Refactoring Dynamic Languages

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ABSTRACT

Typically, beginner programmers do not master the style rules of the programming language they are using and, frequently, do not have yet the logical agility to avoid writing redundant code. As a result, although their programs might be correct, they can also be improved and it is important for the programmer to learn about the improvements that, without changing the meaning of the program, simplify it or transform it according to the style rules of the language. These kinds of transformations are the realm of refactoring tools. However, these tools are typically associated with sophisticated integrated development environments (IDEs) that are excessively complex for beginners.

In this paper, we present a refactoring tool designed for beginner programmers, which we made available in DrRacket, a simple and pedagogical IDE. Our tool provides simple refactoring operations for the typical mistakes made by beginners.

Keywords

Refactoring Tool, Pedagogy, Racket

1. INTRODUCTION

In order to become a proficient programmer, one needs not only to master the syntax and semantics of a programming language, but also the style rules adopted in that language and, more important, the logical rules that allow him to write simple and understandable programs. Given that beginner programmers have insufficient knowledge about these rules, it should not be surprising to verify that their code reveals what more knowledgeable programmers call "poor style," "bad smells," etc. As time goes by, it is usually the case that the beginner programmer will learn those rules and will start producing correct code written in an adequate style. However, that learning process might take a considerable amount of time and, as a result, considerable amounts of poorly-written code might be produced. It is then important to speed up this learning process by showing, from the early learning phases, how a poorly-written fragment of code can be improved.

After learning how to write code in a good style, programmers become critics of their own former code and, whenever they have to work with it, they tend to take advantage of the opportunity to restructure it so that it becomes easier to understand and conforms to the style rules. However, in most cases, these modifications are done without complete knowledge of the requirements and constraints that were considered when the code was originally written and, as result, there is a serious risk that the modifications might introduce bugs. It is thus important to help the programmer in this task so that he can be confident that the code improvements he anticipate are effectively applicable and will not change the meaning of the program. This has been the main goal of *code refactoring*.

Code refactoring is the process of changing a software system in such a way that it does not alter the external behavior of the code yet improves its internal structure [1]. Nowadays, any sophisticated IDE includes an assortment of refactoring tools, e.g., for moving methods along a class hierarchy, to extract interfaces from classes, and to transform anonymous classes into nested classes. It is important to note, however, that these IDEs were designed for advanced programmers, and that the provided refactorings require a level of code sophistication that is not present in programs written by beginners. This makes the refactoring tools inaccessible to beginners.

In this paper, we present a tool that was designed to address the previous problems. In particular, our tool (1) is usable from a pedagogical IDE designed for beginners[2][3], (2) is capable of analyzing the programmer's code and inform him of the presence of the typical mistakes made by beginners, and finally (3) can apply refactoring rules that restructure the program without changing its semantics.

We implemented our idea for a refactoring tool in DrRacket a pedagogical IDE [4][5], used in schools around the world which provides a simple and straightforward interface. DrRacket is tailor made for the Racket programming language, that is well known for its use in introductory programming courses, and currently has only one simple refactoring operation which allows renaming a variable. Our implementation significantly extends the set of refactoring operations available in Drracket.

2. RELATED WORK

Several refactoring tools were analyzed to guide our development. Our focus was on dynamic languages which are

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used in introductory courses or that have similarities with the Racket Language.

2.1 Scheme

A refactoring tool [6] for scheme, implemented in Lisp that uses two forms of information, AST (Abstract Syntax Tree) and PDG (Program Dependence Graph).

The AST represents the abstract syntactic structure of the program. While the PDG explicitly represents the key relationship of dependence between operations in the program. The graph vertices's represent program operations and the edges represent the flow of data and control between operations. However, the PDG only has dependency information of the program and relying only in this information to represent the program could create problems. For example, if two semantically unrelated statements they could be placed arbitrarily with respect to each other. Using the AST as the main representation of the program ensures that statements are not arbitrarily reordered. While the PDG is only used as a notation to prove that transformations preserve the meaning and as a quick way to retrieve needed dependence information. Contours are used with the PDG providing scope information, which is non existent in the PDG, to help reason about transformations in the PDG. With these structures it is possible to have a single formalism to reason effectively about flow dependencies and scope structure.

