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Semidefinite Programming for Geometric Problems in Computer Vision

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I declare that I have work out the presented thesis independently and that I have listed all information sources used in accordance with the Methodical Guidelines about Maintaining Ethical Principles for Writing Academic Theses.					
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Abstract

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List of Symbols and Abbreviations

 $\operatorname{cl}(S)$ Closure of the set S. dom f Domain of the function f. Dom f $\operatorname{cl}(\operatorname{dom} f)$.

1. Introduction

[1]

2. Semidefinite programming

2.1. Preliminaries on semidefinite programs

We introduce here some notation and preliminaries about symmetric matrices and semidefinite programs. We will introduce further notation and preliminaries later on in the text when needed.

At the beginning, let us denote the inner product for two vectors $x, y \in \mathbb{R}^n$

$$\langle x, y \rangle = \sum_{i=1}^{n} x_i y_i \tag{2.1}$$

and the Frobenius inner product for two matrices $X, Y \in \mathbb{R}^{n \times m}$.

$$\langle X, Y \rangle = \sum_{i=1}^{n} \sum_{j=1}^{m} X_{ij} Y_{ij}$$
 (2.2)

2.1.1. Symmetric matrices

Let Sym_n denotes the space of $n \times n$ real symmetric matrices.

For matrix $M \in \operatorname{Sym}_n$, the notation $M \succeq 0$ means that M is positive semidefinite. $M \succeq 0$ if and only if any of the following equivalent properties holds.

- 1. $x^{\top} M x > 0$ for all $x \in \mathbb{R}^n$.
- 2. All eigenvalues of M are nonnegative.

For matrix $M \in \operatorname{Sym}_n$, the notation $M \succ 0$ means that M is positive definite. $M \succ 0$ if and only if any of the following equivalent properties holds.

- 1. $M \succeq 0$ and rank M = n.
- 2. $x^{\top}Mx > 0$ for all $x \in \mathbb{R}^n$.
- 3. All eigenvalues of M are positive.

2.1.2. Semidefinite programs

The standard (primal) form of a semidefinite program in variable $X \in \operatorname{Sym}_n$ is defined as follows:

$$p^* = \sup_{\substack{X \in \operatorname{Sym}_n \\ \text{s.t.}}} \langle C, X \rangle$$
s.t.
$$\langle A_i, X \rangle = b_i \quad (i = 1, \dots, m)$$

$$X \succeq 0$$
 (2.3)

where $C, A_1, \ldots, A_m \in \operatorname{Sym}_n$ and $b \in \mathbb{R}^m$ are given.

The dual form of the primal form is the following program in variable $y \in \mathbb{R}^m$.

$$d^* = \inf_{y \in \mathbb{R}^m} b^\top y$$
s.t.
$$\sum_{i=1}^m A_i y_i - C \succeq 0$$
(2.4)

2.2. State of the art review

2.3. Theoretical background

2.4. Nesterov's approach

In this section, we will follow Chapter 4 of [2] by Y. Nesterov, which is devoted to convex minimization problems. We will extract from it only the minimum, just to be able to introduce a algorithm for SDP programs solving. We will present some basic definitions and theorems, but we will not prove them. For the proofs look into [2].

2.4.1. Self-concordant functions

Definition 2.1 (Self-concordant function in \mathbb{R}). A closed convex function $f: \mathbb{R} \rightarrow$ \mathbb{R} is self-concordant if there exist a constant $M_f \geq 0$ such that the inequality

$$|f'''(x)| \le M_f f''(x)^{3/2}$$
 (2.5)

holds for all $x \in \text{dom } f$.

For better understanding of self-concordant functions, we provide several examples.

Example 2.1.

1. Linear and convex quadratic functions.

$$f'''(x) = 0 \text{ for all } x \tag{2.6}$$

Linear and convex quadratic functions are self-concordant with constant $M_f = 0$.

2. Negative logarithms.

$$f(x) = -\ln(x) \text{ for } x > 0 \tag{2.7}$$

$$f'(x) = -\frac{1}{x} \tag{2.8}$$

$$f'(x) = -\frac{1}{x}$$

$$f''(x) = \frac{1}{x^2}$$

$$f'''(x) = -\frac{2}{x^3}$$
(2.8)
$$f'''(x) = (2.9)$$

$$f'''(x) = -\frac{2}{x^3} \tag{2.10}$$

$$\frac{|f'''(x)|}{f''(x)^{3/2}} = 2 (2.11)$$

Negative logarithms are self-concordant functions with constant $M_f = 2$.

3. Exponential functions.

$$f(x) = e^x (2.12)$$

$$f''(x) = f'''(x) = e^x (2.13)$$

$$f(x) = e^{x}$$

$$f''(x) = f'''(x) = e^{x}$$

$$\frac{|f'''(x)|}{f''(x)^{3/2}} = e^{-x/2} \to \infty \text{ as } x \to -\infty$$
(2.12)
$$(2.13)$$

Exponential functions are not self-concordant functions.

