問題解析:給定一個大小不超過8*8的填字表,並給出可填入的單字,找出可能的組合。

作法:用 dfs 跑主要搜尋,frontier 紀錄發現的點(stack), explored 紀錄跑到的點(stack),每個 node 紀錄該填的位置及填入的字,每個 node 都有自己的 domain。

(所有作法都先用長度刪除 domain, 因此 domain 内只有長度符合的字, 不會有超出格子的問題)

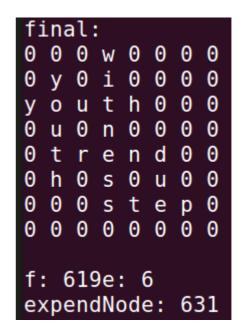
(且 dfs 填字完後, 若發現有無法填的, 即退回上個點, forward checking, 但因為一開始沒發現,不小心寫進去 dfs 了, 所以它關不掉 QQ)

實驗過程:

1. 初始暴力解法:

任意選擇想填入哪個字及内容(在 code 裡是照順序),填入字後判斷剩餘字的 domain(後填的字不可以改動前面填過的字),再任意填,再判斷,直到找出答案為止。(0 是不能填的地方) 實驗結果:

| final: | | | | | | | | | | |
|-----------------|---|---|---|---|---|---|---|--|--|--|
| У | 0 | u | r | 0 | 0 | 0 | 0 | | | |
| 0 | 0 | r | 0 | 0 | 0 | 0 | 0 | | | |
| u | 0 | g | 0 | 0 | 0 | 0 | 0 | | | |
| t | 0 | е | 0 | 0 | 0 | 0 | 0 | | | |
| h | 0 | 0 | 0 | 0 | 0 | 0 | 0 | | | |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | | | |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | | | |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | | | |
| | | | | | | | | | | |
| f: 506e: 5 | | | | | | | | | | |
| expendNode: 514 | | | | | | | | | | |
| | | | | | | | | | | |



```
0 0 y o u r 0 0
 0 o 0 0 a 0
             0
cousin0
             0
e 0 t 0 0 g 0
            0
 0 h u g
         e 0
             0
0 0 0 0 0 0
             0
0 0 0 0 0 0 0
 000000
f: 585e: 6
expendNode: 3204
```

```
w r o n g 0 u 0
o 0 0 0 r u s h
r 0 0 0 a 0 0 a
r e s i d e n t
i 0 u 0 u 0 0 e
e x p l a i n 0
d 0 e 0 t 0 o r
0 u r g e 0 0 0

f: 963e: 12
expendNode: 35854
```

説明:

f 是 frontier 最後剩餘在 stack 的數量

e 是 explored 最後剩餘在 stack 的數量 expendNode 是産生的 node 總數

2. BackTracking Search(逆序)

使用了 Minimum remaining values (MRV) heuristic, 在每次産生新 node 後, 判斷各 doamin 大小, 從剩最少可填的字開始填(但填入哪個字就隨意), 直到找出答案。 實驗結果:

piano0w 0 0 0 v i e w o 0 0 0 e 0 0 i i r h a m a n 0 p 0 1 0 0 u p t 0 ро r 0 0 0 0 k 0 0 е 105e: 12 expendNode: 171

説明:由結果可看出新增的 node 數量比暴力解少許多,且最後一筆測資也明顯比暴力解快, 大約在幾百個 node 之內可以解決。

但若情況是 worst case, 也就是每個 domain 的大小都一樣的時候, 這個方法就會跟暴力解一樣, 需要其他 heuristic 來縮減 domain 才可以優化。像是 degree heuristic, 把關聯最多的字先填, 就可以減少更多不必要的搜尋。

3. BackTracking Search(正序)

將原本從 z~a 的順序改成從 a~z, 驗證看看平均數量 實驗結果:

| final: | | | | | | | | | | |
|-----------------|---|---|---|---|---|---|---|--|--|--|
| a | W | а | У | 0 | 0 | 0 | 0 | | | |
| h | 0 | W | 0 | 0 | 0 | 0 | 0 | | | |
| е | 0 | а | d | 0 | 0 | 0 | 0 | | | |
| а | n | У | 0 | 0 | 0 | 0 | 0 | | | |
| d | 0 | 0 | 0 | 0 | 0 | 0 | 0 | | | |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | | | |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | | | |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | | | |
| | | | | | | | | | | |
| f: 134e: 5 | | | | | | | | | | |
| expendNode: 139 | | | | | | | | | | |

```
final:
0 0 0 0 0 0 0 0 0
0 a 0 v 0 0 0 0
a s s e t 0 0 0
0 i 0 r 0 0 0 0
0 d r a m a 0 0
0 e 0 l 0 c 0 0
0 0 0 l a t e 0
0 0 0 0 0 0 0 0
f: 183e: 6
expendNode: 192
```

b a s i c 0 a 0 a 0 a 0 0 0 a i d e 1 0 0 0 t 0 0 a a a c a d e m i c n 0 d 0 g 0 0 h c o m f o r t 0 e 0 i 0 r 0 o f 0 s t a y 0 0 0 f c 62e: 12 expendNode: 80

