

# Software Engineering for HPC - Notes

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June 2024

## Preface

Each theory section in these notes has been taken from the following sources:

- Course slides. [2]
- SE4HPC Exercise Book. [3]

About:



These notes are an unofficial resource and shouldn't replace the course material or any other book on software engineering. It is not made for commercial purposes. I've made the following notes to help me improve my knowledge and maybe it can be helpful for everyone.

As I have highlighted, a student should choose the teacher's material or a book on the topic. These notes can only be a helpful material.

During the Software Engineering for HPC course, we created two projects:

1. Requirement Engineering and Design Project (section 1 to 5). More information in the following repository (actually private)
  - [SE4HPC\\_project\\_RD](#)
2. DevOps Project (section 6 to 7). More information in the following repositories
  - [SE4HPC\\_project\\_part1](#)
  - [SE4HPC\\_project\\_part2](#)

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## 1 Introduction

### 1.1 The importance of software engineering

Software engineering is so important because it is everywhere. Our society is now totally dependent on software-intensive systems. Think about it. The society could not function without software, for example:

- Transportation systems;
- Energy systems;
- Manufacturing systems.

For these reasons, **software failures cannot be tolerated**.

In the following list, we can see some famous software issues:

- **911 Outage on April 2014.** On 10th April 2014, Washington State had no 911 service for six hours. A software issue causes this event. The software dispatching the calls had a counter used to assign a unique identifier to each call. The counter went over the threshold defined by developers. All calls from that moment on were rejected.

More info is [here](#).

- **Ariane 5, 1996.** On 4th June 1996, forty seconds after take off, Ariane 5 broke up and exploded. The total cost for developing the launcher has been 8000 million dollars. The launcher contained a cluster of satellites for 500 million dollars. Again, the explosion was caused by software failure.

More info is here: [accident tech report](#) and [video](#).

## 1.2 Software engineering: definition

There are some fields of computer science dealing with software systems:

- Large and complex;
- Built by teams;
- It exists in many versions;
- Last many years;
- Undergo changes.

In each field, a software engineer needs to have some skills. In contrast to a *programmer* that has the following abilities:

- They develop a complete program;
- They work on known specifications;
- They work individually.

A **software engineer** has the following **skills**:

- **Identifies requirements** and develops *specifications*;
- **Designs** a component to be combined with other components, developed, maintained, and used by others; component can become part of several systems;
- **Works in a team**.

We can summarize the skills of a software engineer as follows:

- Technical
- Project management
- Cognitive
- Enterprise organization
- Interaction with different cultures
- Domain knowledge

The **main goal** of a software engineer is to **develop software products**. Not only is the product significant, but the **process is also fundamental**. The quality of the process affects the quality of the product.

### 1.3 The software product and the process

The product developed by a software engineer differs from traditional product types. It isn't easy to describe and evaluate because it is intangible. Some aspects affecting the product quality:

- Development technology;
  - Process quality;
  - People quality;
  - Cost, time and schedule.
- 

#### 1.3.1 ISO/IEC 25010

An [ISO](#) (International Organization for Standardization) called **ISO/IEC 25010** comprises the **nine quality characteristics**:

SOFTWARE PRODUCT QUALITY									
FUNCTIONAL SUITABILITY	PERFORMANCE EFFICIENCY	COMPATIBILITY	INTERACTION CAPABILITY	RELIABILITY	SECURITY	Maintainability	FLEXIBILITY	SAFETY	
FUNCTIONAL COMPLETENESS	TIME BEHAVIOUR	CO-EXISTENCE	APPROPRIATENESS RECOGNIZABILITY	FAULTLESSNESS	CONFIDENTIALITY	MODULARITY	ADAPTABILITY	OPERATIONAL CONSTRAINT	
FUNCTIONAL CORRECTNESS	RESOURCE UTILIZATION	INTEROPERABILITY	LEARNABILITY OPERABILITY	AVAILABILITY	INTEGRITY	REUSABILITY	SCALABILITY	RISK IDENTIFICATION	
FUNCTIONAL APPROPRIATENESS	CAPACITY		USER ERROR PROTECTION	FAULT TOLERANCE	NON-REPUDIATION	ANALYSABILITY	INSTALLABILITY	FAIL SAFE	
			USER ENGAGEMENT	RECOVERABILITY	ACCOUNTABILITY	MODIFIABILITY	REPLACEABILITY	HAZARD WARNING	
			INCLUSIVITY		AUTHENTICITY	TESTABILITY		SAFE INTEGRATION	
			USER ASSISTANCE		RESISTANCE				
			SELF-DESCRIPTIVENESS						

Figure 1: [ISO/IEC 25010](#)

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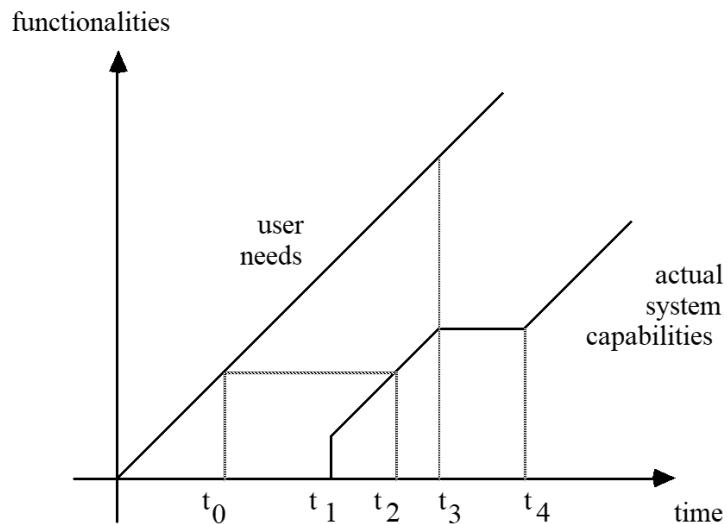
#### 1.3.2 Productivity

A process quality to consider is **productivity** (the **process of producing a product**). The definition can be: “ability to produce a good amount of product”. To **measure it**, we can use **delivered items by unit of effort**, where:

- **Delivered items:** lines of code (and variations) function points;
- **Unit of effort:** person month (note: persons and months cannot be interchanged).

### 1.3.3 Timeliness

Another process quality to consider is timeliness. The definition is: "**the ability to respond to change requests in a timely fashion**".



As you can see by the graph, the “*user needs*” is a linear function (and sometimes can be exponential!). A software engineer should be able to respond to the client’s requests as soon as possible. As the graph shows, a request made on time  $t_0$  is completed on time  $t_2$ ; but another request can be made at that time, and so on. The actual system capabilities can’t grow up always because sometimes there are “brainstorming times” to increase product quality (ISO/IEC 25010).

## 1.4 Software Lifecycle

Initially, no reference model is inside a software lifecycle: code and fix (or refactoring). However, a traditional waterfall model is chosen to react to the many problems that a software engineer faces.

### 1.4.1 Waterfall model

The **waterfall model** is a **breakdown of development activities into linear sequential phases**, meaning they are passed down onto each other, where each phase depends on the deliverables of the previous one and corresponds to a specialization of tasks. [1] Its organization is the following:

- High phases:

- **Feasibility Study**: this is a **cost-benefit analysis**.

The **main goal** is determining whether the project should be started (e.g. buy or make), possible alternatives, and needed resources.

The **outcome** is a **feasibility study document**. This paper provides:

- \* A preliminary problem description;
- \* Some scenarios describing possible solutions;
- \* Costs and schedules for the different alternatives.

- **Requirements Analysis and Specification**: this is an **analysis of the domain in which the application takes place**.

The **main goal** is to **identify requirements and derive specifications for the software**. Note these specifics require a (continuous) interaction with the user and an understanding of the properties of the domain.

The **outcome** is a particular document called **Requirements Analysis and Specification Document** (RASD).

- **Design**: this is the **definition of the software architecture**. There, the definition of components (modules) and the relations/interactions among these components.

The **main goal** is to **support the concurrent development of separate responsibilities**.

The **outcome** is a summary of this info in a **design document**.

- Low phases:

- **Coding and Unit Test**: each module is **implemented using the chosen programming language**. Furthermore, each module is **tested in isolation** by the module developer. Also, the programs should include their documentation.

- **Integration and System Test**: the modules are **integrated into (sub)systems**. The integrated (sub)systems are **tested**. Follows an incremental implementation scheme. A complete system test is needed to verify the overall properties. Note that sometimes we have alpha test and beta test.

- **Deployment**: is the process used to conceive, specify, design, program, document, test, and bug fix to create and maintain applications, frameworks, or other software components.
- **Maintenance**: the maintenance is divided into **two types**:
  - \* **Corrective** deals with the **repair of faults or defects found**.
  - \* **Evolution** is also divided into **three types**:
    - *Adaptive* maintenance: consists of **adapting software to changes in the environment** (the hardware or the operating system, business rules, government policies).
    - *Perfective* maintenance: mainly deals with accommodating new or changed user requirements.
    - *Preventive* maintenance: concerns activities aimed at **increasing the system's maintainability**.

### ⚠ Problems derived from correction and evolution

Note: the **distinction between correction and evolution can be unclear** because specifications often must be completed and clarified. This causes problems because specs are usually part of a developer and customer contract.

- **Early frozen specs** can be problematic because they are more likely to be wrong.
- Another problem is **software evolution** because **it is never anticipated or planned**. Since the software is easy to change, often, under emergency, changes are applied directly to code, and consequently, the state of project documents is inconsistent.

### ✓ Solutions - Best practices

Some good engineering practices exist to solve the evolution problem: **first, modify the design, then change implementation and apply changes consistently in all documents**. Also, the **software must be designed to accommodate changes cost-effectively**. This is one of the *main goals* of software engineering.

### ✓ Flexible processes

We can make the **waterfall model more flexible**. In this case, the **main goal is to adapt to changes** (especially in requirements and specs). The idea is that the stages are **not necessarily sequential**, and processes become **iterative and incremental**.

## 2 HPC Software, Relevant Qualities and Systems Engineering Methods

### 2.1 High Performance Computing Software

There's no single definition of HPC, but it can be explained in a number of ways:

#### Definition 1

The practice of aggregating computing power in a way that delivers much high performance than one could get out of a typical desktop computer or workstation to solve large problems in science, engineering, or business.

Thousands of processors working in parallel to analyze billions of pieces of data in real time, performing calculations thousands of times faster than a normal computer.

The use of parallel processing for running advanced, large-scale application programs efficiently, reliably and very quickly on supercomputer systems.

The platform technology concerned with programming for performance, where performance takes the broad meaning of:

- Speed (reducing time to solution);
- Energy efficiency (doing more with less power);
- Upscaling (handling larger problems such as simulating a wing and then a full plane, or a cell and then an organ);
- High throughput (the ability to handle large volumes of data in near real-time, as required in the financial services industry, telecoms or satellite imagery).

As **Parallel and Distributed Computing (PDC)** exist, it is necessary to explain the difference. The main characteristics of PDC are:

- **Concurrency:** it is a property of software. **A piece of software is also concurrent if it can have more than one active execution context.**
- **Parallelism:** it is a property of software. **The execution of different tasks/pieces of software at the same time.**
- **Distribution.** **The execution of different tasks/pieces of software on physically distinct computing nodes connected through a network, lack of a global clock.**

PDCs are **multi-core machines**, whereas **HPCs** are **quantum computers**. However, both share **parallel machines**, **HPC clusters** and **cloud infrastructures**.

There are **two categories** of HPC software:

- **Compute-intensive applications.** These are **complex calculations** that require a large number of computing resources. They also often require parallel computing.
- **Data-intensive applications.** They **focus on processing, storing and retrieving large amounts of data.** Typically built as distributed systems to ensure reliability and scalability.

## 2.2 Relevant Qualities

For the two categories explained, there are some important characteristics:

- For **Compute-intensive applications**:

– **Correctness**: the software is **correct if it satisfies the specifications**, but be careful! Sometimes, modelling reality into a model (using the specifications) isn't the bigger problem. Instead, it is difficult or impossible to show actual correctness concerning reality. For **example**, imagine you are building a simulator of a planet lander before you have ever visited it.

*How can you fix this issue?* We can **check that the software output fulfils the important desired properties** and **identify and apply a measure of accuracy**.

– **Performance**: it is the **efficient use of resources**. Again, be careful! Is it a good idea? Is performance improvement always a good idea? Because it is **not necessarily if**:

- \* It makes software **more difficult to read and maintain**
- \* It **reduces the portability** of software

– **Portability**.

– **Maintainability**. A system can have this feature if it follows **three principles**:

1. **Operability**: **make it easy for the operations team to run the system and keep it running**. There are a number of things that need to be done to achieve this:

- \* Provide visibility into the runtime behavior and internals of the system, with good monitoring.
- \* Provide good support for automation and integration with standard tools.
- \* Avoidance of dependencies on individual machines (allowing machines to be taken down for maintenance while the system as a whole continues to run uninterrupted).
- \* Provide good documentation and an easy to understand operational model (“If I do X, Y will happen”).
- \* Provide good default behavior, but also give administrators the freedom to override defaults when necessary.
- \* Self-healing when appropriate, but also giving administrators manual control over system state when needed.

2. **Simplicity**: **make it easy for other software engineers to understand the system**. This is necessary because complex systems take more time to understand and increase the cost of maintenance. There are several techniques for doing this:

- \* Reducing *accidental complexity*.
- \* Using abstractions, such as organising the architecture into well-defined components that hide the internal complexity behind a clear and easy-to-use interface; or reusing known solutions.

3. **Evolvability:** make it easy for engineers to change the system as new requirements emerge. There are a number of things that need to be done to achieve this:
  - \* Organize your development process to cope with evolution.
  - \* Keep track of how requirements are mapped to your software structure.
  - \* Update documentation.
  - \* Continue to ensure simplicity and operability.

- For **Data-intensive applications:**

- **Reliability:** can be mathematically defined as **probability of absence of failures for a certain period**. The typical expectations are:
  - \* The application performs the expected function
  - \* It can tolerate mistakes by users
  - \* It prevents unauthorized access and abuse
- **Scalability:** the system ability to cope with increased load. The load unit depends on the product: for web apps can be represented with the number of requests per second; for databases can be the number of read and write operation (or their ratio).
- **Maintainability.** Same as above.

In the software, there can be some errors, but a software engineer should be able to recognize the type of failure, faults or defects:

- A **defect** is an **imperfection or deficiency in a work product** where that work product does not meet its requirements or specifications and needs to be either repaired or replaced.
- A **defect encountered during software execution** is a **fault** (a fault is a subtype of defect, and can be of two types, see below).
- A **system failure** can be:
  - Termination of the ability of a product to perform a required function or its inability to perform within previously specified limits.
  - An event in which a system or system component does not perform a required function within specified limits.

There are some exceptions where systems are **fault-tolerant** or **resilient**. These are systems that can **cope with faults and prevent faults from occurring**. An **advantage of fault-tolerance** is that **reliability is increased**.

The **fault** can be of **two types**:

- **Hardware Faults.**

#### **⚠ Description of the problem**

It is a defect encountered during hardware execution. In a large datacenter these can happen on a daily basis. Different pieces of hardware usually fail independently from each other.

#### **✓ Possible solutions**

The possible solutions are two: **hardware redundancy** and **software fault-tolerance techniques**.

- **Software Faults.**

#### **⚠ Description of the problem**

They result from **software development errors**. Can stay dormant for a long time and appear suddenly. They can **determine failures in multiple components** at the same time.

#### **✓ Possible solutions**

There is no single solution! It is a combination of strategies. So use defensive programming, by testing before release and during operation:

- Reboot the system frequently (rejuvenation)
- Continuous monitoring and alerting in case of possible symptoms
- Deliberately introduce failures to train the fault tolerance machinery (chaos engineering)

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### 2.3 Systems Engineering Methods

There are several systems engineering methodologies required in High Performance Computing:

- Modelling the software structure and checking its properties.
- Performance analysis and improvement.
- Source code management.
- Documentation, standards, support to maintainability.
- Support to scalability.
- Attention to operability and automation.

## 3 Requirement Engineering

### 3.1 Definition

Before the definition, we give a possible scenario to understand what requirement engineering is.

The municipality of Milan says the following: “*The time it takes to make decisions on building permits for residential buildings in the city is too long. We want to develop software that will help us reduce this time*”. So where do we start? How do we identify the most important aspects? How do we make sure that we have understood what our customers want from us?

#### Definition 1

Software measure engineering (**Requirement Engineering**) is the process of discovering the purpose for which the software is intended by identifying stakeholders and their needs, and documenting these in a form suitable for analysis, communication and subsequent implementation.

The questions derived from requirements engineering are:

- Identify stakeholders
- Identify their needs
- Produce documentation
- Analyze, communicate, implement requirements

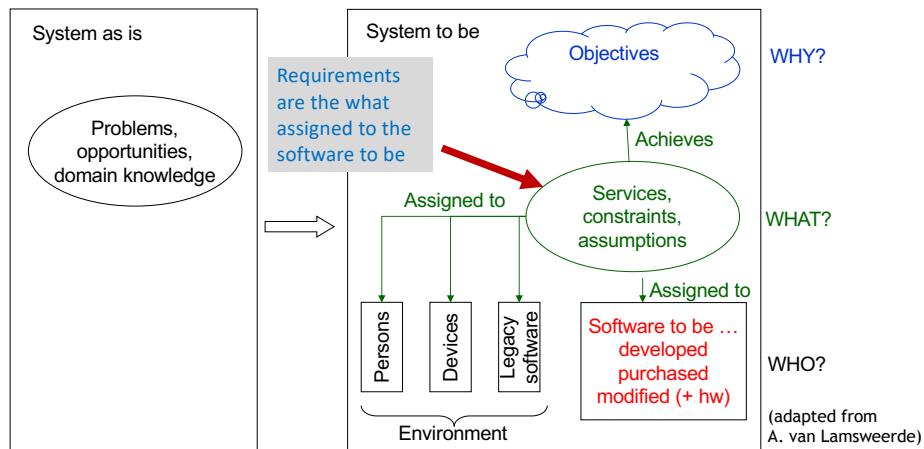


Figure 2: Analyzing the system as is and the system to be.

### 3.2 Studying the interplay between the world and the machine

#### Example 1: ambulance dispatching system

For every urgent call reporting an incident, an ambulance should arrive at the incident location within 14 minutes. For every urgent call, details about the incident are correctly encoded.

When an ambulance is dispatched, it will reach the incident location in the shortest possible time. Accurate ambulance locations are known by GPS. Ambulance crews correctly notify ambulance availability through a mobile data terminal.

Given the previous problem, are you able to extract requirements from this description? Some possible questions might be:

- Should the software system drive the ambulance?
- Who or what will “correctly encode” details about incidents?
- Do terminals already exist or not?

And more in general:

- What are the boundaries of the system? What is inside/outside? What is in-between?
- How do we think about these aspects in a systematic way?

This example is necessary to understand the **phenomena of world and machine**. The **machine** is the part of the system to be developed (typically a software-to-be and a hardware). The **world** (or environment) is the part of the real world that is affected by the machine.

Requirements engineering is concerned with the phenomena that occur in the world. In the previous **example**, RE is concerned with the following phenomena:

- Occurrence of incidents
- Reports of incidents by public calls
- Encoding of call details into dispatching software
- Assignment of an ambulance
- Arrival of an ambulance at the scene of an incident

But RE is also interested in the phenomena that occur inside the machine. In the previous **example**

- The creation of a new object of the class **Incident**
- The updating of a database entry

**Requirements models are models of the world!**

Using the **example** on the previous page, we can show the phenomena we are interested in the world or in the machine set.

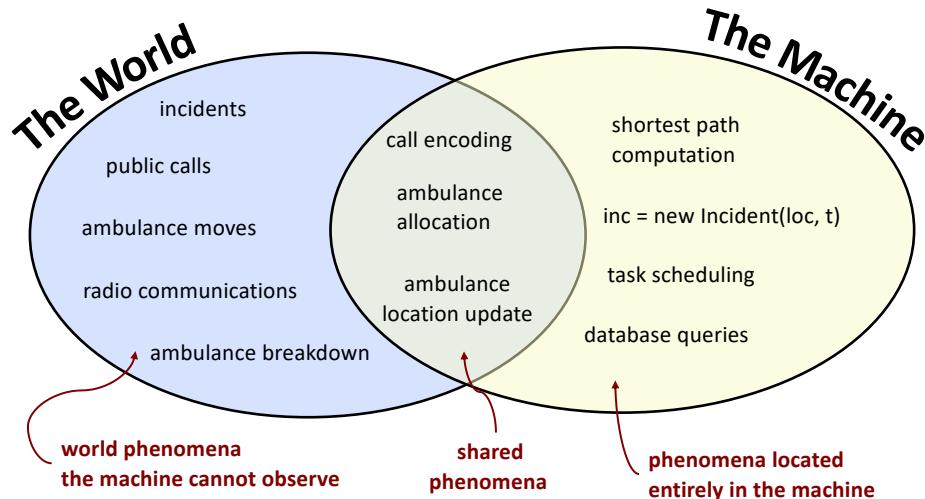


Figure 3: The world and the machine sets, with reference to example on page 17.

More generally, we can divide the machine and the world sets as:

- The world which have goals and domain properties;
- The machine which have computers and programs;
- The requirements which is the bridge between the world and the machine.

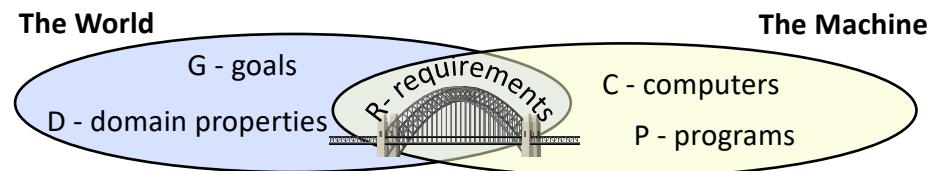


Figure 4: Goals, domain properties, requirements, computers and programs.

We explain more detailed these value inside the two sets:

- **Goals** are prescriptive assertions formulated in terms of world phenomena (not necessarily shared)
- **Domain properties** (or assumptions) are descriptive assertions assumed to hold in the world
- **Requirements** are prescriptive assertions formulated in terms of shared phenomena

Using the **example** on the page 17, we can identify the goal, the domain assumptions and the requirement as follows:

- **Goal:** *For every urgent call reporting an incident, an ambulance should arrive at the incident scene within 14 minutes.*
  - **Domain assumptions:**
    - *For every urgent call, details about the incident are correctly encoded.*
    - *When an ambulance is mobilized, it will reach the incident location in the shortest possible time.*
    - *Accurate ambulances' location are known by GPS.*
    - *Ambulance crews correctly signal ambulance availability through mobile data terminals on board of ambulances.*
  - **Requirement:** *When a call reporting a new incident is encoded, the Automated Dispatching Software should mobilize the nearest available ambulance according to information available from the ambulances' GPS and mobile data terminals.*
- 

### 3.2.1 Completeness of Requirements

Given the set of requirements **R**, goals **G** and domain assumptions **D**.

#### Definition 2

We say that **R** is **complete** if and only if:

- **R** ensures satisfaction of **G** in the context of domain assumptions **D**

$$\mathbf{R} \text{ and } \mathbf{D} \mid = \mathbf{G}$$

We can make an analogy with program correctness. A program **P** running on a particular computer **C** is correct if it satisfies the requirement **R**:  $\mathbf{P} \text{ and } \mathbf{C} \mid = \mathbf{R}$ .

- **G** captures all the **stakeholders' needs**.
- **D** represents **valid properties/assumptions about the world**.

### 3.3 Formulating and classifying requirements

The requirements can be of three types:

- **Functional requirements:** describe the interactions between the system and its environment (independent from implementation). In other words, are the **main (functional) goals the software has to realize**.

For *example*: “*the word processor shall allow users to search for strings in the text*”; “*the system shall allow users to reserve taxis*”.

- **Non-functional requirements (NFRs):** further characterization of user-visible aspects of the system not directly related to functions.

For *example*: “*the response time must be less than 1 second*”; “*the server must be available 24 hours a day*”.

- **Constraints requirements:** imposed by the customer or the environment in which the system operates.

For *example*: “*the implementation language must be Java*”; “*the credit card payment system must be able to be dynamically invoked by other systems relying on it*”.

We make some observations about non-functional requirements. NFRs predicate on **external** non-functional qualities, and these qualities must be **measurable by metrics**. NFRs have some characteristics:

- Constraints on **how functionality must be provided to the end user**.
- The **application domain determines their relevance and their prioritization**.
- Have a **strong influence on the structure of the system to be built**.  
For *example*, a system may require 24/7 availability. As a result, it is likely to be designed as a replicated system (with redundant components).

#### Example 2: are these requirements?

1. “*The user should insert correct information in the enrolment form*”.

