

C & A

Chap. II

Permutation and Combination

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Generating Permutations



- Can we start from the simple thing first?

- The permutation of {1}

1

Induction

High complexity!

Waste of memory!

- The permutation of {1 2} 2 1 2

1 2

2 1

- The permutation of {1 2 3} 3 1 3 2 3

1 2 3

3 2 1

1 3 2

2 3 1

3 2 3 1 3

3 1 2

2 1 3

Generate the permutations of {1,2,...n} based on the permutations of {1,2,...,n-1}

Lexicographic Order

123
132
213

[Eg] For alphabet $\{1,2,3\}$, smaller digits are in the left, so the permutations in lexicographic order are:

123, 132, 213, 231, 312, 321.

A permutation could be regarded as a string, the string could have **prefix** and **suffix**.

The so called next permutation of this one means that there's no more other permutations among this and the next.

This means that this one and next one should have a **common prefix** which is as long as possible, and the changes are limited to **suffixes** as short as possible.

Generating Permutations

单选题 1分

What is the next permutation
for sequence "46739521"
with lexicographical order
_____?

- ☐ A 46791235
- ☐ B 46751238
- ☐ C 46751239
- ☐ D 46759321

What is the next permutation for sequence “35418762” with lexicographical order?

- ☐ A 35428761
- ☐ B 35421678
- ☒ C 35421678
- ☐ D 35421679

Lexicographic Order

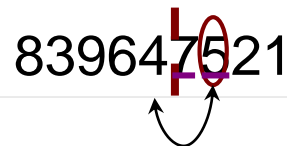
- **[Eg]** Calculate the next permutation of 839647521
- 1 Find the first decrease point from right to left: 4

839647521



- 2. Exchange Find the smallest number larger than 4 in the suffix

839647521



- 3. Overturn the suffix 839651247

- The next permutation is: 839651247

Thinking question: How many permutations are before "839647521" in lexicographical order?

正常使用主观题需2.0以上版本雨课堂

作答

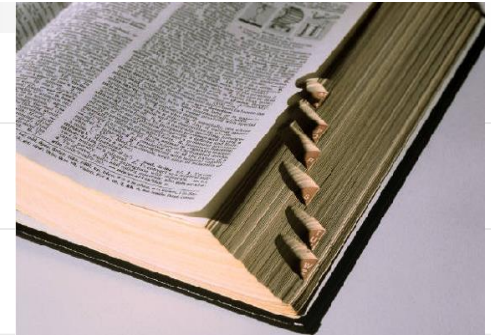
7

Thinking question: How many permutations are before "839647521" in lexicographical order?

$$\begin{aligned} &7 \times 8! + 2 \times 7! + 6 \times 6! + 4 \times 5! + 2 \times 4! + 3 \times 3! + 2 \times 2! + 1 \times 1! \\ &= 297191 \end{aligned}$$

Generating r-combinations

- Lexicographic order (dictionary order)
 - “abz” precedes “acb” in the dictionary
 - 129 precedes 132 in the lexicographic order
- Let A and B be two r-combinations of the set $\{1, 2, \dots, n\}$.
 - *A precedes B* in the lexicographic order provided the smallest integer which is in their union $A \cup B$ but not in their intersection $A \cap B$ is in A.
 - 5-combinations of $\{1, 2, \dots, 8\}$, $A = \{2, 3, 4, 7, 8\}$, $B = \{2, 3, 5, 6, 7\}$.
 - Then the smallest element which is in one but not both of the sets is 4.
 - Hence A precedes B in the lexicographic order.



A theorem

- Consider 5-combinations of $\{1,2,\dots,9\}$. What 5-combination immediately follows 12389?
– 12456
- Let $a_1a_2\dots a_r \neq (n-r+1)(n-r+2)\dots n$ be an r -combination of $\{1, 2, \dots, n\}$.
- Let k be the largest integer such that $a_k < n$ and a_k+1 is different from each of a_1, a_2, \dots, a_r .
- Then the r -combination which is the immediate successor of $a_1a_2\dots a_r$ in the lexicographic ordering is
- $a_1\dots a_{k-1}(a_k+1)(a_k+2)\dots(a_k+r-k+1).$

Consider 5-combinations of $\{1,2,\dots,9\}$. What 5-combination immediately **before** 34579?

☐ A 34569

☒ B 34578

☐ C 34568

☐ D 34597

Algorithm for generating the r-combinations

Begin with the r-combination $a_1a_2\dots a_r=12\dots r$. while $a_1a_2\dots a_r \neq (n-r+1)(n-r+2)\dots n$, do

1) determine the largest integer k such that $a_k+1 \leq n$ and a_k+1 is not one of a_1, a_2, \dots, a_r .

