

# Chapter 7: Entity-Relationship Model

**Database System Concepts, 7th Ed.** 

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### **Chapter 7: Entity-Relationship Model**

- Design Process
- Modeling
- Constraints
- E-R Diagram
- Design Issues
- Weak Entity Sets
- Extended E-R Features
- Design of the Bank Database
- Reduction to Relation Schemas
- Database Design
- UML



#### **Design Phases**

- The initial phase of database design is to characterize fully the data needs of the prospective database users.
- Next, the designer chooses a data model and, by applying the concepts of the chosen data model, translates these requirements into a conceptual schema of the database.
- A fully developed conceptual schema also indicates the functional requirements of the enterprise. In a "specification of functional requirements", users describe the kinds of operations (or transactions) that will be performed on the data.



## **Design Phases (Cont.)**

The process of moving from an abstract data model to the implementation of the database proceeds in two final design phases.

- Logical Design Deciding on the database schema. Database design requires that we find a "good" collection of relation schemas.
  - Business decision What attributes should we record in the database?
  - Computer Science decision What relation schemas should we have and how should the attributes be distributed among the various relation schemas?
- Physical Design Deciding on the physical layout of the database



#### **Design Approaches**

- Entity Relationship Model (covered in this chapter)
  - Models an enterprise as a collection of *entities* and *relationships*
    - 4 Entity: a "thing" or "object" in the enterprise that is distinguishable from other objects
      - Described by a set of attributes
    - 4 Relationship: an association among several entities
  - Represented diagrammatically by an *entity-relationship diagram*:
- Normalization Theory (Chapter 8)
  - Formalize what designs are bad, and test for them



#### **Outline of the ER Model**



## ER model -- Database Modeling

- The ER data mode was developed to facilitate database design by allowing specification of an enterprise schema that represents the overall logical structure of a database.
- The ER model is very useful in mapping the meanings and interactions of real-world enterprises onto a conceptual schema. Because of this usefulness, many database-design tools draw on concepts from the ER model.
- The ER data model employs three basic concepts:
  - entity sets,
  - relationship sets,
  - attributes.
- The ER model also has an associated diagrammatic representation, the ER diagram, which can express the overall logical structure of a database graphically.



### **Entity Sets**

- An **entity** is an object that exists and is distinguishable from other objects.
  - Example: specific person, company, event, plant
- An **entity set** is a set of entities of the same type that share the same properties.
  - Example: set of all persons, companies, trees, holidays
- An entity is represented by a set of attributes; i.e., descriptive properties possessed by all members of an entity set.
  - Example:

```
instructor = (ID, name, street, city, salary )
course= (course id, title, credits)
```

• A subset of the attributes form a **primary key** of the entity set; i.e., uniquely identifying each member of the set.



## Entity Sets -- instructor and student

instructor\_ID instructor\_name

76766	Crick
45565	Katz
10101	Srinivasan
98345	Kim
76543	Singh
22222	Einstein

instructor

student-ID student\_name

98988	Tanaka
12345	Shankar
00128	Zhang
76543	Brown
76653	Aoi
23121	Chavez
44553	Peltier

student



### **Relationship Sets**

• A **relationship** is an association among several entities

Example:

```
44553 (Peltier) <u>advisor</u> 22222 (<u>Einstein</u>) 
student entity relationship set instructor entity
```

• A **relationship set** is a mathematical relation among  $n \ge 2$  entities, each taken from entity sets

$$\{(e_1, e_2, \dots e_n) \mid e_1 \in E_1, e_2 \in E_2, \dots, e_n \in E_n\}$$

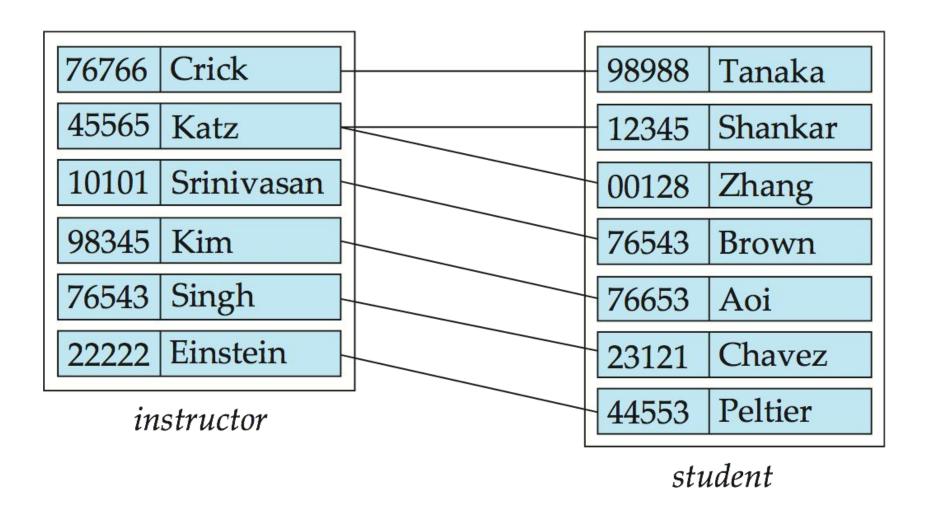
where  $(e_1, e_2, ..., e_n)$  is a relationship