2.2 Haskell

HaRe [7] is a refactoring tool for Haskell that integrates with Emacs and Vim. The HaRe system uses an AST of the program to be refactored in order to reason about the transformations to do. The system has also a token stream in order to preserve the comments and the program layout by keeping information about the source code location and the comments of all tokens. It retrieves scope information from the AST, that allows to have refactoring operations that require binding information of variables. The system also allows the users to design their own refactoring operations using the HaRe API.

2.3 Python

2.3.1 Rope

Rope[8, p. 109] is a Python refactoring tool written in Python, which works like a Python library. In order to make it easier to create refactoring operations, Rope assumes that a Python program only has assignments and functions calls. Thus, by limiting the complexity of the language it reduces the complexity of the refactoring tool.

Rope uses a Static Object Analysis, which analyses the modules or scope to get information about functions. Rope only analyses the scopes when they change and it only analyses the modules when asked by the user, because this approach is time consuming.

Rope also uses a Dynamic Object Analysis that requires running the program in order to work. The Dynamic Object Analysis gathers type information and parameters passed to and returned from functions in order to get all the information needed. It stores the information collected by the analysis in a database. If Rope needs the information and there is nothing on the database the Static object inference starts trying to infer the object information. This approach makes the program run much slower, thus it is only active when the user decides. Rope uses an AST in order to store the syntax information about the programs.

2.3.2 Bicycle Repair Man

Bicycle Repair Man[1] is a refactoring tool for Python written in Python. This refactoring tool can be added to IDEs and editors, such as Emacs, Vi, Eclipse, and Sublime Text. It attempts to create the refactoring browser functionality for Python and has the following refactoring operations: extract method, extract variable, inline variable, move to module, and rename.

The tool has an AST to represent the program and a database that has information about several program entities and dependency information.

2.3.3 Pycharm Edu

Pycharm Educational Edition,[2] or Pycharm Edu, is a IDE for Python created by JetBrains, the creator of IntelliJ. The IDE was specially designed for the educational purpose, for programmers with little or no previous coding experience. Pycharm Edu is a simpler version of Pycharm community which is the free python IDE created by JetBrains. Therefore it is very similar to their normal IDEs and it has interesting features such as code completion, version control tools integration. However, it has a simpler interface when compared with Pycharm Community and other IDEs such as Eclipse or Visual Studio.

It has integrated a python tutorial and the big advantage is the possibility of the teachers creating tasks/tutorials for the students. However, the refactoring tool did not received the same care as the IDE itself. The refactoring operations are exactly the as the Pycharm community IDE which were made for more advanced users. Therefore it does not provide specific refactoring operations to beginners. The embedded refactoring tool uses the AST and def-use-relations in the refactoring operations.

2.4 Javascript

There are few refactoring tools for JavaScript but there is a framework [9] for refactoring JavaScript programs. In order to guarantee the correctness of the refactoring operation, the framework uses preconditions, expressed as query analyses provided by pointer analysis. Queries to the pointer analysis produces over-approximations of sets in a safe way to have correct refactoring operations. For example, while doing a rename operation, it over-approximates the set of expressions that must be modified when a property is renamed in a safe manner. e.g: A set `L` of object labels over-approximates a set `O` of runtime objects if every object $o \in O$ is represented by some $l \in L$.

To prove the concept, three refactoring operations were implemented, namely rename, encapsulate property, and extract module. By using over-approximations it is possible to be sure when a refactoring operation is valid. However, this approach has the disadvantage of not applying every possible refactoring operation, because the refactoring operations for which the framework cannot guarantee behavior preservation are prevented. The wrongly prevented operations accounts for 6.2% of all rejections.

3. ARCHITECTURE

[1]https://pypi.python.org/pypi/bicyclerepair/0.7.1
[2]https://www.jetbrains.com/pycharm-edu/

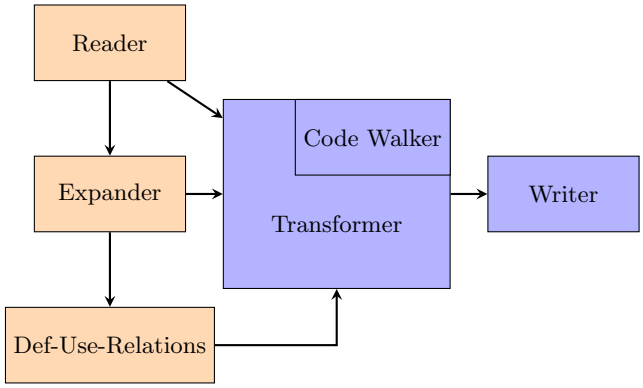

Figure 1: Main modules and information flow between modules. Unlabeled arrows represent information flow between modules.

