

Introduction machine learning on graphs

Prof. O-Joun Lee

Dept. of Artificial Intelligence,
The Catholic University of Korea
ojlee@catholic.ac.kr

Contents

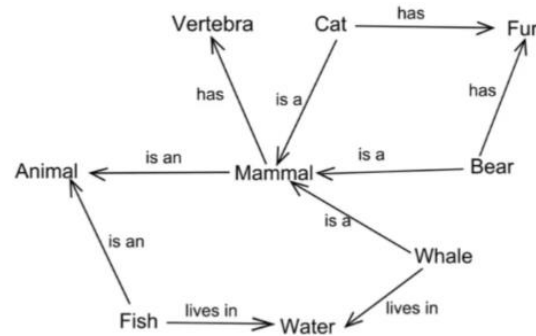


- The overview of Machine Learning on Graphs
- Graph Terminology
- Graph Characterization
 - Centrality measurements
 - Community
- Sample code

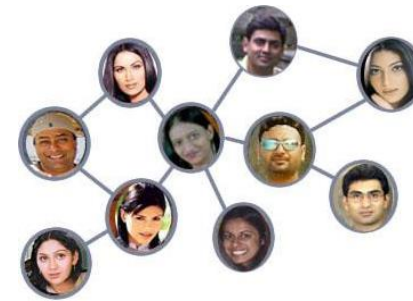
Networks are a general language for describing and modeling complex systems



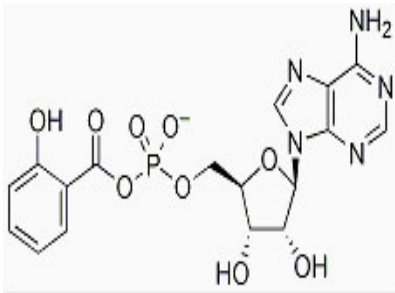
Street network



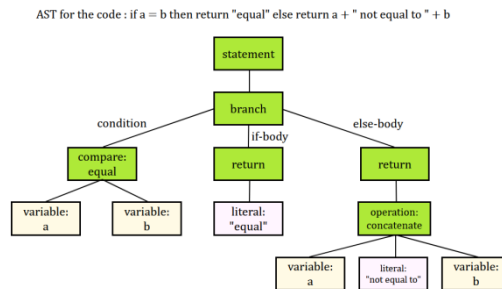
Ecological network



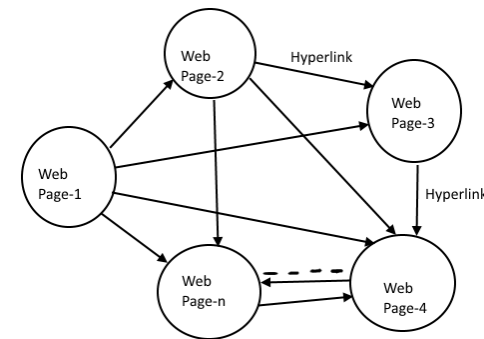
Social media



Chemical network



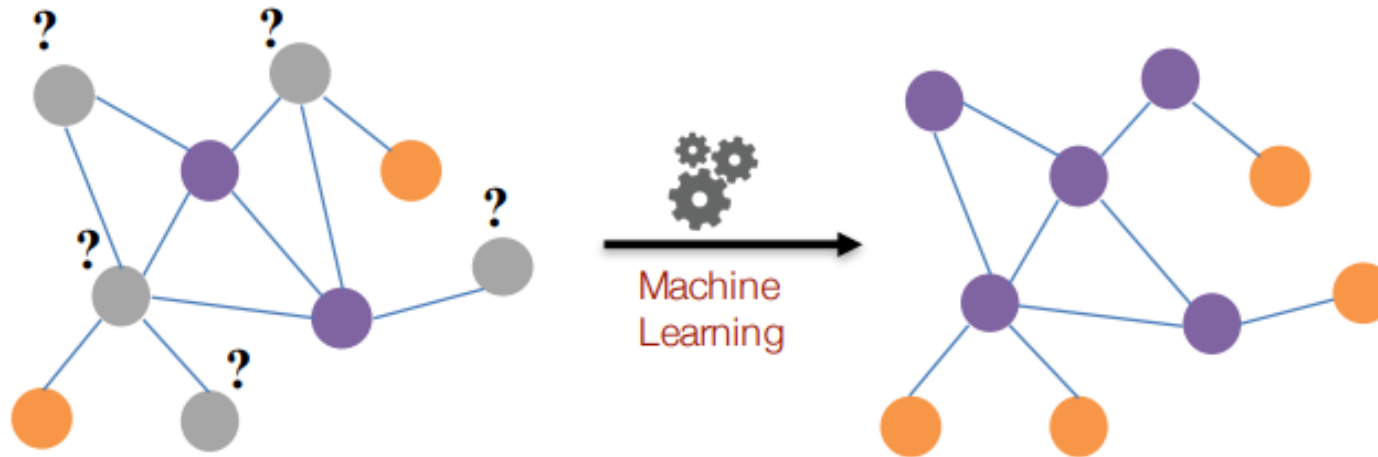
Program flow



Web graph

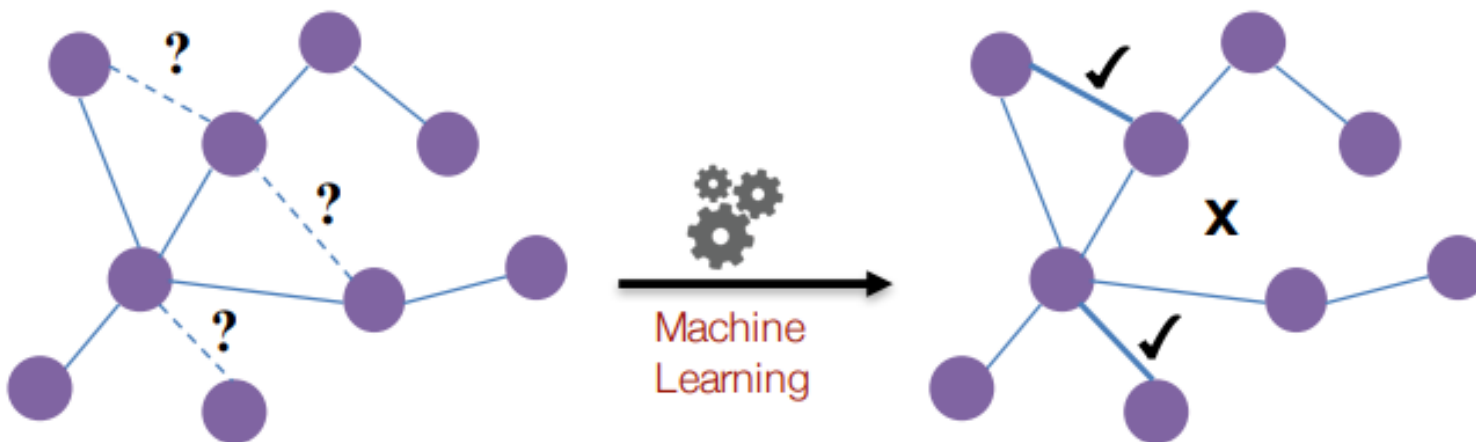
- Universal language for describing complex data
 - Chemical compounds (Cheminformatics)
 - Protein structures, biological pathways/networks (Bioinformatics)
 - Program control flow, traffic flow, and workflow analysis
- Data availability (+computational challenges)
 - Web/mobile, bio,health, and medical data
- Shared vocabulary between fields:
 - Computer science, Social science, Physics, Statistics, Biology
- Impact:
 - Social networking, social media, Drug design