This is not a requirement! How can the software prevent a user from entering incorrect information? Specifically, is a domain assumption!

2. “*The system should check whether fiscal code are well formed*”.

Yes, the software can do this! So it is a requirement.

#### Example 3: types of requirements

##### Example of **functional requirements**:

- “*The system shall allow users to reserve taxis*”.
- “*The system should never allowe non-registered users to see the list of other users willing to share a taxi*”.

- “*The system should guarantee that the reserved taxi picks the user up*”.

But attention! There is **unfeasible** (from the perspective of the software to be) **functional requirements**:

- “*The system should guarantee that the reserved taxi picks the user up*”.

This is because the software cannot guarantee this feature!

Example of **non-functional requirements**:

- “*The system has to provide a feedback in 5 seconds*”.
- “*The system should be available 24/7*”.

Example of **technical requirements**:

- “*The system should be implemented in Java*”.
- “*The search for the available taxi should be implemented in class Controller*”.

#### Example 4: bad requirements

1. “*The system shall process all mouse clicks very fast to ensure users do not have to wait*”.

The problem here is that it **cannot be verified (tested)**, because what does “fast” mean? Do we have a metric? Can you quantify it?

2. “*The user must have Adobe Acrobat installed*”.

The problem here is that it **cannot be achieved by the software system itself**. It is not something that the system has to do. But it could be expressed as a domain assumption, so it is not a functional requirement for our software.

### 3.4 Eliciting requirements

The **Requirements Elicitation** is the practice of researching and discovering the requirements of a system from users, customers, and other stakeholders. The goal of requirements elicitation is to ensure that the software development process is based on a clear and comprehensive understanding of the customer's needs and requirements. To do that, exist a simple and effective tool called **scenarios**.

#### Definition 3

A **scenario** is a concrete, focused, informal description of a single feature of the system to be.

#### Example 5: warehouse on fire

*Bob driving down main street in his patrol car notices smoke coming out of a warehouse. His partner, Alice, reports the emergency from her car.*

*Alice enters the address of the building, a brief description of its location (i.e. north west corner), and an emergency level. In addition to a fire unit, she requests several paramedic units on the scene given that area appears to be relatively busy. She confirms her input and waits for an acknowledgment.*

*John, the Dispatcher, is alerted to the emergency by a beep of his workstation. He reviews the information submitted by Alice and acknowledges the report. He allocates a fire unit and two paramedic units to the incident site and sends their estimated time of arrival (ETA) to Alice.*

*Alice received the acknowledgment and the ETA.*

There are heuristics for finding scenarios, such as asking the customer a few questions:

- Which user groups are supported by the system to perform their work?
- What are the primary tasks that the system needs to perform?
- What data will the actor create, store, change, remove or add in the system?
- What external changes does the system need to know about?
- What changes or events will the actor of the system need to be informed about?

However, it's very important not to rely on questionnaires alone! **Insist on task observation** (if possible), ask to **speak to the end user**, not just the software contractor, and expect resistance, but try to overcome it.

Scenarios provide a nice summary of what the requirements analysis team can derive from observation, interviews, analysis of documentation. By generalizing the scenarios, we can get **Use Cases**.

To specify a use case, it's very important to follow the following scheme.

**Definition 4: Use Cases Schema**

- **Name of Use Case**
- **Actors**
  - *Description of Actors involved in use case.*
- **Entry condition**
  - *"When this use case starts the following condition is true..."*.
- **Flow of Events**
  - *Free form, informal natural language.*
- **Exit condition**
  - *"This use case terminates when the following condition holds..."*.
- **Exceptions**
  - *Describe what happens if things go wrong.*
- **Special Requirements**
  - *Nonfunctional Requirements, Constraints.*

The following **suggestions** are useful in defining an appropriate use case:

- Use cases named with verbs that indicate what the user is trying to accomplish
- Actors named with nouns
- Use cases steps in active voice
- The causal relationship between steps should be clear
- A use case per user transaction
- Separate description of exceptions
- Keep use cases small (no more than two/three pages)
- The steps accomplished by actors and those accomplished by the system should be clearly distinguished (this helps us in identifying the boundaries of the system)

First of all, we present an example of a **bad use case**.

#### Example 6: bad use case

**Example** of a bad use case referring to the ambulance dispatching example on page 17:

- Use case name: Accident
- Participating Actors:
  - Field Officer
- Flow of Events:
  1. The Field Officer reports the accident
  2. An ambulance is dispatched
  3. The Dispatcher is notified when the ambulance arrives on site

The errors are as follows:

- In the *use case name* field, the **word is a noun**. It's better to use a verb that indicates what the user is trying to achieve.
- The Dispatcher actor is **not declared** in the *Participating Actors* field, but is mentioned in the *Flow of Events* field.
- There are two main errors in the *Flow of Events* section: the first is in the sentence “*An ambulance is dispatched*”. But **who sends it?** The second is in the third sentence, because **who notifies the Dispatcher?**

Now we present an example of a *well composed* use case.

#### Example 7: use case ReportEmergency with reference to the example on page 22

There are two **actors** involved:

- Field Officer (Bob and Alice in the Scenario)
- Dispatcher (John in the Scenario)

The **Entry Condition** is always true because an emergency can be reported at any time. The **sequence of events** is as follows:

- The **FieldOfficer** activates the Report Emergency function of her terminal.
- **Friend** (the system to be developed) responds by presenting a form to the officer.
- The FieldOfficer fills the form, by selecting the emergency level, type, location, and brief description of the situation. The FieldOff-

ficer also describes possible responses to the emergency situation. Once the form is completed, the FieldOfficer submits the form.

- At which point, the **Dispatcher** is notified.
- The Dispatcher reviews the submitted information and allocates resources by invoking the **AllocateResources** use case. The Dispatcher selects a response and acknowledges the emergency report.

The **Exit Condition** is the following: the FieldOfficer has received the acknowledgment and the selected response.

There are two possible **exceptions**:

- The FieldOfficer is notified immediately if the connection between her terminal and the control room is lost.
- The Dispatcher is notified immediately if the connection between any logged in FieldOfficer and the control room is lost.

And the **special requirements** are:

- The FieldOfficer's report is acknowledged within 30 seconds;
- The selected response arrives no later than 30 seconds after it is sent by the Dispatcher.

#### Example 8: use case **AllocateResources** with reference to the example on page 22

- Use case name: **AllocateResources**
- Participating Actors:
  - Dispatcher (John in the Scenario. The Dispatcher allocates resources to an Emergency if the resource is available; of course, he also updates and removes Emergency Incidents, Actions, and Requests in the system)
  - Resource Allocator (the Resource Allocator is responsible for allocating resources in case they are scarce)
  - Resources (the Resources that are allocated to the Emergency)
- Entry Condition:
  - An Incident has been opened
- Flow of Events:
  - The Dispatcher selects the types and number of Resources that are needed for the incident.

- Friend replies with a list of Resources that fulfill the Dispatcher's request.
  - The Dispatcher selects the Resources from the list and allocates them for the incident.
  - Friend automatically notifies the Resources.
  - The Resources send a confirmation.
- Exit Condition:
    - The use case terminates when the resource is committed.
    - The selected Resource is now unavailable to any other Emergencies or Resource Requests.
- Exceptions:
    - If the list of Resources provided by Friend is insufficient to fulfill the needs of the emergency, the Dispatcher informs the Resource Allocator.
    - The Resource Allocator analyzes the situation and selects new Resources by decommitting them from their previous work.
    - Friend automatically notifies the Resources and the Dispatcher.
    - The Resources send a confirmation.

## 4 Software Design

### 4.1 Software Architecture

#### Definition 1

The **Software Architecture (SA)** of a system is the **set of structures** needed to reason about the system. These structures comprise software **elements, relations** among them, and **properties** of both.

The software architecture is a tool for thinking about systems, and it's made up of a set of structures.

The Architecture is so important because it is the vehicle for communication: internal (different teams) and external (teams and stakeholders). The Architecture manifests the first set of design decisions and is a portable abstraction of a system.

## 4.2 Architecture and multiple structures

There is a set of structures relevant to the software:

- **Component-and-connector (C&C)** structures. Describe how the system is structured as a **set of elements** that have **runtime behavior** (called components) and interactions (called connectors).

- The **components** are the principal units of computation (for **example** the clients, servers, services, etc.)
- The **connectors** represent communication (for **example** request-response mechanisms, pipes, asynchronous messages, etc.)

The **purpose** of these structures is to **enable us to answer questions** such as:

- *What are the major executing components and how do they interact at runtime?*
- *What are the major shared data stores?*
- *Which parts of the system are replicated?*
- *How does data progress through the system?*
- *Which parts of the system can run in parallel?*
- *How does the system's structure evolve during execution?*

Also, **allow us to study runtime properties** such as availability and performance.

- **Module** structures. Show how a system is structured as a **set of code or data units** that have to be procured or constructed, **together with their relations**. An **example** of modules: packages, classes, functions, libraries, layers, database tables, etc.

Modules constitute **implementation units** that can be used as the basis for work splitting (identifying functional areas of responsibility). Typical relations among modules are: uses, is-a (generalization), is-part-of.

The **purpose** of these structures is to **allow us to answer questions** such as:

- *What is the primary functional responsibility assigned to each module?*
- *What other software elements is a module allowed to use?*
- *What other software does it actually use and depend on?*
- *What modules are related to other modules by generalization or specialization (i.e. inheritance) relationships?*

Can also be used to answer questions about the **impact on the system when the responsibilities assigned to each module change**.

- **Allocation** structures. Define **how the elements** from component-and-connector or module structures **map** onto things that are not software. For **example** hardware (possibly virtualized), file systems, teams. Some typical allocation structures include deployment structure, implementation structure, work assignment structure.

### 4.3 Software design descriptions and UML

In the following pages we present some diagrams that are fundamental to creating appropriate documentation. The following diagrams show how to create some UML diagrams.

#### 4.3.1 Component Diagram (C&C structure)

A **Component Diagram** breaks down the actual **system** under development into **various high levels of functionality**. Each component is responsible for one clear aim within the entire system and only interacts with other essential elements on a need-to-know basis.

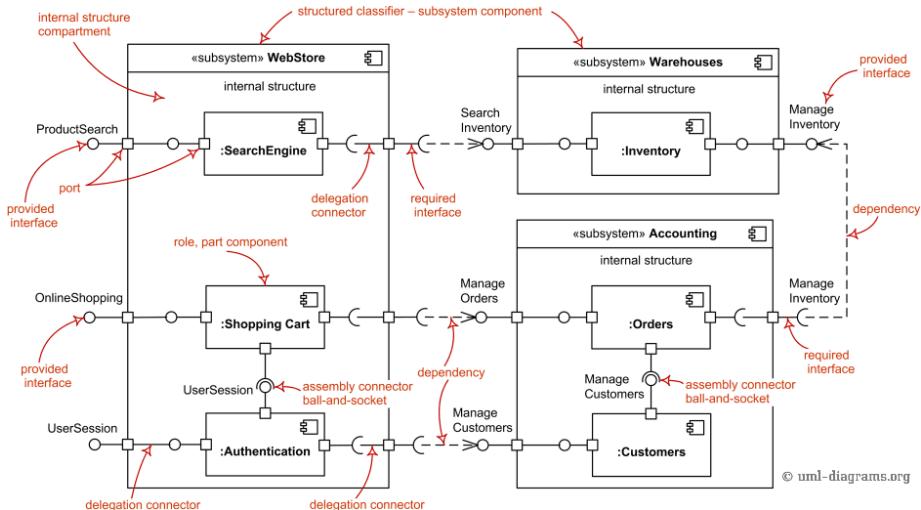


Figure 5: Component Diagram.

To view the component diagram in high resolution, scan (or click) the QR code below.



A complete guide can be found on the following page: [What is Component Diagram?](#)

### 4.3.2 Sequence Diagram (C&C structure)

**Sequence Diagram** show elements as they interact over time and they are organized according to object (horizontally) and time (vertically).

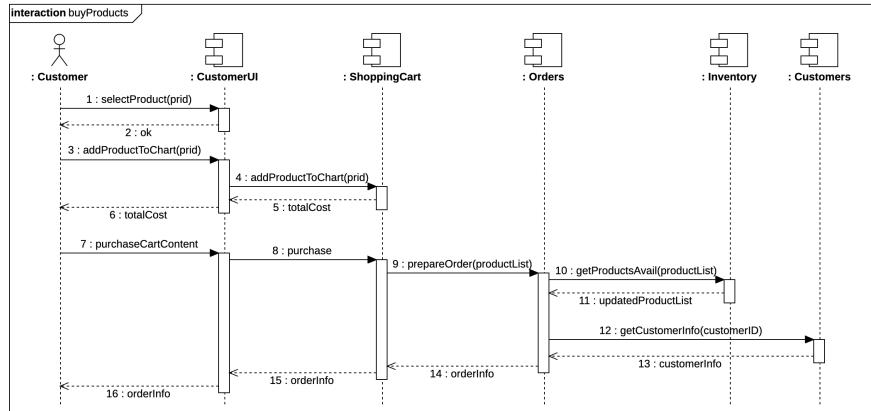


Figure 6: Sequence Diagram.

To view the sequence diagram in high resolution, scan (or click) the QR code below.



A complete guide can be found on the following page: [What is Sequence Diagram?](#)

### 4.3.3 Class Diagram (module structure)

A **Class Diagram** is a type of static structure diagram that describes the structure of a system by showing the system's classes, their attributes, operations (or methods), and the relationships among objects.

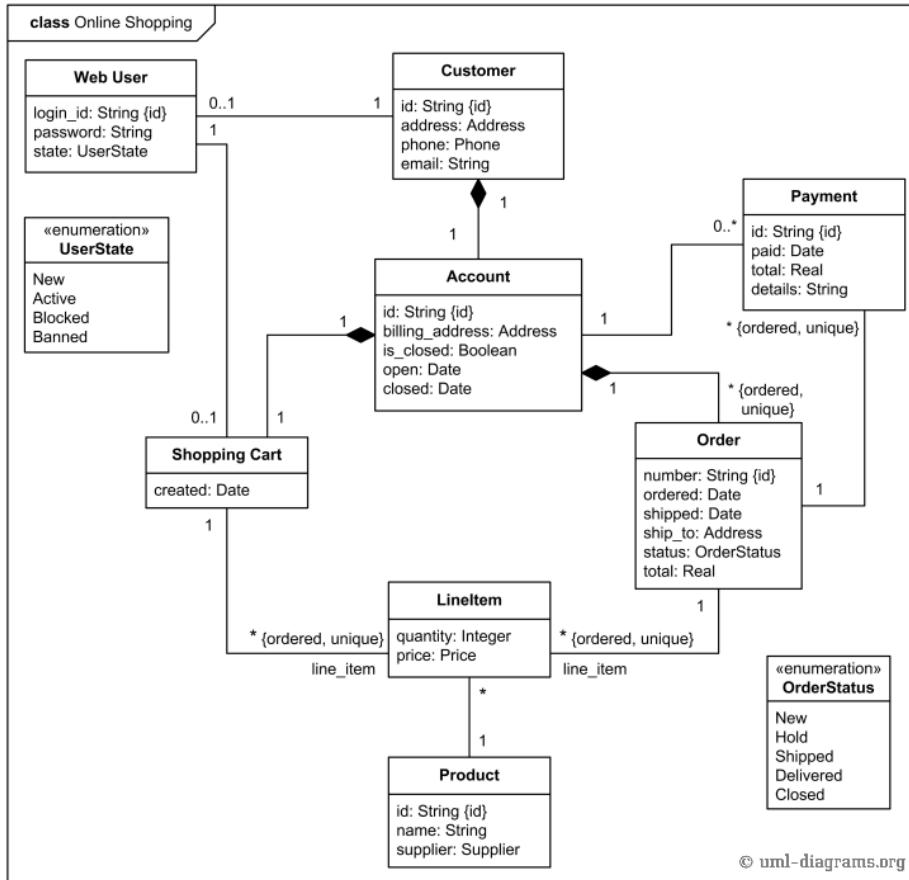


Figure 7: Class Diagram.

To view the sequence diagram in high resolution, scan (or click) the QR code below.



A complete guide can be found on the following page: [What is Class Diagram?](#)

#### 4.3.4 Package Diagram (module structure)

A **Package Diagram**, a kind of structural diagram, shows the **arrangement and organization of model elements in middle to large scale project**. Package diagram can show both structure and dependencies between sub-systems or modules, showing different views of a system, for example, as multi-layered (aka multi-tiered) application - multi-layered application model.

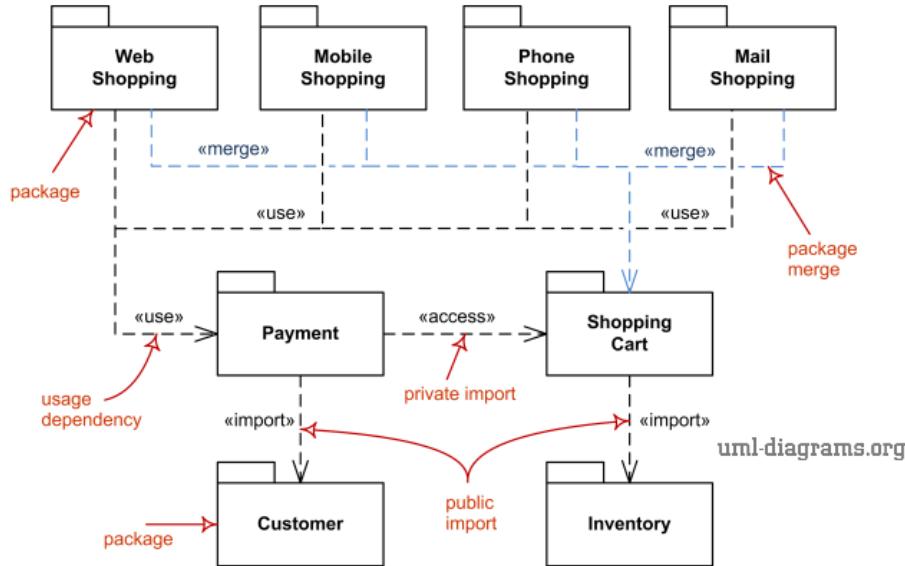


Figure 8: Package Diagram.

To view the sequence diagram in high resolution, scan (or click) the QR code below.



A complete guide can be found on the following page: [What is Package Diagram?](#)

#### 4.3.5 Deployment Diagram (allocation structure)

A **Deployment Diagram** is a diagram that shows the **configuration of runtime processing nodes and the components that live on them**. Deployment diagrams is a kind of structure diagram used in modeling the physical aspects of an object-oriented system. They are often be used to model the static deployment view of a system (topology of the hardware).

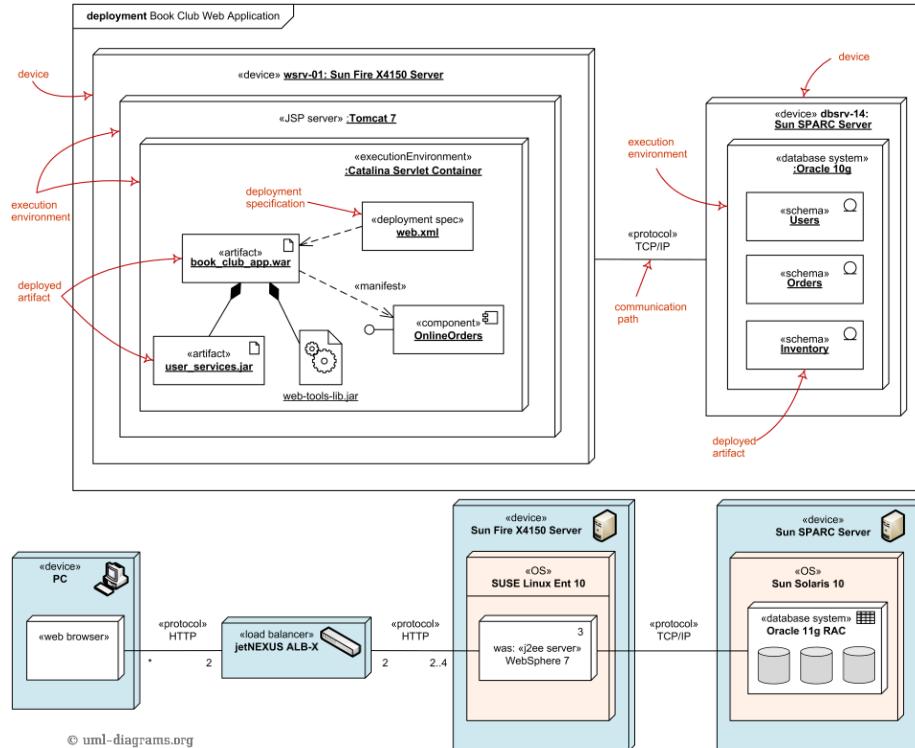
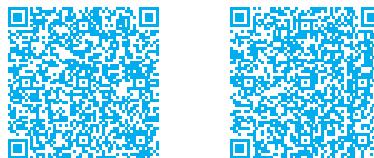


Figure 9: Deployment Diagram.

To view the sequence diagram in high resolution, scan (or click) the QR code below.



A complete guide can be found on the following page: [What is Deployment Diagram?](#)

## 4.4 Design principles

The following is a list of important design principles in software engineering.

### Definition 2

1. **Divide et impera** (also called *divide and conquer*)
2. **Keep the level of abstraction as high as possible**
3. **Increase cohesion where possible**
4. **Reduce coupling where possible**
5. **Design for reusability**
6. **Reuse existing designs and code**
7. **Design for flexibility**
8. **Anticipate obsolescence**
9. **Design for portability**
10. **Design for testability**
11. **Design defensively**

### 1. Divide et impera (*divide and conquer*)

**Divide and Conquer** is a **problem-solving strategy** that involves **breaking down a complex problem into smaller, more manageable parts**, solving each part individually, and then combining the solutions to solve the original problem.

### 2. Keep the level of abstraction as high as possible

Ensure that your designs allow you to **hide or defer consideration of details, thus reducing complexity**. A good abstraction is said to provide **information hiding**. Abstractions allow you to understand the essence of a subsystem without having to know unnecessary details.

### 3. Increase cohesion where possible

In general, a file, module, class or whatever should contain the same logical methods. For example, in the following class we have two functions with two different purposes (error!).

```

1 Class Utility {
2     ComputeAverageScore(Student s[])
3     ReduceImage(Image i)
4 }
```

#### 4. Reduce coupling where possible

**Coupling** is the **degree of interdependence between software modules**; a measure of how closely connected two routines or modules are; the strength of the relationships between modules. There are different **types of couplings**:

- **Content coupling** is said to occur when **one module uses the code of another module**, for instance a branch. This violates *information hiding* (2nd design principle).
- **Communication coupling** is said to occur when one **module sends too many messages to another module**. The creation of a message can be optimized and the number of messages sent between these two modules can be reduced.
- **Control coupling** is one **module controlling the flow of another**, by passing it information on what to do. For **example**, passing a what-to-do flag or the following code:

```

1 class b {
2     func(flag f) {
3         if(f == flag1) do this
4         else if(f == flag2) do that
5         else...
6     }
7 }
```

Other types can be viewed [here](#).

#### 5. Design for reusability

Design the various aspects of your system so that they can be **used again in other contexts**. To do this, you need to follow these rules:

- Generalize your design as much as possible;
- Simplify your design as much as possible;
- Follow the preceding all other design principles;
- Design your system to be extensible.

#### 6. Reuse existing designs and code

**Design with reuse is complementary to design for reusability.** Take advantage of the investment you or others have made in reusable components. Note: cloning should not be seen as a form of reuse.

## 7. Design for flexibility

Actively anticipate changes that a design may have to undergo in the future, and prepare for them. To do this, you need to follow these rules:

- Reduce coupling and increase cohesion;
- Create abstractions;
- Use reusable code and make code reusable;
- Do not hard-code anything.

## 8. Anticipate obsolescence

Plan for changes in the technology or environment so the software will continue to run or can be easily changed. So do not rush using early releases of technology. If possible:

- Avoid using software libraries that are specific to particular environments;
- Avoid using undocumented features or little-used features of software libraries;
- Avoid using software or special hardware from companies that are less likely to provide long-term support;
- Use standard languages and technologies that are supported by multiple vendors.

## 9. Design for portability

Have the software run on as many platforms as possible. Avoid, if possible, the use of facilities that are specific to one particular environment (e.g. a library only available in Microsoft Windows).

## 10. Design for testability

Take steps to make testing easier. Design a program to automatically test the software:

- Ensure that all the functionality of the code can be driven by an external program, bypassing a graphical user interface;
- Create proper code to exercise the other methods/functions;
- Use unit test automation frameworks.

## 11. Design defensively

Be careful when you trust how others will try to use a component you are designing. Handle all cases where other code might attempt to use your component inappropriately. Check that all of the inputs to your component are valid: the preconditions. Unfortunately, over-zealous defensive design can result in unnecessarily repetitive checking.

