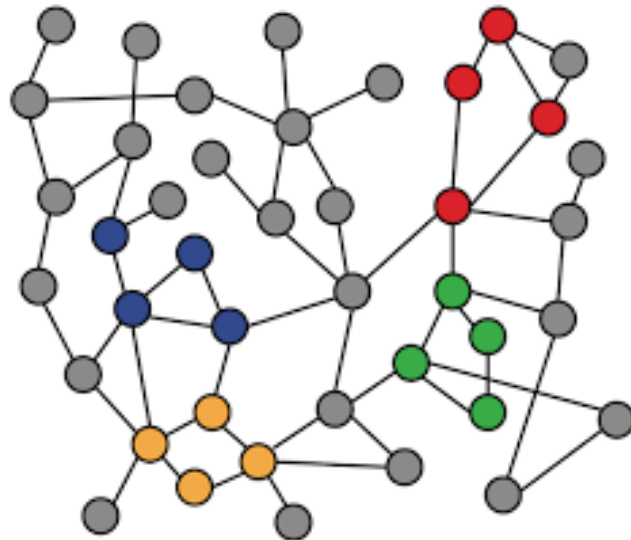
(b)
Testing



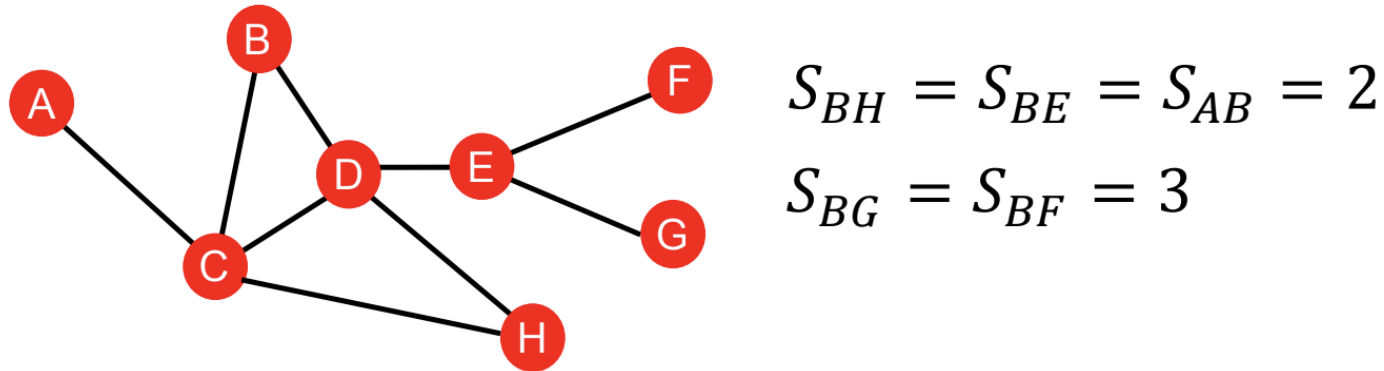
- Select top k pairs as new edges
- $\text{Edge_feature}(x, y, e) \rightarrow \text{new edge_feature}$



- Goal: Characterize the structure and connectivity of nodes and edges in the network:
- Distance-based features
- Local neighborhood overlap
- Global neighborhood overlap



- Shortest-path distance between two nodes
- For example:



- However, this does not capture the degree of neighborhood overlap:
- Node pair (B, H) has 2 shared neighboring nodes, while pairs (B, E) and (A, B) only have 1 such node.

➤ Captures # neighboring nodes shared between two nodes v_1 and v_2 :

➤ Common neighbors: $|N(v_1) \cap N(v_2)|$

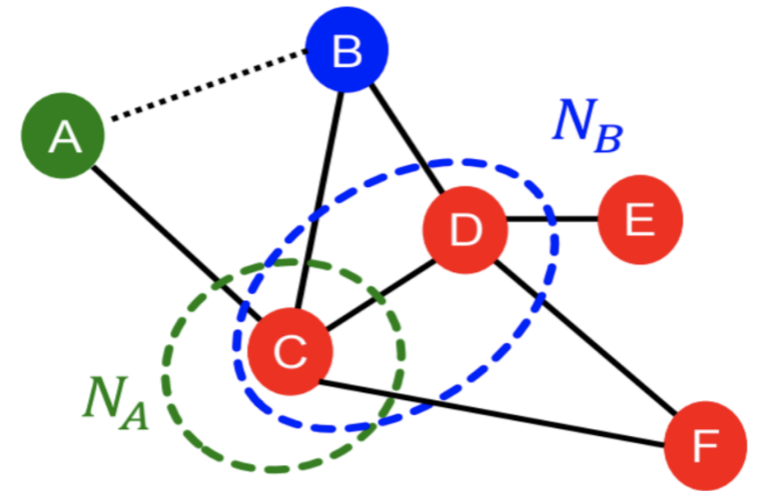
■ Example: $|N(A) \cap N(B)| = |\{C\}| = 1$

