## The problem description is very short:

The equilibrium index of a sequence is an index such that the sum of elements at lower indexes is equal to the sum of elements at higher indexes. For example, in a sequence A:

3 is an equilibrium index, because:

$$A[0]+A[1]+A[2]=A[4]+A[5]+A[6]$$

6 is also an equilibrium index, because:

$$A[0]+A[1]+A[2]+A[3]+A[4]+A[5]=0$$

(The sum of zero elements is zero) 7 is not an equilibrium index - because it is not a valid index of sequence A. If you still have doubts, here is a precise definition: The integer k is an equilibrium index of a sequence A[0],A[1]..,A[n-1] if and only if 0 <= k and sum(A[0..(k-1)]) = sum(A[(k+1)..(n-1)]). Assume the sum of zero elements is equal to zero.

Write a function

```
int equi(int A[], int n)
```

that, given a sequence, returns its equilibrium index (any) or -1 if no equilibrium index exists. Assume that the sequence may be very long.

The problem can be solved by using various approaches, the most common being simply to follow the equilibrium definition:

```
int equi ( int A[], int n ) { 
    int k, m, lsum, rsum; 
    for(k = 0; k < n; ++k) { 
        lsum = 0; rsum = 0; 
        for(m = 0; m < k; ++m) lsum += A[m]; 
        for(m = k + 1; m < n; ++m) rsum += A[m]; 
        if (lsum == rsum) return k; 
    } 
    return -1; 
}
```

Unfortunately, this approach has two disadvantages:

- it fails on large input data sets, since the time complexity is  $O(n^2)$
- it fails on large input values (for example if the input array contains values like MIN/MAX\_INT) due to the arithmetic overflows

The solution analysis will detect such problems in submitted code:



We can fix the first problem with a better algorithm, and the second problem with a better data-type (for example, using long long type instead of int for sum computations). The key observation for better running time is to update the left/right sums in constant time instead of recomputing them from the scratch.

```
int equi(int arr[], int n) {
    if (n==0) return -1;
    long long sum = 0;
    int i;
    for(i=0;i<n;i++) sum+=(long long) arr[i];

long long sum_left = 0;
    for(i=0;i<n;i++) {
        long long sum_right = sum - sum_left - (long long) arr[i];
        if (sum_left == sum_right) return i;
        sum_left += (long long) arr[i];
    }
    return -1;
}</pre>
```

Using this solution, you can obtain a perfect score:

Detected time complexity:  O(n)		
example Test from the task description	0.001 s.	ок
extreme_empty Empty array	0.004 s.	ок
extreme_first	0.001 s.	ОК
extreme_large_numbers Sequence with extremity large numbers testing arithmetic overflow.	0.001 s.	ок
extreme_last	0.001 s.	ок
extreme_single_zero	0.001 s.	ОК
extreme_sum_0 sequence with sum=0	0.001 s.	ок
simple	0.004 s.	ок
single_non_zero	0.004 s.	ок
combinations_of_two multiple runs, all combinations of {-1,0,1}^2	0.001 s.	ок
combinations_of_three multiple runs, all combinations of {-1,0,1}^3	0.004 s.	ок
small_pyramid	0.004 s.	ОК
large_long_sequence_of_ones	0.012 s.	ОК
large_long_sequence_of_minus_ones	0.020 s.	ОК
medium_pyramid	0.012 s.	ОК
large_pyramid Large performance test, O(n^2) solutions should fall.	0.048 s.	ок