## 5 Architectural styles

### 5.1 Definition

#### Definition 1

An architectural style determines the **vocabulary** of **components** and **connectors** that can be used in instances of that style, together with a set of **constraints** on how they can be combined.

These can include topological constraints on architectural descriptions (e.g., no cycles). Other constraints - say, having to do with execution semantics - might also be part of the style definition.

## 5.2 Client-Server

A **Client-Server Architecture** is a **network-based computing structure** where responsibilities and operations get **distributed between clients and servers**. Client-Server Architecture is widely used for network applications such as email, web, online banking, e-commerce, etc.

### ✓ When to use it

The three most common cases are:

- When **multiple users** need to access a **single resource** (e.g. database).
- When there is a preexisting software and we must **access remotely** (e.g. email server).
- When it is convenient to organize the system around a **shared piece of functionality used by multiple components** (e.g. authentication or authorization server).

### ⚠ Technical issues

With this architecture, it's necessary to **design** and **document** proper **interfaces** for our server. It is also necessary to ensure that the server can **handle multiple simultaneous requests**.

#### 5.2.1 Interface design

An **interface design** is a **boundary** across which components interact. Proper definition of interfaces is an architectural concern (affects maintainability, usability, testability, performance, integrability). There are two important **guiding principles** for interface design: **information hiding** and **low coupling**. An interface should encapsulate a component implementation so that it can be changed without affecting other components.

There are several aspects to interface design that need to be considered:

- **Contract principle**: any resource (operation, data) added to an interface implies a **commitment to maintaining** it.
- **Least surprise principle**: interfaces should **behave** consistently **with expectations**.
- **Small interfaces principle**: interfaces should limit the **exposed resources to the minimum**.

There are also some important elements to define: **interaction style** (e.g. sockets, RPC, REST); **representation** and structure of exchanged data (affecting expressiveness, interoperability, performance and transparency); **error handling**.

### 5.2.2 Error handling, multiple interfaces and interface evolution

Sometimes there may be some problems, for example: an operation is called with invalid parameters and consequently the call doesn't return anything. This simple example can provoke some scenarios: the component cannot handle the request in its current state; or hardware/software errors prevent successful execution; or there is a misconfiguration issue (e.g. the server is not correctly connected to the database).

There are three possible **solutions**: **raising an exception**; **return an error code** (common); **log the problem**. There's no single solution, but we can choose several (e.g. error code and log the problem).

A **server** can offer **multiple interfaces** at the same time. This enables separation of concerns, different levels of access rights and support of **interface evolution**.

**Interface evolution** occurs for many reasons (e.g. to support new requirements). Several **strategies** are needed to support continuity:

- **Deprecation**: declare well in advance that an interface version will be retired by a certain date;
- **Versioning**: maintain multiple active versions of the interface;
- **Extension**: a new version extends the previous one.

### 5.2.3 Handling multiple requests

The server must be able to receive and process requests from multiple clients. There are two main approaches to this: *forking* and *worker pooling*.

#### Forking

The **forking** approach is the same as that used by the Apache Web Server: **one process per request or per client**.

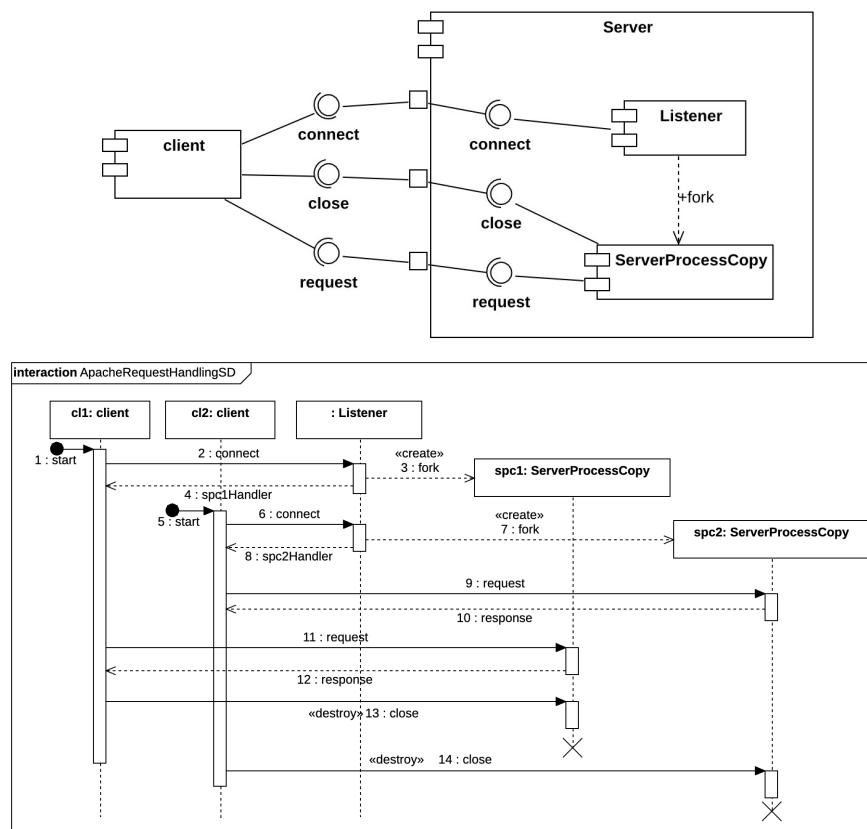


Figure 10: Forking diagrams.

#### ✓ Forking Advantages

- Architectural **simplicity**.
- **Isolation** and **protection** given by the one-connection-per-process model.  
Note: slow processes do not affect other incoming connections.
- **Simple to program**.

**A Forking Issues**

- Growth of the WWW over the last 20 years (number of users and weight of web pages).
- The number of **active processes** at time  $t$  is **difficult to predict** and may **saturate resources**.
- **Expensive** fork-kill operations for each **incoming connection**.

## Worker pooling

It is an alternative approach adopted by NGINX Web Server. It is **designed for high concurrency** but has to deal with **scalability issues**.

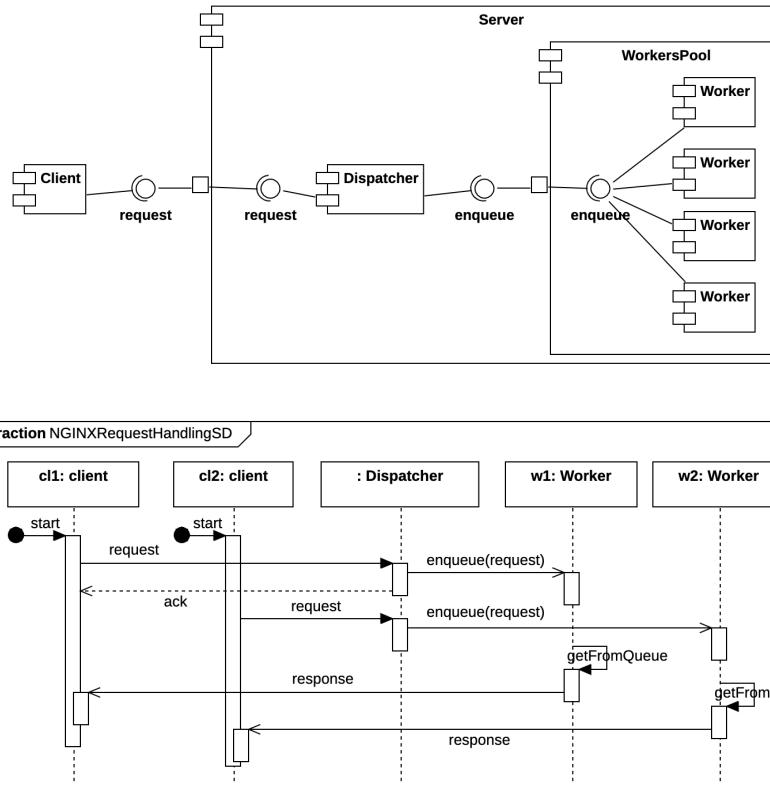


Figure 11: Worker pooling diagrams.

Despite the well-known problem of this architecture (scalability), NGINX addresses the previous problems by introducing a new **architectural tactic**. A *tactic* is a **design decision that affects the control of one or more quality attributes**.

### ✓ Worker Pooling Advantages (quality attribute trade-offs)

- Number of **workers** is fixed, so they **do not saturate available resources**.
- **Each worker** has a **queue**.
- When **queues** are **full** the **dispatcher** drops the incoming requests to keep high performance (**optimize scalability and performance by sacrificing availability**).
- Dispatcher **balances** the **workload** among available workers **according to specific policies**.

### 5.3 Three-Tier Architecture

The following is a summary of the [IBM guide](#).

**Three-tier architecture** is a well-established software application architecture that **organizes applications into three logical and physical computing tiers**:

- The **presentation** tier, or user interface;
- The **application** tier, where data is processed;
- The **data** tier, where application data is stored and managed.

#### ✓ Benefits

The chief benefit of three-tier architecture is its **logical and physical separation** of functionality. Each tier can run on a separate operating system and server platform - for example, web server, application server, database server - that best fits its functional requirements. And each tier runs on at least one dedicated server hardware or virtual server, so the services of **each tier can be customized and optimized without impacting the other tiers**. Other benefits include:

- **Faster development:** Because *each tier can be developed simultaneously by different teams*, an organization can bring the application to market faster. And programmers can use the latest and best languages and tools for each tier.
- **Improved scalability:** *Any tier can be scaled independently* of the others as needed.
- **Improved reliability:** An outage in one tier is less likely to impact the availability or performance of the other tiers.
- **Improved security:** Because the *presentation tier and data tier can't communicate directly*, a well-designed application tier can function as an internal firewall, preventing SQL injections and other malicious exploits.

#### 5.3.1 N-tier architecture

**N-tier architecture** (also called or multitier architecture) refers to any application architecture with **more than one tier**. But applications with more than three layers are rare because extra layers offer **few benefits** and can make the **application slower, harder to manage and more expensive to run**. As a result, n-tier architecture and multitier architecture are usually synonyms for three-tier architecture.

## 5.4 Microservice architectural style

The microservice architectural style is an approach to developing a single application as a suite of **small services**, each running in its own process and communicating **lightweight mechanisms**, often an HTTP resource API.

### ✓ Benefits

There are two main benefits:

- **Technology heterogeneity. Each service uses its own technology stack.** The technology stack can be selected to fit the task best (e.g. data analysis vs video streaming). The teams can experiment with new technologies within a single microservice (e.g. we can deploy two versions and do A/B testing). Also, no unnecessary dependencies or libraries for each service.
- **Scaling. Each microservice can be scaled independently.** Also, identified bottlenecks can be addressed directly. Parts of the system that do not represent bottlenecks can remain simple and unscaled.

## 5.5 Event-Driven Architecture

An **Event-Driven Architecture** uses events to trigger and communicate between decoupled services and is common in modern applications built with microservices. An event is a change in state, or an update, like an item being placed in a shopping cart on an e-commerce website. Events can either carry the state (the item purchased, its price, and a delivery address) or events can be identifiers (a notification that an order was shipped).

Often it's called **publish-subscribe** (publish is the event generation, and subscribe is the declaration of the interest).

### ✓ Benefits

- **Very common in modern development practices** (e.g. continuous integration and deployment, such as GitHub Actions).
- **Easy addition/deletion of components** (publishers and subscribers are decoupled; the event dispatcher handles this dynamic set).

### ⚠ Problems

- **Potential scalability problems** (the event dispatcher may become a bottleneck under high workload).
- **Ordering of events** (not guaranteed, not straightforward).

Other characteristics of this architecture:

- The messages and the events are **asynchronous**.
- Computation is **reactive** (driven by receipt of message).
- **Destination of messages determined by receiver**, not sender (location/identity abstraction).
- **Loose coupling** (senders and receivers added without reconfiguration).
- **Flexible** communication means (one-to-many, many-to-one, many-to-many).

Some **examples** of relevant technologies are: [Apache Kafka](#) and [RabbitMQ](#).

### 5.5.1 Apache Kafka

**Kafka** is a **framework** for the event-driven paradigm:

- Includes primitives to create **event produces** and **consumers** and a runtime infrastructure to handle **event transfer** from producers to consumers.
- **Stores events** durably and reliably.
- Allow consumers to **process events** as they occur or retrospectively.

These services are offered in a distributed, highly scalable, elastic, fault-tolerant, and secure manner.

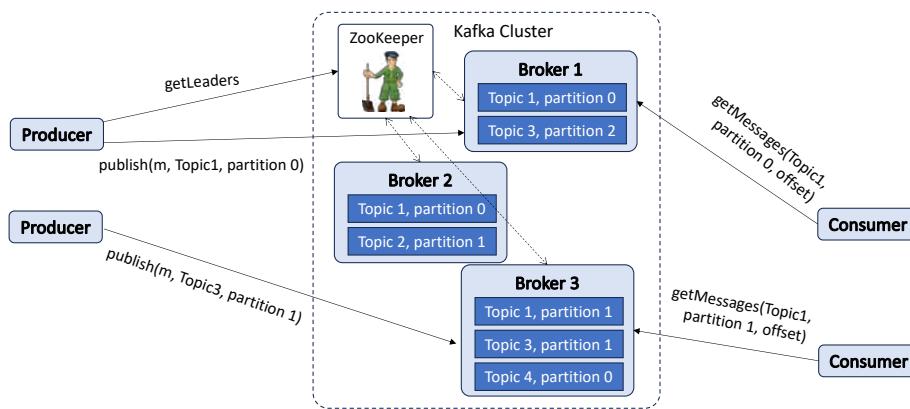


Figure 12: Kafka architecture (the ZooKeeper is a “health manager”).

Some important features:

- Each **broker** handles a set of **topics** and **topic partitions**, parts including sets of messages on the topic.
- The *partitions* are independent from each other and can be **replicated** on multiple brokers for fault tolerance.
- There is **one leading broker per partition**. The other brokers containing the same partition are **followers**.
- The **producers** know the available leading brokers and send messages to them.
- **Messages in the same topic** are organized in **batches** at the producers’ side and then sent to the broker when the batch size overcomes a certain threshold.
- **Consumers** adopt a **pull approach**. They *receive in a single batch all messages* belonging to a certain partition starting from a specified offset.
- **Messages remain available** at the brokers’ side **for a specified period** and can be **read multiple times** in this period.

- The leader keeps track of the **in-synch followers**.
- **ZooKeeper** is used to monitor the correct operation of the cluster. All brokers send heartbeats to ZooKeeper. ZooKeeper will replace a failed broker by electing a new leader for all partitions that the failed broker was leading. It can also start/restart brokers.

### Message delivery

#### Producer

1. Brokers commit messages by storing them in the corresponding partition;
2. Leader adds the message to followers (replicas) if available.

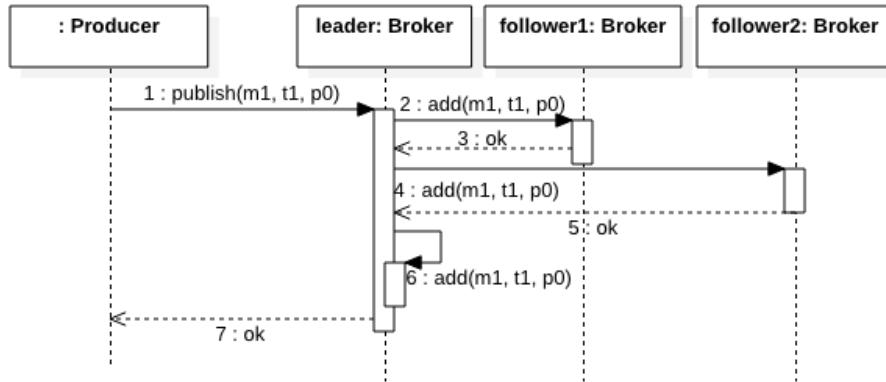


Figure 13: Sequence diagram Kafka producer.

A possible **issue**: in case of failure, the **producer may not get the response** (message number 7 in figure). In this case, the producer has to resend the message and kafka brokers can identify and eliminate duplicates.

Synchronization with replicas can be transactional and it's possible to choose between the following options:

- **Exactly-once** semantics is possible but long waiting time. So **replicas are not allowed**, but the problem is that Kafka spent a **long time trying to guarantee uniqueness**.
- **At-least-once** can be chosen by excluding duplicates' management.
- **At-most-once** can be chosen by publishing messages asynchronously.

#### Consumer

Each **consumer** can rely on a **persistent log** to keep track of the **offset** so that it is not lost in case of failure.

**Issue** case: if the consumer fails after having elaborated messages and before storing the new offset in the log, the same messages will be retrieved again (**at-least-once semantics**). Note that the delivery semantics can be changed if

the new offset is stored before the elaboration and we can choose **at-most-once semantics** because, if failing after storing the offset, the effect of the received messages does not materialize. Finally, transactional management of the log also allows for **exactly-once semantics**.

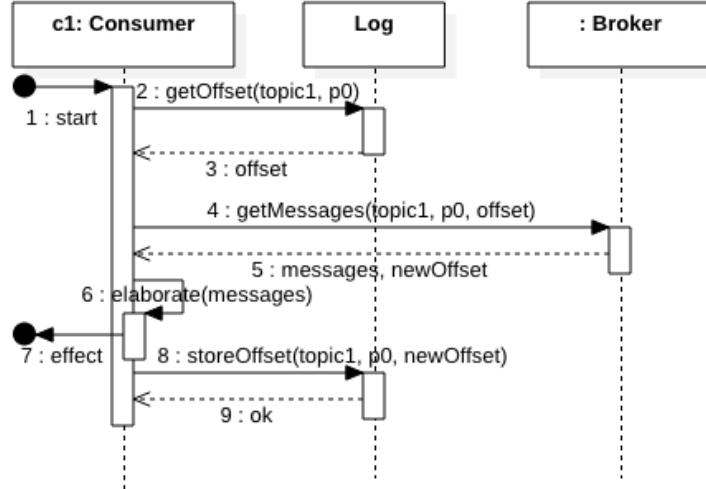


Figure 14: Sequence diagram Kafka consumer.

## Kafka architectural tactics

There are some tactics used to improve some features of Kafka. In the following section we can see scalability and fault tolerance.

### Improve Scalability

By **creating multiple partitions and multiple brokers**, we can create the ability to distribute producers/consumers to different partitions handled by different brokers. We can also **scale the operations** because Kafka supports the **creation of clusters of brokers**. Consider that each cluster contains up to a hundred brokers capable of handling trillions of messages per day.

### Improve Fault Tolerance

By **creating partitions**, we use the **persistence** of the partitions. **Replication** also reduces the risk of data loss. Finally, cluster management takes care of restarting brokers and setting leaders as needed.

## 5.6 Data-Intensive applications

Before we introduce the architectural styles for data-intensive applications, we explain the difference between batch and stream processing.

**Batch processing** is a method of running software programs called jobs in batches automatically. While users are required to submit the jobs, no other interaction by the user is required to process the batch.

**Stream processing** (also known as event stream processing, data stream processing, or distributed stream processing) is a programming paradigm which views streams, or sequences of events in time, as the central input and output objects of computation.

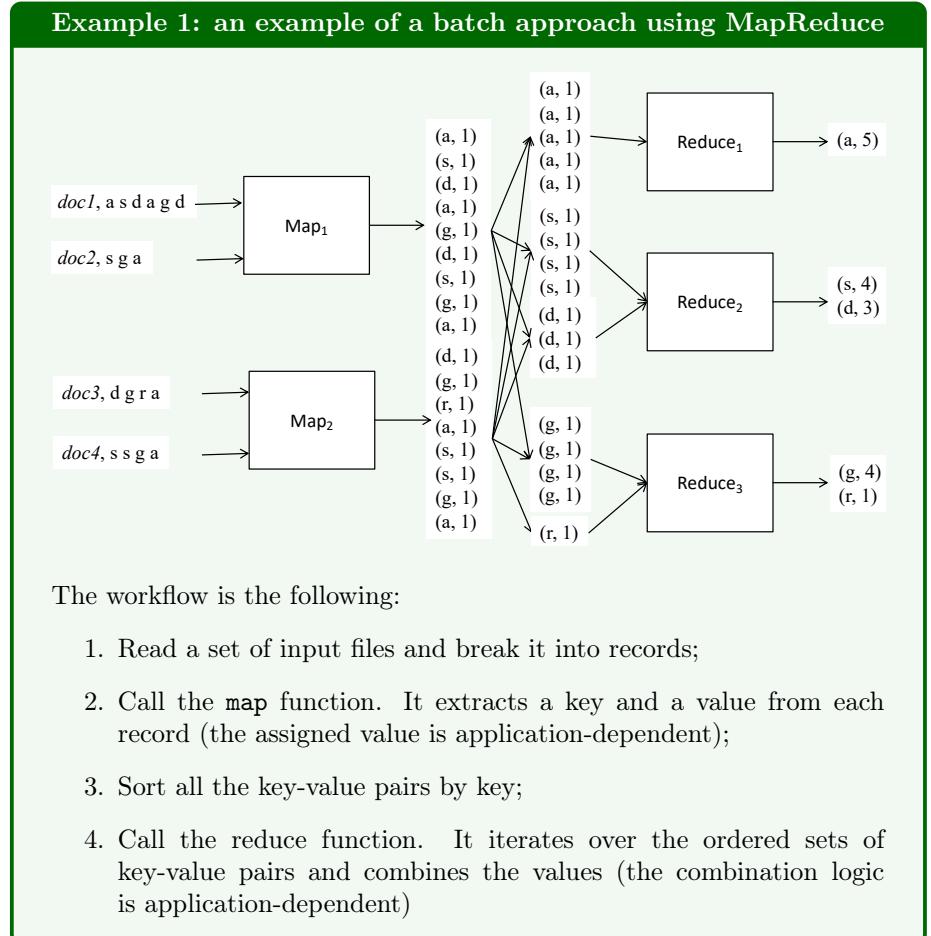
Batch	Stream
Has access to all data.	Computes a function of one data element, or a smallish window of recent data.
Might compute something big and complex.	Computes something relatively simple.
Is generally more concerned with throughput than latency of individual components of the computation.	Needs to complete each computation in near-real-time - probably seconds at most.
Has latency measured in minutes or more.	Computations are generally independent.
	Asynchronous - source of data doesn't interact with the stream processing directly, like by waiting for an answer.

Table 1: Batch vs Stream processing.

### 5.6.1 Batch approach: MapReduce

**MapReduce** is a **programming architecture** and an associated implementation for processing and generating big data sets with a parallel, distributed algorithm on a cluster.

A MapReduce is composed of a **map procedure**, which performs filtering and sorting (such as sorting students by first name into queues, one queue for each name), and a **reduce method**, which performs a summary operation (such as counting the number of students in each queue, yielding name frequencies). The “MapReduce System” (also called “infrastructure” or “framework”) orchestrates the processing by marshalling the distributed servers, running the various tasks in parallel, managing all communications and data transfers between the various parts of the system, and providing for redundancy and fault tolerance.



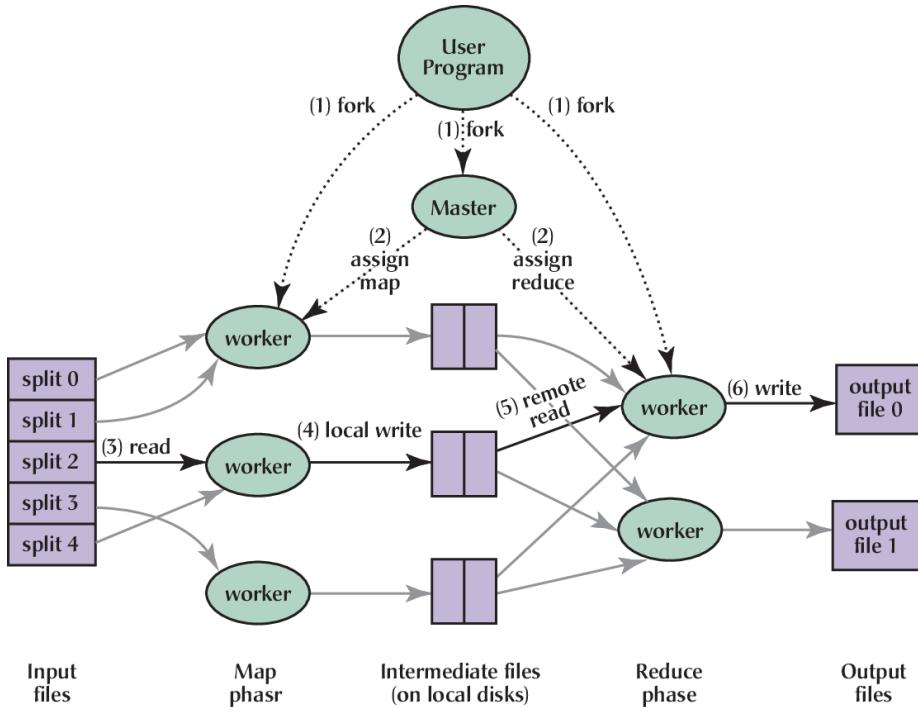


Figure 15: MapReduce architecture.