In order to create correct refactoring operations, the refactoring tool uses two sources of information, the def-use-relations and the AST of the program. The def-use-relations represent the definition of an identifier and its usage, it is visually represented in a form of arrows in DrRacket.

The AST is represented by a list of syntax-objecs which composes the racket program.

Figure 1 summarizes the work flow of the refactoring tool where the Reader produces the non expanded AST of the program while the Expander expands the AST produced by the Reader. In order to produce the def-use-Relations it is necessary to use the expanded AST produced by the Expander because it has the correct dependency information. The Transformer uses the Code Walker to parse the ASTs and the information of the Def-Use-Relations to correctly perform the refactoring operations. Then it goes to the Writing module to produce the output in DrRacket's definitions area.

3.1 Syntax Expressions

The syntax-object list represent the AST, which provides information about the structure of the program. The syntax-object list is already being produced and used by the Racket language and in DrRacket in order to provide error information to the user. DrRacket already provides functions which computes the program's syntax-object list and uses some of those functions in the Background Check Syntax and in the Check Syntax button callback.

3.1.1 Syntax Expression tree forms

DrRacket provides functions to compute the syntax-object list in two different formats. One format is the expanded program, which computes the program with all the macros expanded. The other format is the non-expanded program and computes the program with the macros unexpanded.

The expanded program has the macros expanded and the identifier information correctly computed. However, it is harder to extract the relevant information when compared with the non expanded program.

For example, the following program is represented in the expanded program, and in the non expanded program.

Listing 1: Original Code

```
(and alpha beta)
```

Listing 2: Expanded program

```
#<syntax:2:0
 (#%app call-with-values
 (lambda ()
    (if alpha beta (quote #f)))
  print-values)>
```

Listing 3: Non-expanded program

```
#<syntax:2:0 (and alpha beta)>
```

The expanded program transforms the `and`, `or`, `when`, and `unless` forms into `if`s which makes refactoring operations harder to implement.

Racket adds internal representation information to the expanded-program which for most refactoring operations is not necessary. In addition, the expanded program has a format that is likely to change in the future. Racket is an evolving language and the expanded form is a low level and internal form of representation of the program. However, the expanded program has important information regarding the binding information that is not available in the non-expanded form and is rather useful to detect if two identifiers refer to the same binding. Additionally we do not consider macros as part of a code that could be refactored, since the refactoring tool is targeted at unexperienced programmers macros will not be used often and therefore it is not considered part of the scope of this refactoring tool capabilities.

Therefore it is desirable to use the non expanded form for the refactoring operations whenever possible and use the expanded form only for the necessary operations. Nevertheless, if we intended to create a tool that gives support to refactoring macros we would need to use the expanded program. However, there are no guarantees that would be enough to ensure the correctness of such refactoring operations due to the reflection capabilities of Racket.

3.2 Def-Use-Relations

Def-Use-Relations holds an important information in order to produce correct refactoring operations. They can be used to check whether or not there will be a duplicated name or even to compute the arguments of a function to be extracted.

DrRacket already uses the def-use-relations in the system and they are visually represented by arrows in the GUI. The def-use-relations is computed by the online-compiler that runs in the background. However, it is only computed when a program is syntactically correct.

3.3 Code-walker

The code-walker is used to parse the syntax tree represented by a syntax elements that is a list of syntax-object in Racket. A syntax-object can contain either a symbol, a syntax-pair, a datum (number, boolean or string), or an empty list. While a syntax-pair is a pair containing a syntax object as it first element and either a syntax pair, a syntax element or an empty list as the second argument. Each syntax-object has information about the line where they are defined and this information is to search for the correct elements.

Most of the time using the code-walker we are searching for a specific syntax element and location information contained in the syntax-object is used to skip the syntax blocks that are before the syntax element wanted in the first place.

The Code-walker is a core part of the refactoring tool ensuring that the selected syntax is correctly fed to the refactoring operations.