• Example:

$$(44553,22222) \in advisor$$



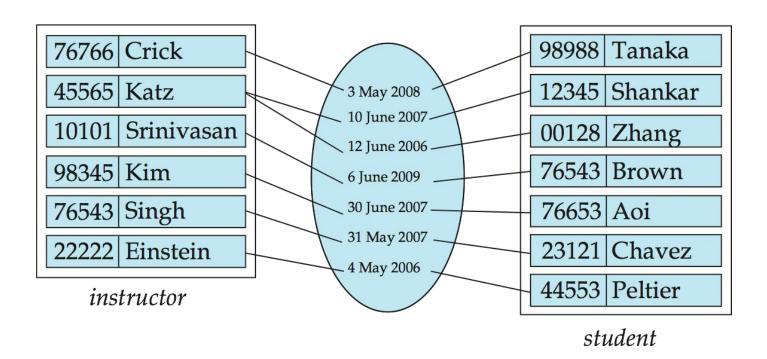
#### Relationship Set advisor





#### **Relationship Sets (Cont.)**

- An attribute can also be associated with a relationship set.
- For instance, the *advisor* relationship set between entity sets *instructor* and *student* may have the attribute *date* which tracks when the student started being associated with the advisor





#### Degree of a Relationship Set

- binary relationship
  - involve two entity sets (or degree two).
  - most relationship sets in a database system are binary.
- Relationships between more than two entity sets are rare. Most relationships are binary. (More on this later.)
  - 4 Example: *students* work on research *projects* under the guidance of an *instructor*.
  - 4 relationship *proj\_guide* is a ternary relationship between *instructor*, *student*, and *project*

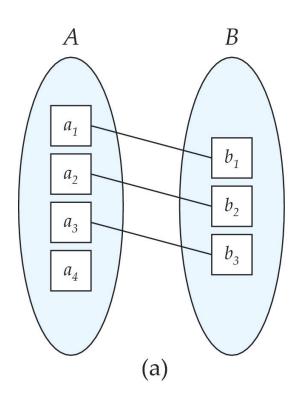


### **Mapping Cardinality Constraints**

- Express the number of entities to which another entity can be associated via a relationship set.
- Most useful in describing binary relationship sets.
- For a binary relationship set the mapping cardinality must be one of the following types:
  - One to one
  - One to many
  - Many to one
  - Many to many



#### **Mapping Cardinalities**



 $\begin{bmatrix} a_1 & b_2 \\ a_2 & b_3 \\ a_3 & b_4 \end{bmatrix}$ (b)

A

B

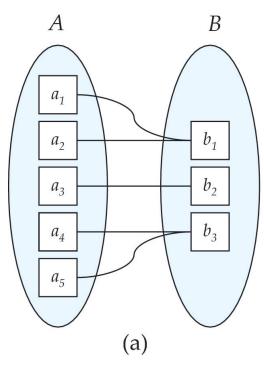
One to one

One to many

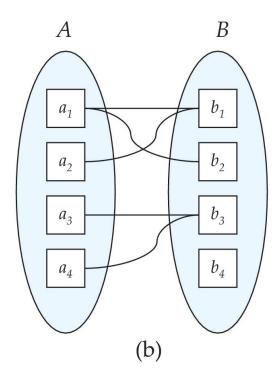
Note: Some elements in A and B may not be mapped to any elements in the other set



#### **Mapping Cardinalities**



Many to one



Many to many

Note: Some elements in A and B may not be mapped to any elements in the other set

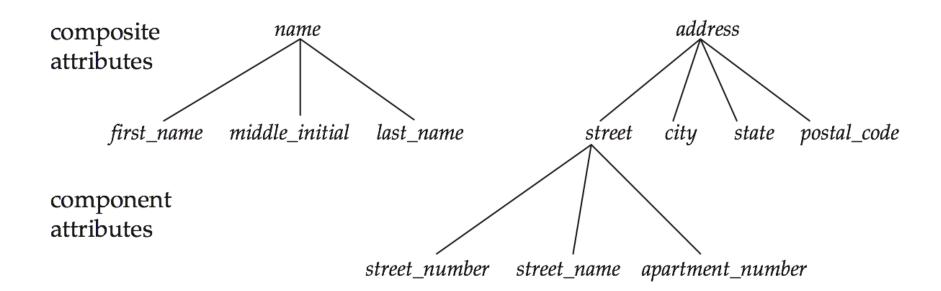


#### **Complex Attributes**

- Attribute types:
  - Simple and composite attributes.
  - Single-valued and multivalued attributes
    - 4 Example: multivalued attribute: *phone numbers*
  - Derived attributes
    - 4 Can be computed from other attributes
    - 4 Example: age, given date of birth
- **Domain** the set of permitted values for each attribute



#### **Composite Attributes**





#### **Redundant Attributes**

- Suppose we have entity sets:
  - *instructor*, with attributes: *ID*, *name*, *dept\_name*, *salary*
  - *department*, with attributes: *dept\_name*, *building*, *budget*
- We model the fact that each instructor has an associated department using a relationship set *inst dept*
- The attribute *dept\_name* appears in both entity sets. Since it is the primary key for the entity set *department*, it replicates information present in the relationship and is therefore redundant in the entity set *instructor* and needs to be removed.
- BUT: when converting back to tables, in some cases the attribute gets reintroduced, as we will see later.