### ✓ Advantages

- Works well on commodity hardware.<sup>1</sup>

### ⚠ Disadvantages

- Implementing a complex processing job is not simple (high level programming model have been built on top of it);
- Reducers have to wait until the preceding Mappers have concluded their job;
- Materialization of intermediate states can be overkilling;
- Sometimes it is not necessary to sort the results of mappers;
- New batch computation approaches supported by frameworks as Spark, Tez, Flink, etc.

<sup>1</sup>Commodity hardware in computing is computers or components that are readily available, inexpensive and easily interchangeable with other commodity hardware. Almost all PCs use commodity hardware.

### 5.6.2 Stream approach: Apache Storm

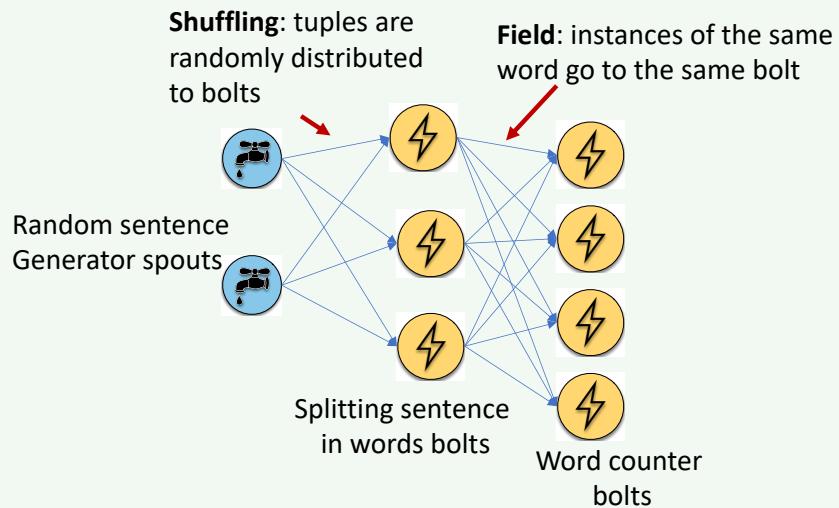
**Apache Storm** is a **distributed stream processing computation framework** written predominantly in the Clojure programming language. Originally created by Nathan Marz and team at BackType, the project was open sourced after being acquired by Twitter. It uses custom created “spouts” and “bolts” to define information sources and manipulations to allow batch, distributed processing of streaming data.

Some features:

- Support stream processing.
- More than 1 million messages per second per node.
- Can scale up to thousands of nodes per cluster.
- Expects and manages failures (fully fault tolerant).
- Provides guaranteed message delivery with exactly once semantics (reliable).

A Storm application is designed as a “topology” in the shape of a **directed acyclic graph** (DAG) with **spouts** (source of streams) and **bolts** (receives messages) acting as the graph vertices. **Edges on the graph are named streams** and direct data from one node to another. Together, the topology acts as a data transformation pipeline. At a superficial level the general topology structure is similar to a MapReduce job, with the main difference being that data is processed in real time as opposed to in individual batches. Additionally, Storm topologies run indefinitely until killed, while a MapReduce job DAG must eventually end.

Stream Grouping	Description
<b>Shuffle</b>	Sends messages to bolts in random, round robin sequence. Use for atomic operations, such as math.
<b>Fields</b>	Sends messages to a bolt based on one or more fields in the tuple. Used to segment an incoming stream and to count tuples of a specified type with a specified value.
<b>All</b>	Sends a single copy of each message to all instances of a receiving bolt. Use to send a signal, such as clear cache or refresh state, to all bolts.
<b>Custom</b>	Customized processing sequence. Use to get maximum flexibility of topology processing based on factors such as data types, load, and seasonality.
<b>Direct</b>	Source decides which bolt receives a message.
<b>Global</b>	Sends messages generated by all instances of a source to a single target instance. Use for global counting operations.

**Example 2: example of topology with different groupings**

### 5.6.3 Combining batch and stream: Lambda Architecture

**Lambda architecture** is a **data-processing architecture** designed to handle massive quantities of data by taking advantage of both batch and stream-processing methods.

This approach to architecture attempts to balance latency, throughput, and fault-tolerance by using batch processing to provide comprehensive and accurate views of batch data, while simultaneously using real-time stream processing to provide views of online data. The two view outputs may be joined before presentation.

The rise of lambda architecture is correlated with the growth of big data, real-time analytics, and the drive to mitigate the latencies of map-reduce.

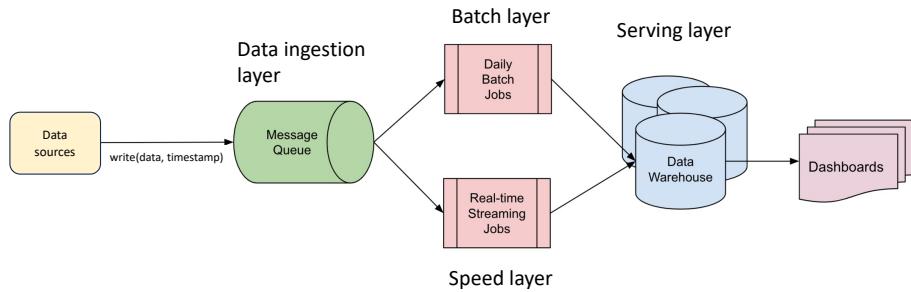


Figure 16: Lambda architecture.

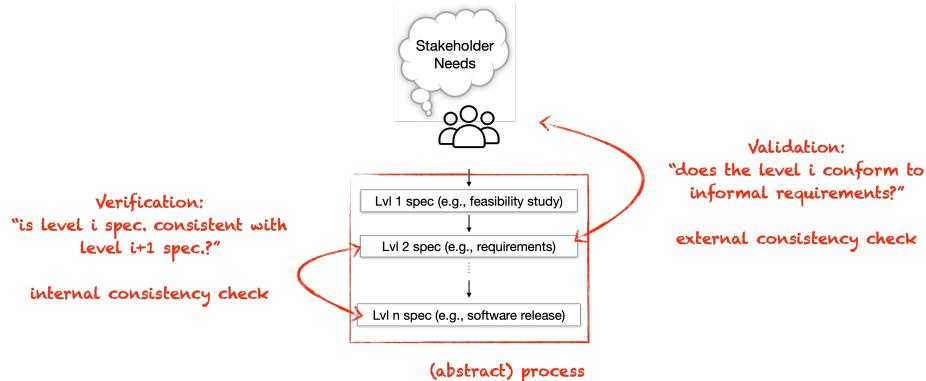
Exist also **Kappa architecture**. Kappa architecture is a software architecture used for processing streaming data with a single technology stack. It is a simplification of Lambda architecture, where the data is processed in batches. Kappa architecture ingests data into a messaging system like Apache Kafka, and performs both real-time and batch processing, especially for analytics, on the same stream. It allows for recomputation on the data by streaming it through the pipeline again.

## 6 Verification and Validation

### 6.1 Terminology

There are some differences between the terms verification and validation.

First of all, the verification is internal. Despite the validation is external. Assuming an abstract process with the following levels:



The **verification** is intended as: “Is level  $i$  consistent with level  $i + 1$ ?” It’s an **internal consistency check**. The **validation** is: “Does level  $i$  conform to needs?” It’s an **external consistency check**.

The [PMBOK guide](#), also adopted by the [IEEE](#) as a standard, defines them as follows in its 4th edition:

- **Validation.** The assurance that a **product, service, or system meets the needs of the customer and other identified stakeholders**. It often involves acceptance and suitability with external customers. Contrast with verification.
- **Verification.** The evaluation of whether or not a **product, service, or system complies with a regulation, requirement, specification, or imposed condition**. It is often an internal process. Contrast with validation.

Another fundamental topic when we speak about verification and validation is **Quality Assurance (QA)**. It **defines the policies and processes to achieve quality**. So it can **judge the quality and find defects**.

A direct **consequence** of the QA is the **improvement of the quality**. With the term “quality”, we refer to an ideal absence of defects (impossible) and an absence of other issues that prevent the fulfilment of non-functional requirements or the degradation of some software qualities.

Since it is impossible to have zero defects, a **periodic quality assurance evaluation is critical**. Ideally, every artefact shall be the subject of QA; even the verification artefacts must be verified!

The **V-model** is a **graphical representation of a systems development lifecycle**. It is used to produce rigorous development lifecycle models and project management models. It describes the activities and the results that must be made during product development.

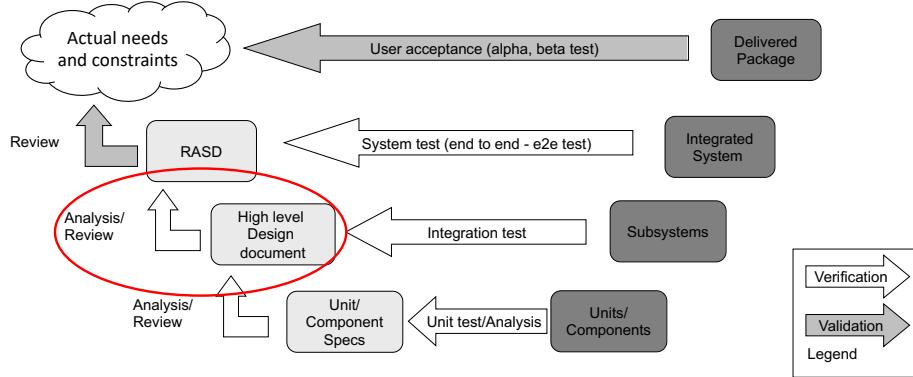


Figure 17: The V-model; verification is emphasized on the left.

The *left side* of the “V” represents the decomposition of requirements and the creation of system specifications. The *right side* of the “V” represents an integration of parts and their validation.

We have presented the V-model to help you understand where the verification can be placed.

Now, the **verification** is concerned with the code and the architecture. Considering the **software** side, it has two possible approaches:

- **Static Analysis.** It is **done using source code or other software artefacts but *without execution*.** Note that the analysis is static, but the properties are dynamic.
- **Dynamic Analysis (Testing).** It is **done by executing the sources.** The analysis is made by **comparing** the actual behaviour and the expected one.

On the other hand, to verify the **architectural level**, it is necessary to consider some aspects:

- The **structure must be consistent**. Some **examples**:
  - For every required interface, a corresponding provided interface exists.
  - Sequence diagrams are consistent with component diagrams and with the defined interfaces.
  - Each component has one or more modules that implement it.
- All **functional requirements must have the possibility to be satisfied**. Some **examples**:

- Each requirement is mapped on one or more components.
  - Each use case event flow is detailed in terms of one or more sequence diagrams.
- **Concurrent use of resources must be correctly defined.** Problems like order violation or a deadlock are expected. Some techniques must be applied to analyze these problems.
  - **Non-functional requirements must have the possibility to be fulfilled.**

#### Definition 1

**DevOps** is a set of practices, tools, and a cultural philosophy that automate and integrate the processes between software development and IT teams. It emphasizes team empowerment, cross-team communication and collaboration, and technology automation.

### 6.1.1 Study concurrent use of resources at architectural level (PT Net)

It is necessary to model distributed systems to study the concurrent use of resources at the architectural level.

A **Petri Net (PT Net or P/T Net)**, a place/transition net (PT net), is one of several **mathematical modelling languages used to describe distributed systems**. Like industry standards such as UML activity diagrams, **Petri nets offer a graphical notation for stepwise processes** that include choice, iteration, and concurrent execution.

The Petri net uses a graphic tool. It is a bipartite-directed graph containing places (circles), transitions (bars), and directed arcs.

A Petri net is a four-tuple:

$$PN = \langle P, T, I, O \rangle \quad (1)$$

- $P$ : a **finite set of places**  $\{p_1, p_2, \dots, p_n\}$
- $T$ : a **finite set of transitions**  $\{t_1, t_2, \dots, t_s\}$
- $I$ : an **input function**  $(T \times P) \rightarrow \{0, 1\}$
- $O$ : an **output function**  $(T \times P) \rightarrow \{0, 1\}$

It's also possible to add another term called  $M^0$ , which is an **initial marking**  $P \rightarrow N$ :

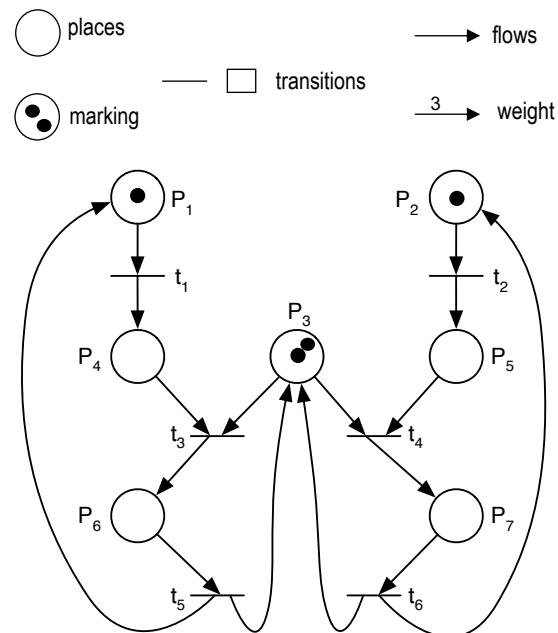
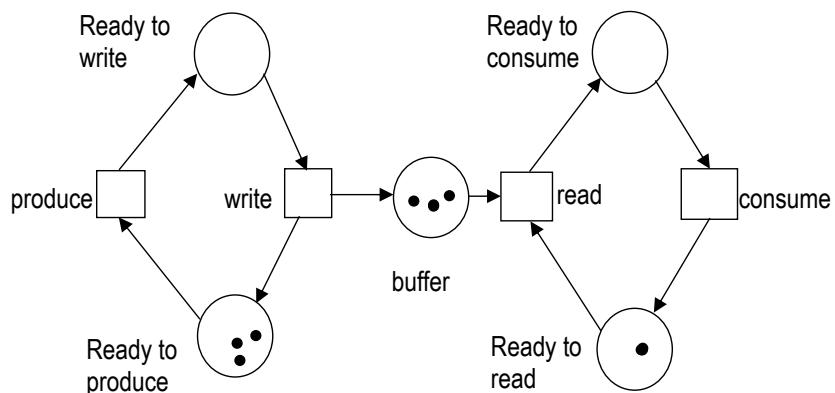
$$PN = \langle P, T, I, O, M^0 \rangle \quad (2)$$

Formula called also **marked Petri net**.

You can find a detailed explanation [here](#). Some observations of the Petri net:

- In a given marking  $M$ , a transition  $t$  can fire only if it is enabled.
- An enabled transition not necessarily fires.
- More than one transition can be enabled in a marking.
- If two transitions are enabled at the same time:
  - Which one fires first is not determined;
  - Petri nets are an intrinsically nondeterministic model;
  - The firing of a transition might disable another enabled transition.

In fact, if two transitions are enabled at the same time, they can fire simultaneously unless the firing of one transition disables the other. Petri nets are **suitable for modelling concurrent systems**. On [this page](#) you can see a live simulation by clicking on the nodes.

Figure 18: **Example** of Petri nets.Figure 19: **Example** of Petri nets of producer-consumer model with unbounded buffer.

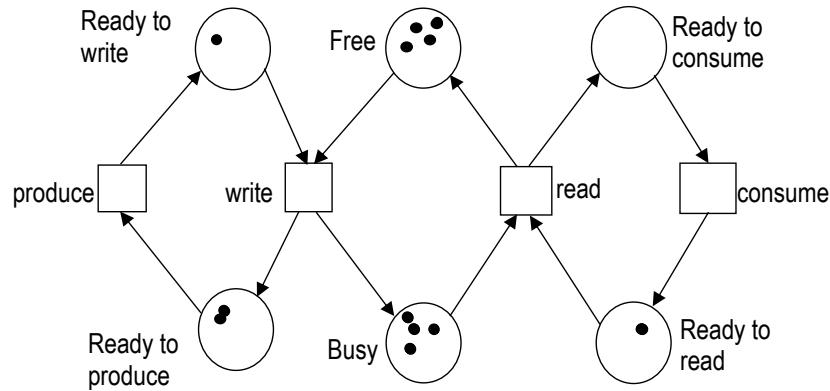


Figure 20: **Example** of Petri nets of producer-consumer model with finite buffer with a parametric number of positions.

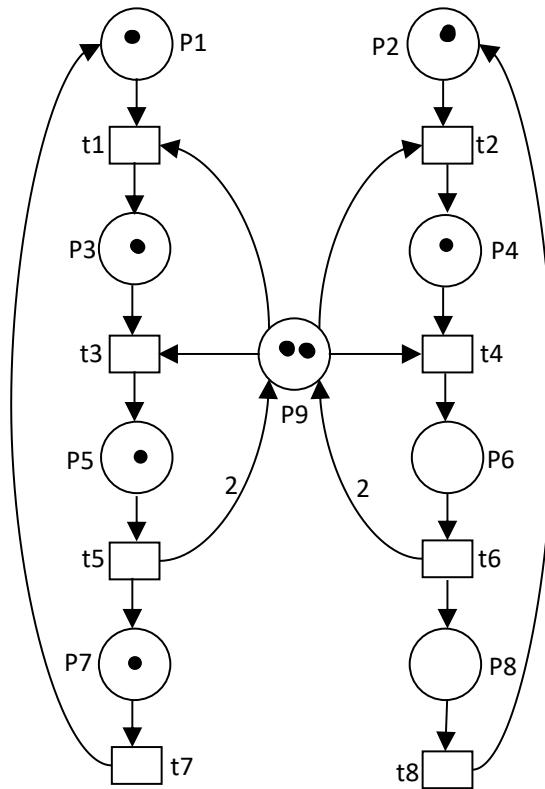


Figure 21: **Example** of Petri nets of deadlock.

### 6.1.2 Quantitative impact of architectural decisions

Architectural choices directly influence several software qualities (e.g., scalability, reliability, availability, usability).

To cope with this, we need metrics to quantify qualities and specific methodologies to analyze the quantitative impact of architectural choices on these qualities. The tactics are also foundational to address the issues.

First, before discussing *how to evaluate the quantitative impact of architectural decisions*, we must introduce the availability concept and explain a system life-cycle to introduce some exciting metrics.

#### 💡 Why is availability so important?

In general, a **service shall be continuously available** to the user, and if it fails after a bit of downtime, it should be a **rapid service recovery**. So the **availability** of a service depends on:

- The **complexity** of the **infrastructure** architecture.
- **Reliability** of the individual components.
- **Ability to respond** quickly and effectively to faults.
- **Quality of the maintenance** by support organizations and suppliers.
- **Quality** and scope of the **operational management** processes.

#### 📝 Study the System Life-Cycle to use availability metrics

The **System Life-Cycle** relates to failures in the following way:

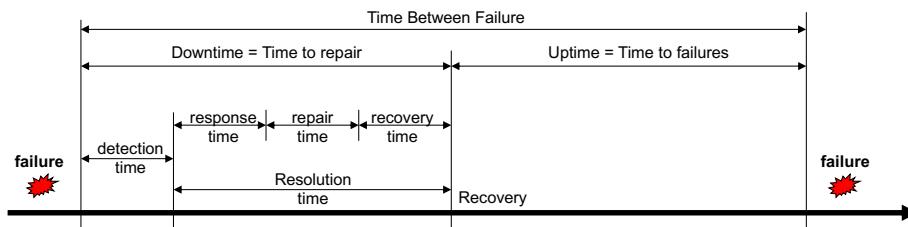


Figure 22: The System Life-Cycle when faults occur.

- **Time of occurrence.** Time at which the *user* becomes aware of the failure.
- **Detection time.** Time at which *operators* become aware of the failure.
- **Response time.** Time required by *operators* to diagnose the issue and respond to users.
- **Repair time.** Time required to fix the service/components that caused the failure.

- **Recovery time.** Time required to restore the system (re-configuration, re-initialization, ...).
- **Mean Time to Repair (MTTR).** Average time between the occurrence of a failure and service recovery, also known as the *downtime*.
- **Mean Time to Failures (MTTF).** Mean time between the recovery from one failure and the occurrence of the next failure, also known as *uptime*.
- **Mean Time Between Failures (MTBF).** Mean time between the occurrences of two consecutive failures.

### Definition 2: Availability Metric

The **availability metric** is the **probability** that a **component works correctly at time  $t$** . As a mathematician term, we can express this definition as the relationship between the Mean Time to Failures (MTTF) and the MTTF plus the Mean Time to Repair (MTTR):

$$A = \frac{\text{MTTF}}{\text{MTTF} + \text{MTTR}} \quad (3)$$

Note that if the Mean Time to Repair (MTTR) is small, then the Mean Time Between Failures (MTBF) is approximately equal to the Mean Time to Failures (MTTF):  $\text{MTBF} \approx \text{MTTF}$ .

### ⌚ Does an easier notation than the availability metric work?

The availability metric is crucial for understanding how a component works at a particular time, but sometimes, we need an easier notation to represent availability. In these cases, we use the **nines notation**.

**Nines** are an informal logarithmic notation for proportions very near to one or, equivalently, percentages very near 100%. Nines are the number of consecutive nines in a percentage such as 99% (two nines). The number of nines of a proportion  $x$  is:

$$\text{nines} = -\log_{10}(1-x) \quad (4)$$

In the computer system availability (our context), a *one nine* (90%) uptime indicates a system that is available 90% of the time or, more commonly described, unavailable 10% of the time.

Availability	Downtime
90% (1-nine)	36.5 days/year
99% (2-nines)	3.65 days/year
99.9% (3-nines)	8.76 hours/year
99.99% (4-nines)	52 minutes/year
99.999% (5-nines)	5 minutes/year

Table 2: Some nines notation and downtime values.

**💡 Now that we have the theory and the tools, what methodology should we use to analyze the impact of the architectural choices?**

The **Analysis Methodology** depends on the system. The Availability is calculated by **modelling the system** as an interconnection of elements in series and parallel:

- **Elements operating in series** mean that if one element fails, the whole combination fails.
- **Elements operating in parallel** mean that if a component fails, the other elements take over the operations of the failed element.

### 📘 Availability in series

The combined system is **operational only if every part is available**. Then, the combined Availability is the **product of the parts' Availability**.

$$A = \prod_{i=1}^n A_i \quad (5)$$

#### Example 1

We assume there is a system composed of two components with the following availability and downtime:

- Component 1 has 99% (2-nines) of availability and 3.65 days/year of downtime.
- Component 2 has 99.999% (5-nines) of availability and 5 minutes/year of downtime.

So the combined availability is 98.999% with 3.65 days/year of downtime.

The downtime is calculated using the following formula:

$$\text{Downtime} = (1 - A) \times 365 \text{ days/year}$$

Note that the  $A$  value is the Availability in terms of simple values and not as percentages (99% become 0.99).

In the previous example, we can see how the low Availability of Component 1 negatively affects the Availability of the entire system. This result means that **a chain is as strong as the weakest link**.

### ■ Availability in parallel

The combined system is **operational if at least one part is available**. Then, the combined Availability is  $1 - P$ , where  $P$  indicates all parts that are not available.

$$A = 1 - \prod_{i=1}^n (1 - A_i) \quad (6)$$

#### Example 2

We assume there is a system composed by two components with the following Availability and downtime:

- Component 1 has 99% (2-nines) of Availability and 3.65 days/year of downtime.
- Component 2 has 99% (2-nines) of Availability and 3.65 days/year of downtime.

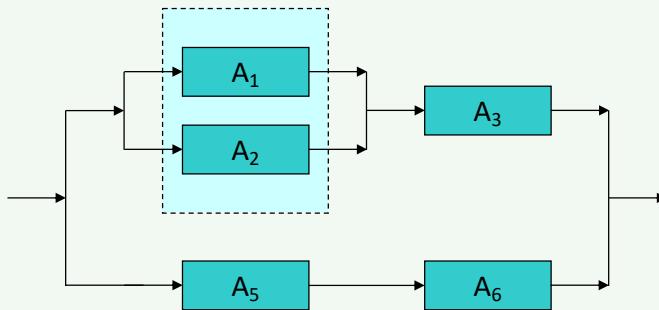
Despite the previous example, the combined availability is 99.99% (4-nines) with 52 minutes/year of downtime.

Even though components with very low Availability are used, the system's overall Availability is much higher than the Availability in series!