3.4 Pretty-printer

Producing correct output is an important part of the refactoring tool. It is necessary to be careful to produce indented code and we decided to use a pretty-printer that is already incorporated in the language. However, this pretty-printer does not follow the convention in the `cond` clauses should be surrounded by `[ ]` parenthesis. This is not considered a problem because Racket supports both representations. One possible solution is to use a different pretty-printer in order to keep the language convention.

3.5 Comments preservation

Preserving the comment information after a refactoring transformation is an important task of the refactoring tool. If the comment in determined place of the program changes its location affecting another structure it could confuse the programmer. However, comment preservation is not implemented yet, making it a limitation of this prototype.

One possible solution is to modify the syntax reader and add a comment node to the AST. While the new node will not be used during refactoring transformations it is used during the output part of the refactoring operation preserving the comment with the correct syntax expression.

3.6 Syntax-Parse

The Syntax-Parse function provided by Racket is rather useful for the refactoring operations regarding mainly syntax information. It provides a wide range of options to help matching the correct syntax with backtracking making it possible to have several rules to be matched in the same syntax parser, which helps to create more sophisticated rules.

4. REFACTORING OPERATIONS

In this section we explain some of the more relevant refactoring operations and some limitations of the refactoring tool.

4.1 Semantic problems

There are some known semantic problems that might occur after doing a refactoring operation. One of them occurs in the refactoring operation that removes redundant `and`s in numeric comparisons, since Racket supports more than two arguments.

Listing 4: And example

```
(and (< 1 (foo 2)) (< (foo 2) 3))
(define (foo arg)
  (displayln "foo")
  arg)
```

The refactoring transforms the code into this:

Listing 5: Refactored code

```
(< 1 (foo 2) 3)
```

Instead of applying the side effect that is displaying the the string "foo" twice it will only display it once. Therefore changing the meaning of the program.

We still kept this refactoring operation because in the vast majority of the cases this refactoring operation do not change the semantic of the program. Furthermore, the possible solution would limit excessively this refactoring operation. Considering Racket's reflection capabilities we would only apply this refactoring operation safely when the arguments of the `<` expressions, in this case the numbers 1, 2, 3, and the function `foo` were datums (number, boolean or string).

Another example of a semantic problem occurs when refactoring the following `if` expression.

Listing 6: Code sample

```
(if ?x
    (begin ?y ...)
    #f)
```

There are two different refactoring transformations possible:

Listing 7: Refactoring option 1

```
(when ?x
      ?y ...)
```

Listing 8: Refactoring option 2

```
(and ?x (begin ?y ...))
```

The first refactoring option changes the meaning of the program, because if the test expression, in this case `?x`, is false the result of the `when` expression is `#<void>`. However, the programmer may still want to choose the first refactoring option if the return value when the `if` is false is not important creating a dilemma, there are two possible refactoring operations applicable. For example, if a programmer is creating a predicate may choose the `and` version, whereas if the programmer is using another control structure and do not care for the result of the expression may prefer the `when` version.

4.2 Extract Function

Extract function is an important refactoring operation that every refactoring tool should have. In order to extract a function it is necessary to compute the arguments needed to the correct use of the function. While giving the name to a function seams quite straightforward it is necessary to check for name duplication in order to produce a correct refactoring. e.g. Having two identifiers with the same name and in the same scope produces an incorrect program and therefore modifying the meaning of the program. Lastly, computing the function body and replacing it by the call should be straightforward.

However, it raises the problem where should the function be extracted to. A function can not be defined in an expression, (e.g inside a `let`) but it could be defined in the top-level or in any other level that is accessible from the current level.

e.g: When extracting the `(+ 1 2)` to a function where should it be defined? Top-level, Level-0, level-1, or in the current level, the level-2?

Listing 9: Extract function levels

```
;;top-level
(define (level-0)
  (define (level-1)
    (define (level-2)
      (+ 1 2))
    (level-2))
  (level-1))
```

The fact is that is extremely difficult to know the answer to this question because it depends on what the user is doing and the user interpretation. Accordingly we decided that the best solution is to let the user decide where the user wants the function defined.

4.2.1 Computing the arguments

In order to compute the arguments we have to know in which scope the variables are being defined, in other words, if the variables are defined inside or outside the extracted function. The variables defined outside the function to be extracted are the candidates to be the arguments of that function. However, imported variables, whether from the language base or from other libraries does not have to be passed as arguments, to solve this we considered two possible solutions:

-Def-use-relations + Text information

-Def-use-relations + AST

The first approach is simpler to implement and more direct than the second one. However, it is less tolerant to future changes and to errors. The second one combines the def-use-relations information with the syntax information to check whether it is imported from the language or from other library.