➤ Jaccard's coefficient: $\frac{|N(v_1) \cap N(v_2)|}{|N(v_1) \cup N(v_2)|}$

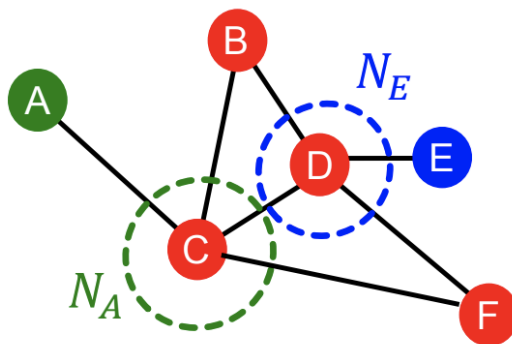
■ Example: $\frac{|N(A) \cap N(B)|}{|N(A) \cup N(B)|} = \frac{|\{C\}|}{|\{C, D\}|} = \frac{1}{2}$

➤ Adamic-Adar index: $\sum_{u \in N(v_1) \cap N(v_2)} \frac{1}{\log(k_u)}$

■ Example: $\frac{1}{\log(k_C)} = \frac{1}{\log 4}$



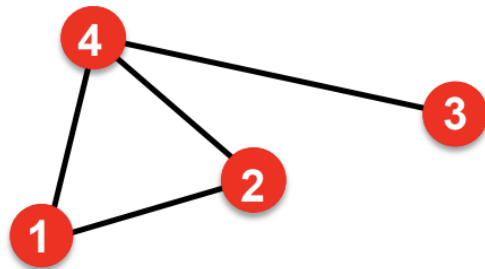
- Limitation of local neighborhood overlapping features:
 - Metric is always zero if the two nodes do not have any neighbors in common.
 - However, the two nodes may still potentially be connected in the future.



$$N_A \cap N_E = \phi$$
$$|N_A \cap N_E| = 0$$

→ Global neighborhood overlap features resolve the limitation by considering the entire graph.

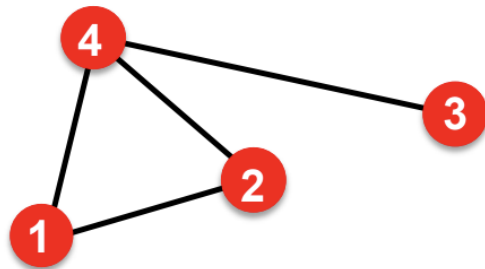
- Katz index:
 - Count the number of walks of all lengths between a given pair of nodes.
- How to compute number of walks?
 - Powers of the graph adjacency matrix



$$A = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

$A_{i,j} = 1$ if node i, j are connected

- So, what is number of walks?
- We want to show: $\mathbf{P}^{(K)} = \mathbf{A}^k$
- where $\mathbf{P}_{uv}^{(K)}$ = #walks of length K between u and v
- $\mathbf{P}_{uv}^{(1)}$ = #walks of length 1 (direct neighborhood) between u and v = \mathbf{A}_{uv}

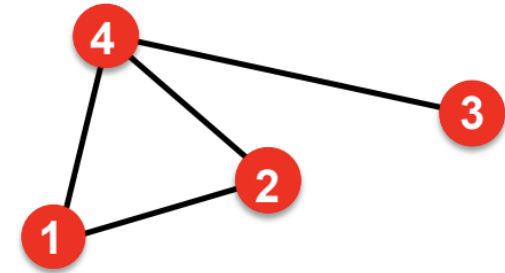


$$\mathbf{P}_{12}^{(1)} = \mathbf{A}_{12}$$

$$\mathbf{A} = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

$\mathbf{A}_{i,j} = 1$ if node i, j are connected

- How to compute $P_{uv}^{(2)}$?
 - Step 1: Compute **#walks** of length 1 **between each of u 's neighbor and v**
 - Step 2: **Sum up** these #walks across u 's neighbors
 - $P_{uv}^{(2)} = \sum_i A_{ui} * P_{iv}^{(1)} = \sum_i A_{ui} * A_{iv} = A_{uv}^2$