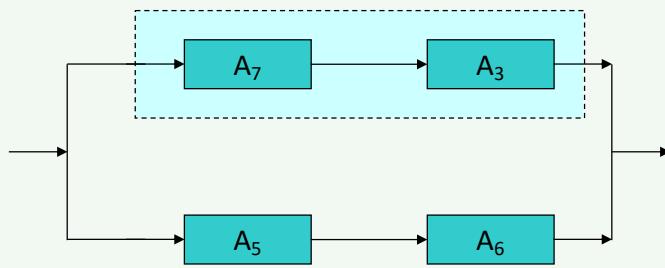
#### Example 3

Some examples of complex systems:

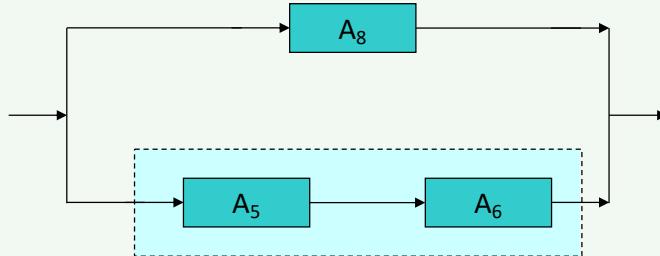
- $A_7 = 1 - (1 - A_1)(1 - A_2)$



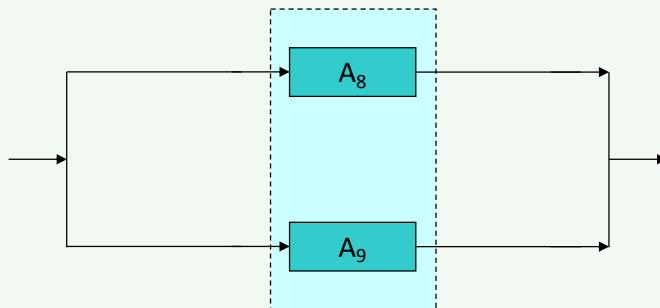
- $A_8 = A_7 A_3$



- $A_9 = A_5 A_6$



- $A = 1 - (1 - A_8)(1 - A_9)$



### ■ Tactics for Availability

As we explained in the past pages, Availability is crucial, but it's also fundamental to use intelligent **tactics to improve the quality of the attributes**.

#### Definition 3

The **Tactics** are design decisions that influence the control of one or more quality attributes.

Some well-known tactics are:

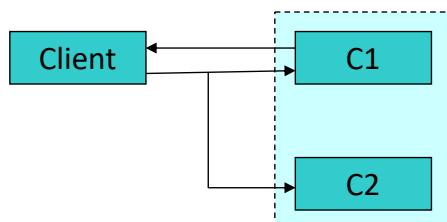
- **Replication**
- **Forward error recovery**
- **Circuit breaker**

### ■ Replication approaches

The **Replication** is very simple to manage in the case of stateless components. The approaches are different:

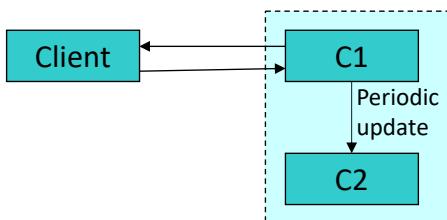
1. **Hot spare**: One component leads, and another is always ready to take over.

In the following example, C1 leads, C2 is always ready to take over.



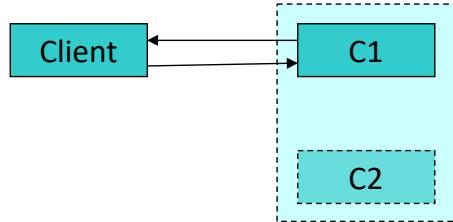
2. **Warm spare**: One component leads and periodically updates another component. If the primary component fails, the second component takes time to update itself fully.

In the following example, C1 leads and periodically updates C2. If C1 fails, some time might be needed to fully update C2.



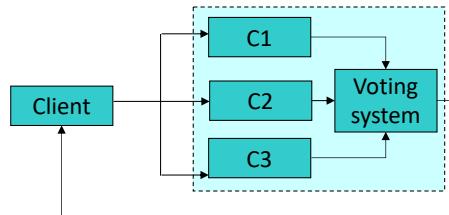
3. **Cold spare:** A second component is dormant, started, and updated only if required.

In the following example, C2 is dormant, started, and updated only if required.



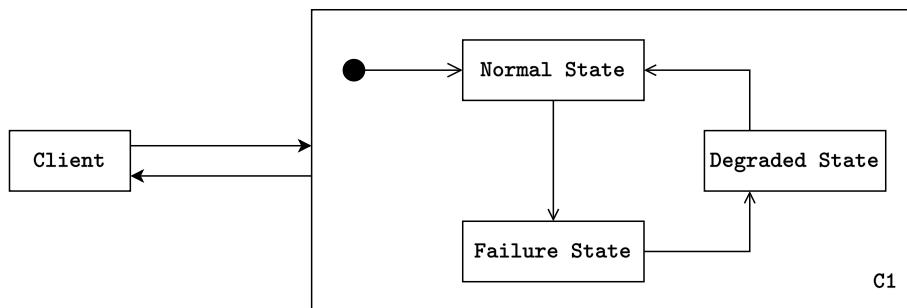
4. **Triple modular redundancy:** Three components are always active, and the result is the one produced by the majority. **This is good when reliability is also important.**

In the following example, C1, C2, and C3 are all active. The result is the one produced by the majority.



### Forward error recovery

**Forward Error Recovery** is a tactic in which a recovery mechanism moves the failed component to a degraded state. In a degraded state, a component continues to be available even if it is not fully functional. Here is an example:



### Circuit breaker

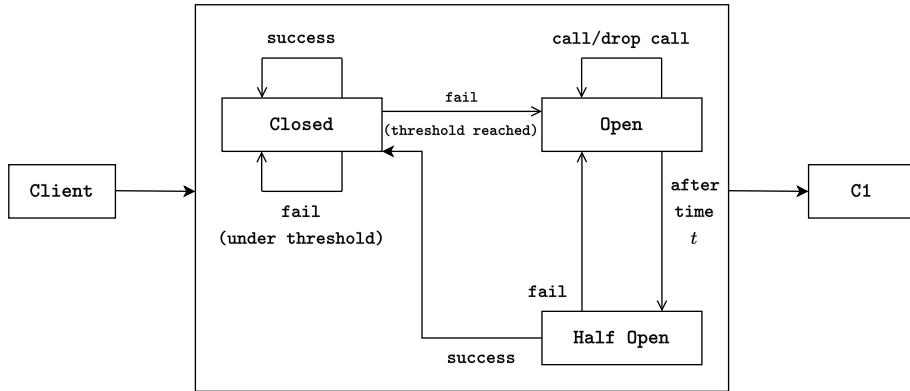
The **Circuit Breaker (CB)** tactic is a client-side resiliency pattern. The CB acts as a proxy for a remote component:

1. A component is called;
2. The CB monitors the call.

But note that there should be possible failures:

- CB receives an error;
- The call takes “too long” (CB kills the call).

If there are too many failures, the circuit breaker inhibits future calls by moving to the open state.



## 6.2 Analysis: Symbolic Execution

As we have discussed in the previous subsection, there are two types of approaches when we speak about verification:

- **Static Analysis.** It analyzes the source code, and each analyzer targets a fixed set of hard-coded (pre-defined, not custom) properties. It is entirely automatic, and the output reports two types of results: safe (no issues) and unsafe (potential problems). Also, the **analysis is made on generic (or symbolic) inputs.**

The properties that we have mentioned are safety properties, such as:

- No *overflow* for integer variables
- No *type errors*
- No *null-pointer* dereferencing
- No *out-of-bound* array accesses
- No *race conditions*
- No *useless assignments*
- No *usage of undefined variables*
- No *execution of specific paths*

- Testing (dynamic analysis) is made at runtime and is related to the software's behavior during execution. The analysis is also made on specific inputs.

Using the static analysis, we can use the symbolic execution.

### Definition 4

**Symbolic Execution** (also symbolic evaluation) **analyses a program to determine what inputs cause each part of a program to execute.**

The **symbolic execution** analyzes actual source code and **reachability** and **path feasibility** properties. It is automatic and may fail to explore all possible paths. Sometimes, it is used to support testing.

The checked properties by the static analysis can be of different types:

- **Reachability.** Does some program execution reach location  $L$  (generic line of code) in  $S$  (source code)? With the reachability property, the symbolic execution tries:

- To **verify** that  $L$  **cannot be reached**;
- Or **spots the condition under which  $L$  can be reached**.

For example, in the following code:

```

1 ...
2 k:      try {
3 k+1:      ...
4 L-1:    } catch (e) {
5 L:        /* error */
6 ...

```

Static analysis checks the reachability properties and verifies that  $L$  cannot be reached, or discovers the condition under which  $L$  can be reached.

- **Path Feasibility.** Is the given path  $p$  feasible? With the path feasibility property, the symbolic execution tries:

- To verify that  $p$  **cannot be executed**;
- Or **spots the condition under which  $p$  can be executed**.

Then  $p$  will be:

$$p = \langle 0, 1, \dots, k, \dots, n \rangle$$

Symbolic execution **executes programs on symbolic values**. Each symbolic value has its **symbolic states**, which keep track of the variables' (symbolic) values. The inputs are initialized with symbolic (generic) values.

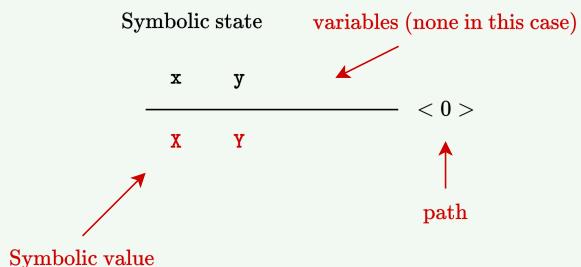
In the following example we can see a complete example of symbolic execution. But before we do, let us introduce some limitations of this methodology.

- The **path conditions may be too complex for constraint solvers**. Because solvers are very good at checking linear constraints, but it is harder for them to reason about non-linear arithmetic, bit-wise operations, string manipulation, etc.
- It is **impossible or difficult to use when the number of paths to be explored is infinite or huge**. For example, unbounded loops give rise to infinite sets of paths. Although the set of paths is finite, checking all loops is expensive and impractical.
- Finally, there may be **external code**. Then the sources are not available, such as a precompiled library, or the behavior is unknown to the solver.

#### Example 4

1. First we introduce the annotation:

```
1 void foo(int x, int y) {
2     ...
3 }
```



2. We introduce a local variable:

```
1 void foo(int x, int y) {
2     int z := x
3 }
```

Symbolic state

$$\frac{x \quad y \quad z}{\quad\quad\quad} < 0, 1 >$$

**X      Y      X**

3. We introduce a condition. A **path condition**  $\pi$  represents a constraint on a path:

```

1 void foo(int x, int y) {
2     int z := x
3     if (z < y)

```

- if condition true

Symbolic state

$$\frac{x \quad y \quad z \quad \pi}{\quad\quad\quad\quad} < 0, 1, 2 >$$

**X      Y      X    X < Y**

- if condition false

Symbolic state

$$\frac{x \quad y \quad z \quad \pi}{\quad\quad\quad\quad} < 0, 1, 2 >$$

**X      Y      X    X \geq Y**

4. **Execution continues along feasible paths.** In this case, the path condition  $\pi$  is satisfiable:

```

1 void foo(int x, int y) {
2     int z := x
3     if (z < y)
4         z := z*2

```

Symbolic state

$$\frac{x \quad y \quad z \quad \pi}{\quad\quad\quad\quad} < 0, 1, 2, 3 >$$

**X      Y      2X    X < Y**

5. Another if condition:

```

1 void foo(int x, int y) {
2     int z := x
3     if (z < y)
4         z := z*2
5     if (x < y && z >= y)

```

- if condition true

Symbolic state

$$\frac{x \quad y \quad z \quad \pi}{x \quad y \quad 2x \quad x < y} < 0, 1, 2, 3, 4 >$$

$$x < y \wedge 2x \geq y$$

- if condition false

Symbolic state

$$\frac{x \quad y \quad z \quad \pi}{x \quad y \quad 2x \quad x < y} < 0, 1, 2, 3, 4 >$$

$$x \geq y \vee 2x < y$$

## 6. Possible outcomes of symbolic execution:

```

1 void foo(int x, int y) {
2     int z := x
3     if (z < y)
4         z := z*2
5     if (x < y && z >= y)
6         print(z)
7 }
```

- (a) **Satisfiable** exit ( $\pi$  is satisfiable): every satisfying assignment to variables in  $\pi$  is an **input that satisfies the given property in a concrete execution**.

Symbolic state

$$\frac{x \quad y \quad z \quad \pi}{x \quad y \quad 2x \quad x < y} < 0, 1, 2, 3, 4, 5 >$$

$$x < y \wedge 2x \geq y$$

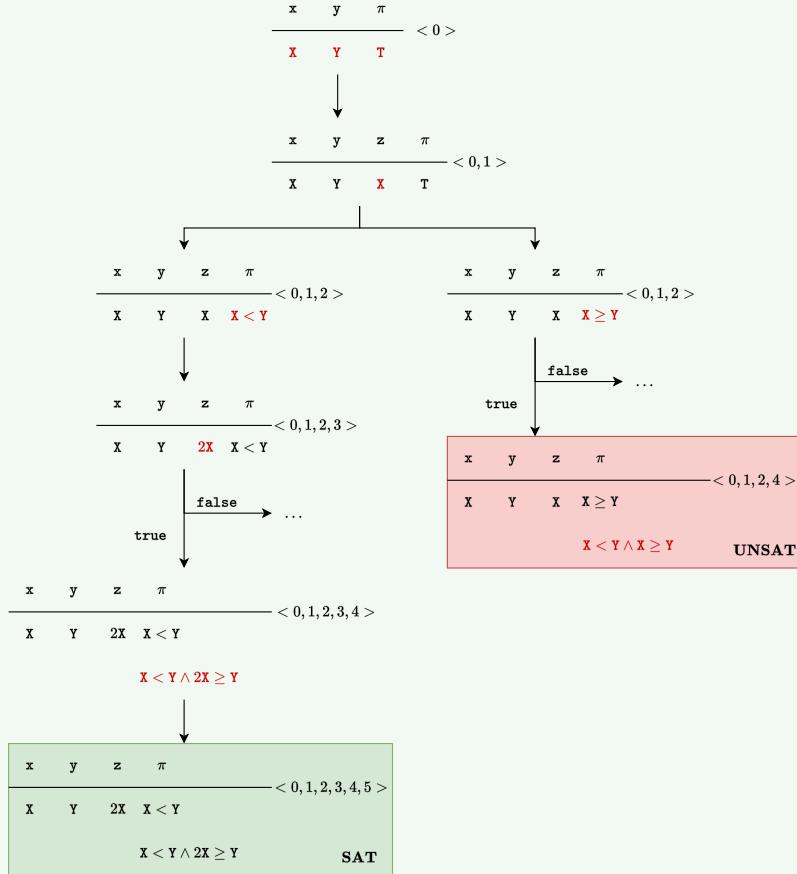
- (b) **Unsatisfiable** exit ( $\pi$  is not satisfiable): the given **property cannot be satisfied by any concrete execution**.

Symbolic state

$$\frac{x \quad y \quad z \quad \pi}{x \quad y \quad 2x \quad x < y} < 0, 1, 2, 4 >$$

$$x < y \wedge x \geq y$$

Finally, we can draw the **Execution Tree**. The execution paths can be collected in an execution tree, where end states are marked as SAT or UNSAT.



To view the tree in high resolution, scan (or click) the QR code below.



### 6.3 Testing: terminology, types of testing activities

**Testing** (dynamic analysis) is an approach to verification. The **main goal of testing is to make programs fail**.

Other *common goals* are:

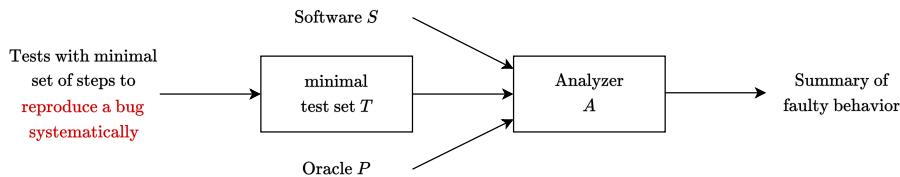
- Exercise different parts of a program to **increase coverage**;
- Make sure the **interaction between components works** (*integration testing*);
- Support **fault localization** and **error removal** (*debugging*);
- Ensure that **bugs introduced in the past do not happen again** (*regression testing*).

The dynamic analysis **analyzes program behavior**. The **properties** are **encoded as executable oracles** representing **expected outputs** and **desired conditions** (assertions).

It can **run only finite test cases**, so it's not exhaustive verification. The failures have **concrete inputs** that **trigger them**, and the **execution is automatic**.

#### ⌚ We have often heard about *debugging*, but what is it?

**Debugging** is a systematic approach to **fault localization** and **error removal**. The output is often used to support debugging.



#### ⌚ ... and *test case*?

##### Definition 5

A **Test Case** is a set of **inputs**, **execution conditions**, and a **pass/fail criterion**.

Running a test case typically involves setup, execution and teardown.

- **Setup.** Bring the program to an **initial state** that fulfils the execution conditions.
- **Execution.** **Run** the program on the actual inputs.
- **Teardown.** **Record** the output, the final state, and any **failure** determined based on the pass/fail criterion.

A **test set**, or **test suite**, can include **multiple test cases**. Finally, a **Test Case Specification** is a **requirement to be satisfied by one or more test cases**. An example of test case specification can be *the input must be a sentence composed of at least two words*, and an example of test case input is *this is a good test case input*.

When discussing test cases, it's necessary to introduce **Unit Testing**. This is **conducted by developers and aims to test small pieces (units) of code in isolation**.

However, when we test in isolation, there should be a **problem**: the **units may depend on other units**. Then, we need to simulate missing units.

The **Integration Testing** (integration of the unit tests) **aims to exercise the interaction between interfaces and components**. The **faults discovered by integration testing are multiple**; some examples:

- **Inconsistent interpretation of parameters** (e.g. mixed units meters or yards)
- **Violations of assumptions about domains** (e.g. buffer overflow)
- **Side effects on parameters or resources** (e.g. conflict on temporary file)
- **Nonfunctional properties** (e.g. unanticipated performance issues)
- **Concurrency-specific problems**

Typically, the integration test is defined by the Design Document. In the Design Document, we can find two types of plans:

- **Build Plan** that establishes the **order of the implementation**;
- A **Test Plan** that defines how to carry out integration testing is needed.

The strategies for the integration test are many:

- **Big Bang**: test only after integrating all modules (not even a real strategy).

#### ✓ Pros

It doesn't require stubs; it only **requires fewer drivers/oracles**.

#### ⚠ Cons

1. Minimum: observability, fault localization/diagnosability, efficacy, feedback;
2. **High cost of repair** (cost of repairing a fault increases as a function of time between the introduction of an error in the code and repair).

- **Iterative and incremental strategies.** The main action is **run after components are released (not just at the end)**. The strategy can be done in three different ways:

– **Hierarchical**. Based on the hierarchical structure of the system. It can be done top-down or bottom-up.

\* **Top-down strategy.** Work **from the top level** (in terms of “use” or “include” relationship) **down to the bottom level**. As modules are completed (according to the building plan), more functionality is testable. We also need to replace some stubs, and we need other stubs for lower levels. **When all modules are incorporated, the whole functionality can be tested.**

#### ✓ Pros

The drivers use the top level interfaces (e.g. REST APIs).

#### ⚠ Cons

This strategy requires stubs of used modules at each step of the process.

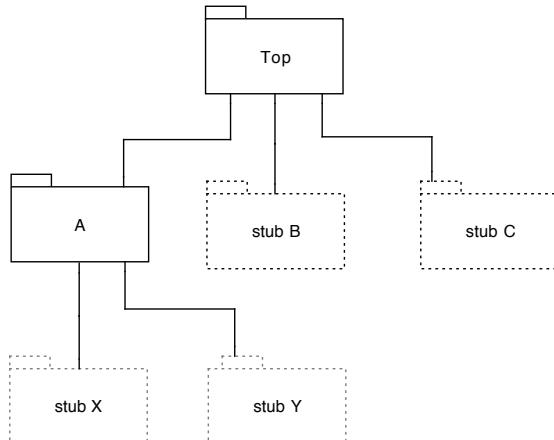


Figure 23: Example of top-down strategy.

#### ❓ What are the stubs?

A **Stub** is dummy code that acts on behalf of the original module without actually calling it during testing. Since all we need is the response, we simply call the dummy code each time, changing the response as needed for testing, and test how our current module behaves with those responses.

Conventionally, we might think of using stubs only when the module is still under development. However, stubs are very important in top-down integration testing, even after the developers have built the module. If we start using the developed module, we may not be able to tell whether the bug is in the submodule or the one we are testing.

- \* **Bottom-up strategy.** Starting from the leaves of the “uses” hierarchy.

✓ **Pros**

An advantage is that it **doesn't require stubs**.

⚠ **Cons**

**Typically requires more drivers** (one for each module, as in unit testing). Can this be a disadvantage? Maybe, because the newly developed module may replace an existing driver, and new modules require new drivers.

Another thing to consider is that **it may create several working subsystems**, and each working subsystem will eventually be integrated into the final one.

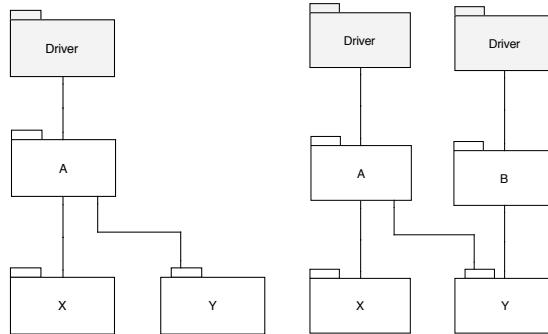


Figure 24: **Example** of bottom-up strategy.

❓ **What are the drivers?**

We work from the bottom up, and we are the smallest module in the first iteration that has no dependencies below it. But we may need the support of the modules above us to confirm that the response this module is sending up is correct.

A **Driver** is a **dummy code that sends the responses and acknowledgments (primarily not necessarily) to the submodules**. It helps us to identify the behavior of the submodule independently and quickly.

- **Threads.** A **thread** is a part of several modules that together provide a user-visible programme function. By using the thread strategy we can have some **advantages**.

✓ **Pros**

- \* We can **maximize the progress visible to the user** (or other stakeholders);
- \* **Reduce drivers and stubs;**
- \* An integration plan is usually more complex.

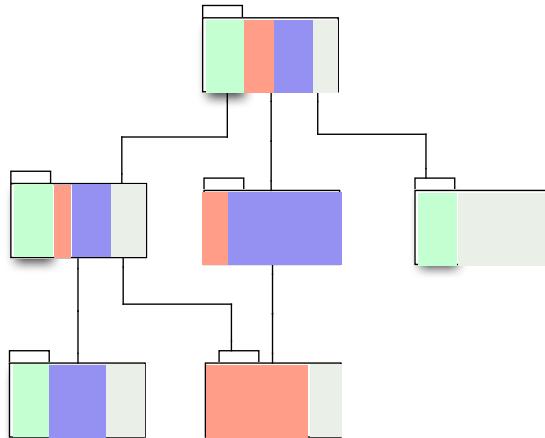


Figure 25: **Example** of threads strategy.

- **Critical.** The critical modules strategy **starts with the highest risk modules**. Risk assessment is a necessary first step. It can include technical risks (e.g. is X feasible?) and process risks (e.g. is the schedule for X realistic?). It may also be similar to a priority process.

The **key point of this strategy is the risk-oriented process**. Integration and testing as a risk mitigation activity, designed to *deliver any bad news as early as possible*.