- Node classification
 - Predict a type of a given node

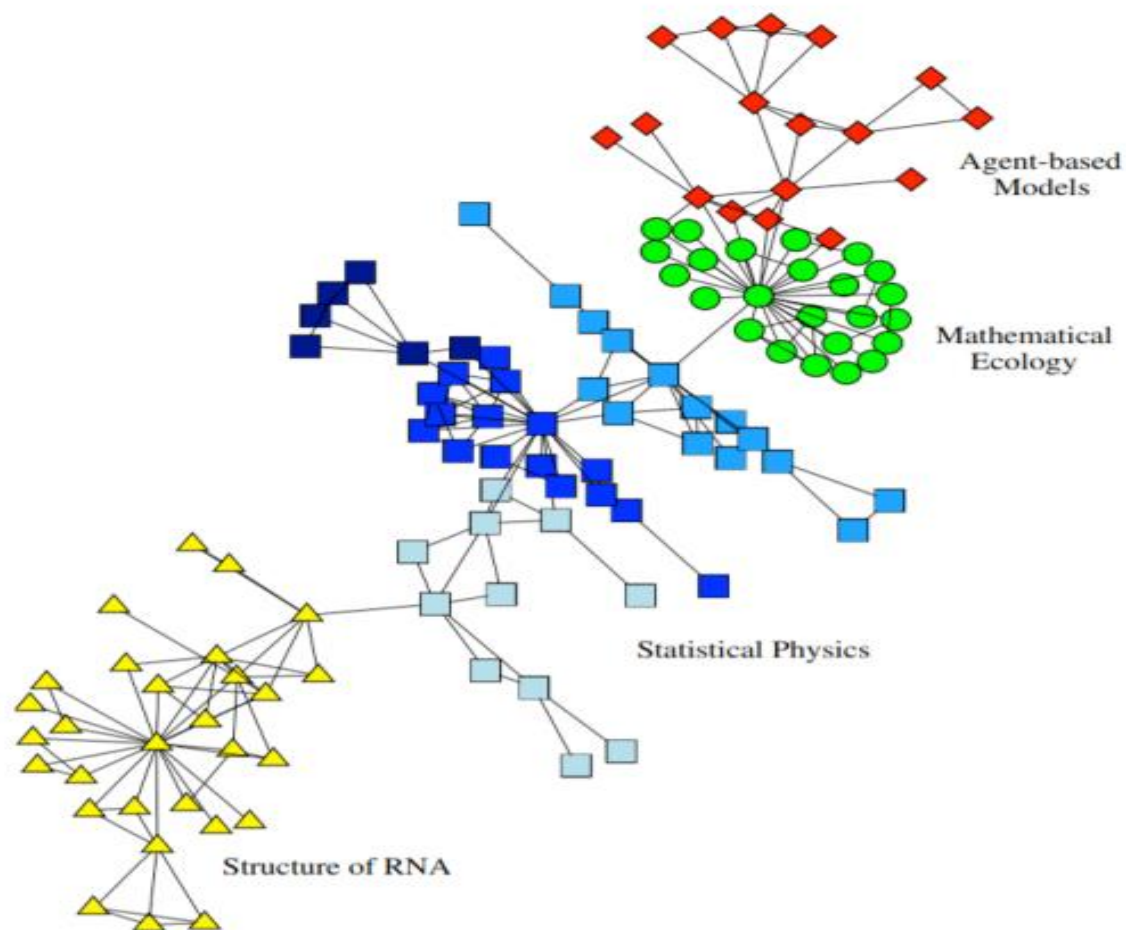


- Many possible ways to create node features:
 - Node degree, PageRank score, motifs

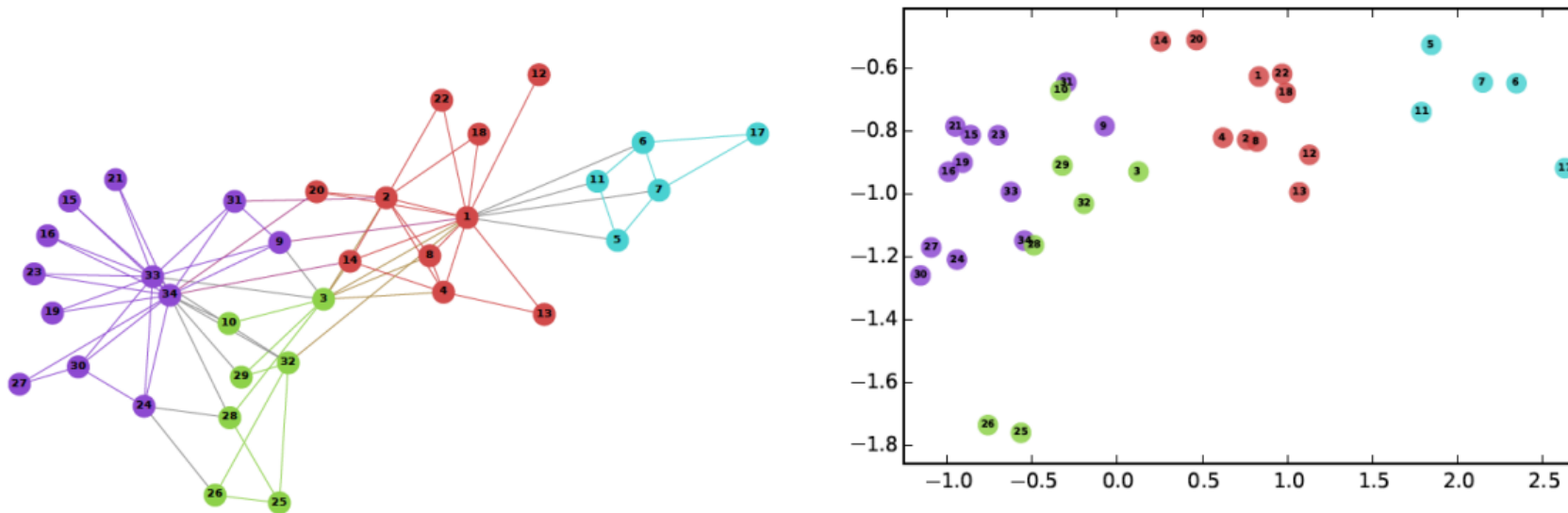
- Link prediction
 - Predict whether two nodes are linked



- Community detection
 - Identify densely linked clusters of nodes



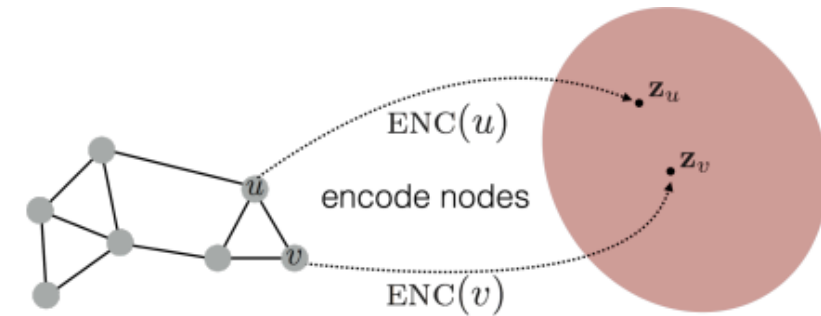
- Goal is to encode nodes so that similarity in the embedding space (e.g., dot product) approximates similarity in the original network.



- Goal is to encode nodes so that similarity in the embedding space (e.g., dot product) approximates similarity in the original network.
- Let z_u be the embedding of node u .
- Goal is to find the encoder function f such that:

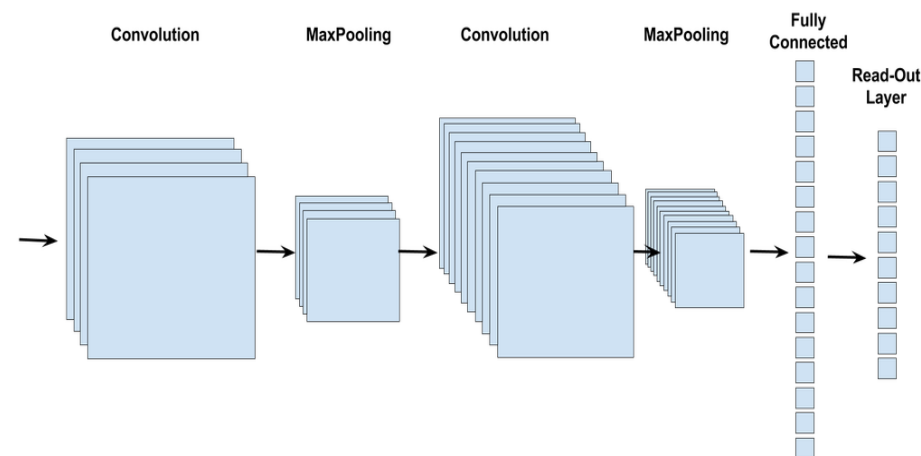
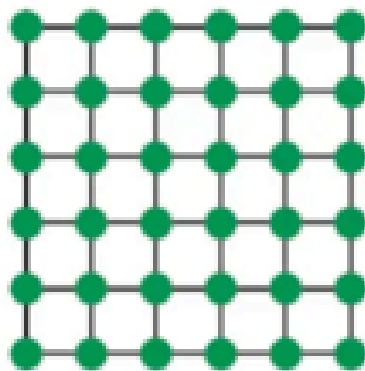
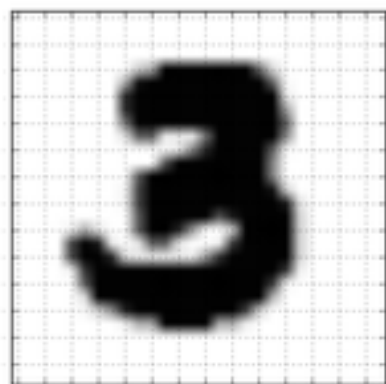
$$\text{similarity}(u, v) \approx z_u^T z_v$$