We choose the second approach in order to provide a more stable solution to compute correctly the arguments of the new function.

4.3 Let to Define Function

We noticed the similarity of a `let` expression with a function, therefore making it easy to choose one of them by mistake. Therefore we decided to provide a refactoring operation that would make that transition simpler.

There are several let forms, but we want to focus in the ones more similar to a function, namely the `let` and the `named let`.

The `let` and the `named let` can be directly mapped to a function, however, the `named let` can be directly mapped to a named function, the `define` keyword, whereas the `let` can only be directly mapped to an anonymous function. However, we did not consider the transformation of a `let` to an anonymous function, `lambda`, making the code simpler and therefore it was not implemented yet.

A refactoring operation that transform a `named let` into a `define` function, could have syntax problems because a `let` form can be used in expressions, but the `define` can not. In the vast majority of cases this refactoring is correct, but when a `let` is used in an expression it is not correct and it changes the meaning of the program, transforming a correct program in a incorrect one. e.g.

Listing 10: Let in an expression

```
(and (let xpto ((a 1)) (< a 2)) (< b c))
```

Modifying this `named let` into a define would raise a syntax error because a define could not be used in an expression context. This could be solved by using the local keyword that is an expression like the let form. However, the `local` keyword is not used very often and might confuse the users. This reason made us keep the refactoring operation without the local keyword that works for most of the cases.

4.4 Wide-Scope Replacement

The Wide-Scope replacement extends the functionality of the extract function refactoring by replacing the duplicated of the extracted code with a function call to the new function.

The Wide-Scope replacement refactoring operation searches for the code that is duplicated of the extracted function and then replaces it for the call of the extracted function and it is divided in two steps:

- Detect duplicated code
- Replace the duplicated code

Replacing the duplicated code is the easy part, however the tool might has to compute the arguments for the duplicated code itself.

4.4.1 Detecting duplicated code

Correctly detect duplicated code is a key part for the correctness of this refactoring. Even the simplest form of duplicated code detection, when it only detects duplicated code when the code is exactly equal, may have some problems regarding the binding information. For example, if the duplicated code is inside a `let` that changes some bindings that must be taken into consideration. In order to solve the binding problem we can use functions already provided in Racket. However, that does not work if we use the program in the not expanded form to do the binding comparisons because there is not enough information for those bindings to work.

Therefore, in order to compute the correct bindings, it is necessary to use the expanded form of the program.

The naive solution is to use the expanded program to detect the duplicated code and then use this information to do the replacing of the duplicated code. However, when expanding the program Racket adds necessary internal information to run the program itself that are not visible for the user. While this does not change the detecting of the duplicated code, it adds unnecessary information that would have to be removed. In order to solve this problem in a simple way we can use the expanded code to detect the correctly duplicated code and use the non expand program to compute which code will be replaced.

However, this detection is a quadratic algorithm which might have some performance problems for bigger programs.

5. FEATURES

This section describes some of the features that improve the utility of this refactoring tool to beginner programmers.

5.1 User FeedBack

It is important to give proper feedback to the user while the user is attempting or preforming refactoring operations. Previewing the outcome of a refactoring operation is an efficient form to help the users understand the result of a refactoring before even applying the refactoring. It works by applying the refactoring operation in a copy version of

Table 1: Data Structures

Name	AST	PDG	Database	Others
Griswold	X	X		
HaRe	X			
Rope	X		X	
Bicycle	X		X	
Pycharm Edu	X		X	
Javascript				X

the AST and displaying those changes to the user. It was also studied the possibility of giving feedback to the user instead of just disabling the refactoring operation button when the tool could not perform the refactoring operation. The tool should provide information about the steps needed in order to be possible to apply that refactoring operation. However, after an analysis it was clear that these situations rarely occur in Racket language and therefore it was not implemented.

5.2 Automatic Suggestions

Beginner programmers usually do not know which refactoring operations exist or which can be applied. By having a automatic suggestion of the possible refactoring operations available the beginner programmer can have an idea what refactoring operations can be applied or not.