### Weak Entity Sets

- Consider a *section* entity, which is uniquely identified by a *course\_id*, *semester*, *year*, and *sec\_id*.
- Clearly, section entities are related to course entities. Suppose we create a relationship set *sec course* between entity sets *section* and *course*.
- Note that the information in *sec\_course* is redundant, since *section* already has an attribute *course\_id*, which identifies the course with which the section is related.
- One option to deal with this redundancy is to get rid of the relationship sec\_course; however, by doing so the relationship between section and course becomes implicit in an attribute, which is not desirable.



### Weak Entity Sets (Cont.)

- An alternative way to deal with this redundancy is to not store the attribute *course\_id* in the *section* entity and to only store the remaining attributes *section\_id*, *year*, and *semester*. However, the entity set *section* then does not have enough attributes to identify a particular *section* entity uniquely; although each *section* entity is distinct, sections for different courses may share the same *section id*, *year*, and *semester*.
- To deal with this problem, we treat the relationship *sec\_course* as a special relationship that provides extra information, in this case, the *course\_id*, required to identify *section* entities uniquely.
- The notion of **weak entity set** formalizes the above intuition. A weak entity set is one whose existence is dependent on another entity, called its **identifying entity**; instead of associating a primary key with a weak entity, we use the identifying entity, along with extra attributes called **discriminator** to uniquely identify a weak entity. An entity set that is not a weak entity set is termed a **strong entity set**.



### Weak Entity Sets (Cont.)

- Every weak entity must be associated with an identifying entity; that is, the weak entity set is said to be **existence dependent** on the identifying entity set. The identifying entity set is said to **own** the weak entity set that it identifies. The relationship associating the weak entity set with the identifying entity set is called the **identifying relationship**.
- Note that the relational schema we eventually create from the entity set *section* does have the attribute *course\_id*, for reasons that will become clear later, even though we have dropped the attribute *course id* from the entity set *section*.



# E-R Diagrams



## **Entity Sets**

- Entities can be represented graphically as follows:
  - Rectangles represent entity sets.
  - Attributes listed inside entity rectangle
  - Underline indicates primary key attributes

instructor

<u>ID</u>

name

salary

student

<u>ID</u>

name

tot\_cred



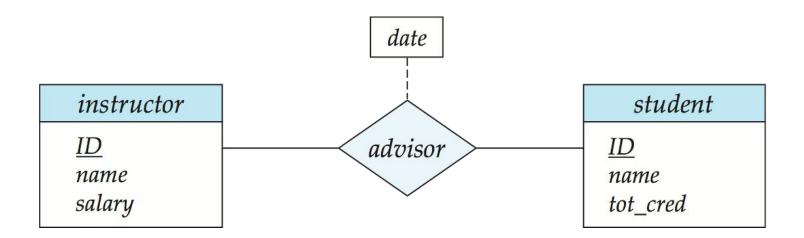
# **Relationship Sets**

• Diamonds represent relationship sets.





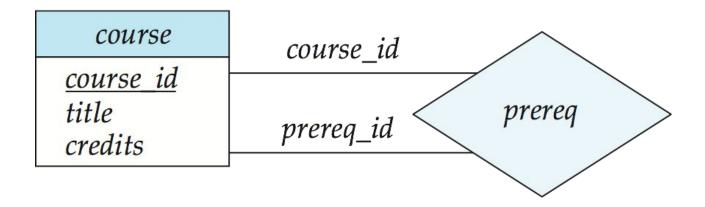
#### **Relationship Sets with Attributes**





#### **Roles**

- Entity sets of a relationship need not be distinct
  - Each occurrence of an entity set plays a "role" in the relationship
- The labels "course\_id" and "prereq\_id" are called roles.





#### **Cardinality Constraints**

- We express cardinality constraints by drawing either a directed line (→), signifying "one," or an undirected line (—), signifying "many," between the relationship set and the entity set.
- One-to-one relationship between an *instructor* and a *student*:
  - A student is associated with at most one *instructor* via the relationship *advisor*
  - A student is associated with at most one department via stud dept





## **One-to-Many Relationship**

- one-to-many relationship between an *instructor* and a *student* 
  - an instructor is associated with several (including 0) students via advisor
  - a student is associated with at most one instructor via advisor,





#### **Many-to-One Relationships**

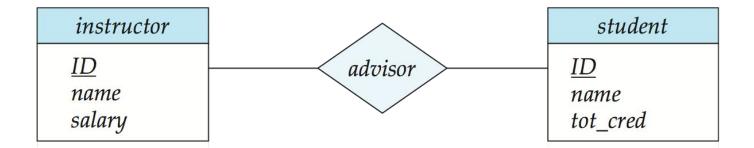
- In a many-to-one relationship between an *instructor* and a *student*,
  - an instructor is associated with at most one student via *advisor*,
  - and a student is associated with several (including 0) instructors via advisor





#### Many-to-Many Relationship

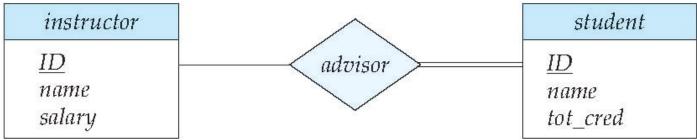
- An instructor is associated with several (possibly 0) students via *advisor*
- A student is associated with several (possibly 0) instructors via *advisor*





#### **Total and Partial Participation**

• Total participation (indicated by double line): every entity in the entity set participates in at least one relationship in the relationship set



participation of student in advisor relation is total

- 4 every *student* must have an associated instructor
- Partial participation: some entities may not participate in any relationship in the relationship set
  - Example: participation of *instructor* in *advisor* is partial



#### **Notation for Expressing More Complex Constraints**

- A line may have an associated minimum and maximum cardinality, shown in the form *l..h*, where *l* is the minimum and *h* the maximum cardinality
  - A minimum value of 1 indicates total participation.
  - A maximum value of 1 indicates that the entity participates in at most one relationship
  - A maximum value of \* indicates no limit.



