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Quantum circuit

- Quantum circuit is a model of computation of Boolean functions which consists of quantum gates.
- Single qubit gate: X gate, Y gate, Z gate, H gate,

$$S := \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \text{ gate } \mathbf{T} := \begin{bmatrix} 1 & 0 \\ 0 & \frac{1+i}{\sqrt{2}} \end{bmatrix} \text{ gate } \mathbf{X}$$

Two qubit gate: CNOT gate _____



• Three qubit gate: Toffoli gate



Spectral norm

For $A \in \mathcal{L}(\mathbb{C}^n, \mathbb{C}^m)$,

$$||A|| := \max_{|\psi\rangle \in \mathbb{C}^n \colon \langle \psi|\psi\rangle = 1} \sqrt{\langle \psi|A^{\dagger}A|\psi\rangle}$$

For the singular value decomposition $A = \sum_i \lambda_i |\psi_i\rangle \langle \varphi_i|$,

$$A^{\dagger}A = \left(\sum_{i} \lambda_{i} |\varphi_{i}\rangle \langle \psi_{i}|\right) \left(\sum_{j} \lambda_{j} |\psi_{j}\rangle \langle \varphi_{j}|\right) = \sum_{i} \lambda_{i}^{2} |\varphi_{i}\rangle \langle \varphi_{i}|$$

Hence,
$$\langle \psi | A^{\dagger} A | \psi \rangle = \sum_{i} \lambda_{i}^{2} |\langle \psi | \varphi_{i} \rangle|^{2} \leq \max_{i} \lambda_{i}^{2}$$

That means ||A|| is the largest singular value of A.

For any unitary matrices
$$U \in \mathcal{L}(\mathbb{C}^m)$$
 and $V \in \mathcal{L}(\mathbb{C}^n)$, $\|UAV\| = \|A\|$.

Theorem (Universality of finite gate set)

For any unitary matrix $U \in \mathcal{L}(\mathbb{C}^{2^n})$ and $\epsilon > 0$, there is a quantum circuit with X, Y, Z, H, S, T, CNOT gates computing \widetilde{U} satisfying $\|U - \widetilde{U}\| < \epsilon$.

Proof.

- Any unitary matrix can be decomposed to a product of two-level unitary matrices.
- 2 Any two-level unitary matrix can be decomposed to a product of controlled-unitary gates.
- **3** Any controlled-untary gate can be decomposed to a product of CNOT and arbitrary single-qubit gates.
- 4 Any single-qubit gate can be approximated by X, Y, Z, H, S and T.

Two-level unitary matrix

There exist $x \neq y \in \{0, 1\}^n$ such that

$$|x\rangle \longmapsto u_{1,1} |x\rangle + u_{1,2} |y\rangle$$

 $|y\rangle \longmapsto u_{2,1} |x\rangle + u_{2,2} |y\rangle$

and for any $z \in \{0,1\}^n \setminus \{x,y\}, |z\rangle \longmapsto |z\rangle$.

Two-level unitary matrix

Theorem (Decomposition to two-level unitary matrices)

For any unitary matrix $U \in \mathcal{L}(\mathbb{C}^d)$, there is a sequence $U_1, U_2, ..., U_m$ of two-level unitary matrices such that $U = U_1 U_2 \cdots U_m$.

Proof.

We will show that there is a sequence $V_1,\ V_2,\ \dots,\ V_m$ of two-level unitary matrices such that

$$V_m V_{m-1} \cdots V_1 U = I$$
.

Since $U_i := V_i^{-1}$ is two-level unitary, this completes a proof. \square

Decomposition to two-level unitary matrix 1/3

$$U = \begin{bmatrix} u_{1,1} & u_{1,2} & u_{1,3} & u_{1,4} \\ u_{2,1} & u_{2,2} & u_{2,3} & u_{2,4} \\ u_{3,1} & u_{3,2} & u_{3,3} & u_{3,4} \\ u_{4,1} & u_{4,2} & u_{4,3} & u_{4,4} \end{bmatrix}$$

If $u_{2,1} = 0$, we skip this step. If $u_{2,1} \neq 0$, apply the two-level unitary matrix

$$V_1 = \frac{1}{z} \begin{bmatrix} u_{1,1}^* & u_{2,1}^* & 0 & 0 \\ u_{2,1} & -u_{1,1} & 0 & 0 \\ 0 & 0 & z & 0 \\ 0 & 0 & 0 & z \end{bmatrix}$$

for
$$z := \sqrt{|u_{1,1}|^2 + |u_{2,1}|^2}$$
.

Decomposition to two-level unitary matrix 2/3

$$V_1 U = \begin{bmatrix} u_{1,1} & u_{1,2} & u_{1,3} & u_{1,4} \\ 0 & u_{2,2} & u_{2,3} & u_{2,4} \\ u_{3,1} & u_{3,2} & u_{3,3} & u_{3,4} \\ u_{4,1} & u_{4,2} & u_{4,3} & u_{4,4} \end{bmatrix}$$

 $u_{i,j}$ s are not equal to those in the previous slide for $i \in \{1, 2\}$.

If $u_{3,1} = 0$, we skip this step.

If $u_{3,1} \neq 0$, apply the two-level unitary matrix

$$V_2 = \frac{1}{z} \begin{bmatrix} u_{1,1}^* & 0 & u_{3,1}^* & 0 \\ 0 & z & 0 & 0 \\ u_{3,1} & 0 & -u_{1,1} & 0 \\ 0 & 0 & 0 & z \end{bmatrix}$$

for
$$z := \sqrt{|u_{1,1}|^2 + |u_{3,1}|^2}$$
.