In order to detect possible refactoring operations it parses the code from the beginning to the end and tries to check if a refactoring is applicable. To do that it tries to match every syntax expression using syntax parse. In other words it uses brute force to check whether a expression can be applied a refactoring operation or not.

To properly display this information it highlights the source-code indicating that there is a possible refactoring. This feature could be improved by having a set of colors for the different types of refactoring operations. Moreover, the color intensity could be proportional to the level of suggestion. (e.g the recommended level to use extract function refactoring increases with the number of duplicated code found)

6. ANALYSIS

The Table 1 summarizes the data structures of the analyzed refactoring tools. It is clear that the AST of a program is an essential part of the refactoring tool information with every refactoring tool having an AST to represent the program. Regarding the PDG and Database it has mainly information about the def-use-relation of the program. The PDG has also control flow information among others.

HaRe only uses the AST as a source of information of the program. Thus, by not having the def-use-relation or a PDG it has less information to perform the refactoring operations. However, because HaRe is for the Haskell program language that is purely-functional programming language that extra information is not necessary to perform a good set of refactoring operations correctly.

Our implementation uses the same data structures, the AST and def-use-relations. The def-use-relation is often represented as a database, some refactoring tools annotate that information in the AST, some tools extract the information from the AST itself when it is possible and it is possible to extract that information from the PDG.

Some tools have more information about the program, either because they need that information to perform the refactoring operation or because they need to prove that the refactoring is correct.

Some tools like the one build by Griswold focus on the correctness of the refactoring operations. Others, focus in offering refactoring operations for professional or advanced users. However, the goal of our refactoring tool is to provide refactoring operations designed for beginners. Therefore we are not interested in refactoring operations formerly proven correct or provide refactoring operations only used in advanced and complex use cases. We intend to have simple, useful, and correct refactoring operations for the usual use cases a beginner would use. With this set of scope we exclude macros usage, classes and other complex structures not used by beginners.

7. EVALUATION

In this section we present some code examples written by beginner programmers in their final project of an introductory course and their possible improvements using the refactoring operations available in our refactoring tool. The examples show the usage of some of the refactoring operations previous presented and here is explained the motivation for their existence.

The first example is a very typical error made beginner programmers.

```
(if (>= n_plays 35)
      #T
      #F)
```

It is rather a simple refactoring operation, but nevertheless it improves the code.

```
(>= n_plays 35)
```

The next example is related with the conditional expressions, namely the `and` or `or` expressions. We decided to choose the `and` expression to exemplify a rather typical usage of this expression.

```
(and
  (and
   (eq? #t (correct-moviment? player play))
   (eq? #t (player-piece? player play)))
  (and
   (eq? #t (empty-destination? play))
   (eq? #t (empty-start? play))))
```

Transforming the code by removing the redundant `and` expression makes the code cleaner and simpler to understand.

```
(and (eq? #t (correct-moviment? player play))
     (eq? #t (player-piece? player play))
     (eq? #t (empty-destination? play))
     (eq? #t (empty-start? play)))
```

However, this code could still be improved, the `(eq? #t ?x)` is a redundant way of simple writing `?x`.

```
(and (correct-moviment? player play)
     (player-piece? player play)
     (empty-destination? play)
     (empty-start? play))
```

While a student is still learning and experiencing it is easier to confuse whether or not an expression needs the `begin` expression or not. The `when`, `cond` and `let` expressions have a implicit `begin` expression and as a result it is not necessary to add the `begin` expression. Moreover, sometimes they still kept the `begin` keyword because often use error and trial approach in writing code. Our refactoring tool checks for those mistakes and corrects them.

```
(if (odd? line-value)
    (let ((internal-column (sub1 (/ column 2))))
      (begin
        (if (integer? internal-column)
            internal-column
            #f)))
    (let ((internal-column (/ (sub1 column) 2)))
      (begin
        (if (integer? internal-column)
            internal-column
            #f))))
```

It is a simple refactoring operation and it does not have a big impact, however it makes the code clear and helps the beginner programmer to learn that a `let` does not need a `begin`.

```
(if (odd? line-value)
      (let ((internal-column (sub1 (/ column 2))))
        (or (integer? internal-column)
            internal-column))
      (let ((internal-column (/ (sub1 column) 2)))
        (or (integer? internal-column)
            internal-column)))
```