説明:由實驗結果可看出 expend 出來的 node 數量差距不大,且整體字母比較前面(不過最後一種情況應該是巧合,剛好造成地圖較大,卻比較快找到的結果)

4. BackTracking Search

再額外新增 degree heuristic, 當要選擇填哪個位置時, 若遇到剩餘 domain 大小一樣情況, 根據當前地圖, 選擇影響最大的位置, 例如 domain 都是 25 的話, 選擇有三個交叉點的填而不選兩個。

實驗結果:

```
final:
s h a r p 0 a 0
a 0 0 0 l a s t
t 0 0 0 e 0 0 e
i n s t a n c e
s 0 e 0 s 0 0 n
f o r m u l a 0
y 0 v 0 r 0 d o
0 g e n e 0 0 0

f: 67e: 12
expendNode: 101
```

説明:上圖是第四種跑出來的結果,其餘三種皆跟上一種一樣,而且這個反而跑更多 node...,有可能是因為照順序跑剛好跑到有解,而增加 degree heuristic 雖然平均起來較快,但也因此避開了巧合解,所以 node 數量變多

其他想法:

- 1. 如果使用 AC-3 演算法的話, 就可以直接從 binary constraint 下手, 不用一個一個嘗試, expend 的 node 數量感覺也會明顯下降, 而且甚至不會避掉巧合(畢竟巧合也在 binary constraint 下), 不過實做上我沒什麼想法…
- 2. 至於跑出全部的結果我覺得是有可能的, 只要每輸出一次結果, 就把剛填過的最後一個字從該行 (列)的 domain 拿掉, 再繼續跑應該就可以了>< (等於是假裝這個結果不行繼續跑)
- 3. 每次都隨機輸出的話我想到的只有用上面 MRV heuristic 的方式, 在 push node 到 frontier 上時, 亂數 push, 這樣就有可能每次都不一樣

總結:

我覺得填字問題,對變數的限制數最高只有 2,用 AC-3 最為適合,因為 binary constraint 是這個問題最麻煩的事,且數量跟其他 constraint 比起來也少了許多。(但感覺好難做不出來)

以我實驗的幾種方法來看,在單字數量到 3000 的情況下, MRV+forward checking 就可以達到不錯的效果,但當數量過於龐大的時候,加入 degree heuristic 穩定 expendNode 的數量似乎比較好。

