CLASS 01

Class Atendence

16h10

submission of work done at the end of class

IMPORTANT

1st project delivery — 9th April

- mathematical formulation
- pseudocode

IMPORTANT

2nd project delivery — 21st May

- implementation
- analysis

Other evaluation points:

- Exam (Theoretical + Practical components)
- Kahoot! (Theoretical Questions)
- HackerRank (Programming Exercises/Competition)

TP Class 1

1.

```
Input: T, selected[3]
for i ← 1 to T do
    for j ← 1 to i do
        for K ← 1 to j do
            if i+j+K == T then
                selected[0] ← i
                selected[1] ← j
                selected[2] ← K
                return true
            end if
return false
```

7.

a) Consider:

set $T = \{a_1, a_2, \ldots, a_n\}$ of tasks

a_i requires t_i units to be completed

c_i completion time of task a_i

Restrictions:

one task at a time

Objective:

$$\min \frac{1}{n} \sum_{i=1}^{n} c_i$$

b) proof (non rigorous):

Consider an ordered set of tasks $T = \{a_1, \ldots, a_k\}$

$$C_{k=0} = 0 \; ; \; C_k = t_k + C_{k-1}$$

Example: $K \in \{0, 3\}$

$$K=0 \rightarrow C_{K=0} = 0$$
$$K=1 \rightarrow C_{K=1} = t_{K=1}$$
$$K=2 \rightarrow C_{K=2} = t_{K=2} + C_{K=1} = t_{K=2} + t_{K=1} \quad \text{(substitution)}$$
$$K=3 \rightarrow C_{K=3} = t_{K=3} + C_{K=2} = t_{K=3} + t_{K=2} + t_{K=1}$$

$$\min (1/n) \times \sum_{i=1}^{n} c_i = \min \sum_{i=1}^{n} c_i \quad \text{(minimize sum)}$$
$$= \min (0 + t_{K=1} + t_{K=2} + t_{K=1} + t_{K=3} + t_{K=2} + t_{K=1})$$

$0 = C_{K=0}$, $t_{K=1} = C_{K=1}$, $t_{K=2} + t_{K=1} = C_{K=2}$, $t_{K=3} + t_{K=2} + t_{K=1} = C_{K=3}$

It is trivial to see that we need to choose tasks, a_i, with the lowest t_i first since their terms will be propagated through the summation.
A greedy algorithm will give an optimal solution, with selection criteria lowest t_i.