Decomposition to two-level unitary matrix 3/3

$$V_3V_2V_1U = \begin{bmatrix} u_{1,1} & u_{1,2} & u_{1,3} & u_{1,4} \\ 0 & u_{2,2} & u_{2,3} & u_{2,4} \\ 0 & u_{3,2} & u_{3,3} & u_{3,4} \\ 0 & u_{4,2} & u_{4,3} & u_{4,4} \end{bmatrix} = \begin{bmatrix} u_{1,1} & 0 & 0 & 0 \\ 0 & u_{2,2} & u_{2,3} & u_{2,4} \\ 0 & u_{3,2} & u_{3,3} & u_{3,4} \\ 0 & u_{4,2} & u_{4,3} & u_{4,4} \end{bmatrix}$$

 $u_{1,1} = 1$ unless $u_{2,1}$, $u_{3,1}$, $u_{4,1}$ are originally 0. In this case, apply one-level unitary for making $u_{1,1} = 1$.

Arbitrary $d \times d$ unitary matrix can be decomposed to a product of at most d(d-1)/2 two-level unitary matrices for d > 2.

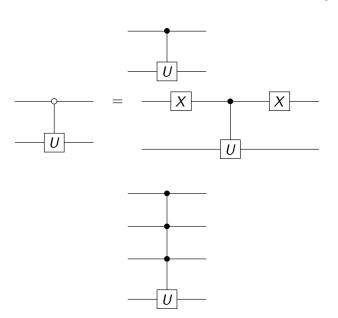
Theorem (Universality of finite gate set)

For any unitary matrix $U \in \mathcal{L}(\mathbb{C}^{2^n})$ and $\epsilon > 0$, there is a quantum circuit with X, Y, Z, H, S, T, CNOT gates computing \widetilde{U} satisfying $\|U - \widetilde{U}\| < \epsilon$.

Proof.

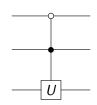
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Controlled-unitary



Special cases

0	0	0	0	0	0	٦0
1	0	0	0	0	0	0
0		0	0	0	$u_{1,2}$	0
		1	0	0		0
0	0	0		0	0	0
	0	0	0	1	0	0
0	$u_{2,1}$	0	0	0	<i>u</i> _{2,2}	0
		0	0	0	0	$1 \rfloor$
	1 0 0 0 0	1 0	1 0 0 0 $u_{1,1}$ 0 0 0 1 0 0 0 0 0 0 0 $u_{2,1}$ 0	$ \begin{array}{cccccccccccccccccccccccccccccccccccc$	1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	$ \begin{array}{cccccccccccccccccccccccccccccccccccc$



Lemma

Any $2^n \times 2^n$ two-level unitary matrix can be decomposed to a product of controlled-unitary gates.

Proof.

Assume that the two-level unitary matrix acts on a 2-dimentional subspace span($\{|x\rangle, |y\rangle\}$) for $x \neq y \in \{0, 1\}^n$.

Assume that for $i \in \{1, 2, ..., n\}$, $x_i = 1$ and $y_i = 0$. Apply at most

n-1 CNOT gates such that

$$|x\rangle \longmapsto |y \oplus e_i\rangle$$

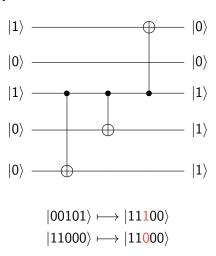
$$|y\rangle \longmapsto |y\rangle$$

$$\forall z \neq x, y \qquad \exists \tilde{z} \neq x, y \qquad |z\rangle \longmapsto |\tilde{z}\rangle,$$

Then, apply "controlled unitary" and reverse the permutation of the basis.

The first part

Let x = 00101, y = 11000.



Controlled-unitary

$$|x\rangle = |00101\rangle \longmapsto |11100\rangle$$

$$|y\rangle = |11000\rangle \longmapsto |11000\rangle$$

$$|1\rangle \longrightarrow 0$$

$$|u\rangle \longrightarrow 0$$

$$|0\rangle \longrightarrow 0$$

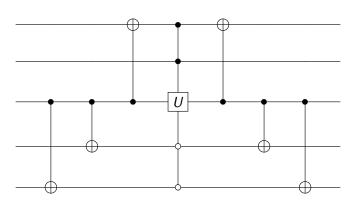
Finally, reverse the basis

$$|11100\rangle \longmapsto |00101\rangle = |x\rangle$$

 $|11000\rangle \longmapsto |11000\rangle = |y\rangle$

Whole quantum circuit

Let x = 00101, y = 11000.



$$\begin{array}{c} |00101\rangle \longmapsto |11100\rangle \longmapsto |00101\rangle \\ |11000\rangle \longmapsto |11000\rangle \longmapsto |11000\rangle \end{array}$$

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Assignments

1 Show a decomposition of

$$\frac{1}{2} \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ 1 & -1 & 1 & -1 \\ 1 & -i & -1 & i \end{bmatrix}$$

into a product of two-level unitary matrices.

2 Show a decomposition of two-level unitary

$$\begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & a & 0 & 0 & 0 & 0 & c \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & b & 0 & 0 & 0 & 0 & d \end{bmatrix}$$

into a product of controlled-unitary gates.