Node 1's neighbors

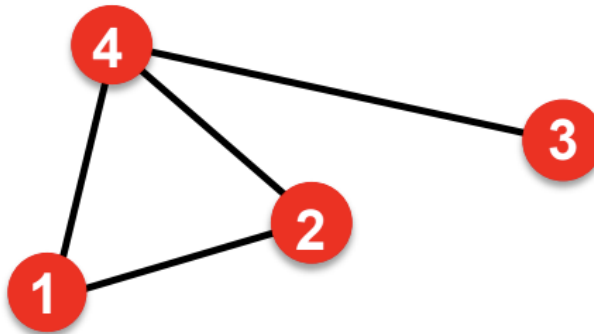
#walks of length 1 between Node 1's neighbors and Node 2

$P_{12}^{(2)} = A_{12}^2$

Power of adjacency

$$A^2 = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix} \times \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix} = \begin{pmatrix} 2 & 1 & 1 & 1 \\ 1 & 2 & 1 & 1 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 3 \end{pmatrix}$$

- Using powers of adjacency matrix, we can calculate number of walks of all lengths between a pair of nodes.
- A_{uv} specifies #walks of length 1 (direct neighborhood) between u and v.
- A_{uv}^2 specifies #walks of length 2 (neighbor of neighbor) between u and v.
- A_{uv}^l specifies #walks of length l.



- **Katz index** between v_1 and v_2 is calculated as **sum over all walk lengths**.

$$S_{v_1 v_2} = \sum_{l=1}^{\infty} \boxed{\beta^l} \boxed{A_{v_1 v_2}^l} \quad \begin{array}{l} \text{\#walks of length } l \\ \text{between } v_1 \text{ and } v_2 \end{array}$$

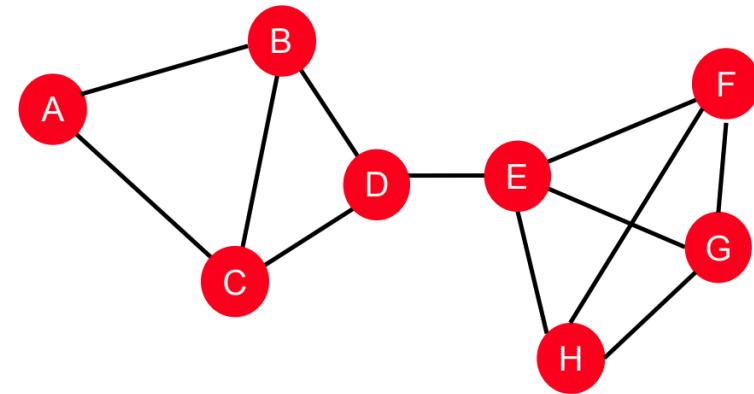
$0 < \beta < 1$: discount factor

- **Katz index matrix** is computed in closed-form:

$$S = \sum_{i=1}^{\infty} \beta^i A^i = (I - \beta A)^{-1} - I$$

- Distance-based features:
 - Calculates shortest path between 2 nodes, but cannot capture overlapping neighborhoods.
- Local neighborhood overlap:
 - Captures number of sharing neighborhoods between 2 nodes.
 - Only focuses nodes within 2-hop.
- Global neighborhood overlap:
 - Katz index uses entire graph structure to score 2 nodes.
 - It can capture the structure globally.

- Goal: We want features that characterize the structure of an entire graph.
- Similar graphs has similar features
- Graph Kernel Methods (kernel methods are widely-used in traditional ML for graph-level prediction.):
- Graphlet Kernel
- Weisfeiler-Lehman Kernel



$f(\text{graph}) = ?$

➤ A quick introduction to Kernels:

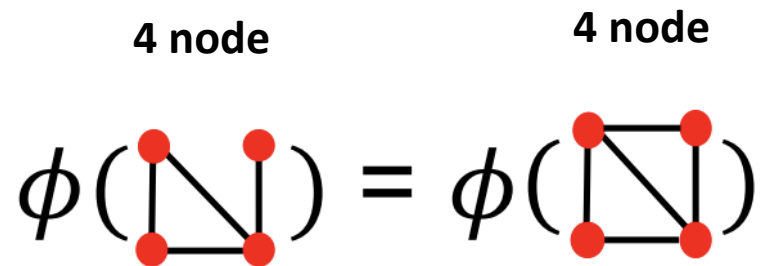
- Kernel $K(G, G') \in \mathbb{R}$ measures similarity between two graphs (data)
- There exists a feature representation $\phi(\cdot)$ such that:

$$K(G, G') = \phi(G)^T \phi(G')$$

- Kernel enables vectors to be operated in higher dimension
- Once the kernel is defined, off-the-shelf ML model, such as kernel SVM, can be used to make predictions.

$$K = (\begin{array}{c} \text{Graph 1} \\ \text{Graph 2} \end{array} , \begin{array}{c} \text{Graph 3} \\ \text{Graph 4} \end{array}) \rightarrow \phi (\begin{array}{c} \text{Graph 5} \\ \text{Graph 6} \end{array})$$

- Goal: Design graph kernel $\phi(G)$
- Key Idea: Bag-of-Words (BoW) for a graph
 - In NLP, BoW counts the word's frequency in a document as feature
 - Simplest way on graph: **Regard nodes as words.**
 - For example, we found the features of 2 graphs are same.



- Problem: Bag of node counts doesn't work well...
- What if we use Bag of Node Degrees?

Deg1: ● Deg2: ● Deg3: ●

$$\phi(\text{Graph 1}) = \text{count}(\text{Graph 2}) = [1, 2, 1]$$

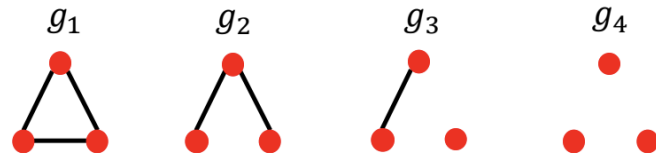
$$\phi(\text{Graph 3}) = \text{count}(\text{Graph 4}) = [0, 2, 2]$$

Obtains different features for different graphs!

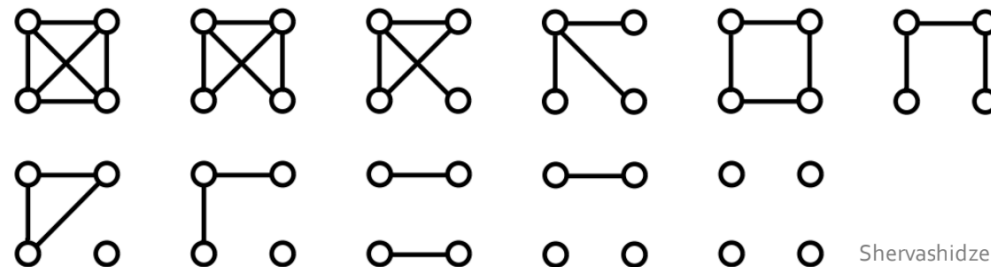
- Both Graph kernel and Weisfeiler-Lehman kernel use Bag-of-* representation of graph, but * is more sophisticated than node degrees!

- Key Idea: Count the number of different graphlets in a graph
- Graphlet Differences:
 - Nodes do not need to be connected.
 - Graphlets are not rooted.

- For $k = 3$, there are 4 graphlets.