Instructor can advise 0 or more students. A student must have 1 advisor; cannot have multiple advisors



#### **Notation to Express Entity with Complex Attributes**

#### instructor

```
ID
name
  first_name
   middle_initial
   last_name
address
   street
     street_number
     street_name
     apt_number
   city
   state
   zip
{ phone_number }
date_of_birth
age()
```



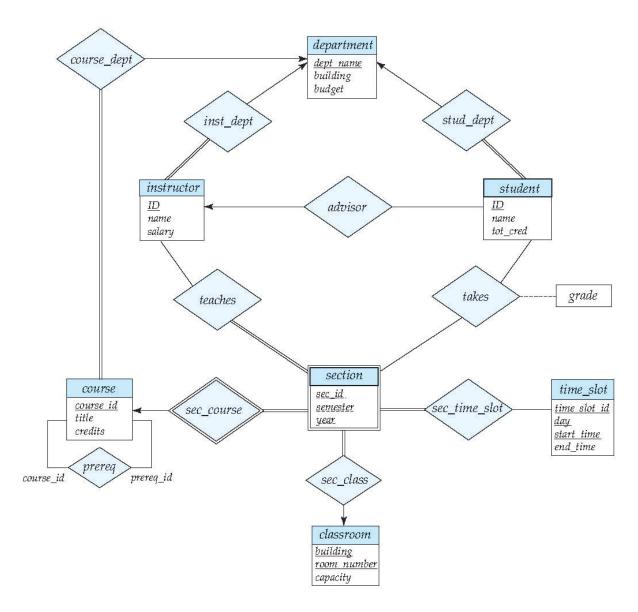
#### **Expressing Weak Entity Sets**

- In E-R diagrams, a weak entity set is depicted via a double rectangle.
- We underline the discriminator of a weak entity set with a dashed line.
- The relationship set connecting the weak entity set to the identifying strong entity set is depicted by a double diamond.
- Primary key for section (course\_id, sec\_id, semester, year)





# E-R Diagram for a University Enterprise





## **Reduction to Relation Schemas**



### **Reduction to Relation Schemas**

- Entity sets and relationship sets can be expressed uniformly as *relation schemas* that represent the contents of the database.
- A database which conforms to an E-R diagram can be represented by a collection of schemas.
- For each entity set and relationship set there is a unique schema that is assigned the name of the corresponding entity set or relationship set.
- Each schema has a number of columns (generally corresponding to attributes), which have unique names.



# **Representing Entity Sets**

• A strong entity set reduces to a schema with the same attributes student(<u>ID</u>, name, tot cred)

 A weak entity set becomes a table that includes a column for the primary key of the identifying strong entity set

section (course id, sec id, sem, year)

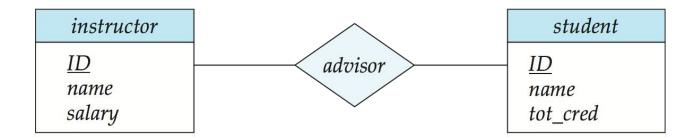




# Representing Relationship Sets

- A many-to-many relationship set is represented as a schema with attributes for the primary keys of the two participating entity sets, and any descriptive attributes of the relationship set.
- Example: schema for relationship set *advisor*

$$advisor = (\underline{s} \ id, i \ id)$$





### Representation of Entity Sets with Composite Attributes

### instructor

```
ID
name
  first_name
   middle initial
   last name
address
   street
     street number
     street_name
     apt_number
   city
   state
   zip
{ phone_number }
date_of_birth
age()
```

- Composite attributes are flattened out by creating a separate attribute for each component attribute
  - Example: given entity set *instructor* with composite attribute *name* with component attributes *first\_name* and *last\_name* the schema corresponding to the entity set has two attributes *name\_first\_name* and *name\_last\_name* 
    - 4 Prefix omitted if there is no ambiguity (name\_first\_name could be first\_name)
- Ignoring multivalued attributes, extended instructor schema is
  - instructor(ID, first\_name, middle\_initial, last\_name, street\_number, street\_name, apt\_number, city, state, zip\_code, date\_of\_birth)



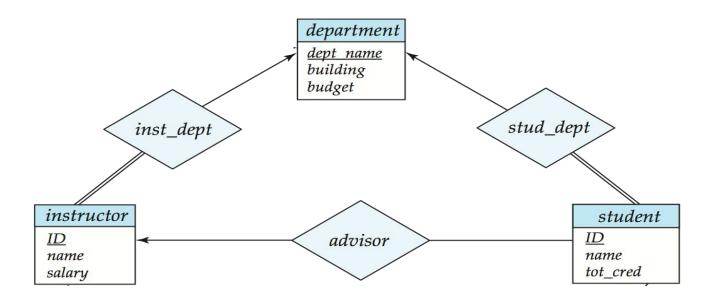
### Representation of Entity Sets with Multivalued Attributes