- Learning node embedding:
 - Define an encoder .
 - Define a node similarity function.
 - Optimize the parameters of the encoder so that:
 $\text{similarity}(u, v) \approx z_u^T z_v$

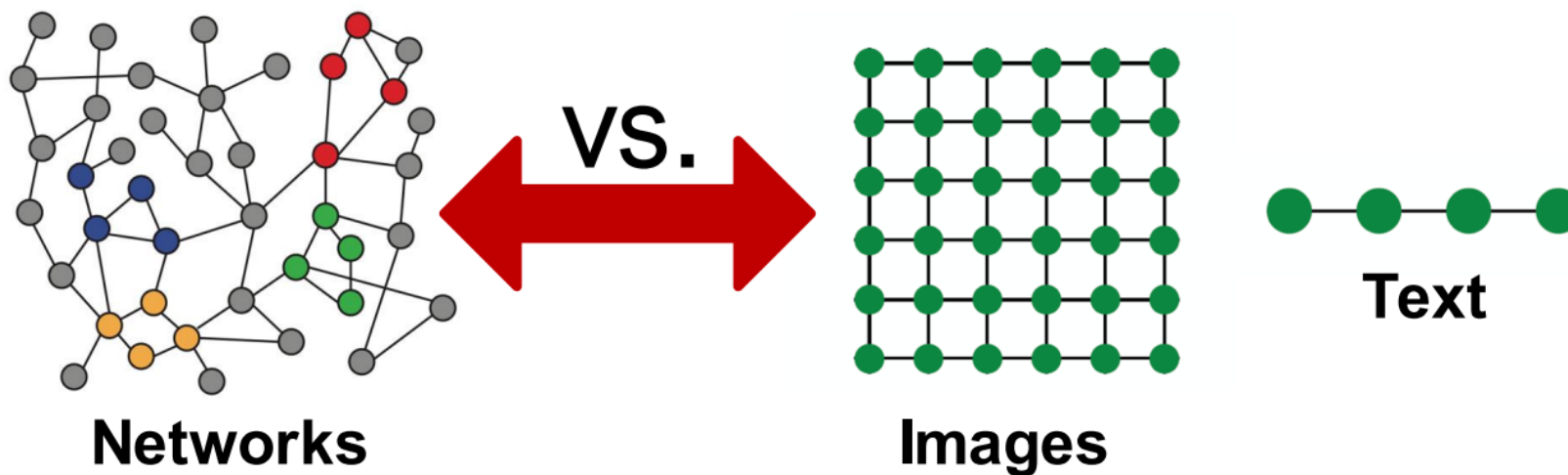


- The goal is to map each node into a low-dimensional space
 - Distributed representation for nodes
 - Similarity between nodes indicates link strength
 - Encodes network information and generate node representation

- Graph data is so complex that it's created a lot of challenges for existing machine learning algorithms.
- Images with the same structure and size can be considered as fixed-size grid graphs.
- Text and speech are sequences, so they can be considered as line graphs. (text and speech have linear 1D structure)



- Graphs have arbitrary size and complex topological structure.
- In graphs, there is no fixed node ordering or reference point.
- Graphs are often dynamic and have multimodal features.

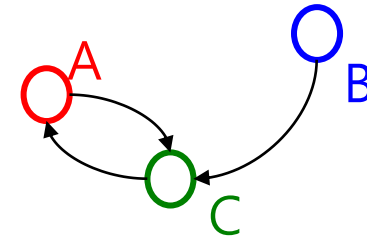


- How can we develop neural networks that are much more broadly applicable?
- Feature learning for networks:
 - “Linearizing” the graphs:
 - Create a “sentence” for each node using random walks (node2vec)
 - Graph neural networks:
 - Propagate information between the nodes in graphs (message passing)

A graph is a pair: $G = (V, E)$:

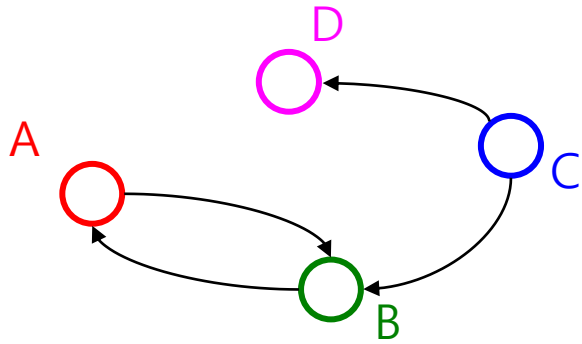
- A set of nodes, also known as nodes: $V = \{v_1, v_2, \dots, v_n\}$
- A set of edges $E = \{e_1, e_2, \dots, e_m\}$
 - Each edge e_i is a pair of nodes (v_j, v_k)
 - An edge "connects" the nodes

Graphs can be *directed* or *undirected*

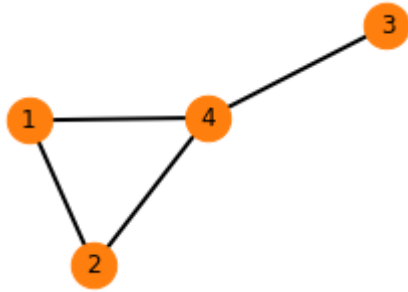


$$V = \{A, B, C\}$$
$$E = \{(B, C), (A, C), (C, A)\}$$

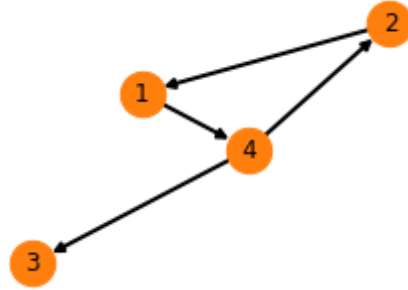
➤ Adjacency Matrix



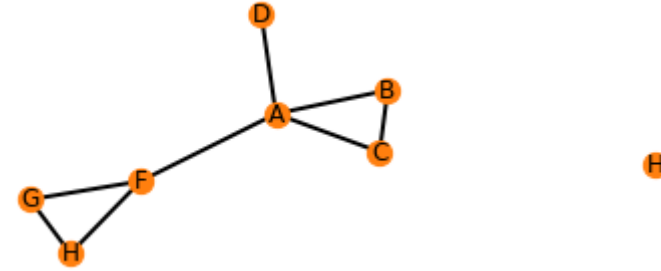
	A	B	C	D
A	0	1	0	0
B	1	0	0	0
C	0	1	0	1
D	0	0	0	0



