

Advanced Logics – Assignment 12

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Due: Thursday, 14. July 2022, 9:20

The stars indicate the difficulty of an exercise (no star - easiest, one star - harder, two stars - hardest).

Exercise 12.1 [FOL-sentences for star-free languages]

Let $\Sigma = \{a, b\}$. Provide an FOL-sentence ϕ over $\text{Voc}_{\Sigma, \leq}$ that defines the language over $\Sigma = \{a, b\}$ given by the following star-free expression:

$$\overline{b\emptyset a\emptyset}$$

Exercise 12.2 [From FOL to star-free and regular expressions]

Let $\Sigma = \{a, b\}$. Provide star-free and regular expressions for the star-free languages defined by the following FOL-sentences ϕ_i over $\text{Voc}_{\Sigma, \leq}$:

$$\phi_1 = \forall x \exists y. (P_a(x) \wedge P_b(y))$$

$$\phi_2 = \forall x. (P_a(x) \rightarrow \exists y \exists z. (E(y, x) \wedge E(z, y) \wedge P_b(z)))$$

Exercise 12.3 [Star-freeness of a regular language?]

Is the regular language given by the regular expression

$$(ab)^*$$

star-free? Justify your answer.

Exercise 12.4 [Regular language associated with an MSO-formula]

Provide a characterization of $\mathcal{L}(\phi)$ for the MSO-formula over $\text{Voc}_{\Sigma, \text{graph}}$ for $\Sigma = \{a, b\}$

$$\phi(z, Y) = \psi_1 \wedge \psi_2, \text{ where}$$

$$\psi_1 = P_a(z) \rightarrow z \notin Y$$

$$\psi_2 = \exists X. (\text{first} \in X \wedge z \in X \wedge \forall y \forall y'. (E(y, y') \rightarrow (y \in X \leftrightarrow y' \notin X)))$$

and depict an NFA for $\mathcal{L}(\phi)$.

Exercise 12.5 [REG2EMSO (*)]

Provide an alternative proof for the EMSO-definability of regular languages by defining a direct transformation from regular expressions to EMSO-sentences without using the equivalence of regular expressions and finite automata.

Exercise 12.6 [Hamilton path problem and MSO/ESOL-definability]

- (a) Show that there is no MSO-sentence ϕ over $\text{Voc}_{\text{graph}}$ such that for all finite undirected graphs \mathcal{G} :

$\mathcal{G} \models \phi$ if and only if \mathcal{G} has a Hamilton path.

Remind: A Hamilton path in a graph $\mathcal{G} = (V, E)$ is a path $v_0 v_1 \dots v_n$ s.t. $v_i \neq v_j$ for $0 \leq i < j \leq n$ and $V = \{v_0, v_1, \dots, v_n\}$

- (b) Does there exists an ESOL-sentence ϕ over $\text{Voc}_{\text{graph}}$ s.t. for all finite undirected graphs \mathcal{G} :

$\mathcal{G} \models \phi$ if and only if \mathcal{G} has a Hamilton path?