Definition 2.2 (Self-concordant function in \mathbb{R}^n). A closed convex function f: $\mathbb{R}^n \to \mathbb{R}$ is self-concordant if function

$$g(t) = f(x+tv) (2.15)$$

is self-concordant for all $x \in \text{dom } f$ and all $v \in \mathbb{R}^n$.

Now, let us focus on the main properties of self-concordant functions.

Theorem 2.1. Let functions f_i be self-concordant with constants M_i and let $\alpha_i > 0$, i = 1, 2. Then the function

$$f(x) = \alpha_1 f_1(x) + \alpha_2 f_2(x) \tag{2.16}$$

is self-concordant with constant

$$M_f = \max\left\{\frac{1}{\sqrt{\alpha_1}}M_1, \frac{1}{\sqrt{\alpha_2}}M_2\right\} \tag{2.17}$$

and

$$dom f = dom f_1 \cap dom f_2. (2.18)$$

Corollary 2.1. Let function f be self-concordant with some constant M_f and $\alpha > 0$. Then the function $\phi(x) = \alpha f(x)$ is also self-concordant with the constant $M_{\phi} = \frac{1}{\sqrt{\alpha}} M_f$.

We call function f(x) as the standard self-concordant function if f(x) is some self-concordant function with the constant $M_f = 2$. Using Corollary 2.1, we can see that any self-concordant function can be transformed into the standard self-concordant function by scaling.

Theorem 2.2. Let function f be self-concordant. If dom f contains no straight line, then the Hessian f''(x) is nondegenerate at any x from dom f.

For some self-concordant function f(x), for which we assume, that dom f contains no straight line (which implies that all f''(x) are nondegenerate, see Theorem 2.2), we denote two local norms as

$$||u||_x = \sqrt{u^\top f''(x)u} \tag{2.19}$$

$$||u||_x^* = \sqrt{u^\top f''(x)^{-1}u}.$$
 (2.20)

Consider following minimization problem

$$x^* = \arg\min_{x \in \text{dom } f} f(x) \tag{2.21}$$

as the minimization of the self-concordant function f(x). Algorithm 1 is describing the iterative process of solving optimization problem (2.21). The algorithm is divided into two stages by the value of $||f'(x_k)||_{x_k}^*$. The splitting parameter β guarantees quadratic convergence rate for the second part of the algorithm. The parameter β is chosen from interval $(0, \bar{\lambda})$, where

$$\bar{\lambda} = \frac{3 - \sqrt{5}}{2},\tag{2.22}$$

Algorithm 1 Newton method for minimization of self-concordant functions

Input:

f a self-concordant function to minimize $x_0 \in \text{dom } f$ a starting point $\beta \in (0, \bar{\lambda})$ a parameter of size of the region of quadratic convergence

 ε a precision

Output:

 x^* an optimal solution to the minimization problem (2.21)

```
1: k \leftarrow 0

2: while ||f'(x_k)||_{x_k}^* \ge \beta \operatorname{do}

3: x_{k+1} \leftarrow x_k - \frac{1}{1+||f'(x_k)||_{x_k}^*} f''(x_k)^{-1} f'(x_k)

4: k \leftarrow k+1

5: end while

6: while ||f'(x_k)||_{x_k}^* > \varepsilon \operatorname{do}

7: x_{k+1} \leftarrow x_k - f''(x_k)^{-1} f'(x_k)

8: k \leftarrow k+1

9: end while

10: return x^* \leftarrow x_k
```

which is a root of the equation

$$\frac{\lambda}{(1-\lambda)^2} = 1. (2.23)$$

The first while loop (lines 2-5) represents damped Newton method, where at each iteration we have

$$f(x_k) - f(x_{k+1}) \ge \beta - \ln(1+\beta) \text{ for } k \ge 0,$$
 (2.24)

where

$$\beta - \ln(1+\beta) > 0 \text{ for } \beta > 0, \tag{2.25}$$

and therefore the global convergence of the algorithm is ensured. It can be shown that the local convergence rate of the damped Newton method is also quadratic, but the presented switching strategy is preferred as it gives better complexity bounds.

The second while loop of the algorithm (lines 6-9) is the standard Newton method with quadratic convergence rate.

The algorithm terminates when the required precision ε is reached.

2.4.2. Self-concordant barriers

To be able to introduce self-concordant barriers, let us denote $\operatorname{Dom} f = \operatorname{cl}(\operatorname{dom} f)$.

Definition 2.3 (Self-concordant barrier). Let F(x) be a standard self-concordant function. We call it a ν -self-concordant barrier for set Dom F, if

$$\sup_{u \in \mathbb{R}^n} \left(2u^\top F'(x) - u^\top F''(x)u \right) \le \nu \tag{2.26}$$

for all $x \in \text{dom } F$. The value ν is called the parameter of the barrier.

The inequality (2.26) can be rewritten into the following equivalent matrix notation:

$$F''(x) \succeq \frac{1}{\nu} F'(x) F'(x)^{\top}.$$
 (2.27)

In Definition 2.3, the hessian F''(x) is not required to be nondegenerate. However, in case that F''(x) is nondegenerate, the inequality (2.26) is equivalent to

$$F'^{\top}(x)F''(x)^{-1}F'(x) \leq \nu. \tag{2.28}$$

Let us explore, which basic functions are self-concordant barriers.