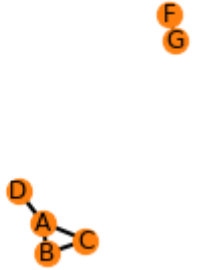
Undirected



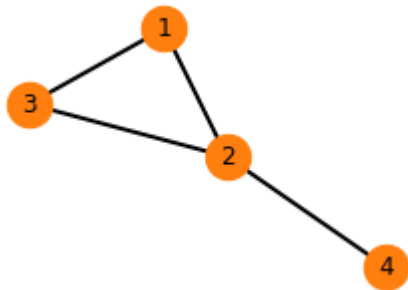
Directed



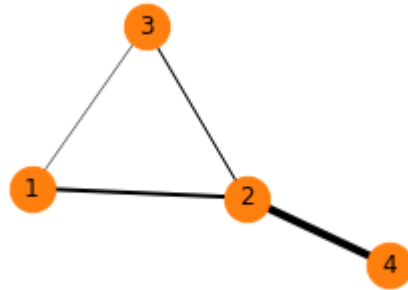
Connected



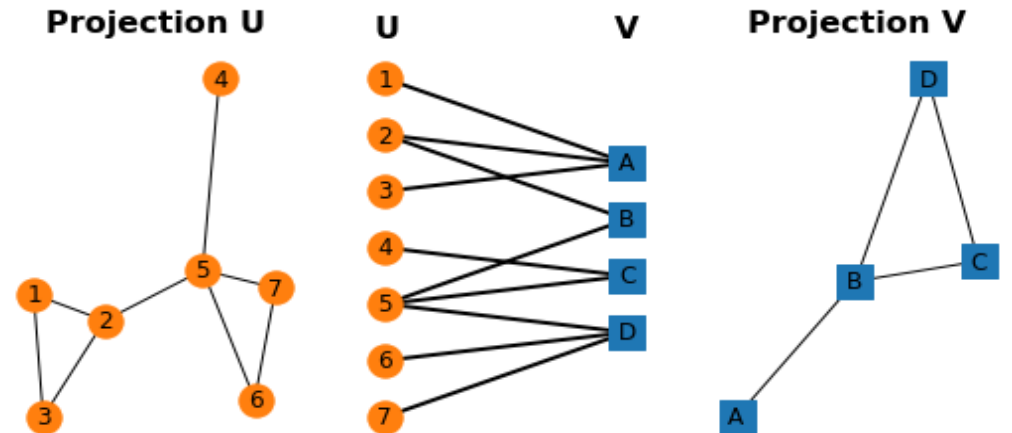
Disconnected



Unweighted

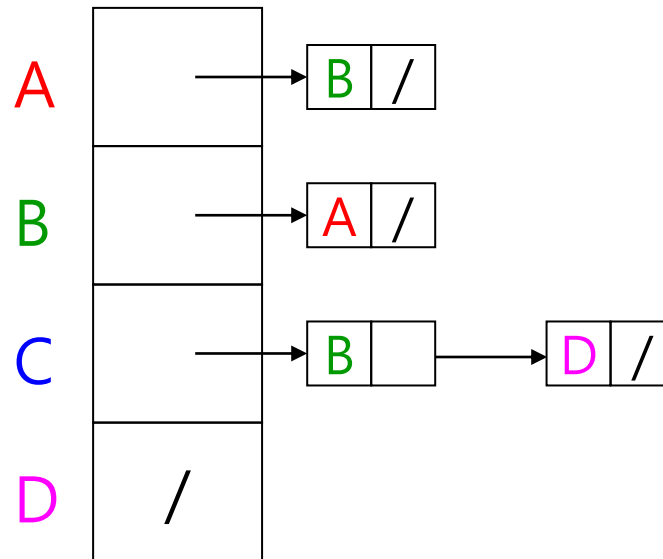
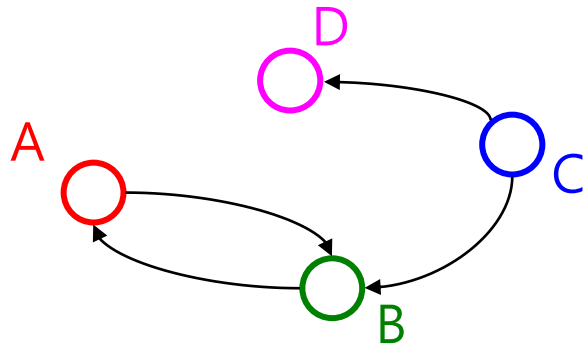


Weighted



Folded/Projected Bipartite Graphs

➤ Adjacency List



- Running time to:
 - Get a vertex's out-edges: $O(d)$ where d is out-degree of vertex
 - Get a vertex's in-edges: $O(|E|)$ (could keep a second adjacency list for this!)
 - Decide if some edge exists: $O(d)$ where d is out-degree of source
 - Insert an edge: $O(1)$ (unless you need to check if it's already there)
 - Delete an edge: $O(d)$ where d is out-degree of source
- Space requirements: $O(|V|+|E|)$
- Best for sparse or dense graphs? **sparse**

- Knowing the network structure, we can calculate various useful quantities or measures that capture features of network topology
- Centrality measures represent the most important nodes in graphs:
 - The most influential person in a social network.
 - The most critical nodes in a infrastructure.
 - The highest spreaders of disease.
- Several common measurements:
 - Degree centrality
 - Betweenness centrality
 - Closeness centrality
 - Eigenvector centrality
 - PageRank

- Using Freeman's general formula for centralization (which ranges from 0 to 1):

$$C_D(G) = \frac{\sum_{i=1}^n [C_D(v^*) - C_D(v_i)]}{(n-1)(n-2)},$$

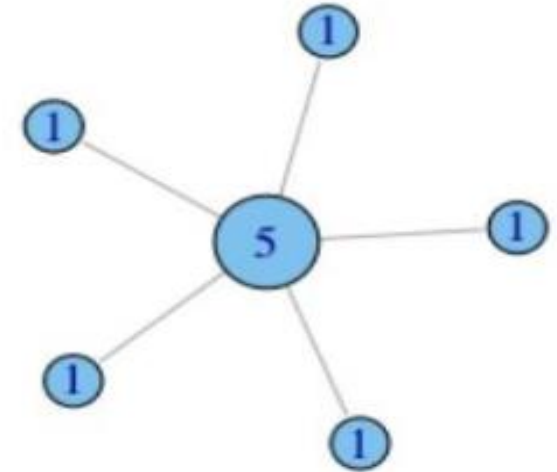
where :

v^* : the node with the highest degree in G



$$C_D = 0.167$$

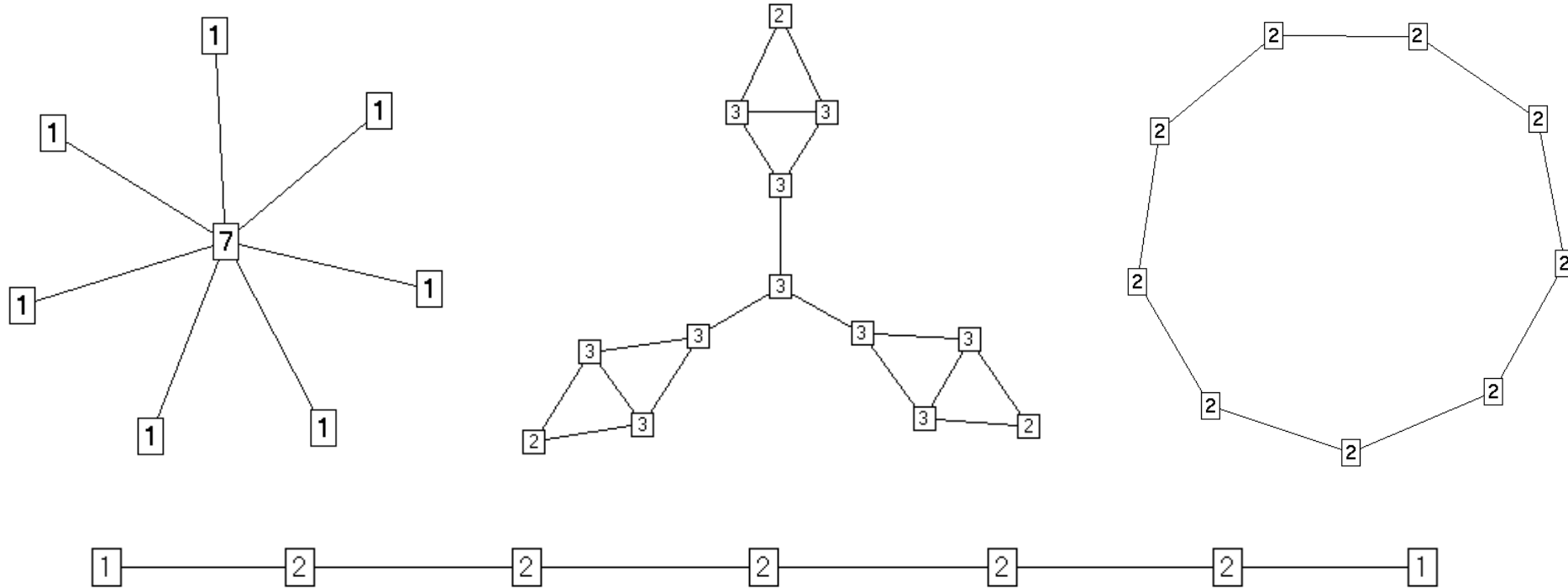
$$C_D(G) = \frac{(2-1) + (2-0) + \dots + (2-1)}{(5-1)(5-2)} = 0.167$$