- A multivalued attribute M of an entity E is represented by a separate schema EM
- Schema *EM* has attributes corresponding to the primary key of *E* and an attribute corresponding to multivalued attribute *M*
- Example: Multivalued attribute *phone\_number* of *instructor* is represented by a schema:
  - inst\_phone= ( <u>ID</u>, <u>phone\_number</u>)
- Each value of the multivalued attribute maps to a separate tuple of the relation on schema *EM* 
  - For example, an *instructor* entity with primary key 22222 and phone numbers 456-7890 and 123-4567 maps to two tuples: (22222, 456-7890) and (22222, 123-4567)



# **Redundancy of Schemas**

- Many-to-one and one-to-many relationship sets that are total on the many-side can be represented by adding an extra attribute to the "many" side, containing the primary key of the "one" side
- Example: Instead of creating a schema for relationship set *inst\_dept*, add an attribute *dept\_name* to the schema arising from entity set *instructor*





# **Redundancy of Schemas (Cont.)**

- For one-to-one relationship sets, either side can be chosen to act as the "many" side
  - That is, an extra attribute can be added to either of the tables corresponding to the two entity sets
- If participation is *partial* on the "many" side, replacing a schema by an extra attribute in the schema corresponding to the "many" side could result in null values



# **Redundancy of Schemas (Cont.)**

- The schema corresponding to a relationship set linking a weak entity set to its identifying strong entity set is redundant.
- Example: The *section* schema already contains the attributes that would appear in the *sec course* schema



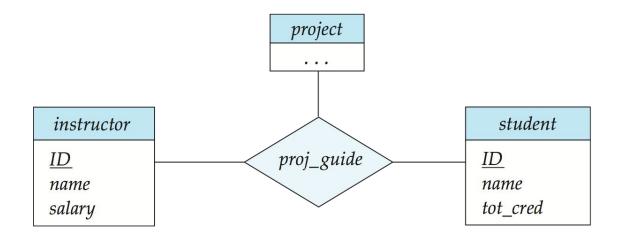


# **Advanced Topics**



## **Non-binary Relationship Sets**

- Most relationship sets are binary
- There are occasions when it is more convenient to represent relationships as non-binary.
- E-R Diagram with a Ternary Relationship





# **Cardinality Constraints on Ternary Relationship**

- We allow at most one arrow out of a ternary (or greater degree) relationship to indicate a cardinality constraint
- For exampe, an arrow from *proj\_guide* to *instructor* indicates each student has at most one guide for a project
- If there is more than one arrow, there are two ways of defining the meaning.
  - For example, a ternary relationship *R* between *A*, *B* and *C* with arrows to *B* and *C* could mean
    - 1. Each A entity is associated with a unique entity from B and C or
    - 2. Each pair of entities from (A, B) is associated with a unique C entity, and each pair (A, C) is associated with a unique B
  - Each alternative has been used in different formalisms
  - To avoid confusion we outlaw more than one arrow



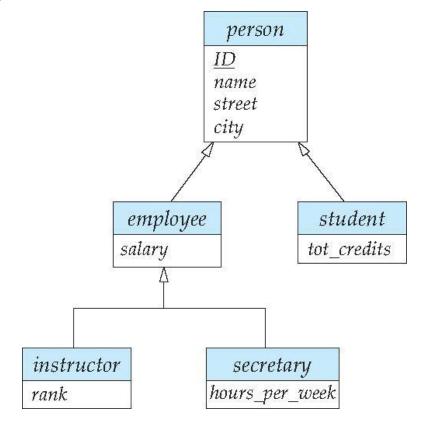
# **Specialization**

- Top-down design process; we designate sub-groupings within an entity set that are distinctive from other entities in the set.
- These sub-groupings become lower-level entity sets that have attributes or participate in relationships that do not apply to the higher-level entity set.
- Depicted by a *triangle* component labeled ISA (e.g., *instructor* "is a" *person*).
- **Attribute inheritance** a lower-level entity set inherits all the attributes and relationship participation of the higher-level entity set to which it is linked.



# **Specialization Example**

- Overlapping employee and student
- **Disjoint** *instructor* and *secretary*
- Total and partial





## Representing Specialization via Schemas

#### Method 1:

- Form a schema for the higher-level entity
- Form a schema for each lower-level entity set, include primary key of higher-level entity set and local attributes

schema a	ttributes
person	ID, name, street, city
student	ID, tot_cred
employee	ID, tot_cred ID, salary

• Drawback: getting information about, an *employee* requires accessing two relations, the one corresponding to the low-level schema and the one corresponding to the high-level schema



## Representing Specialization as Schemas (Cont.)

### • Method 2:

 Form a schema for each entity set with all local and inherited attributes

ttributes
ID, name, street, city
ID, name, street, city, tot_cred
ID, name, street, city, salary

• Drawback: *name*, *street* and *city* may be stored redundantly for people who are both students and employees



### Generalization

- A bottom-up design process combine a number of entity sets that share the same features into a higher-level entity set.
- Specialization and generalization are simple inversions of each other; they are represented in an E-R diagram in the same way.
- The terms specialization and generalization are used interchangeably.