The next example shows a nested `if`. Nested `if`s are difficult to understand, error prone, and a debugging nightmare. It is much simpler to have a `cond` expression. In addition, every true branch of this nested if contains `if` expressions that are or expressions and by refactoring those `if` expressions to `or`s it makes the code simpler to understand.

```
(define (search-aux? board line column piece)
  (if
    (> column 8)
    #f
    (if
      (= line 1)
      (if
        (eq? (house-board
             board 1 column) piece)
        #t
        (search-aux? board line
          (+ 2 column) piece))
      (if (= line 2)
        (if
          (eq? (house-board
            board 2 column) piece)
          #t
          (search-aux? board line
            (+ 2 column) piece))
        (if (= line 3)
          (if
            (eq? (house-board
              board 3 column) piece)
            #t
            (search-aux? board line
              (+ 2 column) piece))
          (if (= line 4)
            (if
              (eq? (house-board
                board 4 column) piece)
              #t
              (search-aux? board line
                (+ 2 column) piece))
            (if (= line 5)
              (if
                (eq? (house-board
                  board 5 column) piece)
                #t
                (search-aux? board line
                  (+ 2 column) piece))
              (if (= line 6)
                (if (eq? (house-board
                      board 6 column) piece)
                  #t
                  (search-aux? board line
                         (+ 2 column) piece))
                (if (= line 7)
                  (if (eq? (house-board
                        board 7 column) piece)
                    #t
                    (search-aux? board line
                      (+ 2 column) piece))
                  (if (= line 8)
                    (if (eq? (house-board
                      board 8 column) piece)
                      #t
                     (search-aux? board
                       line (+ 2 column) piece))
                  null)))))))))))
```

This transformations drastically reduced the lines of code, from 59 to 28 that accounts for a 52% reduction of the lines of code.

```
(define (search-aux? board line column piece)
 (cond
 [(> column 8) #f]
 [(= line 1)
  (or (eq? (house-board board 1 column) piece)
        (search-aux? board line (+ 2 column) piece))]
 [(= line 2)
  (or (eq? (house-board board 2 column) piece)
        (search-aux? board line (+ 2 column) piece))]
 [(= line 3)
  (or (eq? (house-board board 3 column) piece)
        (search-aux? board line (+ 2 column) piece))]
 [(= line 4)
  (or (eq? (house-board board 4 column) piece)
        (search-aux? board line (+ 2 column) piece))]
 [(= line 5)
  (or (eq? (house-board board 5 column) piece)
        (search-aux? board line (+ 2 column) piece))]
 [(= line 6)
  (or (eq? (house-board board 6 column) piece)
        (search-aux? board line (+ 2 column) piece))]
 [(= line 7)
  (or (eq? (house-board board 7 column) piece)
        (search-aux? board line (+ 2 column) piece))]
```

```
  [(= line 8)
   (or (eq? (house-board board 8 column) piece)
         (search-aux? board line (+ 2 column) piece))]
  [else null]))
```

However, this code could still be further improved by refactoring it into a `case`.

The examples presented above appear repeatedly in almost every code submission of this final project supporting the need to proper support to beginner programmers.

8. CONCLUSION

A refactoring tool designed for beginner programmers would benefit them by providing a tool to restructure the programs and improve what more knowledgeable programmers call "poor style," "bad smells," etc. In order to help those users it is necessary to be usable from a pedagogical IDE, to inform the programmer of the presence of the typical mistakes made by beginners, and to correctly apply refactoring operations preserving semantics.

Our solution tries to help those users to improve their programs by using the AST of the program and the def-use-relations to create refactoring operations that do not change the program's semantics. This structure is then used to analyze the code to detect typical mistakes using automatic suggestions and correct them using the refactoring operations provided.

There are still some improvements that we consider important for the user. Firstly, the detection of duplicated code is still very naive and improving to understand if two variables represent the same even if the names are different or even if the order of some commutative expressions is not the same would make a huge improvement on the automatic suggestion. Then it is possible to improve the automatic suggestion of refactoring operations by having different colors for different types of refactoring operations. With a lower intensity for low "priority" refactoring operations and a high intensity for higher "priority". Thus giving the user a better knowledge of what is a better way to solve a problem or what is a strongly recommendation to change the code. Lastly it would be interesting to detect when a developer is refactoring in order to help the developer finish the refactoring by doing it automatically [10].

9. ACKNOWLEDGMENTS

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