

- ★ - Design LRU cache 146
- Design LFU cache 460
- ★ - Copy Linked List with random pointer 138
 - ↳ with extra space
 - ↳ without extra space

① LRU cache {Least Recently Used}

Eg: whatsapp chat ranking \rightarrow latest chat first.

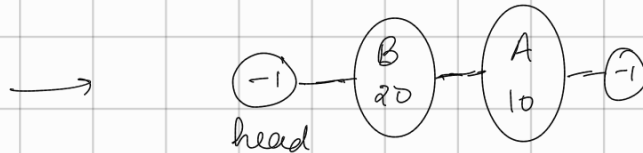
Removal: Can be from any point \rightarrow Array $\times O(n) \rightarrow$ LL ✓

Insertion: Only from head \rightarrow LL $\rightarrow O(1)$



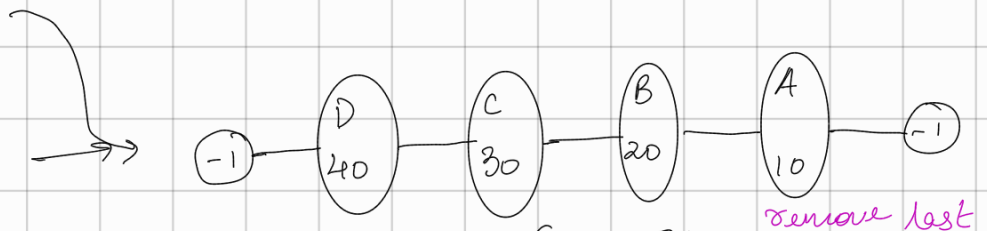
① put (A, 10)

② put (B, 20)

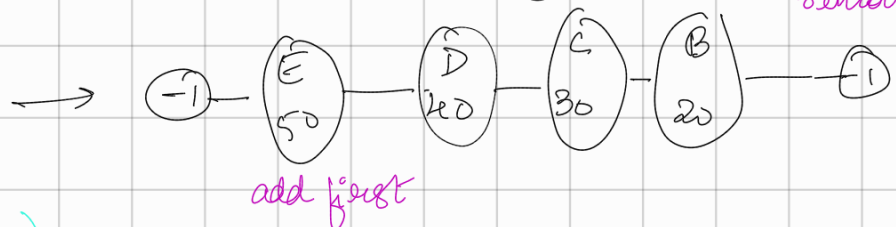


③ put (C, 30)

④ put (D, 40)



⑤ put (E, 50)



(put = addFirst)

⑥ get(A) \Rightarrow -1 (Not in cache) but with LL, gives $O(n)$
 So to improve complexity, use HashMap.
 $\rightarrow O(1)$

⑦ $get(1)$: Process: Check in Hash Map & if present return value.

If we use HashSet, this operation will take $O(\text{capacity})$ because we need to traverse list which is not acceptable for current problem. So use Hash Map.

Once node is found, remove Node & add First the remove Node. \rightarrow this will also take $O(n)$ in SLL

So use DLL (Doubly Linked List)

⑧ $put(D, 50) \rightarrow$ check if D is present in Hash Map $O(1)$ \swarrow
if present, remove from its current position and add Front (D, 50). $\rightarrow \underline{O(1)}$

⑨ $put(E, 60) \rightarrow$ check if E is present in Hash Map $O(1)$,
if present, remove from its current position and add Front (E, 60). $\rightarrow \underline{O(1)}$

⑩ $put(F, 60) \rightarrow$ check if F is present in Hash Map $O(1)$
if not present, then $removeLast() \rightarrow O(1)$
and $addFirst(F, 60) \rightarrow O(1)$

To remove last node, No need to update tail node,
just do $curr.prev.next = curr.next$
 $curr.next.prev = curr.prev$

Hashmap \rightarrow Insert	\rightarrow avg case $\rightarrow O(1)$
	\rightarrow worst case $\rightarrow O(n)$

UNORDERED MAP

Page / Frame Not found in cache \rightarrow Page Fault or
Cache miss
" FOUND " \rightarrow Cache hit

* DLL Operations

Remove first

curr = head.next

curr.prev.next = curr.next

curr.next.prev = curr.prev

Remove last

curr = tail.prev

curr.prev.next = curr.next

curr.next.prev = curr.prev

Remove at

curr.prev.next = curr.next

curr.next.prev = curr.prev

Add first

Node curr = new Node(val);

curr.prev = head;

curr.next = head.next;

head.next = curr;

curr.next.prev = curr;

Add last

Node curr = new Node(val)

curr.next = tail

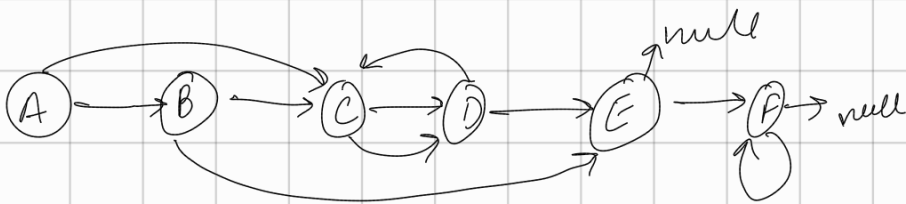
curr.prev = tail.prev

curr.next.prev = curr

curr.prev.next = curr

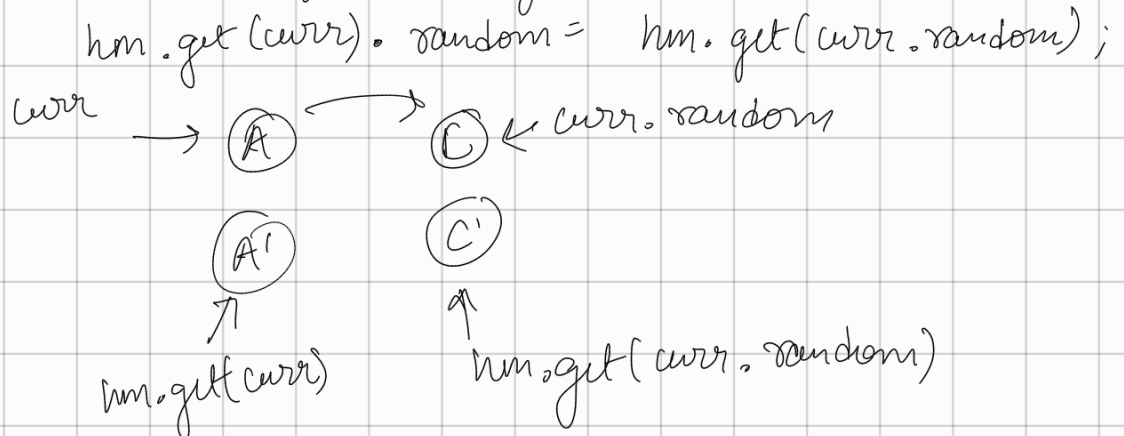
② Copy Linked list with random pointer Deep Copy

If no random pointer: Simple traversal & add next $\rightarrow O(n)$



Use hashmap $\langle \text{Node}, \text{Node} \rangle$
old node \uparrow cloned node \uparrow

for each node after cloning without random pointers



Constant Extra Space:

Insert duplicate node just after it's original one