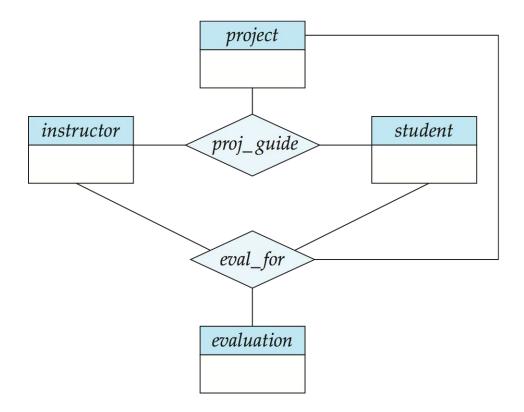
### Design Constraints on a Specialization/Generalization

- Completeness constraint -- specifies whether or not an entity in the higher-level entity set must belong to at least one of the lower-level entity sets within a generalization.
  - total: an entity must belong to one of the lower-level entity sets
  - **partial**: an entity need not belong to one of the lower-level entity sets
- Partial generalization is the default. We can specify total generalization in an ER diagram by adding the keyword **total** in the diagram and drawing a dashed line from the keyword to the corresponding hollow arrow-head to which it applies (for a total generalization), or to the set of hollow arrow-heads to which it applies (for an overlapping generalization).
- The *student* generalization is total: All student entities must be either graduate or undergraduate. Because the higher-level entity set arrived at through generalization is generally composed of only those entities in the lower-level entity sets, the completeness constraint for a generalized higher-level entity set is usually total



# **Aggregation**

- Consider the ternary relationship *proj\_guide*, which we saw earlier
- Suppose we want to record evaluations of a student by a guide on a project





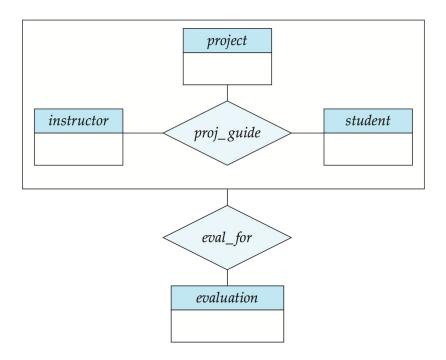
# **Aggregation (Cont.)**

- Relationship sets *eval\_for* and *proj\_guide* represent overlapping information
  - Every eval\_for relationship corresponds to a proj\_guide relationship
  - However, some *proj\_guide* relationships may not correspond to any *eval\_for* relationships
    - 4 So we can't discard the *proj\_guide* relationship
- Eliminate this redundancy via *aggregation* 
  - Treat relationship as an abstract entity
  - Allows relationships between relationships
  - Abstraction of relationship into new entity



# **Aggregation (Cont.)**

- Eliminate this redundancy via *aggregation* without introducing redundancy, the following diagram represents:
  - A student is guided by a particular instructor on a particular project
  - A student, instructor, project combination may have an associated evaluation





## Representing Aggregation via Schemas

- To represent aggregation, create a schema containing
  - Primary key of the aggregated relationship,
  - The primary key of the associated entity set
  - Any descriptive attributes
- In our example:
  - The schema eval\_for is:

    eval for (s ID, project id, i ID, evaluation id)
  - The schema *proj guide* is redundant.

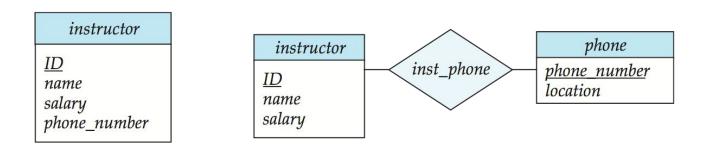


# **Design Issues**



## Entities vs. Attributes

• Use of entity sets vs. attributes



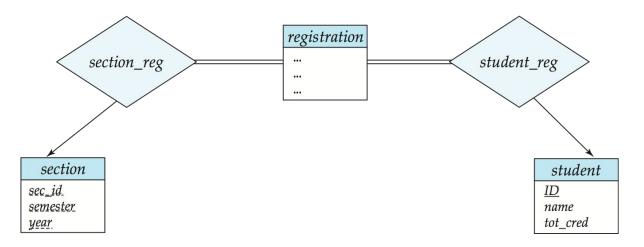
• Use of phone as an entity allows extra information about phone numbers (plus multiple phone numbers)



# Entities vs. Relationship sets

### • Use of entity sets vs. relationship sets

Possible guideline is to designate a relationship set to describe an action that occurs between entities



### Placement of relationship attributes

For example, attribute date as attribute of advisor or as attribute of student



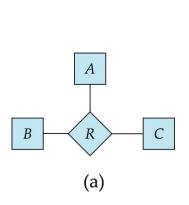
# Binary Vs. Non-Binary Relationships

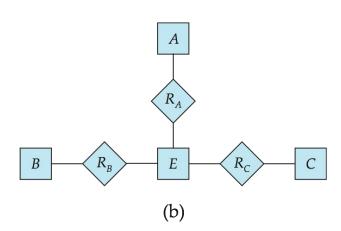
- Although it is possible to replace any non-binary (n-ary, for n > 2) relationship set by a number of distinct binary relationship sets, a n-ary relationship set shows more clearly that several entities participate in a single relationship.
- Some relationships that appear to be non-binary may be better represented using binary relationships
  - For example, a ternary relationship *parents*, relating a child to his/her father and mother, is best replaced by two binary relationships, *father* and *mother* 
    - 4 Using two binary relationships allows partial information (e.g., only mother being known)
  - But there are some relationships that are naturally non-binary
    - 4 Example: proj guide