$$C_D = 1.0$$

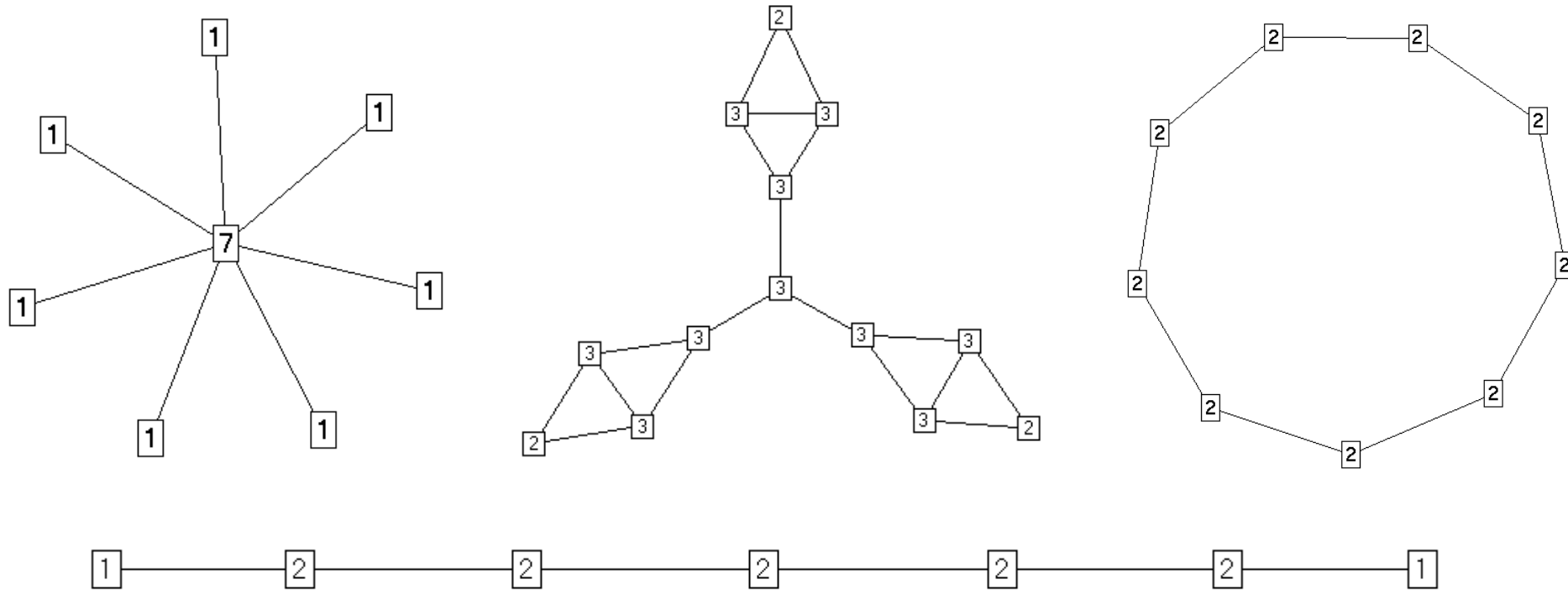
$$C_D(G) = \frac{(5-1) + \dots + (5-1)}{(6-1)(6-2)} = \frac{20}{20} = 1$$

- The most intuitive notion of centrality focuses on degree:
 - The actor with the most ties is the most important:



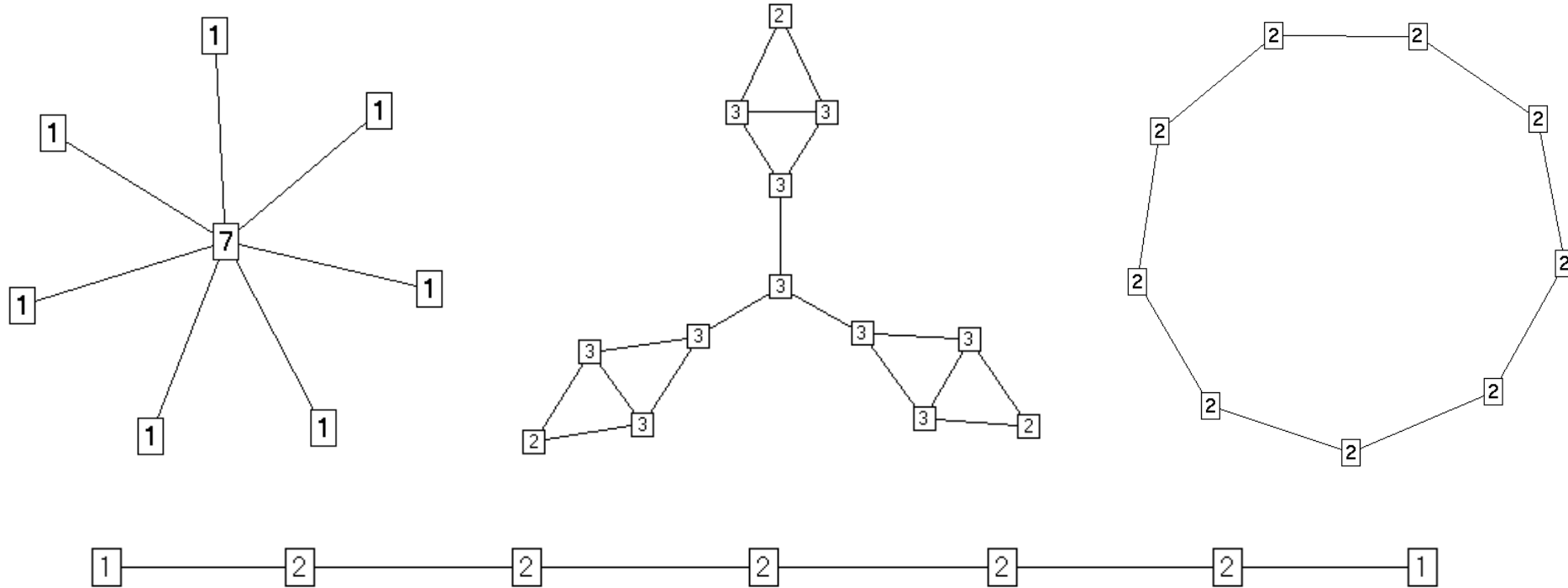
$$C_D(v_i) = d(v_i) = \sum_{j=1}^n A_{ij}$$

- The most intuitive notion of centrality focuses on degree:
 - The actor with the most ties is the most important:



$$C_D(v_i) = d(v_i) = \sum_{j=1}^n A_{ij}$$

- The most intuitive notion of centrality focuses on degree:
 - The actor with the most ties is the most important:



$$C_D(v_i) = d(v_i) = \sum_{j=1}^n A_{ij}$$

- Betweenness centrality of node v_i :

$$B(v_i) = \sum_{v_j, v_k \in G} \left| SPD_{v_j \rightarrow v_k}(v_i) \right|$$

The number of shortest paths between v_j and v_k that pass through the vertex v_i

- Usually normalized by:

$$\bar{B}(v_i) = B(v_i) / [(n-1)(n-2) / 2]$$

No. pairs of nodes excluding the node itself

- The closeness is defined so that if a vertex is close to every other vertex, then the value is larger than if the vertex is not close to everything else.

SPD: The number of nodes in the shortest path between 2 nodes

$$C(v_i) = \left[\sum_{j=1}^n \left| \text{SPD}(v_i, v_j) \right| \right]^{-1}$$

- Normalized Closeness Centrality:


$$\bar{C}(v_i) = C(v_i) / (n-1)$$

- Define the centrality x'_i of i recursively in terms of the centrality of its neighbors:

$$x'_i = \sum_{v_j \in N(v_i)} A_{ij} x_j \quad \text{with the initial node centrality } x_j = 1, \forall j$$

- That is equivalent to:

$$x_i(t) = \sum_{v_j \in N(v_i)} A_{ij} x_j(t-1) \quad \text{with the centrality at time } t=0 \text{ being } x_j(0) = 1, \forall j$$



The centrality of nodes x_i and x_j at time t and $(t-1)$, respectively.

- Katz centrality computes the centrality for a node based on the centrality of its neighbours. It is a generalization of the eigenvector centrality.
- The Katz centrality for node v_i is:

$$x_i = \alpha \sum_j A_{ij} x_j + \beta,$$

where:

α is a constant called damping factor, and β is a bias constant,

A is the adjacency matrix.

- When $\alpha = 1 / \lambda_{\max}$, $\beta = 0$, Katz centrality is the same as eigenvector centrality

PageRank is a numeric value that represents how important a page is on the web.

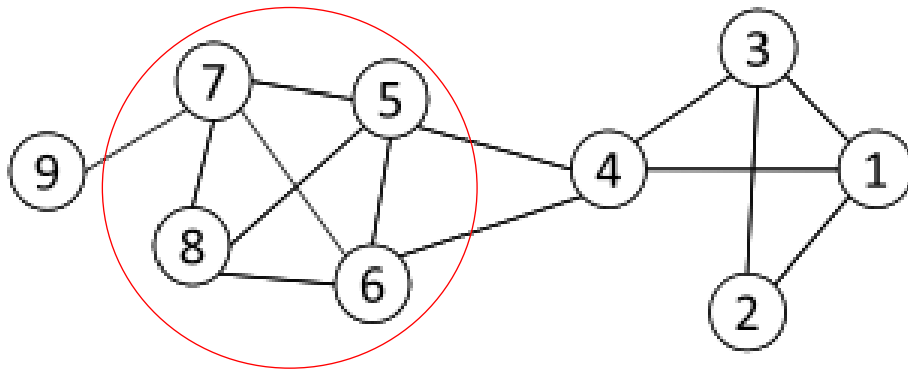
- Webpage importance
 - One page links to another page = A vote for the other page A link from page A to page B is a vote on A to B.
 - If page A is more important itself, then the vote of A to B should carry more weight.
 - More votes = More important the page must be
- How can we model this importance?