2) replace $a_1a_2\dots a_r$ with the r-combination
 $a_1a_2\dots a_{k-1}(a_k+1)(a_k+2)\dots(a_k+r-k+1)$.

Generate the 4-combinations of $S = \{1,2,3,4,5,6\}$.

1234;	1245;	1345;	1456;	2356;
1235;	1246;	1346;	2345;	2456;
1236;	1256;	1356;	2346;	3456.

Generating Permutations



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Waste of memory!

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- The permutation of {1 2 3} 3 1 3 2 3

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3 2 1

1 3 2

2 3 1

3 2 3 1 3

3 1 2

2 1 3

Generate the permutations of {1,2,...n} based on the permutations of {1,2,...,n-1}

Generating Permutations

– The permutation of {1 2 3 4}

1	2	3↔4	3	1	2↔4	2	3	1↔4	
1	2↔4	3	3	1↔4	2	2	3↔4	1	
1↔4	2	3	3	4↔1	2	2↔4	3	1	
4	1	2↔3	4	3	1↔2	4	2	3↔1	
4↔1	3	2	4↔3	2	1	4↔2	1	3	
1	4↔3	2	3	4↔2	1	2	4↔1	3	
1	3	4↔2	3	2	4↔1	2	1	4↔3	
1↔3	2	4	3↔2	1	4	2	1	3	4

• Each permutation other than the first one is obtained from the preceding one by **switching two adjacent numbers**.

- Move the largest number 4 between two ends
- When the end is met, the largest number 4 will be fixed for a step and the switching of the second largest number 3 makes 4 movable again.

Mobile Integer

- Given an integer k we assign a direction to it by writing an arrow above it pointing to the left or to the right. Consider a permutation of $\{1, 2, \dots, n\}$ in which each of the integers is given a direction.

- The integer k is called *mobile* if its arrow points to a smaller integer adjacent to it.

$\overrightarrow{2} \overrightarrow{6} \overrightarrow{3} \overrightarrow{1} \overrightarrow{5} \overrightarrow{4}$ only 3, 5, and 6 are mobile.

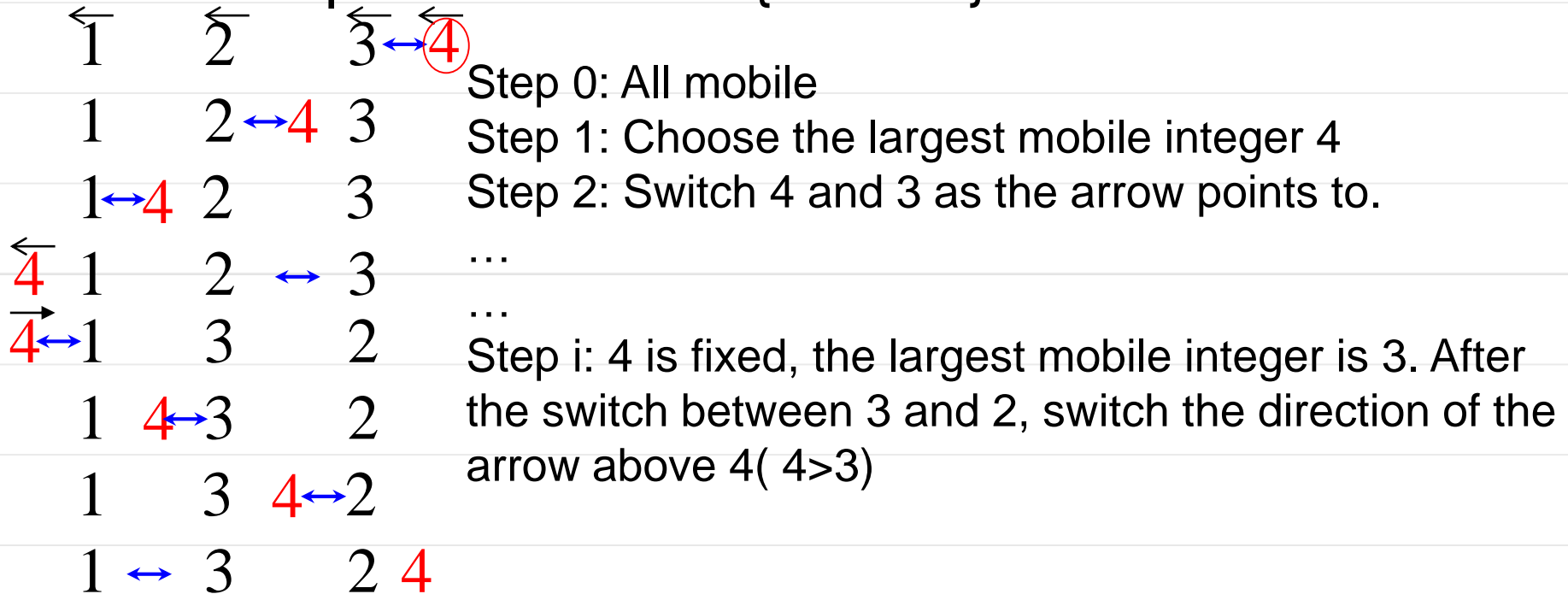
- Integer 1 can never be mobile since there is no integer smaller than 1.
- The integer n is always mobile, except in two cases:
 - n is the first integer and its arrow points to the left.
 - n is the last integer and its arrow points to the right.