### **Converting Non-Binary Relationships to Binary Form**

- In general, any non-binary relationship can be represented using binary relationships by creating an artificial entity set.
  - Replace *R* between entity sets A, B and C by an entity set *E*, and three relationship sets:
  - 1.  $R_A$ , relating E and A 2.  $R_B$ , relating E and B 3.  $R_C$ , relating E and C
  - Create an identifying attribute for *E* and add any attributes of *R* to *E*
  - For each relationship  $(a_i, b_i, c_i)$  in R, create
    - 1. a new entity  $e_i$  in the entity set E 2. add  $(e_i, a_i)$  to  $R_A$
    - 3. add  $(e_i, b_i)$  to  $R_B$
- 4. add  $(e_i, c_i)$  to  $R_C$







## **Converting Non-Binary Relationships (Cont.)**

- Also need to translate constraints
  - Translating all constraints may not be possible
  - There may be instances in the translated schema that cannot correspond to any instance of *R* 
    - 4 Exercise: add constraints to the relationships  $R_A$ ,  $R_B$  and  $R_C$  to ensure that a newly created entity corresponds to exactly one entity in each of entity sets A, B and C
  - We can avoid creating an identifying attribute by making E a weak entity set (described shortly) identified by the three relationship sets

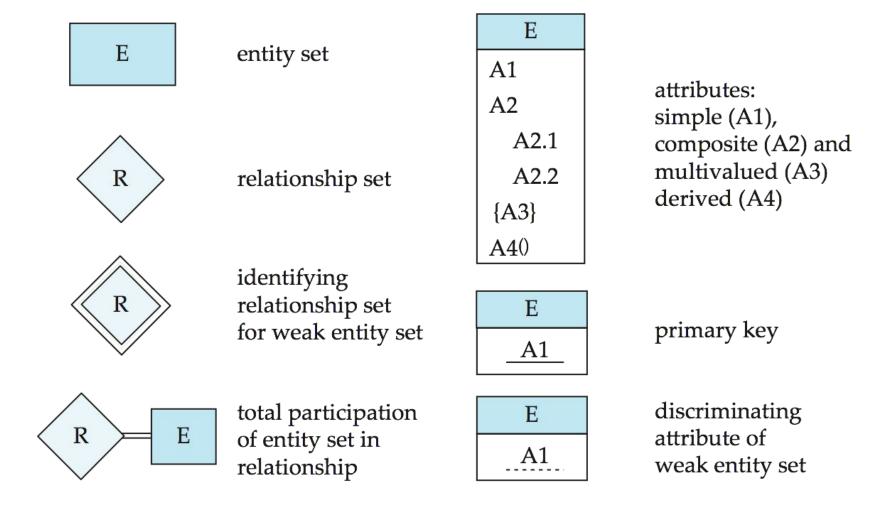


# **E-R Design Decisions**

- The use of an attribute or entity set to represent an object.
- Whether a real-world concept is best expressed by an entity set or a relationship set.
- The use of a ternary relationship versus a pair of binary relationships.
- The use of a strong or weak entity set.
- The use of specialization/generalization contributes to modularity in the design.
- The use of aggregation can treat the aggregate entity set as a single unit without concern for the details of its internal structure.

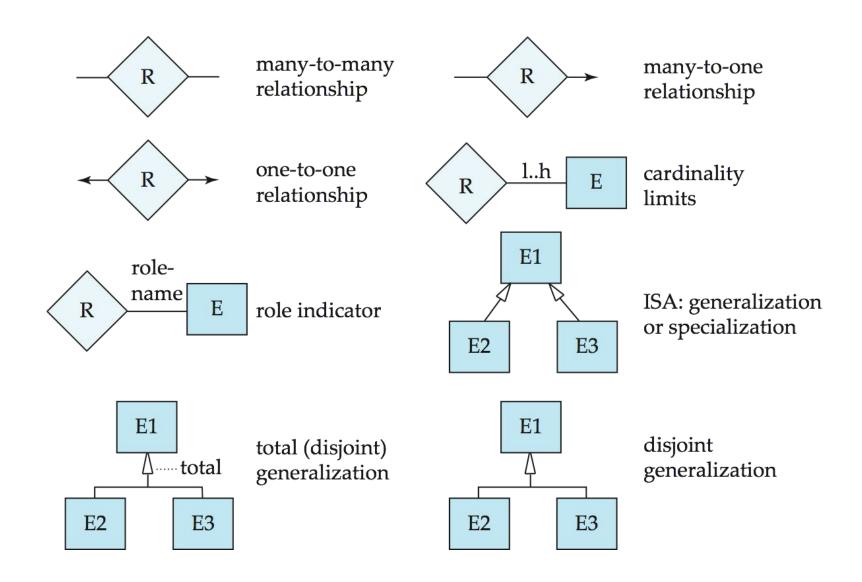


## **Summary of Symbols Used in E-R Notation**





# Symbols Used in E-R Notation (Cont.)

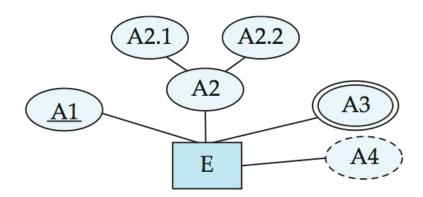




### **Alternative ER Notations**

• Chen, IDE1FX, ...