Example 2.2.

1. Linear functions.

$$F(x) = \alpha + a^{\top} x, \text{ dom } F = \mathbb{R}^n$$
 (2.29)

$$F''(x) = 0 (2.30)$$

From (2.27) and for $a \neq 0$ follows, that linear functions are not self-concordant barriers.

2. Convex quadratic functions.

For $A = A^{\top} \succ 0$:

$$F(x) = \alpha + a^{\mathsf{T}}x + \frac{1}{2}x^{\mathsf{T}}Ax, \ \operatorname{dom} F = \mathbb{R}^n$$
 (2.31)

$$F'(x) = a + Ax (2.32)$$

$$F''(x) = A (2.33)$$

After substitution into (2.28) we obtain

$$(a+Ax)^{\top}A^{-1}(a+Ax) = a^{\top}A^{-1}a + 2a^{\top}x + x^{\top}Ax,$$
 (2.34)

which is unbounded from above on \mathbb{R}^n . Therefore, quadratic functions are not self-concordant barriers.

3. Logarithmic barrier for a ray.

$$F(x) = -\ln x, \text{ dom } F = \{x \in \mathbb{R} \mid x > 0\}$$
 (2.35)

$$F'(x) = -\frac{1}{x} \tag{2.36}$$

$$F''(x) = \frac{1}{x^2} \tag{2.37}$$

From (2.28), when F'(x) and F''(x) are both scalars, we get

$$\frac{F'(x)^2}{F''(x)} = \frac{x^2}{x^2} = 1. {(2.38)}$$

Therefore, the logarithmic barrier for a ray is a self-concordant barrier with parameter $\nu = 1$ on domain $\{x \in \mathbb{R} \mid x > 0\}$.

Now, let us focus on the main properties of self-concordant barriers.

Theorem 2.3. Let F(x) be a self-concordant barrier. Then the function $c^{\top}x + F(x)$ is a self-concordant function on dom F.

Theorem 2.4. Let F_i be ν_i -self-concordant barriers, i=1,2. Then the function

$$F(x) = F_1(x) + F_2(x) (2.39)$$

is a self-concordant barrier for convex set

$$Dom F = Dom F_1 \cap Dom F_2 \tag{2.40}$$

with the parameter

$$\nu = \nu_1 + \nu_2. \tag{2.41}$$

Theorem 2.5. Let F(x) be a ν -self-concordant barrier. Then for any $x \in \text{Dom } F$ and $y \in \text{Dom } F$ such that

$$(y-x)^{\top} F'(x) \geq 0, \tag{2.42}$$

we have

$$||y - x||_x \le \nu + 2\sqrt{\nu}.$$
 (2.43)

There is one special point of convex set, which is important for solving convex minimization problem. It is called analytic center of convex set and we will focus on its properties.

Definition 2.4. Let F(x) be a ν -self-concordant barrier for the set Dom F. The point

$$x_F^* = \arg\min_{x \in \text{Dom } F} F(x) \tag{2.44}$$

is called the analytic center of convex set Dom F, generated by the barrier F(x).

Theorem 2.6. Assume that the analytic center of a ν -self-concordant barrier F(x) exists. Then for any $x \in \text{Dom } F$ we have

$$||x - x_F^*||_{x_F^*} \le \nu + 2\sqrt{\nu}.$$
 (2.45)

This property clearly follows from Theorem 2.5 and the fact, that $F'(x_F^*) = 0$.

Thus, if Dom F contains no straight line, then the existence of x_F^* (which leads to nondegenerate $F''(x_F^*)$) implies, that the set Dom F is bounded.

Now, we describe the algorithm and its properties for obtaining an approximation to the analytic center. To find the analytic center, we need to solve the minimization problem (2.44). For that, we will use the standard implementation of the damped Newton method with termination condition

$$||F'(x_k)||_{x_k}^* \le \beta, \text{ for } \beta \in (0,1).$$
 (2.46)

The pseudocode of the whole minimization process is shown in Algorithm 2.

2.5. Implementation details

2.6. Comparison with the state of the art methods

Algorithm 2 Damped Newton method for analytic centers

Input:

 ${\cal F}$ a self-concordant barrier

$$x_0 \in \operatorname{Dom} F$$
 a starting point

$$\beta \in (0,1)$$

Output:

 x_F^* an approximation of the analytic center of the set $\operatorname{Dom} F$

2: while
$$||F'(x_k)||_{x_k}^* > \beta$$
 do
3: $x_{k+1} \leftarrow x_k - \frac{1}{1+||F'(x_k)||_{x_k}^*} F''(x_k)^{-1} F'(x_k)$
4: $k \leftarrow k+1$

- 5: end while
- 6: **return** $x_F^* \leftarrow x_k$

3. Conclusion

A. Contents of the enclosed CD

Bibliography

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