Code:

```
#includesbits/stdc++.h>
#includesbits/stdc-+.h
#includesbits/stdc-+.h
#includesbits/stdc-+.h
#includesbits/stdc-+.h
#includesbits/stdc-+.h
#includesbits/stdc-+.h
#includ
```

```
void loadMap(vector < vector <int> > &position, vector < vector <char> > &map, int depth){
    for(int i=0; i<depth; i++){
        if(position[i][3] == 0){
            for(int j=0; j<sposition[i][2]; j++){
                map[position[i][1]][position[i][0]] = '1';
        }
    }
} }

// set the initial map, '1' for blank that can fill word, '0' for the place can't fill word

void loadMap2(vector < vector <int> > &position, vector < vector <char> > &map, int depth){
    for(int i=0; i<depth; i++){
        if [position[i][3] == 0){
            for(int j=0; j<position[i][2]; j++){
                map[position[i][1]][position[i][0]]++;
            }
        }
} }else{
    for(int j=0; j<position[i][2]; j++){
            map[position[i][1]]; j++){
            map[position[i][2]; j++){
```

```
inf findMin(vector < vector <int> > &position, vector < int > &domainSize, vector < vector <char> > &map,

vector <br/>
vector <br/>
**vector <br/>
**v
```

```
void dfsSearch(vector < vector < int> > &position, vector < vint > & domain, stack < node > & explored, stack < node > & frontier, vector < vector < char> > & map, vector < int > & domainSize, vector < book > visited, let on the vector < vector < vector < vector < vint > & domainSize, vector < vector <
```

```
visited[index] = true;
// use for dfs to remember which variable has ran

newMode = 0;
// use to memorize how many node will expend if this current node push

for (int i=domainSize[index]-1; i>=0; i--){
    node temp;
    temp.variable = index;
    temp.value = domain[index][i];
    frontier.push(temp);
    newMode++;
    expendNode++;

// expend node

white(!frontier.empty() && explored.size()!=depth){
    if(newNode = 0) {
        explored.pop();
        loadMap(position, map, depth);
        stack < node > temp = explored;
        white(!temp.empty()){
        inf tempindex = temp.top().variable;
        if(position[tempindex][3] == 0){
            for(int j=0; sposition[tempindex][0]+j] = temp.top().value[j];
        }
    }
} else{
    for(int j=0; sposition[tempindex][2]; j++){
        map[position[tempindex][1]|position[tempindex][0]] = temp.top().value[j];
    }
} else{
    for(int j=0; sposition[tempindex][2]; j++){
        map[position[tempindex][1]+j][position[tempindex][0]] = temp.top().value[j];
    }
} temp.pop();
} return;
} // newNode == 0 means the frontier it expend is all pop, i need a new explore and return to it's domain
```

```
int draw = frontier.top().variable;
if (position(fraw)[3] = 0}{
    for(int j=0; j<position(draw)[2]; j++){
        map[position[draw][2]; j++){
        map[position[draw][2]; j++){
        map[position[draw][2]; j++){
        map[position[draw][1]+j][position[draw][0]] = frontier.top().value[j];
    }
}

vector { int j=0; j<position[draw][1]+j][position[draw][0]] = frontier.top().value[j];
}

//draw the word on the map

vector < vector < string> > newDomain(depth);
vector < int > newDomainslze(depth);
//rew domain and domainslze for dfs

for(int i=0; i<depth; i++){
            newDomainslze = domainslze;

// newDomainslize = domainslze;

setConstraint(position, newDomain, newDomainslze, map, visited, depth);
//rebuilt each domain

makeSize(newDomain, newDomainslze, depth); //rebuilt each size

explored.push(frontier.top());
rankOde++;
//let the word in the explored

frontier.pop();
//let the word out of frontier
newNode--;

dfsSearch(position, newDomain, explored, frontier, map, newDomainSize, visited, binaryCons, dictionary, depth, newNode);
}

// setting dfsSearch(position, newDomain, explored, frontier, map, newDomainSize, visited, binaryCons, dictionary, depth, newNode);
}
</pre>
```

```
int main(){
  vector < vector < int> > position;
  vector < vector < vector < int> > position;
  vector < 
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                  vector <string> dictionary;
                 stack < node > explored;
//data structure to store the node we explored, in detail, which blank we fill in what word
                  stack < node > frontier;
//data_structure to store the node we can explored, in detail, which will be the next to refill
                  ifstream fin("puzzle.txt", ios::in);
string str;
                  loadWord(dictionary);
                 getline(fin,str);
getline(fin,str);
getline(fin,str);
getline(fin,str);
                 position.clear();
if(!explored.empty()){
                         explored.pop();
                 }
if(!frontier.empty()){
frontier.pop();
                  }
// initialize
                  vector < vector <char> > map(8, vector<char> (8, '0'));
//to store current drawing on the map, '1' for blank, '0' can't fill
                  vector < vector < char> > map2(8, vector< char> (8, '0'));
                    int i=0;
int depth=0;
                  //how many variable to 
istringstream iss(str);
vector <int> tmp;
char temp;
while(iss >> temp){
   if(i<4){
      if(temp == 'A'){
      tmp.push_back(0);
   }
}</pre>
                                    i++;

else if(temp == 'D'){

tmp.push_back(1);
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                             i++;
}else{
                                    tmp.push_back(temp-'0');
i++;
                        }
}
if(i==4){
  position.push_back(tmp);
  tmp.clear();
  i = 0;
  depth++;
                    int newNode=0;
                   vector <bool> visited(depth, false);
// to recgonize which variable have filled
                  vector < vector <string> > domain(depth);
                   loadDomain(domain, position, dictionary, visited, depth);
                    vector < int > domainSize(depth, 0);
                   vector < int > domainSize(depth, 0);
                  vector < int > binaryCons(depth, 0);
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                  makeSize(domain, domainSize, depth);
                  loadMap(position, map, depth);
                  loadMap2(position, map2, depth);
                  loadBinaryCons(position, map2, binaryCons, depth);
                  dfsSearch(position, domain, explored, frontier, map, domainSize, visited, binaryCons, dictionary, depth, newNode);
                 cout << "final:" << endl;
for(int i=0; i<8; i++){
   for(int j=0; j<8; j++){
      cout << map[i][j] << " ";</pre>
                        cout << endl;
                  cout << endl;
cout << "f: " << frontier.size() << "e: " << explored.size() << endl;
cout << "expendNode: " << expendNode << endl;
//cout << "ranNode: " << ranNode << endl;</pre>
 359
360
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```