entity set E with simple attribute A1, composite attribute A2, multivalued attribute A3, derived attribute A4, and primary key A1



weak entity set



generalization



total generalization

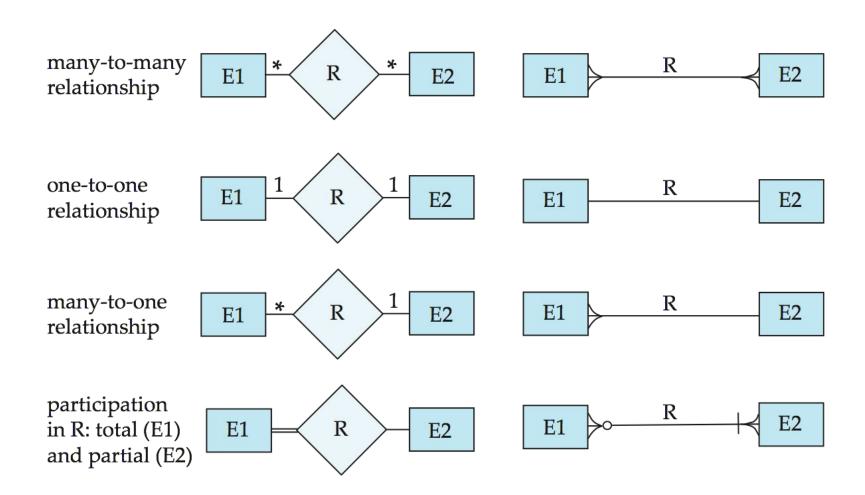




### **Alternative ER Notations**



### **IDE1FX (Crows feet notation)**





### **UML**

- **UML**: Unified Modeling Language
- UML has many components to graphically model different aspects of an entire software system
- UML Class Diagrams correspond to E-R Diagram, but several differences.



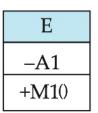
# ER vs. UML Class Diagrams

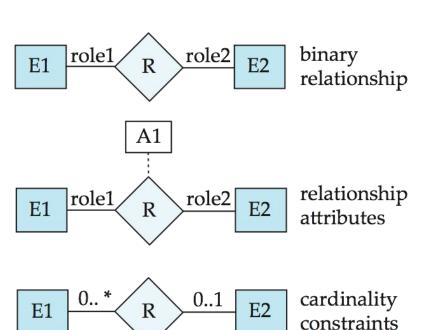
#### **ER Diagram Notation**

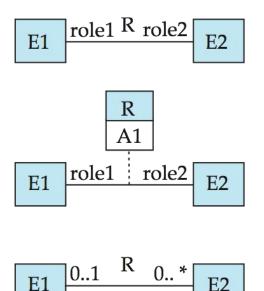


entity with attributes (simple, composite, multivalued, derived)

### **Equivalent in UML**



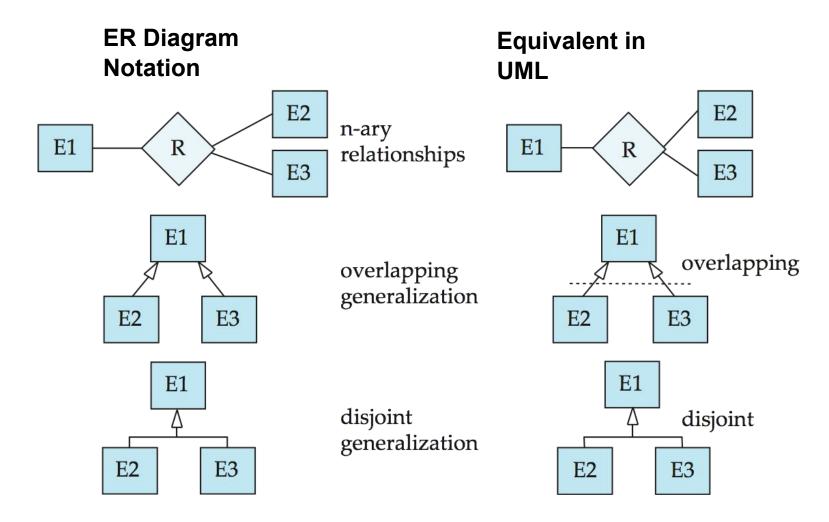




<sup>\*</sup>Note reversal of position in cardinality constraint depiction



# ER vs. UML Class Diagrams



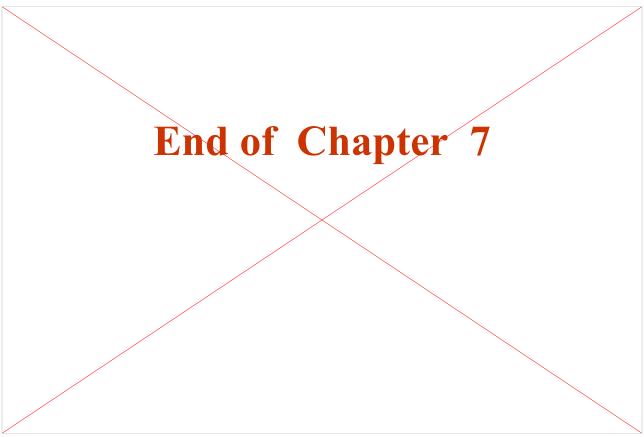
<sup>\*</sup>Generalization can use merged or separate arrows independent of disjoint/overlapping



# **UML Class Diagrams (Cont.)**

- Binary relationship sets are represented in UML by just drawing a line connecting the entity sets. The relationship set name is written adjacent to the line.
- The role played by an entity set in a relationship set may also be specified by writing the role name on the line, adjacent to the entity set.
- The relationship set name may alternatively be written in a box, along with attributes of the relationship set, and the box is connected, using a dotted line, to the line depicting the relationship set.





### **Database System Concepts, 7th Ed.**

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