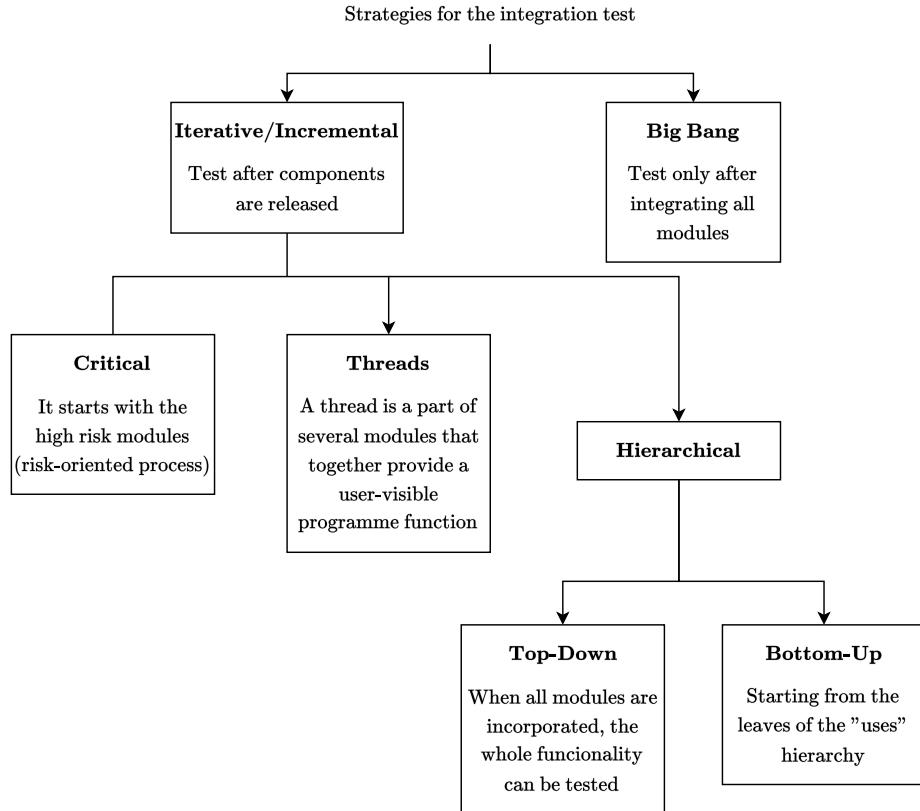


Figure 26: Summary of integration test strategies.

### ❸ Which one should we choose?

Given the three strategies above, *which one should we choose?* Well, the structural strategies (bottom-up or top-down) are simpler, but thread and critical modules provide better external visibility of progress (especially in complex systems).

So the **best choice** should be a **combination of different strategies**: - Use **top-down/bottom-up** for relatively **small components and subsystems**; - Combinations of **thread** and **critical module integration testing** for larger **subsystems**.

### 6.3.1 E2E Testing

#### Definition 6

**End-to-end (E2E)** testing is a software testing methodology to test a functional and data application flow consisting of several sub-systems working together from start to end.<sup>a</sup>

<sup>a</sup>[Engineering Fundamentals Playbook](#)

At times, these systems are developed in different technologies by different teams or organizations. Finally, they come together to form a functional business application. Hence, testing a single system would not suffice. Therefore, end-to-end testing verifies the application from start to end putting all its components together.

The following is a list of **common types of tests** that use the E2E system:

- **Functional Testing**

#### ❑ Purpose

Check whether the **software meets the functional requirements**.

#### ❓ How?

Use the software as described by use cases in the RASD (pag. 9), check whether requirements are fulfilled.

- **Performance Testing**

#### ❑ Purpose

1. Detect **bottlenecks** affecting response time, utilization, throughput
2. Detect **inefficient algorithms**
3. Detect **hardware/network issues**
4. Identify **optimization possibilities**

#### ❓ How?

Load the system with the expected workload and measure and compare acceptable performance.

- **Load Testing**

**❑ Purpose**

1. **Expose bugs** such as memory leaks, mismanagement of memory, buffer overflows
2. Identify **upper limits of components**
3. **Compare alternative architectural options**

**❑ How?**

Test the system at increasing workload until it can support it, and load the system for a long period.

- **Stress Testing**

**❑ Purpose**

Make sure that the **system recovers gracefully after failure**.

**❑ How?**

Trying to break the system under the test by overwhelming its resources or by reducing resources.

For **example** double the baseline number for concurrent users/HTTP connections, or randomly shut down and restart ports on the network switches/routers that connect servers.

## 6.4 Test case generation

### 6.4.1 Introduction

**Testing Workflow** is a type of software testing that verifies that each software workflow accurately reflects the given business process. A workflow is a series of tasks to produce a desired result, usually involving several stages or steps. For any business process, testing these sequential steps is defined as “workflow testing”.

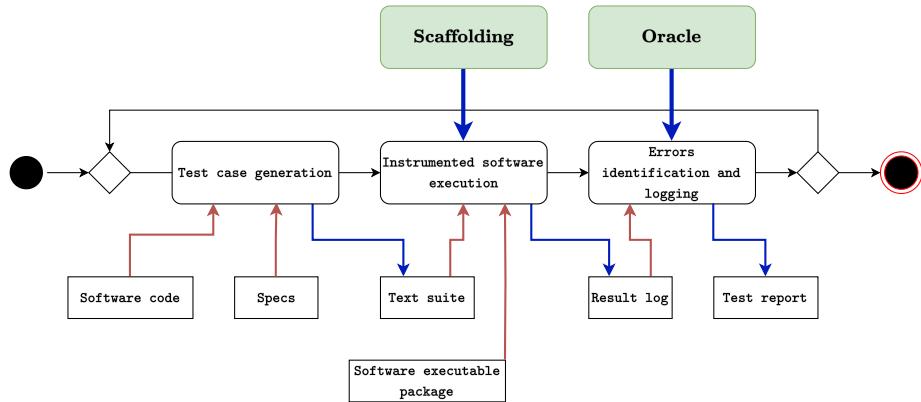


Figure 27: Testing Workflow.

To view the component diagram in high resolution, scan (or click) the QR code below.



As we can see in Figure 27, test cases can be generated in a **black-box** or **white-box** manner. The **White Box** is a generation based on code features. Meanwhile, the **Black Box** is a generation based on specification features.

Test case generation can be done manually (no need to explain) or automatically. Automatic generation can be done in several ways:

- **Combinatorial** testing. It enumerates all possible inputs according to some policy (e.g. smaller to larger).
- **Concolic Execution** (analyzed in the section 6.4.2 on page 83). It's a pseudo-random generation of inputs guided by symbolic path properties.
- **Fuzz** testing (*fuzzing*). It's a pseudo-random generation of inputs, including invalid, unexpected inputs.
- **Search-based** testing. It explores the space of valid inputs, looking for those that improve some metric (e.g. coverage, diversity, fault inducing capability).
- **Metamorphic** testing. Generates new test cases based on some metamorphic relationships and other previously defined test cases.

### 6.4.2 Concolic Execution

**Concolic Execution** (concrete-symbolic execution) is an **automatic generation of test cases**. It's a pseudo-random generation of inputs guided by symbolic path properties.

In other words, the concolic execution **performs symbolic execution** (definition on page 69) **alongside** concrete execution (**concrete inputs**). Under the hood, in concolic execution a **state** combines a symbolic part and a concrete part, used as needed to make progress in the exploration.

The steps are then as follows:

- Concrete  $\xrightarrow{\text{to}}$  Symbolic, derive conditions to explore new paths.
- Symbolic  $\xrightarrow{\text{to}}$  Concrete, simplifying conditions to generate concrete inputs.

Let's take an example to clarify the explanation.

#### Example 5

See the code below:

```

1 def m2(x: int, y: int):
2     z: int = bb(y) # black-box function
3     if z == x:
4         z = y + 10
5         if x <= z:
6             print("Log message.")
7 # end

```

Let's explore all the paths of the `m2` method, starting with a **(random) concrete input** and at the *same time* building the **symbolic condition of the explored path**. Unfortunately, in some cases we will not be able to solve the symbolic execution. For example, the behavior of the first if-condition (`z == x`) is unknown in the code. For this reason, we execute it with the identified input cases: given `y = 7`, run `bb(7)` and return 14. With this arrangement, the condition can be solved.

The annotation used will be the same as in the Symbolic Execution example on page 70.

1. Let's start with defined parameters and no condition state:

$$\begin{array}{cccc}
 x & y & z & \pi \\
 \hline
 X & Y & T & \\
 22 & 7 & &
 \end{array} < 0 >$$

2. Call the function with the parameter  $y$ , i.e. value 7. The return value of the called function will be 14.

$$\frac{\begin{array}{cccc} \mathbf{x} & \mathbf{y} & \mathbf{z} & \pi \\ \mathbf{X} & \mathbf{Y} & \mathbf{bb(Y)} & \mathbf{T} \end{array}}{\begin{array}{ccc} 22 & 7 & 14 \end{array}} < 0, 1 >$$

3. The if condition  $14 == 22$  is obviously false, so the path will be 6 and 7. By the way, the condition can also be seen as a logical condition  $\neg(bb(Y) \neq X)$ . But it can't be solved, so the concolic execution goes from symbolic to concrete  $\neg(14 \neq X) \equiv 14 = X$ .

$$\frac{\begin{array}{cccc} \mathbf{x} & \mathbf{y} & \mathbf{z} & \pi \\ \mathbf{X} & \mathbf{Y} & \mathbf{bb(Y)} & \mathbf{bb(Y) \neq X} \end{array}}{\begin{array}{ccc} 22 & 7 & 14 \end{array}} < 0, 1, 2, 6, 7 >$$

4. After a random concrete input, we can solve the constraint and start a new exploration. The parameter value of  $X$  is now 14. This is because we do not want the condition from the previous step to be *not equal*.

$$\frac{\begin{array}{cccc} \mathbf{x} & \mathbf{y} & \mathbf{z} & \pi \\ \mathbf{X} & \mathbf{Y} & & \mathbf{T} \end{array}}{\begin{array}{ccc} 14 & 7 \end{array}} < 0 >$$

5. The result of the function  $bb(7)$  is 14.

$$\frac{\begin{array}{cccc} \mathbf{x} & \mathbf{y} & \mathbf{z} & \pi \\ \mathbf{X} & \mathbf{Y} & \mathbf{bb(Y)} & \mathbf{T} \end{array}}{\begin{array}{ccc} 14 & 7 & 14 \end{array}} < 0, 1 >$$

6. Unlike before, the **if** condition is now true!

$$\frac{\begin{array}{cccc} \mathbf{x} & \mathbf{y} & \mathbf{z} & \pi \\ \mathbf{X} & \mathbf{Y} & \mathbf{bb(Y)} & \mathbf{bb(Y)=X} \end{array}}{\begin{array}{ccc} 14 & 7 & 14 \end{array}} < 0, 1, 2 >$$

7. The evaluation of the third line of code needs no explanation.

$$\frac{\begin{array}{cccc} \mathbf{x} & \mathbf{y} & \mathbf{z} & \pi \\ \text{X} & \text{Y} & \text{Y+10} & \mathbf{bb(Y)=X} \\ 14 & 7 & 17 & \end{array}}{<0,1,2,3>} \quad \mathbf{bb(Y)=X}$$

8. The path just explored explores the whole code. To try another path, we can negate the condition, but be careful! We only need to negate the second condition, because we already negated the first condition in the first exploration.

$$\mathbf{bb(Y)=X \wedge X \leq Y+10 \implies bb(Y)=X \wedge X > Y+10}$$

$$\frac{\begin{array}{cccc} \mathbf{x} & \mathbf{y} & \mathbf{z} & \pi \\ \text{X} & \text{Y} & \text{Y+10} & \mathbf{bb(Y)=X} \\ 14 & 7 & 17 & \mathbf{X \leq Y+10} \end{array}}{<0,1,2,3,4,5,6,7>} \quad \mathbf{bb(Y)=X}$$

- 
9. Explore the new path using two (random) values that satisfy the previous logical condition.

$$\frac{\begin{array}{cccc} \mathbf{x} & \mathbf{y} & \mathbf{z} & \pi \\ \text{X} & \text{Y} & & \mathbf{T} \\ 34 & 17 & & \end{array}}{<0>} \quad \mathbf{T}$$

10. The result of the `bb` function is 34.

$$\frac{\begin{array}{cccc} \mathbf{x} & \mathbf{y} & \mathbf{z} & \pi \\ \text{X} & \text{Y} & \mathbf{bb(Y)} & \mathbf{T} \\ 34 & 17 & & 34 \end{array}}{<0,1>} \quad \mathbf{bb(Y)=T}$$

11. The `if` condition is obviously true.

$$\frac{\begin{array}{cccc} \mathbf{x} & \mathbf{y} & \mathbf{z} & \pi \\ \text{X} & \text{Y} & \mathbf{bb(Y)} & \mathbf{bb(Y)=X} \\ 34 & 17 & 34 & \end{array}}{<0,1,2>} \quad \mathbf{bb(Y)=X}$$

12. We add 10 to the variable Y.

x	y	z	$\pi$	
X	Y	<b>Y+10</b>	$\text{bb}(Y)=X$	$< 0, 1, 2, 3 >$
34	17	<b>27</b>		

13. In this case the **if** condition is false and we complete the exercise because we have explored every possible path!

x	y	z	$\pi$	
X	Y	Y+10	$\text{bb}(Y)=X$	$< 0, 1, 2, 3, 4, 6, 7 >$
34	17	27	<b>X &gt; Y+10</b>	

### ✓ Advantages

The concolic execution has two main advantages:

1. Can handle **black-box functions in path conditions** (not possible with symbolic execution!).
2. Can automatically generate concrete test cases according to a **code coverage criterion**.

### ⚠ Limitations

And the disadvantages are:

1. It will only **find one input example per path**. And this could be a problem, because typically the errors only occur with certain inputs. What's more, if the errors are infrequent events, it's difficult to detect them with concolic execution.
2. The number of paths explodes due to complex nested conditions, then it **requires a large search space**.
3. It **doesn't guide the exploration**, it just explores possible paths one by one, as long as we have the budget (e.g. time, number of runs).

### 6.4.3 Concurrent systems testing

There are many difficulties in testing concurrent software. For example, the **concurrency bugs are non-deterministic** and can only manifest themselves within certain *interleavings*. Furthermore, the interleavings depend on execution conditions that are not under the direct control of the program.

#### Example 6

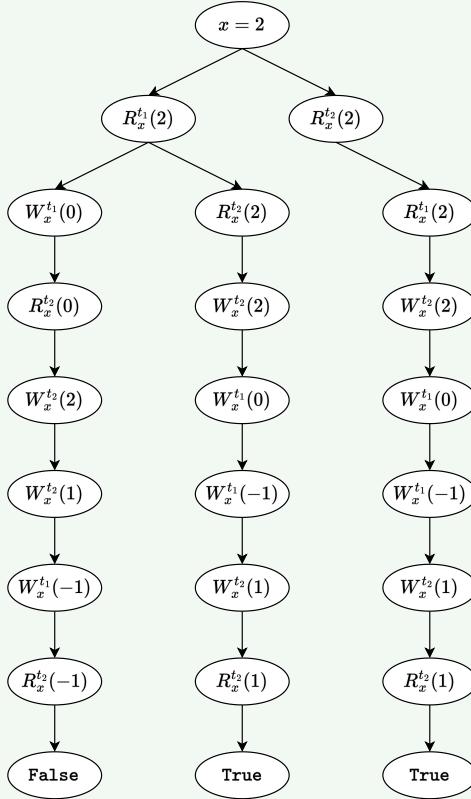
The following code:

```

1 int x = 2;
2 begin t_1           begin t_2:
3 if (x >= 0)         if (x >= 0)
4   x = 0;             x = 2;
5 x = -1;             x = 1;
6 end t_1             result = array[x];
7                     end t_2
8                     assert(result != null)

```

Can have the following possible execution behaviors:



**Example 7: an example of data race**

```

1 begin f_1           begin f_2
2 x = 'AAAAAAA'     x = 'BBBBBBB'
3 end f_1            end f_2

```

**Example 8: an example of atomicity violation**

```

1 begin f_1           begin f_2
2 if (x > 0)         x = 0
3   x = x - 1        end f_2
4 end f_1

```

**Example 9: an example of deadlock**

```

1 begin f_1           begin f_2
2 acquire(L1)         acquire(L2)
3   acquire(L2)       acquire(L1)
4   ...
5   release(L2)      ...
6   release(L1)      release(L1)
7 end f_1             end f_2

```

**⌚ So how do you test concurrent programmes?**

**Metamorphic Testing (MT)** is a property-based software testing technique that can be an effective approach to solving the test oracle problem<sup>2</sup> and the test case generation problem (this chapter). The test oracle problem is the difficulty in determining the expected results of selected test cases, or whether the actual results match the expected results.

The MT technique is applicable to any type of software and has recently been adapted to detect ***data race*** problems in concurrent software.

The basic idea of general MT is to derive **Metamorphic Relations (MRs)** between multiple program inputs and corresponding outputs:

1. New test cases from existing ones
2. An oracle for the program

**Example 10**

Consider a program  $P$  computing the shortest path in an undirected graph  $G$ . We may express the following properties (MRs):

**MR1** The length of the shortest path between two points in the graph

<sup>2</sup>In software testing, a **Test Oracle** (or simply oracle) is a provider of information that describes the correct output based on the input of a test case.

doesn't vary if we swap the two points:

$$|P(G, a, b)| = |P(G, b, a)|$$

**MR2** If  $c$  belongs to the shortest path between  $a$  and  $b$ :

$$|P(G, a, c)| + |P(G, c, b)| = |P(G, a, b)|$$

Now we can:

1. Pick a source test case  $\langle a1, b1 \rangle$  selected with any technique;
2. Run the software and obtain the result, that is, the shortest path and its length
3. Based on MR1, generate a follow up test case  $\langle b1, a1 \rangle$
4. Run the software with the new test case
5. Check if MR holds:

$$|P(G, a1, b1)| = |P(G, b1, a1)|$$

If not, then  $P$  is failing.

The **test case** for concurrent programmes can be represented as follows:

$$C = (X, S)$$

- Where  $X$  is a set of values for the input parameters.
- $S$  is a series of read and write events on the variable shared between different threads/processes.

The Metamorphic Relations (MRs) are relations between several programme inputs + series of read/write events and corresponding outputs.

#### Example 11: example of a test case that raises some doubts

An example of a test case  $p_{\text{array}}$ :

$$\begin{aligned} C = ( & \{2\}, \\ & \langle R_x^{t1}(2), W_x^{t1}(0), R_x^{t2}(0), \\ & \quad W_x^{t2}(2), W_x^{t1}(-1), W_x^{t2}(1), R_x^{t2}(1) \rangle ) \end{aligned}$$

For the following concurrent program:

```

1 int x = 2;
2 begin t_1      begin t_2:
3 if (x >= 0)    if (x >= 0)
4   x = 0;        x = 2;
5 x = -1;        x = 1;
6 end t_1        result = array[x];
7                      end t_2
8                      assert(result != null)

```

Let's assume that this is the source test case. Intuitively, our MR is that the output remains the same for the same inputs and different read/write sequences.

So we define follow-on test cases based on the MR and run the concurrent program, making sure that the selected read/write sequences occur. If the output changes, the test fails.

From the previous example, there are a number of questions:

- *How to select the read/write sequences for the subsequent test cases?*

For example, we can select those that may raise specific issues.

Data Access Pattern	Description
1. $R_u(l) W_{u'}(l) W_u(l)$	Value read is stale by the time an update is made in $u$ .
2. $R_u(l) W_{u'}(l) R_u(l)$	Two reads of the same location yield different values in $u$ .
3. $W_u(l) R_{u'}(l) W_u(l)$	An intermediate state is observed by $u'$ .
4. $W_u(l) W_{u'}(l) R_u(l)$	Value read is not the same as the one written last in $u$ .
5. $W_u(l) W_{u'}(l) W_u(l)$	Value written by $u'$ is lost.
6. $W_u(l_1) W_{u'}(l_1) W_{u'}(l_2) W_u(l_2)$	Memory is left in an inconsistent state.
7. $W_u(l_1) W_{u'}(l_2) W_{u'}(l_1) W_u(l_2)$	Same as above.
8. $W_u(l_1) W_{u'}(l_2) W_u(l_2) W_{u'}(l_1)$	Same as above.
9. $W_u(l_1) R_{u'}(l_1) R_{u'}(l_2) W_u(l_2)$	State observed is inconsistent.
10. $W_u(l_1) R_{u'}(l_2) R_{u'}(l_1) W_u(l_2)$	Same as above.
11. $R_u(l_1) W_{u'}(l_1) W_{u'}(l_2) R_u(l_2)$	Same as above.
12. $R_u(l_1) W_{u'}(l_2) W_{u'}(l_1) R_u(l_2)$	Same as above.
13. $R_u(l_1) W_{u'}(l_2) R_u(l_2) W_{u'}(l_1)$	Same as above.
14. $W_u(l_1) R_{u'}(l_2) W_u(l_2) R_{u'}(l_1)$	Same as above.

- How to control that the program runs according to the selected nesting?

Instrumenting the bytecode:

```
1 int x = 2;
2 begin t_1:
3 Scheduler.beforeRead();
4 if (x >= 0)
5     Scheduler.beforeWrite();
6     x = 0;
7 Scheduler.beforeWrite();
8 x = -1;
9 end t_1
10
11 begin t_2:
Scheduler.beforeRead();
if (x >= 0)
    Scheduler.beforeWrite();
    x = 2;
Scheduler.beforeWrite();
x = 1;
Scheduler.beforeRead();
y = array[x];
assert(y != null)
```

## 6.5 Testing frameworks

In software development, we typically use **unit testing frameworks** such as the xUnit frameworks (for example, JUnit and NUnit), which allow us to run unit tests to determine whether different parts of the code behave as expected under different circumstances.

The elements of unit test frameworks:

- **Test Runner:** it's a *component that orchestrates the execution of tests and delivers the result to the user*. The runner can use a graphical interface, a textual interface or return a special value to indicate the results of the execution of the tests.
- **Test Case:** in most cases this *is a class* from which our application-specific code inherits.
- **Test Fixture:** it represents the *preparation needed to set up the initial state required for a test case before the test, and to return to the original state after the test*.
- **Test Suite:** this is a *collection of test cases that share the same fixture*.
- **Assertions:** the *functions/macros that check the state or output of the system under test* (oracles).

Some examples of frameworks: for Java are [JUnit](#) ([here is an example](#)) and [Apache JMeter](#), for C++ are [Google Test](#) ([here](#) is an example of a factorial function), [CppUnit](#), [CxxTest](#), [Microsoft Unit Testing Framework for C++](#).

## 7 Software Configuration Management

### 7.1 Introduction

**Configuration Management (CM)** is a systems engineering process for establishing consistency of a product's attributes throughout its life. In the technology world, configuration management is an IT management process that tracks individual configuration items of an IT system. IT systems are composed of IT assets that vary in granularity. An IT asset may represent a piece of software, or a server, or a cluster of servers. The following focuses on configuration management as it directly applies to IT software assets and software asset CI/CD.

**Software Configuration Management** is a systems engineering process that tracks and monitors changes to a software systems configuration metadata. In software development, configuration management is commonly used alongside version control and CI/CD infrastructure (explained later).

The basic **approach to using a decentralized CM** is to have a repository (project) on the server side.

- When we want to work on the project, we clone the repository on the local PC. This workflow is used because we can work offline and on the local project without making critical changes to the repository server side.
- The local changes can be saved using the commit command and when we are ready to publish our changes, we use the push command to update the repository on the server side.
- After a push, anyone who has a local copy should make a pull command to update the local project.

This workflow can be done using git commands, and a good cheat sheet can be found [here](#).

## 7.2 Workflows

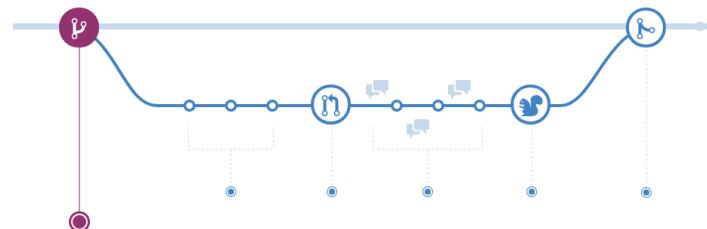
In Configuration Management (CM), it's important to define a workflow. A workflow then defines how an organization uses CM for a specific project.

The **main principles** are:

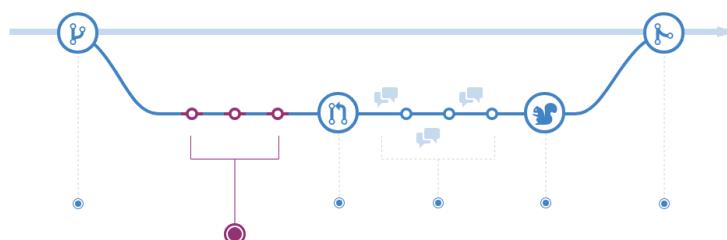
- Keep the master repository **as clean as possible**.
- The distinction between local and central repositories is not sufficient in the case of multiple contributors.
- On the central repository we can create a master repository and as many branches as we want.
- A new commit only becomes part of the master **if the whole team agrees**.

In detail, the [GitFlow](#) is:

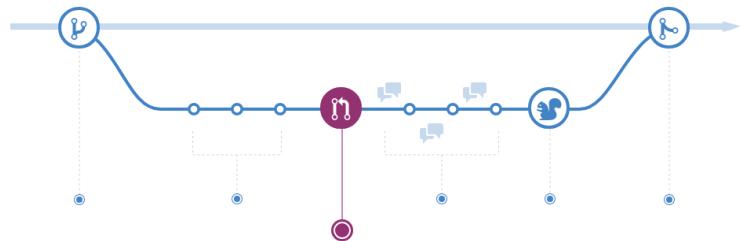
1. If we want to develop a new feature or idea, **create a new branch**. A branch can be thought of as a new, independent sub-repository, where we can experiment without affecting what is in the master.