- Criteria vary depending on the tasks.
- Roughly, community detection methods can be divided into 4 categories (not exclusive):
 - 1. Node-Centric Community
 - Each node in a group satisfies certain properties
 - 2. Group-Centric Community
 - Consider the connections within a group as a whole. The group has to satisfy certain properties without zooming into node-level
 - 3. Network-Centric Community
 - Partition the whole network into several disjoint sets
 - 4. Hierarchy-Centric Community
 - Construct a hierarchical structure of communities

- Nodes satisfy different properties
 - Complete Mutuality
 - cliques
 - Reachability of members
 - k-clique, k-clan, k-club
 - Nodal degrees
 - k-plex, k-core
 - Relative frequency of Within-Outside Ties
 - LS sets, Lambda sets
- Commonly used in traditional social network analysis

We discuss some representative ones

- Clique: a maximum complete subgraph in which all nodes are adjacent to each other

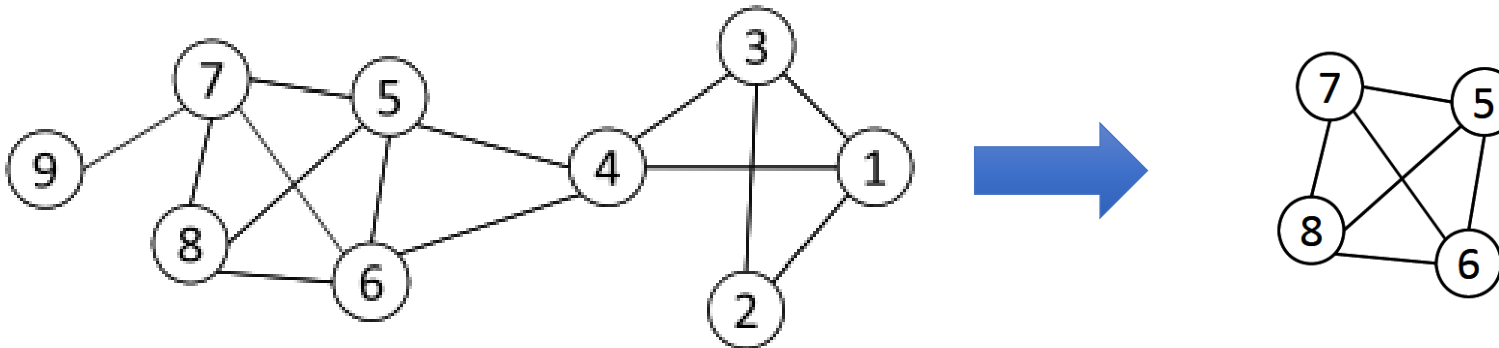


Nodes 5, 6, 7 and 8 form a clique

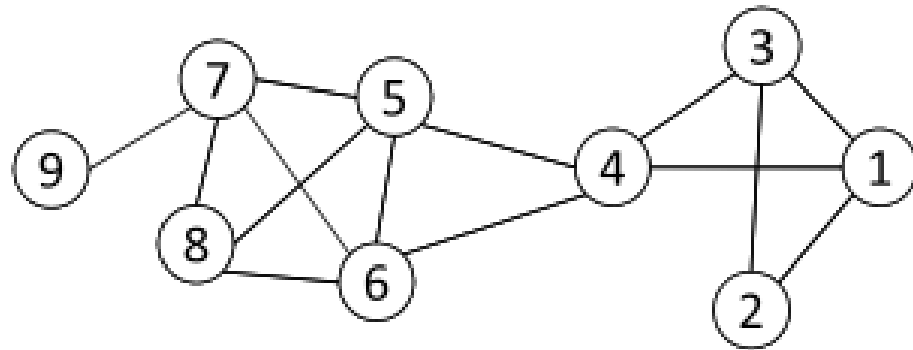
- NP-hard to find the maximum clique in a network
- Straightforward implementation to find cliques is very expensive in time complexity

- In a clique of size k , each node maintains degree $\geq k-1$
- Nodes with degree $< k-1$ will not be included in the maximum clique
- Recursively apply the following pruning procedure:
 - Sample a sub-network from the given network, and find a clique in the sub-network, say, by a greedy approach
 - Suppose the clique above is size k , in order to find out a larger clique, all nodes with degree $\leq k-1$ should be removed.
- Repeat until the network is small enough
- Many nodes will be pruned as social media networks follow a power law distribution for node degrees

- Suppose we sample a sub-network with nodes {1-5} and find a clique {1, 2, 3} of size 3
- In order to find a clique >3 , remove all nodes with degree $\leq 3 - 1 = 2$
 - Remove nodes 2 and 9
 - Remove nodes 1 and 3
 - Remove node 4



- Clique is a very strict definition, unstable
- Normally use cliques as a core or a seed to find larger communities
- CPM is such a method to find overlapping communities
- Input
 - A parameter k , and a network
- Procedure
 - Find out all cliques of size k in a given network
 - Construct a clique graph. Two cliques are adjacent if they share $k-1$ nodes
 - Each connected components in the clique graph form a community

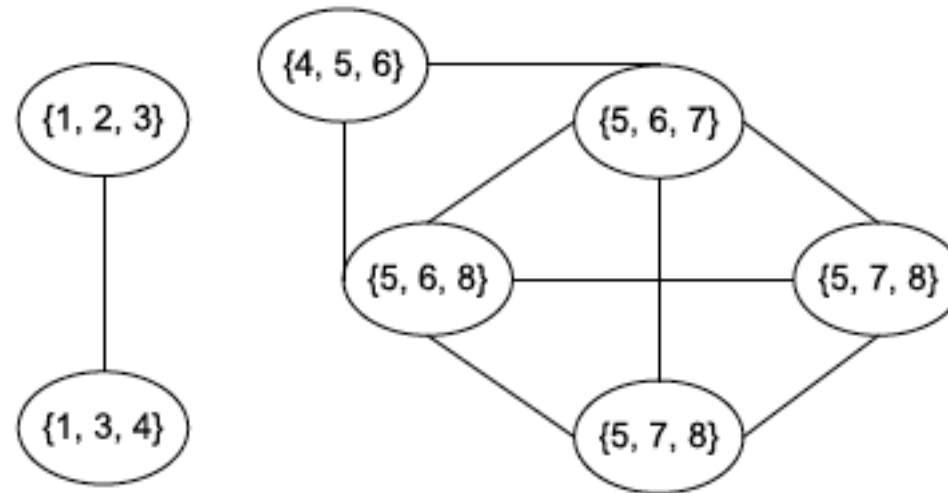


Cliques of size 3:

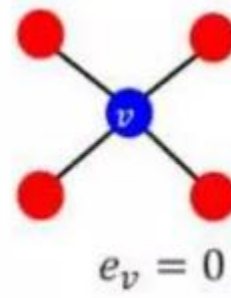
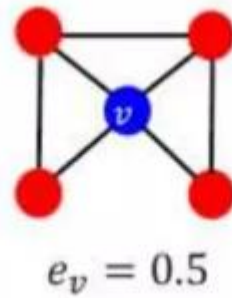
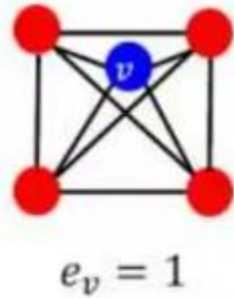
$\{1, 2, 3\}$, $\{1, 3, 4\}$, $\{4, 5, 6\}$,
 $\{5, 6, 7\}$, $\{5, 6, 8\}$, $\{5, 7, 8\}$,
 $\{6, 7, 8\}$

Communities:

$\{1, 2, 3, 4\}$
 $\{4, 5, 6, 7, 8\}$

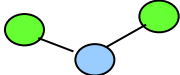
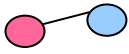
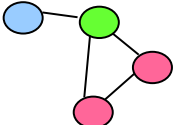


