



**Jawahar Education Societys Annasaheb Chudaman Patil College of  
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**SUBJECT: Analysis of Algorithms Lab**

**Implement a C program for All Pair Shortest Path using Dynamic Programming Approach. Using Dynamic Programming**

**EXPERIMENT: 07**

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Aim:- TO implement All pair shortest path using dynamic programming Approach.

Objectives:- TO implement complexity of All pair shortest path [Floyd warshall's Algorithm.]

Outcomes:- Student will be able to compute the complexity of All pair shortest path.

Hardware / Software Required: Turbo C++

Theory

- Dynamic programming:  
Dynamic programming is an algorithm design method that can be used when the solution to the problem can be viewed as the result of a sequence of decisions. It avoids recomputing solution that have already been computed. DP uses "principle of optimality."

The principle of optimality states that "in an optimal sequence of decisions or choice, each subsequence must also be optimal."

- Problem Definition:-  
Let  $G = (V, E)$  be a directed graph with  $n$  vertices. Let  $A$  be a cost adjacency matrix for  $G$  such that  $A(i, i) = 0$  for all  $i = 1$  to  $n$ .  $A(i, j)$  is a length of edge  $(i, j)$  if  $(i, j) \in E(G)$  and  $A(i, j) = \infty$  if  $i = j$  and  $(i, j) \notin E(G)$ . The all pair shortest path problem is to determine a matrix  $A(i, j)$  is length of shortest path from  $i$  to  $j$ .

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The all pair shortest path algorithm is also known as Floyd-Warshall algorithm. It is used to find all pair shortest path problem from a given weight graph. As a result of this algorithm, it will generate a matrix, which will represent the minimum distance from any node to all other nodes in the graph.

Procedure / Algorithm:

Algorithm App (W, n)

- 1)  $A \leftarrow W$  // initialize D array to W
- 2)  $P \leftarrow 0$  // initialize P array to 0
- 3) for  $k \leftarrow 1$  to  $n$   
    // computing  $D'$  from  $D$   
    4) do for  $i \leftarrow 1$  to  $n$   
        5) do for  $j \leftarrow 1$  to  $n$   
            6) if  $(A[i, j] > A[i, k] + A[k, j])$   
                7) Then  $A'[i, j] \leftarrow A[i, k] + A[k, j]$   
                8)  $P[i, j] \leftarrow k$   
            9) else  $A'[i, j] \leftarrow A[i, j]$   
    10) move  $A'$  to  $A$

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- Analysis :

$$T(n) = \sum_{k=1}^n \sum_{j=1}^n \sum_{i=1}^n$$

$$= O(n^3)$$

Conclusion:

Thus, it observed that the complexity of All pair Shortest path problem is  $O(n^3)$ .

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**Input:**

```
1 #include<stdio.h>
2 #include<conio.h>
3 int min(int,int);
4 void floyds(int p[10][10],int n)
5 {
6     int i,j,k;
7     for(k=1;k<=n;k++)
8         for(i=1;i<=n;i++)
9             for(j=1;j<=n;j++)
10                 if(i==j)
11                     p[i][j]=0;
12                 else
13                     p[i][j]=min(p[i][j],p[i][k]+p[k][j]);
14 }
15 int min(int a,int b)
16 {
17     if(a<b)
18         return(a);
19     else
20         return(b);
21 }
22 void main()
23 {
24     int p[10][10],w,n,e,u,v,i,j;
25     printf("\n Enter the number of vertices:");
26     scanf("%d",&n);
27     printf("\n Enter the number of edges:\n");
28     scanf("%d",&e);
29     for(i=1;i<=n;i++)
30     {
31         for(j=1;j<=n;j++)
32             p[i][j]=999;
33     }
34     for(i=1;i<=e;i++)
35     {
36         printf("\n Enter the end vertices of edge%d with its weight \n",i);
37         scanf("%d%d%d",&u,&v,&w);
38         p[u][v]=w;
39     }
40     printf("\n Matrix of input data:\n");
41     for(i=1;i<=n;i++)
42     {
43         for(j=1;j<=n;j++)
44             printf("%d\t",p[i][j]);
45         printf("\n");
46     }
47     floyds(p,n);
48     printf("\n Transitive closure:\n");
49     for(i=1;i<=n;i++)
50     {
51         for(j=1;j<=n;j++)
52             printf("%d\t",p[i][j]);
53         printf("\n");
54     }
55     printf("\n The shortest paths are:\n");
56     for(i=1;i<=n;i++)
57     {
58         for(j=1;j<=n;j++)
59             if(i!=j)
60                 printf("\n <%d,%d>=%d",i,j,p[i][j]);
61     }
62     getch();
63 }
```

## Implement a C program for All Pair Shortest Path using Dynamic Programming Approach. Using Dynamic Programming

### Output:

```
C:\Users\Rupesh\Documents\OS\AOA07\bin\Debug\AOA07.exe

Enter the number of vertices:3

Enter the number of edges:
4

Enter the end vertices of edge1 with its weight
1 2 3 4

Enter the end vertices of edge2 with its weight
1 5 6

Enter the end vertices of edge3 with its weight
7 8 9

Enter the end vertices of edge4 with its weight
1 2 3

Matrix of input data:
999    3    999
999    999    4
999    999    999

Transitive closure:
999    3    7
999    0    4
999    999    0

The shortest paths are:
<1,2>=3
<1,3>=7
<2,1>=999
<2,3>=4
<3,1>=999
<3,2>=999
```

**Conclusion:** Thus it is observed that **All Pair Shortest Path Problem** is  $O(n^3)$ .