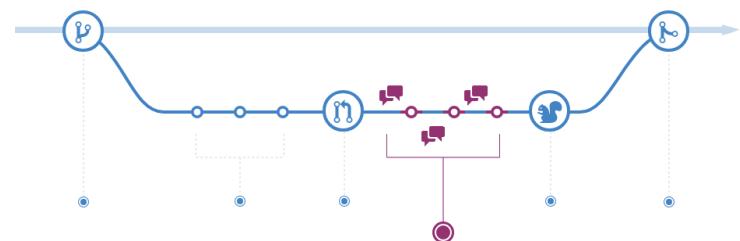
2. New **versions of software can be produced** in the branch.



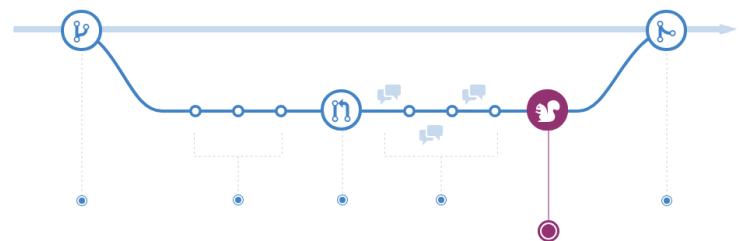
3. When the work is done, we can **make a pull request**. Then the sub-teams ask the whole team to review the changes and decide whether or not to accept them. The sub-team has already merged any local changes into the server side of the branch.



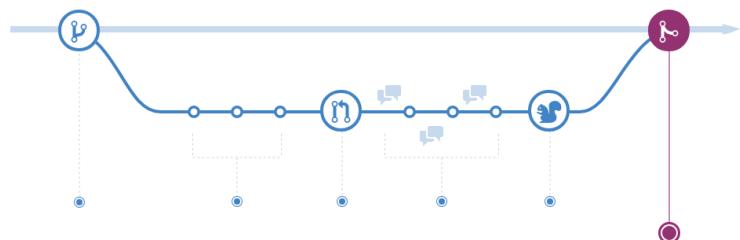
4. A **discussion starts**, and new versions can be created in the branch during the discussion.



5. Finally, the **code is deployed and tested**.



6. And the **branch is merged into the master**.



### 7.3 Continuous Integration and Continuous Delivery

**Continuous Integration (CI)** is the practice of integrating changes into the main codebase as often as possible. It relies on automation to ensure that proper checks are performed during integration. Typically, commits or pull requests trigger compilation or build, unit and integration testing, code quality checking, and merging into the main repository.

However, **Continuous Delivery (CD)** is different from continuous deployment. The former is the automatic preparation and tracking of a release into production. Instead, continuous deployment automatically deploys into the operational environment.

This theory can be put into practice using [GitHub Actions](#). A **workflow** is a configurable automated process that runs one or more jobs. An **event** triggers the execution of workflows, and a **job** is a set of steps in a workflow. A **step** can be an action (reusable in different contexts) or a shell script. The **runner** is a [Virtual Machine](#) in which a job is executed.

To create the workflow, we'll need to create a directory called `.github` and a folder called `workflows` within it. Some [pre-written examples](#) are:

- **Example** of CI workflow using `cmake`: [cmake-single-platform.yml](#)
- **Example** of CI workflow using `Maven`: [maven.yml](#)

A note about YAML. YAML (YAML Ain't Markup Language) is a data serialization language designed to be human friendly.

## 7.4 Containers and Schedulers

In general, software should run correctly on multiple environments, such as a developer's laptop. Unfortunately, there are some problems, such as different versions of libraries and middleware.

An optimal solution is containers. A **Container** consists of an **entire runtime environment**: an application plus all its dependencies, libraries and other binaries and configuration files needed to run it, **bundled into one package**.

A similar concept is the **Virtual Machine**. A Virtual Machine, commonly shortened to just VM, is no different than any other physical computer like a laptop, smart phone, or server. It has a CPU, memory, disks to store your files, and can connect to the internet if needed. While the parts that make up your computer (called hardware) are physical and tangible, VMs are often thought of as virtual computers or software-defined computers within physical servers, existing only as code.<sup>3</sup>

The differences between containers and virtual machines are as follows:

- **Containers**

- Rely on the underlying operating system.
- A container image is in the order of MBs.
- Are available almost instantly.

- **VMs**

- Each has its own operating system.
- A VM image is in the order of GBs.
- Require several minutes to bootstrap.

The most common container frameworks are Docker, which is widely used in classical software development, Singularity, which is used in HPC software, and **OCI** (Open Container Initiative).

### Terminology

The terminology used when talking about containers is as follows:

- **Image**. This is a read-only **template that contains instructions for creating a container** (such as Docker). An image can be based on another image, with some customization.
- **Container**. An **executable instance of an image**. Containers can be created, started, stopped, moved or deleted. They can be connected to one or more networks and have storage attached to them.
- **Dockerfile** is an example of one of the frameworks and contains the instructions for creating and running a new image.

---

<sup>3</sup>[What is a virtual machine \(VM\)?](#)

#### 7.4.1 Singularity

**Singularity** is designed to **run complex applications on HPC clusters** in a simple, portable and reproducible way.

This framework has the following features:

- Verifiable reproducibility and security, using cryptographic signatures, an immutable container image format, and in-memory decryption.
- Integration over isolation by default. Easily make use of GPUs, high speed networks, parallel filesystems on a cluster or server by default.
- Mobility of compute. The single file SIF container format is easy to transport and share.
- A simple, effective security model. You are the same user inside a container as outside, and cannot gain additional privilege on the host system by default.

However, the previous features are detailed explained [here](#).

---

#### 7.4.2 Slurm

Along with Singularity, Slurm is also commonly used. **Simple Linux Utility for Resource Management** (SLURM) is a free and open source **job scheduler** (see definition below) for Linux and Unix-like kernels, used by many of the world's supercomputers and computer clusters.

Slurm is the workload manager on approximately 60% of the TOP500 supercomputers.

A **Resource Manager** co-ordinates the actions of all other components in the batch system by maintaining a database of all resources, submitted requests and running jobs.

In addition, a **Job Scheduler** takes the node and job information from the resource manager and creates a list, sorted by job priority, that tells the resource manager when and where to run each job.

Slurm was originally designed as a **simple resource manager**, capable only of allocating whole nodes to jobs. It has evolved into a **comprehensive workload scheduler capable of managing the most demanding workflows** on many of the world's largest computers. Its design goals are to be open source, portable, scalable, fault tolerant, secure and sys admin friendly.

## Terminology

Some of the main concepts of Slurm:

- **Job**: resource allocation request, two types:
  1. Batch script to be queued for later execution.
  2. Interactive job for which the user awaits resource allocation and then makes real-time use of it.
- **Job Step**: a set of parallel tasks, typically an MPI application.
- **Node**: computational element.
- **Cluster**: a collection of nodes sharing the same network.
- **Federation**: a collection of clusters sharing a standard configuration DB.
- **Job Partition**: It is a job queue. Note that a job can be partitioned into multiple parts.
- **Account**: an access group that gives users access to resources in a cluster.
- **Slurm Association**: a combination of the cluster, account, username, and (optional) partition name.

## Commands

And here is a list of the most common Slurm commands ([summary of commands](#)):

- **slurmd**: controls the execution of jobs and job steps on a node
- **slurmctld**: monitors the status of resources, decides when and where to initiate jobs and job steps, and processes almost all user commands.
- **slurmdbd**: interfaces with a DBMS such as MariaDB or MySQL, keeps accounting information and manages some configuration information centrally (many limits, fair share information, QOS, licences, etc.).
- **slurmrestd**: interfaces to Slurm via REST.
- **sbatch**: submits a job script for later execution.
- **srun**: submits a job for execution or initiate job steps in real time.
- **salloc**: allocates resources for a job in real time.
- **squeue**: reports the state of jobs or job steps.
- **sinfo**: reports the state of partitions and nodes managed by Slurm.

The default scheduling approach is FIFO. But with the Multifactor Priority plugin it is ordered by priority. There is a detailed explanation in the following [link](#) (these are slides).

## 8 Exercises

These exercises are taken from the course workbook [3]. In the following sections you can see how we approached the exercises.

### 8.1 Availability

#### 8.1.1 Collection and analysis

##### Problem description

Consider the system in Figure 28. Assume that the participating components offer the following availability:

- DataCollector: 99%
- MessageQueue: 99.99%
- DataAnalyzer: 99.5%

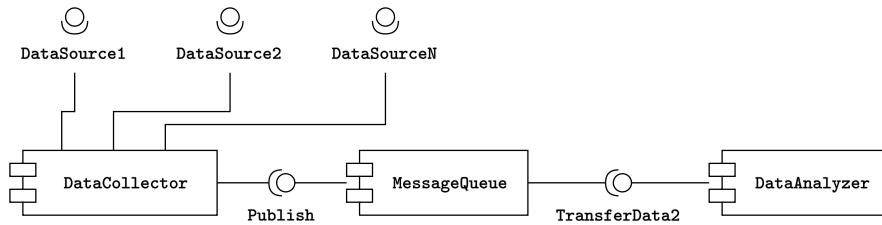


Figure 28: Collection and analysis system.

##### Questions

1. Provide an estimate of your system's total availability (a rough estimation is acceptable without complete calculations). Assume you want to improve this total availability through replication, which component(s) would you choose to replicate? Explain your reasoning.
2. How would such replication impact the way the system works and is designed?

##### Solution

**Solution 1.** We can use Equation 5 on page 63 to estimate our system's total availability. Then, we can multiply the availability of each item in the system:

$$\begin{aligned}\text{Availability} &= \text{DataCollector} \times \text{MessageQueue} \times \text{DataAnalyzer} \\ &= (99 \div 100) \times (99.99 \div 100) \times (99.5 \div 100) \\ &= 0.99 \times 0.9999 \times 0.995 \\ &= 0.984951495 \approx 0.985\end{aligned}$$

Now, we want to improve the system's accessibility. As the exercise suggests, we use the replication technique. Since the component to replicate is the worst, we replicate the **DataAnalyzer** (99.5% availability).

So, how can we calculate accessibility? Well, when we replicate a component, we take advantage of the *parallel* concept. Then, we use Equation 6 (page 64) on these components. Finally, the availability becomes:

$$\begin{aligned}\text{Improved Avail.} &= 0.99 \times 0.9999 \times [1 - (1 - 0.995) \times (1 - 0.995)] \\ &= 0.989901 \times [1 - 0.000025] \\ &= 0.989901 \times 0.999975 \\ &= 0.989876252475 \approx 0.989 \approx 0.99\end{aligned}$$

**Solution 2.** We have improved the availability system with replication, but how are two parallel components managed? We can choose between three available tactics (page 66): Hot Spare, Warm Spare, and Cold Spare (*Triple Modular Redundancy* cannot be applied because there are only two parallel modules).

- *Hot Spare* approach: **DataCollector1** and **DataCollector2** are updated at the same time from the **MessageQueue**. One component leads, and another is always ready to take over.

However, the main problem with this choice is the *data duplication*. Because when the two components receive the same data, both process and return the information at the source **MessageQueue**. To manage this issue, we can code a system to avoid duplicates on the queue.

Another issue when two sources answer simultaneously is *concurrency* (a race condition problem). This problem can be solved by implementing mutual exclusion mechanisms.

- *Warm or Cold Spare*: In the first, **DataCollector1** leads and periodically updates the **DataCollector2**. If the primary **DataCollector** fails, the second **DataCollector** takes time to update itself fully. In the second, **DataCollector2** is dormant, and it's started and updated only if required.

### 8.1.2 Train Ticket

#### Problem description

Consider a microservices application called **TrainTicket** composed of 3 domain microservices (**search**, **reserve**, **buy**) and 1 additional microservice that acts as the API gateway. **TrainTicket** supports two basic operations invoked using the exposed RESTful APIs.

- Search request: `/APIv1/search/{args}`
- Reserve request: `/APIv1/reserve/{args}`

Requests (search and reserve) are received and dispatched by the API gateway. In particular, Figure 29 shows how requests propagate from the gateway to internal microservices. Note that in this example, reserve also includes the purchase of reserved items.

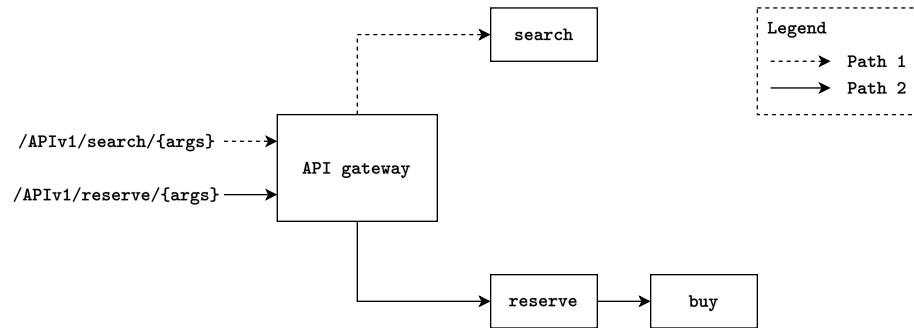


Figure 29: Train ticket system.

Microservices run in units deployed onto 2 different Virtual Machines (VMs), VM1 and VM2 as shown in the following UML deployment diagram.

The available VMs have Computational Resources (CRs) that can be allocated to run microservices. Each VM has a maximum number of CRs and each microservice requires a certain number of CRs, according to the executed artifact. As shown in the schema, available CRs are as follows.

- VM1: 20 CRs
- VM2: 22 CRs

The mapping between microservices and required CRs is as follows.

- API gateway: 2 CRs
- search: 5 CRs
- reserve: 4 CRs
- buy: 5 CRs

The deployment diagram in Figure 30 shows that each microservice can be replicated to have redundant business-critical components. In the latter case, requests are directed to all the replicas rather than to an individual instance, and the first answer received from a replica is returned to the caller, while the others are simply ignored. The number of replicas for each microservice shall be defined so that the following nonfunctional requirement is satisfied and the deployment constraints defined in the deployment diagram and above are fulfilled.

*R<sub>1</sub>: “Both search and reserve services exposed through API gateway shall have availability greater than or equal to 0.99.”*

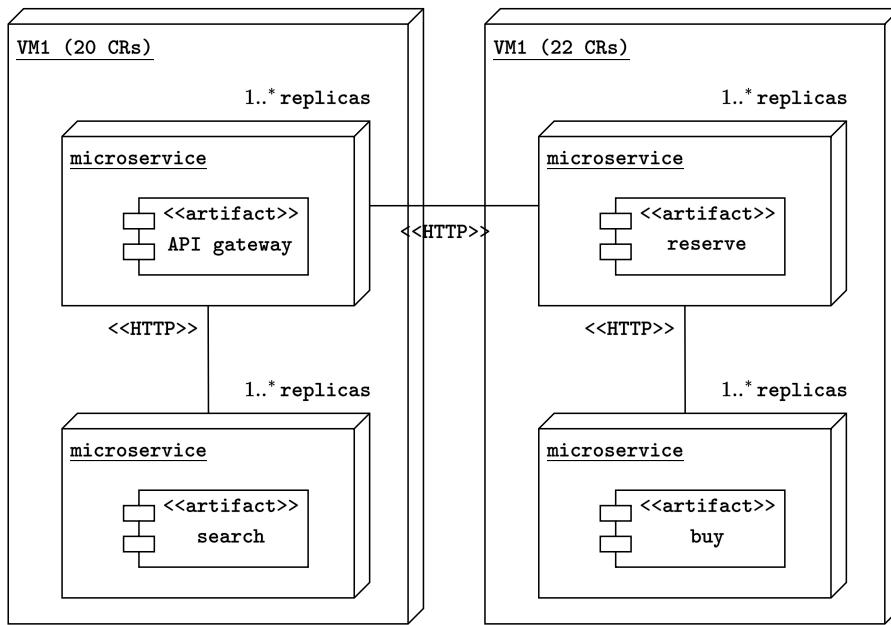


Figure 30: Train ticket deployment diagram.

## ❷ Questions

- Considering the constraints of the execution environment represented, determine whether requirement  $R_1$  can be satisfied or not assuming the following availability estimates for each microservice.
  - API gateway: 0.99
  - search: 0.98
  - reserve: 0.95
  - buy: 0.9
- Consider the problem of resource allocation taking into account the operational profile, that is, the behavior of the users. Assume the following workload in terms of average number of concurrent users for each request.

- (a) search: 50 users
- (b) reserve: 90 users

Assume that for reserve, only 20% of users complete the purchase at reservation time. This means that 20% of reserve requests get through and reach the buy microservice, while 80% of them terminate the execution without calling buy. After a preliminary analysis, we realize that availability depends on the workload according to the estimates presented in Table 3. Does the execution environment have enough computational resources to support the workload defined above still fulfilling requirement *R1* and the defined constraints? Justify your answer.

LOW workload		HIGH workload	
0-60 concurrent users		60-150 concurrent users	
Microservice	Availability	Microservice	Availability
API gateway	0.99	API gateway	0.98
search	0.98	search	0.95
reserve	0.95	reserve	0.93
buy	0.91	buy	0.90

Table 3: Microservices availability under different workloads.

### ✓ Solution

**Solution 1.** The requirement *R1* can be satisfied if both virtual machines respect the constraints about the computational resources (CR) and the constraints about the availability percentage.

To analyze the constraint about the CRs, we need to write two simple inequations:

$$\begin{cases} 2x + 5y \leq 20 \\ 4u + 5z \leq 22 \end{cases}$$

Where the first one refers to VM1 and the second refers to VM2. To the right side, we have the computational resource limit; to the left, we have a sum of the microservices inside each Virtual Machine. The unknown variables indicate “*how many resources are requested*”. For example, if  $x$  will be 5, the API gateway has a request equal to 10 ( $5 \times 2$ ).

With the same logic, we can write about the inequations in availability. We use the same unknown variables because the computational resources directly affect the availability calculation.

$$\begin{cases} (1 - (1 - 0.99)^x) \cdot (1 - (1 - 0.98)^y) \geq 0.99 \\ (1 - (1 - 0.99)^x) \cdot (1 - (1 - 0.95)^u) \cdot (1 - (1 - 0.91)^z) \geq 0.99 \end{cases}$$

To the right side, we have the value accepted by the request *R1*; as a symbol, we use  $\geq$  because the request says “[...] availability greater than or equal to 0.99.”; finally, to the left, we have any component written using the parallelization

formula. The reason is that if we have multiple computational resources at a time, we need a duplication. An example of a valid assignment can be:

$$x = 2, y = 2, u = 3, z = 2$$

**Solution 2.** On time  $t$ , we have 50 users who request the search microservice and 90 who request the reserve service. Then, the API gateway module will receive 140 ( $50+90$ ) concurrent users, and its workload will be HIGH (determined by looking at the Table 3).

The search module will receive 50 concurrent requests, and looking at the table; its workload will be LOW.

The reserve module must manage 90 concurrent requests, and its workload will be HIGH.

Finally, the buy module will receive only 18 concurrent requests (20% of 90 requests), and its workload will be LOW. Then summarizing:

- API gateway: 140 users (HIGH)  $\rightarrow$  0.98 availability
- search: 50 users (LOW)  $\rightarrow$  0.98 availability
- reserve: 90 users (HIGH)  $\rightarrow$  0.93 availability
- buy: 18 users (LOW)  $\rightarrow$  0.91 availability

The availability inequations written in the previous solution it is modified as follows:

$$\begin{cases} (1 - (1 - 0.98)^x) \cdot (1 - (1 - 0.98)^y) \geq 0.99 \\ (1 - (1 - 0.98)^x) \cdot (1 - (1 - 0.93)^u) \cdot (1 - (1 - 0.91)^z) \geq 0.99 \end{cases}$$

An optimal resource allocation is represented by the assignment:

$$x = 2, y = 2, u = 3, z = 2$$

That is again feasible according to environment constraints.

## 8.2 Petri Nets

### 8.2.1 Logical And

#### Problem description

Define a Petri Net (P/T net) model for calculating the logical conjunction of two variables,  $x$ , and  $y$ , each of which takes the values *true* or *false*, with each variable's value being independent of the other.

#### Questions

1. Draw the diagram describing the aforementioned P/T ned model.
2. Define a P/T net that computes the negation of  $x$ , which again takes only the values *true* and *false*.
3. Define a P/T net that computes  $!(x \wedge y)$  by composing the previous models.

#### Solution

**Solution 1.** We use a Petri Net (section 6.1.1 at page 58) to describe distributed systems. In other words, it is a soft UML graph with an exact mathematical definition of its execution semantics.

The exercise requests that we design a P/T net that describes the logic and operator (conjunction of two variables):  $x \wedge y$ . The values of the variables are their domain, *true* and *false*, and then they are boolean values. The values to try depend on the logic table:

$x$	$y$	$\wedge$
True	True	True
True	False	False
False	True	False
False	False	False

The Petri Net model is exposed in the Figure 31.

**Solution 2.** The model of the negation  $x$  depends on the negation logic table:

$x$	$\neg x$
True	False
False	True

The Petri Net model is exposed in the Figure 32.

**Solution 3.** To combine the previous two models, we can extend the first model (Figure 31) with the second (Figure 32) in a very simple way. See the final Petri Net model in the Figure 33.

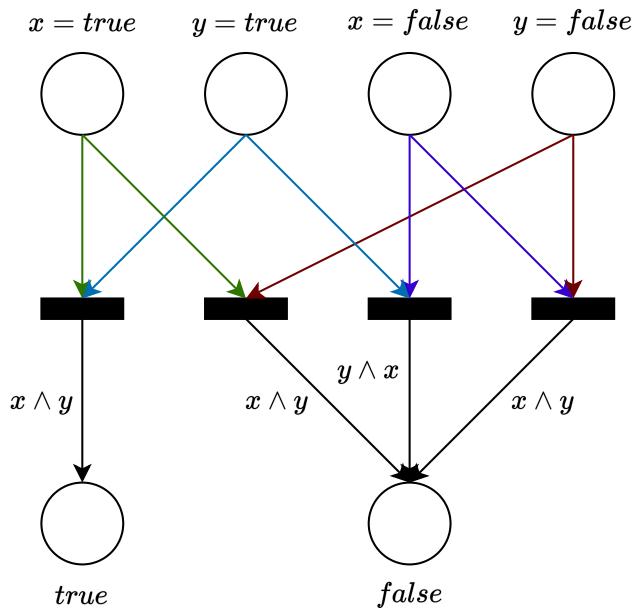


Figure 31: Petri Net model of logical and.

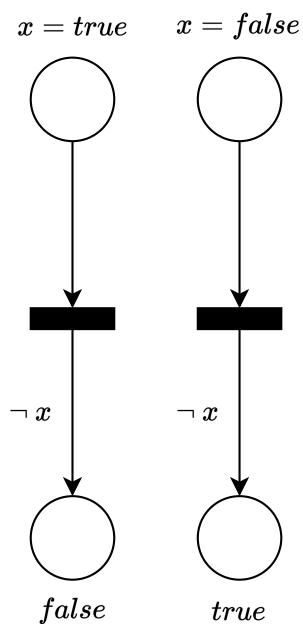
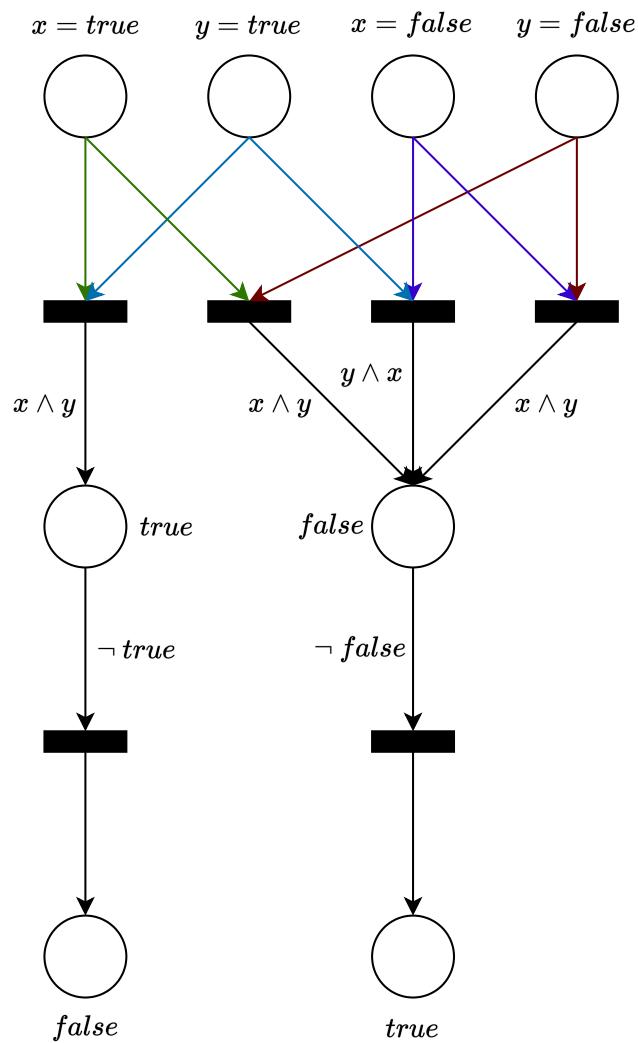


Figure 32: Petri Net model of logical not.