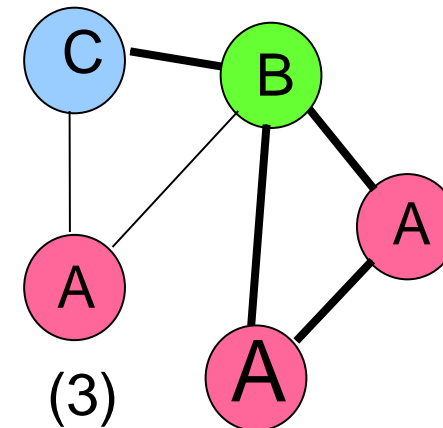
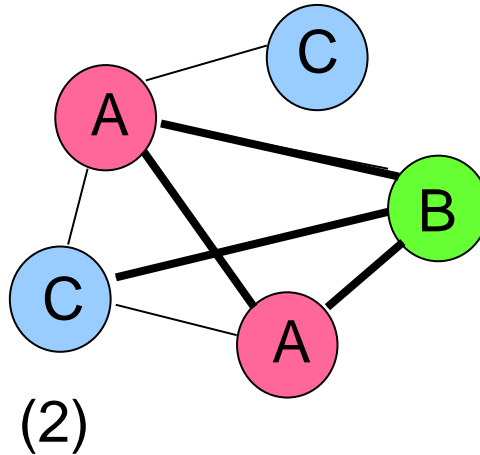
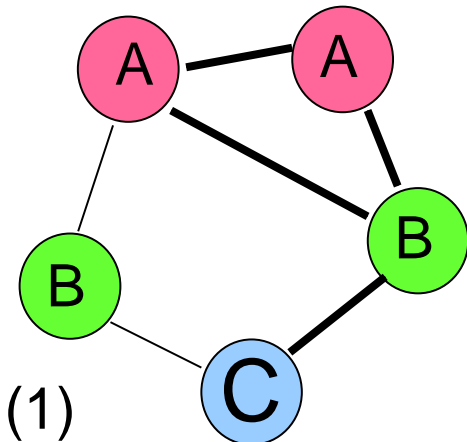
- Degree of nodes
- Clustering Coefficient
 - Measures how connected neighboring nodes are
 - E.g., The number of edges among neighboring nodes



➤ Frequent subgraphs

- A (sub)graph is frequent if its support (occurrence frequency) in a given dataset is no less than a minimum support threshold
- Suppose $t = 2$, the frequent subgraphs are (only edge labels)
 - a, b, c
 - a-a, a-c, b-c, c-c
 - a-c-a ...

Support	1	3	3
Subgraph			



- Graph kernels based on **bags of patterns**:
 - Extraction of a set of patterns from graphs
 - Comparison between patterns
 - Comparison between bags of patterns

Deg1: ● Deg2: ● Deg3: ●

$$\phi(\text{graph1}) = \phi(\text{graph2}) \Rightarrow$$

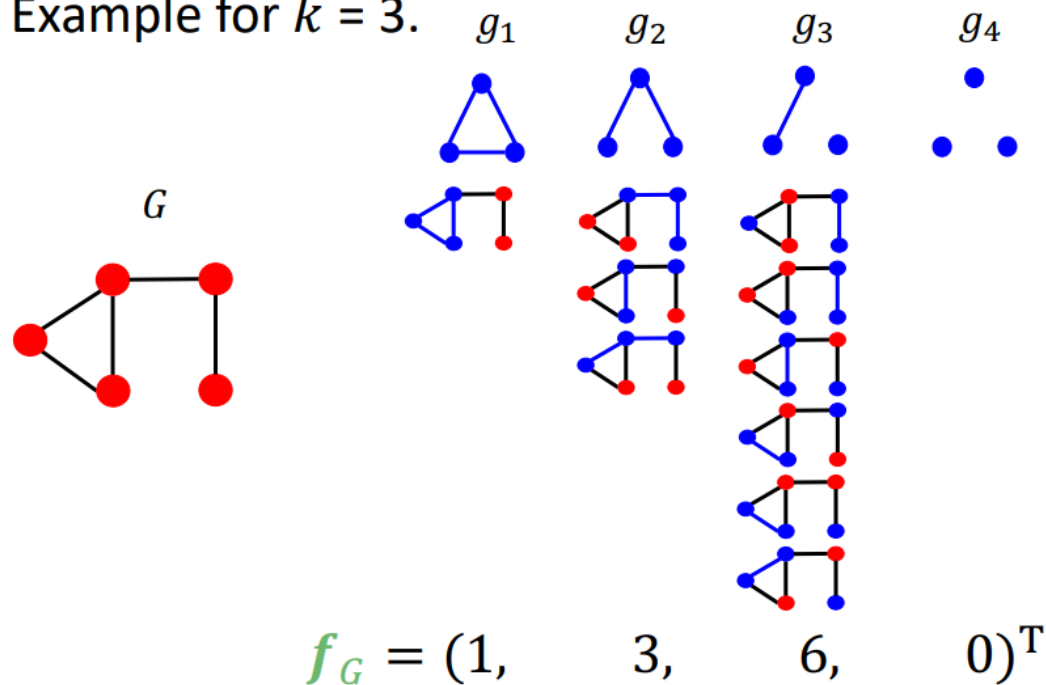
$$\phi(\text{graph1}) = \text{count}(\text{graph3}) = [1, 2, 1]$$

$$\phi(\text{graph2}) = \text{count}(\text{graph4}) = [0, 2, 2]$$

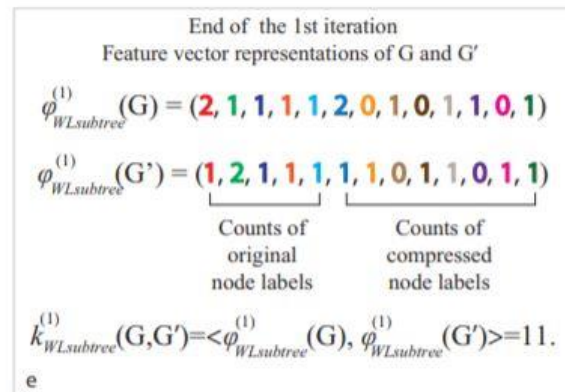
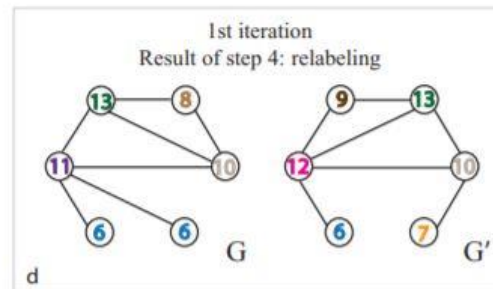
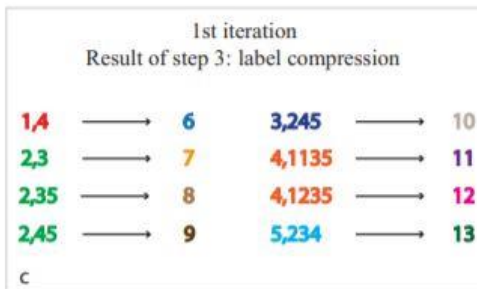
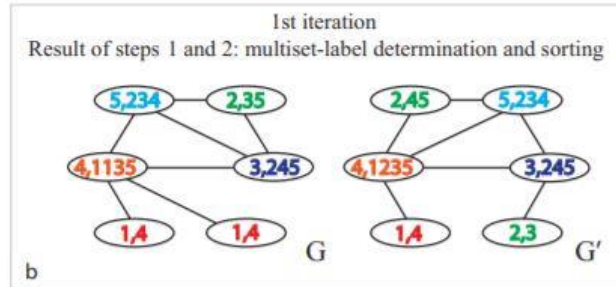
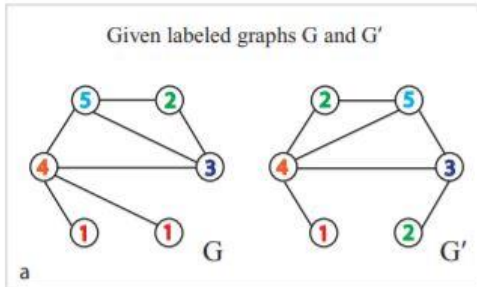
Obta for di

- Graphlet Kernel (B., Petri, et al., MLG 2007)
- Count subgraphs of limited size 3:

■ Example for $k = 3$.