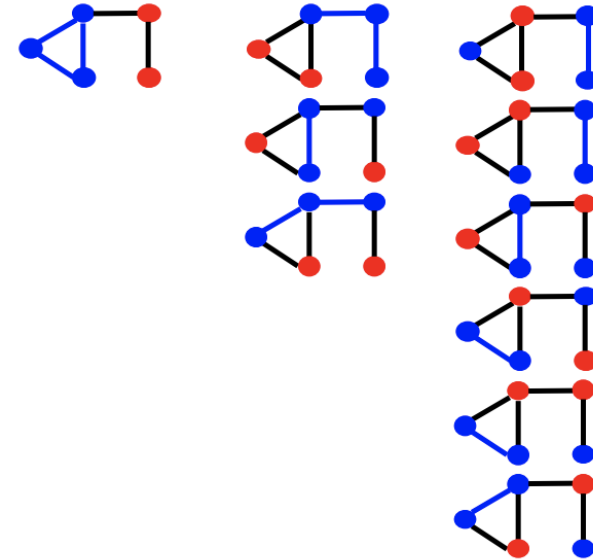
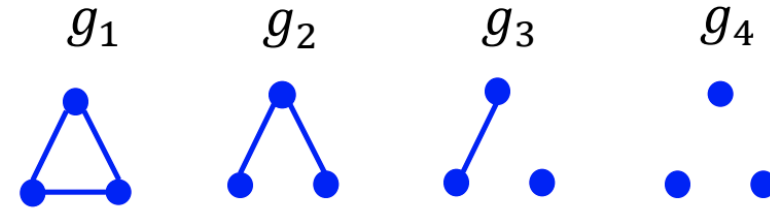
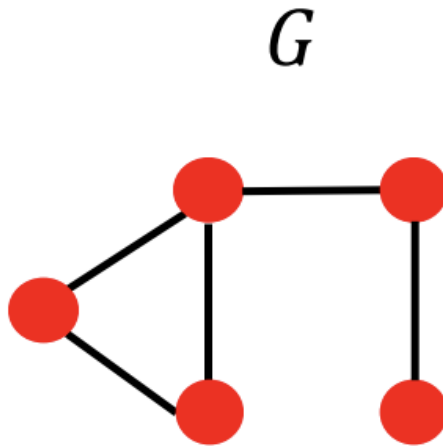
- For $k = 4$, there are 11 graphlets.



Shervashidze et al., AISTATS 2011

➤ For example:

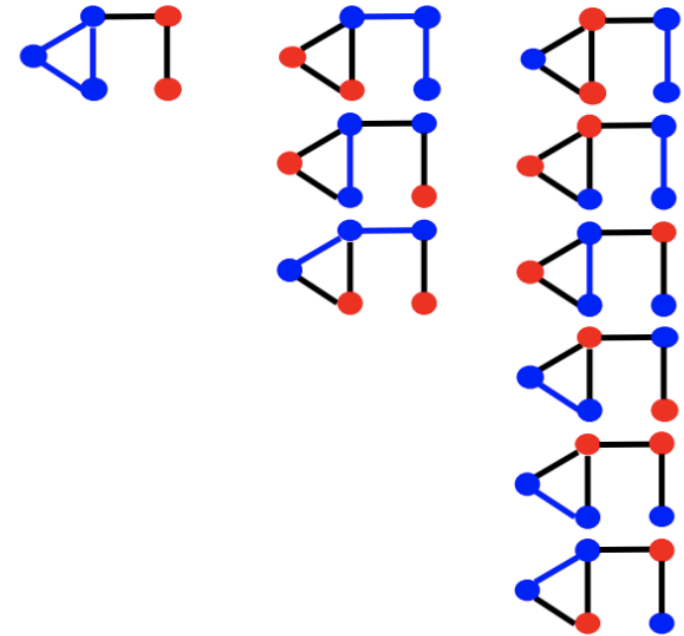
➤ For $k = 3$:



$$\mathbf{f}_G = (1, 3, 6, 0)^T$$

- Counting graphlets is expensive!
- Counting size k graphlets for graph with n nodes by enumeration takes n^k time.
- If our graph changes with time, we have to recalculate the features again.

→ **Can we design a more efficient graph kernel?**



- Goal: Design an efficient graph feature descriptor $\phi(G)$
- Key Idea:
 - Iteratively use neighborhood structure to describe node's neighboring topology.
 - Generalized version of Bag of node degrees (node degree only contains one-hop neighborhood information).

→ **Color Refinement Algorithm**

- For a graph G with nodes V :
- Assign initial color $c^{(0)}(v)$ to each node v
- Iteratively refine node colors by:

$$c^{(k+1)}(v) = \text{HASH} \left(\left\{ c^{(k)}(v), \{c^{(k)}(u)\}_{u \in N(v)} \right\} \right)$$

where HASH maps different inputs to different colors.

- After K steps of iteration, $c(K)(v)$ represents the structure of the K -hop neighborhood.

- WL kernel benefits:
 - Computationally efficient (color refinement need $\#(\text{edges})$ steps)
 - $\#(\text{colors})$ depends on total number of nodes.
 - Can be used to check graph isomorphism

- Graphlet Kernel:
 - Bag of graphlets
 - Computationally expensive
- Weisfeiler-Lehman Kernel:
 - Bag of colors
 - Capture graph structure within K-hop
 - Computationally efficient
 - Closely related to graph neural networks!



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