Figure 33: Petri Net model of logical  $\wedge$  and  $\neg$ .

### 8.2.2 Reachability graph

#### Problem description

Consider the P/T Net  $(N, M_0)$  presented in Figure 34.

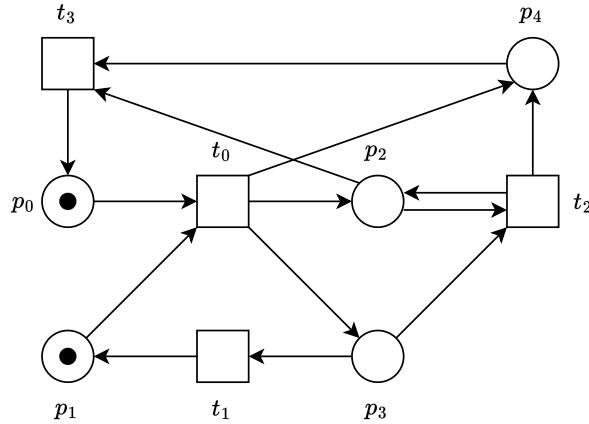


Figure 34: An example of P/T Net model.

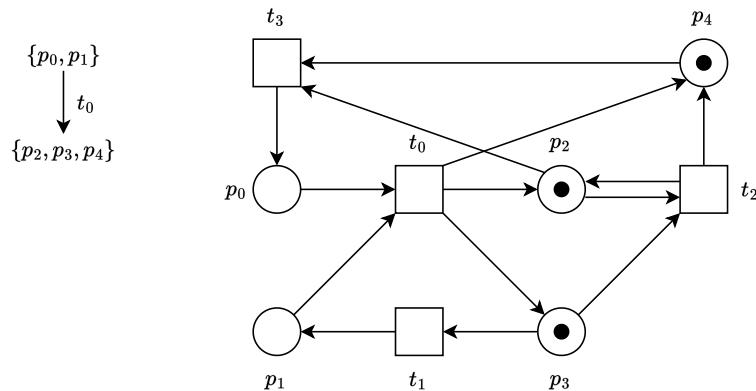
#### Questions

1. Define the reachability graph of  $(N, M_0)$ .

#### Solution

**Solution 1.** The reachability graph of the Petri Net:

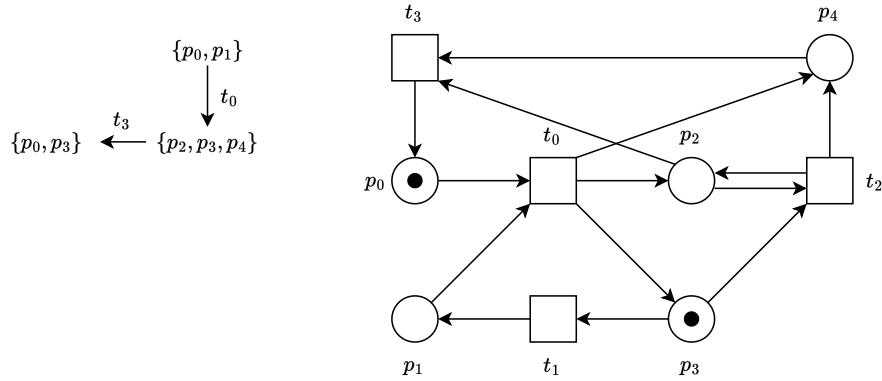
1. The initial state of the model has the state  $p_0$  and  $p_1$  active. This is the start of the reachability graph:  $\{p_0, p_1\}$ .
2. The arcs leave states  $p_0$  and  $p_1$  and go to transition 0 ( $t_0$ ). The destinations are  $p_2$ ,  $p_3$  and  $p_4$ . This is the reachability graph:



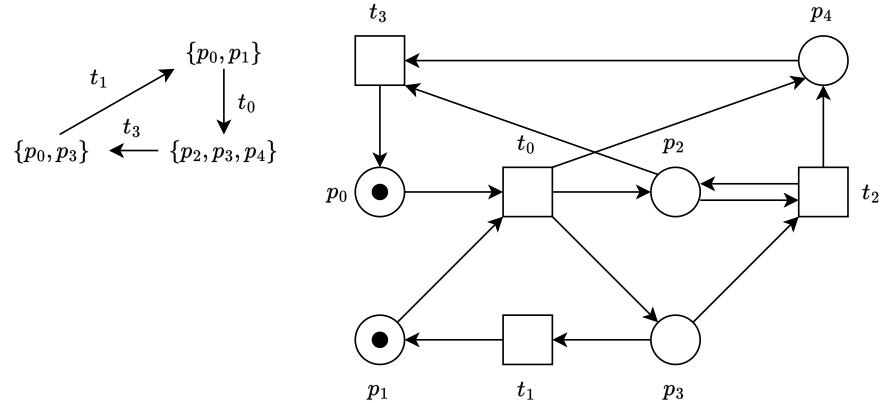
3. We are evaluating three possible options:

- $t_3$

(a) The transition  $t_3$  requires the states  $p_3$  and  $p_4$ . The target will be  $p_0$ :

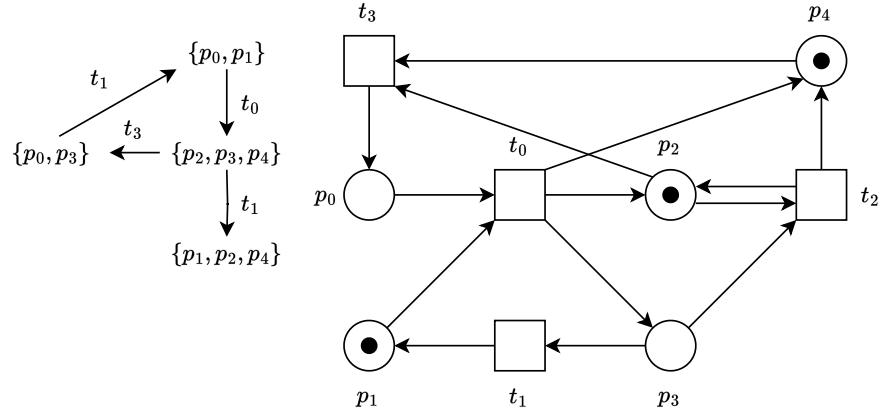


(b) The transition  $t_0$  can't be made by the  $p_0$  state because it needs the  $p_1$  state. Then the  $p_3$  state makes the transition  $t_1$  towards the  $p_1$  state. At this point we can see that it's the same system as the original one.

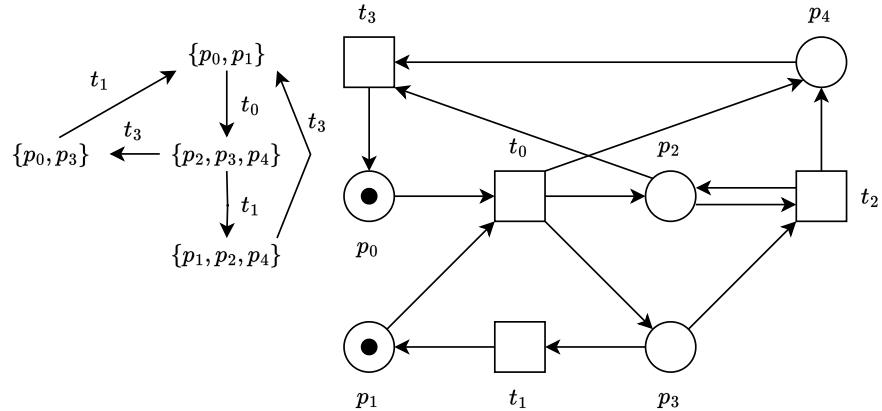


- $t_1$

(a) The transition  $t_1$  is made by the state  $p_3$ . So the destination is  $p_1$ .

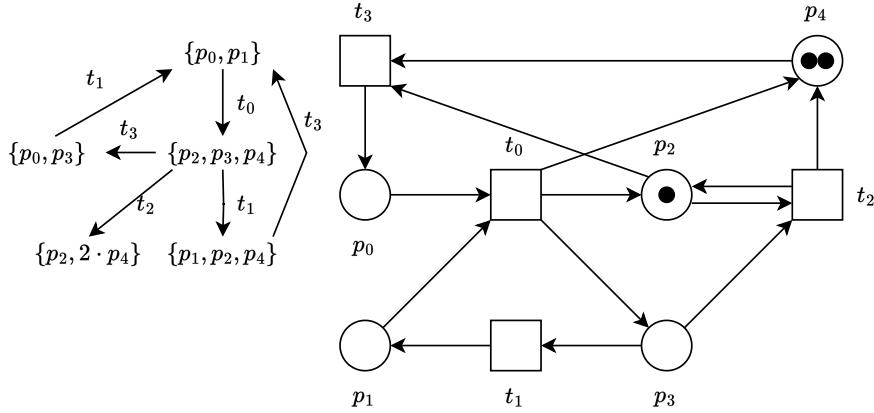


(b) At this point we can see that the transition  $t_0$  cannot be performed by the state  $p_1$  because  $t_0$  needs the state  $p_0$ . This state can be performed by the transition  $t_3$ . Again, this situation brings us back to the initial state.

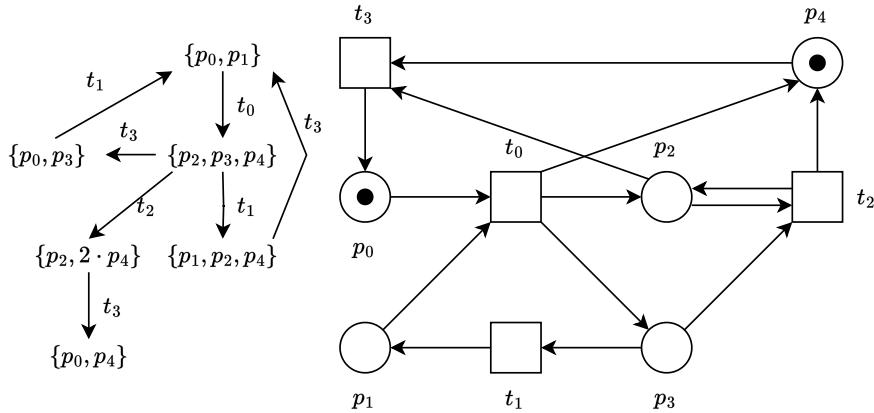


- $t_2$

(a) The transition  $t_2$  is made by the state  $p_3$  and the state  $p_2$ . As we can see, the state  $p_4$  gets another point and the state  $p_2$  doesn't lose its point because it has an internal loop.



(b) Finally, we make the transition  $t_3$  with  $p_2$  and  $p_4$ .



## 8.3 Symbolic Execution

### 8.3.1 Simple Computation

#### Problem description

Consider the following fragment of code:

```

0 int computation(int a[], int n) {
1     int i, count, prod;
2     if (n < 2)
3         printf("%d", n);
4     i = 0;
5     count = 0;
6     prod = 1;
7     while (i < n) {
8         if (a[i] == 0)
9             count++;
10        prod = prod * a[i];
11        i++;
12    }
13    if (count < 2 || prod == 0)
14        return -1;
15    else
16        return prod;
17 }
```

#### Questions

- Derive the path condition corresponding to the execution of path:

$<1, 2, 4, 5, 6, 7, 8, 10, 11, 12, 7, 8, 9, 10, 11, 12, 7, 13, 14>$

- Derive the path condition corresponding to the execution of path:

$<1, 2, 4, 5, 6, 7, 13, 15>$

#### Solution

The symbolic execution analyses a program to determine what inputs cause each part of a program to execute. The full explanation can be found in section 6.2 on page 69.

The answer to both questions:

- Path: 0

$$\begin{array}{c}
 a \qquad n \\
 \hline
 A \qquad N
 \end{array} <0>$$

- Path: 0, 1

$$\frac{a \quad n \quad i \quad count \quad prod}{A \quad N \quad I \quad COUNT \quad PROD} < 0, 1 >$$

- Path: 0, 1, 2

$$\frac{a \quad n \quad i \quad count \quad prod \quad \pi}{A \quad N \quad I \quad COUNT \quad PROD \quad N \geq 2} < 0, 1, 2 >$$

- Path: 0, 1, 2, 4

$< 0, 1, 2, 4 >$

$$\frac{a \quad n \quad i \quad count \quad prod \quad \pi}{A \quad N \quad 0 \quad COUNT \quad PROD \quad N \geq 2}$$

- Path: 0, 1, 2, 4, 5, 6

$< 0, 1, 2, 4, 5, 6 >$

$$\frac{a \quad n \quad i \quad count \quad prod \quad \pi}{A \quad N \quad 0 \quad 0 \quad 1 \quad N \geq 2}$$

Here is the answer to the first question ([solution 1](#)):

- Path: 0, 1, 2, 4, 5, 6, 7

$< 0, 1, 2, 4, 5, 6, 7 >$

$a$	$n$	$i$	$count$	$prod$	$\pi$
$A$	$N$	0	0	1	$N \geq 2$
$N \geq 2 \wedge I < N$					
↓					
$N \geq 2 \leftarrow N \geq 2 \wedge 0 < N$					

- Path: 0, 1, 2, 4, 5, 6, 7, 8

$< 0, 1, 2, 4, 5, 6, 7, 8 >$

$a[0]$	$a$	$n$	$i$	$count$	$prod$	$\pi$
$A[0]$	$A$	$N$	0	0	1	$N \geq 2 \wedge$
$A[0] \neq 0$						

- Path: 0, 1, 2, 4, 5, 6, 7, 8, 10

$< 0, 1, 2, 4, 5, 6, 7, 8, 10 >$

$a[0]$	$a$	$n$	$i$	$count$	$prod$	$\pi$
$A[0]$	$A$	$N$	0	0	$1 \cdot A[0]$	$N \geq 2 \wedge$
$A[0] \neq 0$						

- Path: 0, 1, 2, 4, 5, 6, 7, 8, 10, 11, 12

$< 0, 1, 2, 4, 5, 6, 7, 8, 10, 11, 12 >$

$a[0]$	$a$	$n$	$i$	$count$	$prod$	$\pi$
$A[0]$	$A$	$N$	$0 + 1$	0	$1 \cdot A[0]$	$N \geq 2 \wedge$
						$A[0] \neq 0$

- Path: 0, 1, 2, 4, 5, 6, 7, 8, 10, 11, 12, 7

$< 0, 1, 2, 4, 5, 6, 7, 8, 10, 11, 12, 7 >$

$a[0]$	$a$	$n$	$i$	$count$	$prod$	$\pi$
$A[0]$	$A$	$N$	1	0	$1 \cdot A[0]$	$N \geq 2 \wedge$
						$A[0] \neq 0 \wedge$
						1 < $N$

- Path: 0, 1, 2, 4, 5, 6, 7, 8, 10, 11, 12, 7, 8, 9

$< 0, 1, 2, 4, 5, 6, 7, 8, 10, 11, 12, 7, 8, 9 >$

$a[1]$	$a[0]$	$a$	$n$	$i$	$count$	$prod$	$\pi$
$A[1]$	$A[0]$	$A$	$N$	1	1	$1 \cdot A[0]$	$N \geq 2 \wedge$
							$A[0] \neq 0 \wedge$
							$A[1] = 0$

- Path: 0, 1, 2, 4, 5, 6, 7, 8, 10, 11, 12, 7, 8, 9, 10

$$<0, 1, 2, 4, 5, 6, 7, 8, 10, 11, 12, 7, 8, 9, 10>$$

$a[1]$	$a[0]$	$a$	$n$	$i$	$count$	$prod$	$\pi$
$A[1]$	$A[0]$	$A$	$N$	1	1	$A[0] \cdot A[1]$	$N \geq 2 \wedge$
							$A[0] \neq 0 \wedge$
							$A[1] = 0$

- Path: 0, 1, 2, 4, 5, 6, 7, 8, 10, 11, 12, 7, 8, 9, 10, 11, 12

$$<0, 1, 2, 4, 5, 6, 7, 8, 10, 11, 12, 7, 8, 9, 10, 11, 12>$$

$a[1]$	$a[0]$	$a$	$n$	$i$	$count$	$prod$	$\pi$
$A[1]$	$A[0]$	$A$	$N$	2	1	$A[0] \cdot A[1]$	$N \geq 2 \wedge$
							$A[0] \neq 0 \wedge$
							$A[1] = 0$

- Path: 0, 1, 2, 4, 5, 6, 7, 8, 10, 11, 12, 7, 8, 9, 10, 11, 12, 7

$$<0, 1, 2, 4, 5, 6, 7, 8, 10, 11, 12, 7, 8, 9, 10, 11, 12, 7>$$

$a[1]$	$a[0]$	$a$	$n$	$i$	$count$	$prod$	$\pi$
$A[1]$	$A[0]$	$A$	$N$	2	1	$A[0] \cdot A[1]$	$N \geq 2 \wedge$
							$A[0] \neq 0 \wedge$
							$A[1] = 0 \wedge$
							$N = 2 \quad \leftarrow \quad 2 \geq N \quad \leftarrow \quad \neg(2 < N)$

- Path: 0, 1, 2, 4, 5, 6, 7, 8, 10, 11, 12, 7, 8, 9, 10, 11, 12, 7, 13, 14

$< 0, 1, 2, 4, 5, 6, 7, 8, 10, 11, 12, 7, 8, 9, 10, 11, 12, 7, 13, 14 >$

$a[1]$	$a[0]$	$a$	$n$	$i$	$count$	$prod$	$\pi$
$A[1]$	$A[0]$	$A$	$N$	2	1	$A[0] \cdot A[1]$	$N = 2 \wedge$
							$A[0] \neq 0 \wedge$
							$A[1] = 0 \wedge$
							$True! \leftarrow (1 < 2 \vee A[0] \cdot A[1] = 0)$

This path is *satisfiable*. An unsatisfiable case is solution number 2:

- Path: 0, 1, 2, 4, 5, 6, 7

$a$	$n$	$i$	$count$	$prod$	$\pi$
					$< 0, 1, 2, 4, 5, 6, 7 >$
$A$	$N$	0	0	1	$N \geq 2 \wedge$
					$N \leq 0$

It's unsatisfiable because the logical condition is a contradiction:  $N \geq 2 \wedge N \leq 0$ .

## 8.4 Strategies for test case identification

### 8.4.1 Triangles

#### ☒ Problem description

Consider a function with the following signature:

```
bool isTriangle(float a, float b, float c)
```

The function receives three integers representing the lengths of three segments. The function returns true if the three segments can constitute the sides of a triangle and false otherwise.

#### ❓ Questions

Assume we have the following test cases and corresponding output:

- $a = 10, b = 12, c = 9$ , result **true**
- $a = 10, b = 12, c = 22$ , result **false**

Identify one or more metamorphic relations between inputs and outputs and some corresponding test cases.

#### ✓ Solution

We can identify at least the following two **Metamorphic Relations** predicated on the sides of a triangle:

1. *Relation between the sides of a triangle*: the length of each side of a triangle is less than the sum of the lengths of the other two sides.
2. *Similarity*: two triangles, A and B, are said to be similar if their side pairs are proportional.

The corresponding test cases are:

1. By exploiting the first relation:
  - $a = 12, b = 9, c = 10$ , result **true**
  - $a = 12, b = 22, c = 10$ , result **false**
2. By exploiting the similarity:
  - $a = 5, b = 6, c = 4.5$ , result **true**
  - $a = 100, b = 120, c = 90$ , result **true**

## 8.5 Integration Testing Plan Definition

### 8.5.1 Microservices

#### Problem description

Consider the microservice-based architecture shown in Figure 35. The architecture is organized into eight stateless microservices collaborating to fulfill requests  $R1$  and  $R2$ .  $S1$  is the front-end service that receives both requests. The fulfillment of request  $R1$  requires the interaction with services  $S2$  and  $S3$  (through sub-requests  $R1.1$  and  $R1.2$ , respectively), which, in turn, need to interact with other services. In particular,  $S2$  interacts with  $S4$  and  $S5$  and  $S3$  with  $S5$  and  $S6$ . The fulfillment of  $R2$  requires that  $S1$  interacts with  $S8$ , which, in turn, interacts with  $S6$  and  $S7$ .

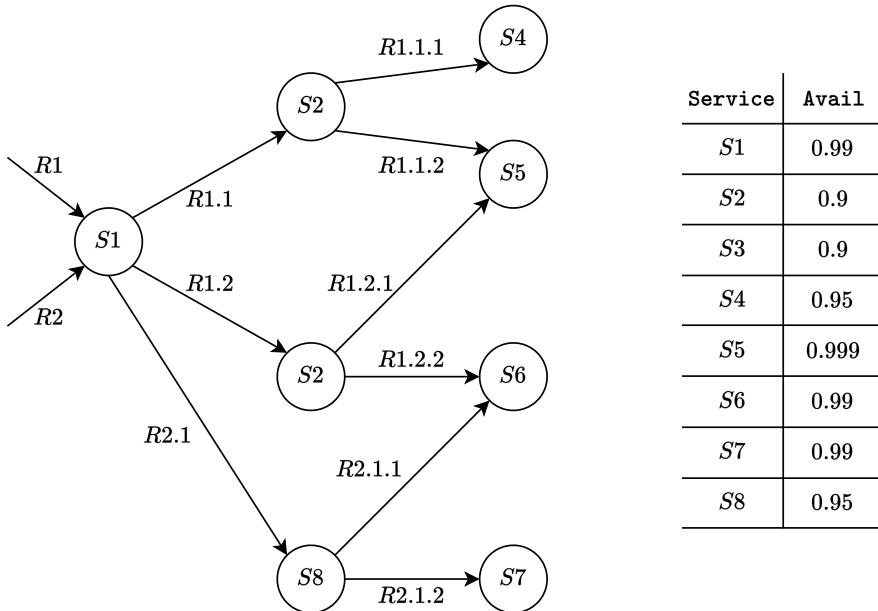


Figure 35: Microservice-based architecture.

#### Questions

1. Define an integration and test plan for the architecture, motivating your choice.
2. Assume that we are asked to test whether the availability improvement theoretically obtained with the duplication of services is actually achieved. How would we proceed to perform such a test?

#### Solution

**Solution 1.** The strategy used for integration testing (a summary of the strategies can be found in Figure 26 on page 79) is the iterative/incremental strategy.

In particular, we use the *Threads* strategy to test some of the modules. This is possible because the set of modules together provide a user-visible programme function.

We also use a *hierarchical* strategy with a *top-down* approach because the modules are integrated and all functionality can be tested.

The test plan can be divided into  $R1$  and  $R2$  for the given architecture. Given the request  $R1$ , we can test a part of the architecture. Meanwhile, given the request  $R2$ , we can test the remaining part of the architecture. Then, the test plan is the following:

1. Integration of the  $S1$  (front-end) and testing by emulating two requests called  $R1$  and  $R2$ . We need the following stubs to test this service:  $R1.1$ ,  $R1.2$  and  $R2.1$ .
2. Once completed, we temporarily leave aside  $R2$  and implement the  $R1$  tree. We start from the  $S2$  and  $S3$  by creating the requests  $R1.1.1$ ,  $R1.1.2$ ,  $R1.2.1$  and  $R1.2.2$  as stubs. Now we can also test the triangles  $S1$ ,  $S2$  and  $S3$ .
3. Finally, we implement  $S4$ ,  $S5$  and  $S6$ . We do not use stubs to test these services because the other modules have already been implemented ( $S2$ ,  $S3$ ). At this point, the request  $R1$  has been thoroughly tested and integrated.
4. We can now integrate the  $R2$  tree. We implement/test the service  $S8$  by using  $R2.1.2$  as stubs. The  $S6$  was already implemented (previous step), so we can use that module directly.
5. Finally, we implement  $S7$  and test the whole architecture because all the services are implemented.

**Solution 2.** To measure the quantitative impact of architectural decisions, we use the availability metric (Formula 3 on page 62).

The main values required are MTTR (Mean Time To Repair) and MTTF (Mean Time To Failures). If we want to do some tests to compare the two architectures before and after duplication, we can run the application for a long time and simulate a failure of one of the services. We can then calculate the time taken by the architecture to restore the service. The requests can be made sequentially or in parallel to study how the architecture responds.

Using the availability metric, we can understand if the chosen replication makes sense and if the availability is better than the architecture without duplication of services.

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