➤ Weisfeiler-Lehman Isomorphism Testing:



Algorithm 1: WL-1 algorithm (Weisfeiler & Lehmann, 1968)

Input: Initial node coloring $(h_1^{(0)}, h_2^{(0)}, \dots, h_N^{(0)})$

Output: Final node coloring $(h_1^{(T)}, h_2^{(T)}, \dots, h_N^{(T)})$

$t \leftarrow 0$;

repeat

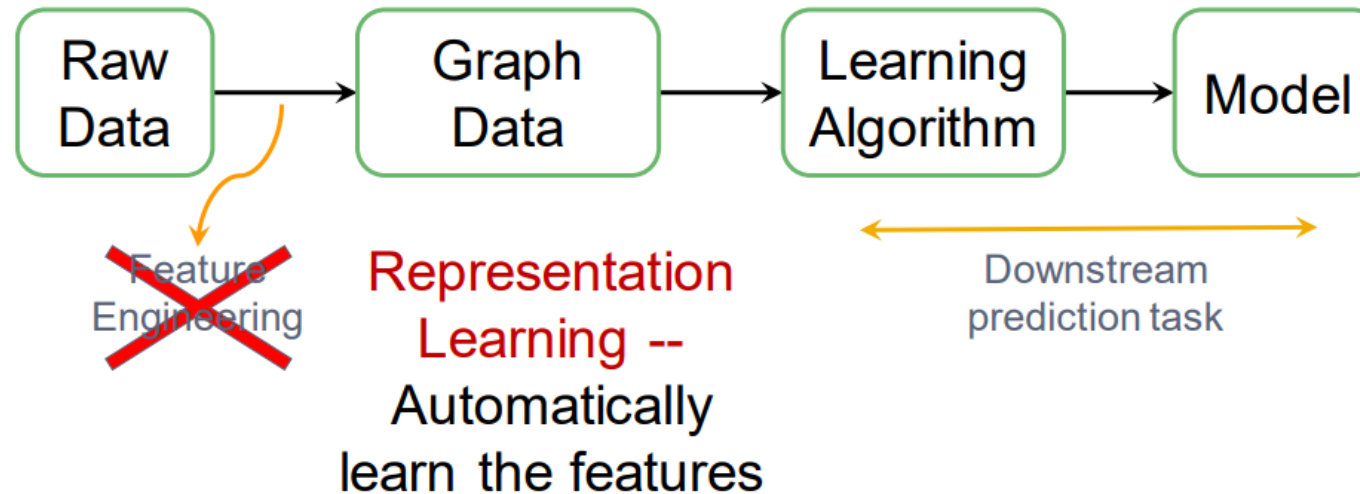
for $v_i \in \mathcal{V}$ **do**

$h_i^{(t+1)} \leftarrow \text{hash} \left(\sum_{j \in \mathcal{N}_i} h_j^{(t)} \right)$;

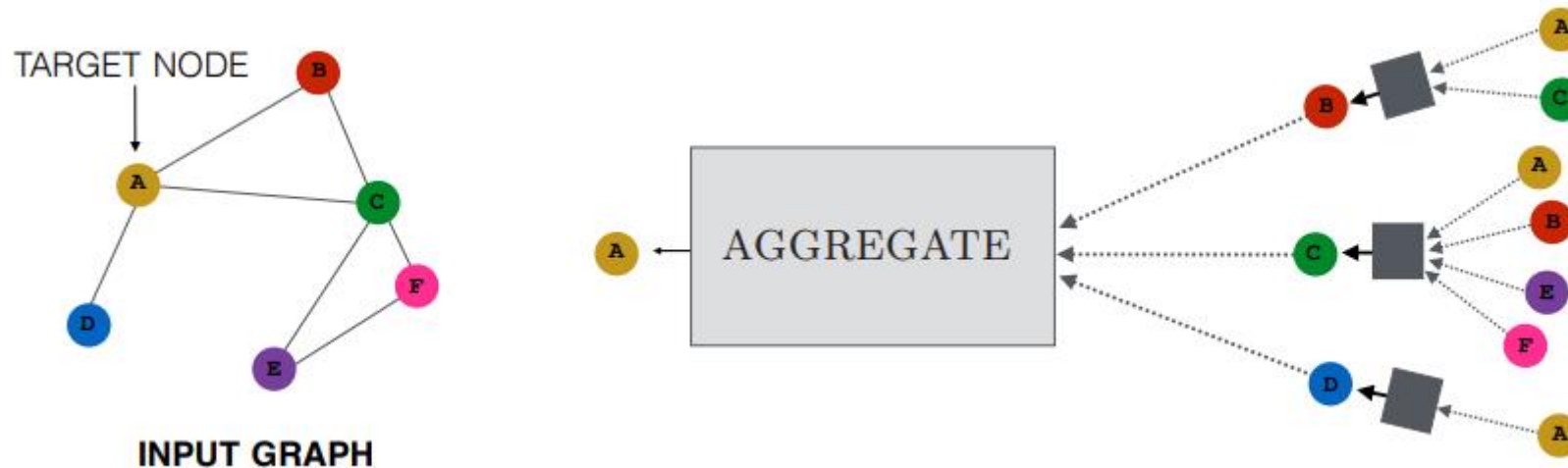
$t \leftarrow t + 1$;

until *stable node coloring is reached*;

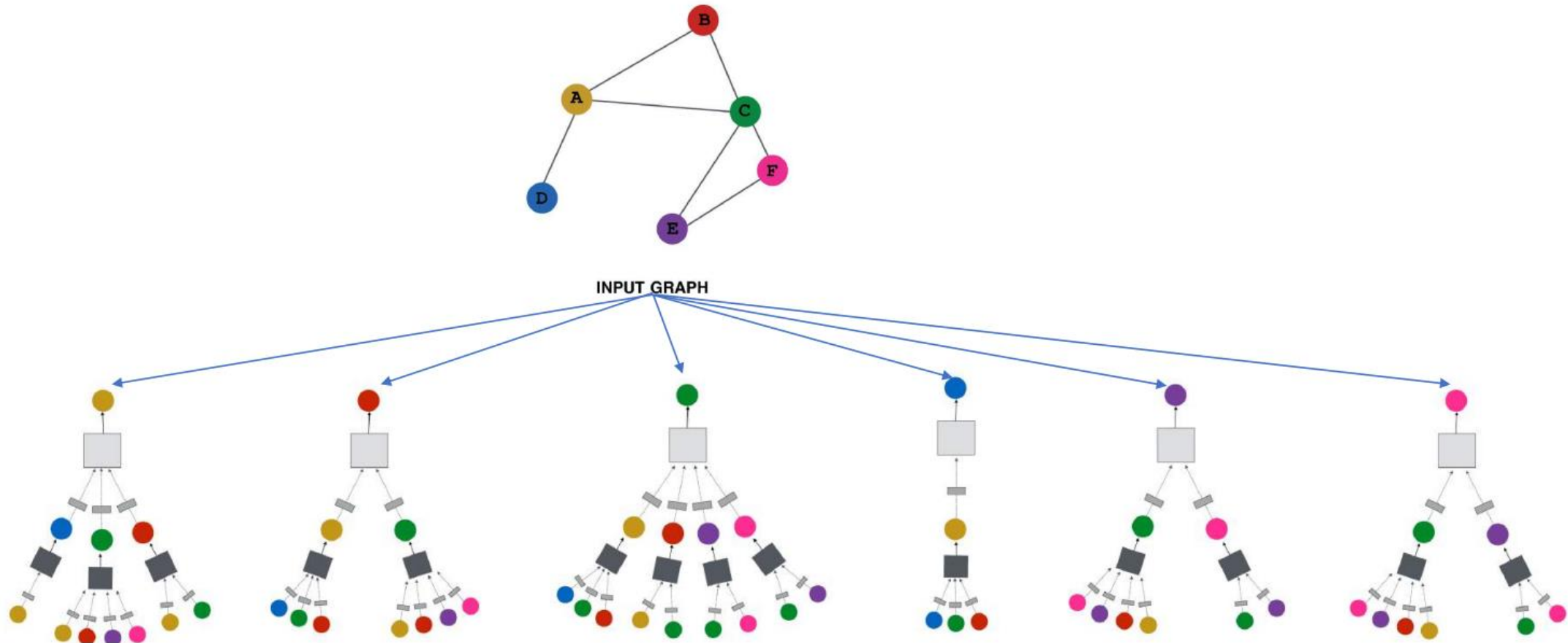
- (Supervised) Machine Learning Lifecycle
- This feature, that feature. Every single time!



- The idea is to generate node embeddings based on local neighborhoods
- The intuition is nodes aggregate information from their neighbors using neural networks.



- Network neighborhood defines a computation graph





네트워크 과학연구실
NETWORK SCIENCE LAB



가톨릭대학교
THE CATHOLIC UNIVERSITY OF KOREA