$\overleftarrow{4} \overleftarrow{1} \overleftarrow{2} \overleftarrow{3}$

$\overrightarrow{3} \overleftarrow{2} \overleftarrow{1} \overrightarrow{4}$

Generating Permutations

– The permutation of {1 2 3 4}



• Each permutation other than the first one is obtained from the preceding one by **switching two adjacent numbers**.

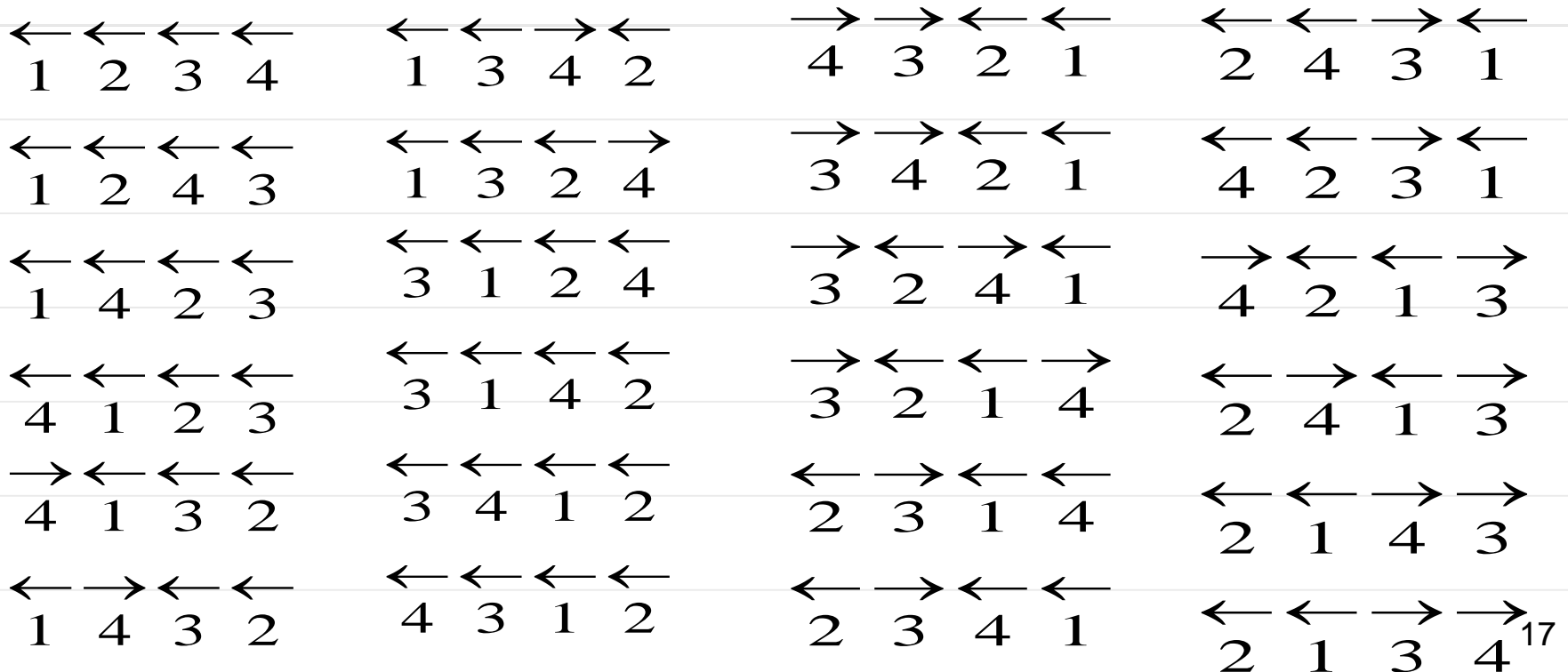
- Move the largest number 4 between two ends
- When the end is met, the largest number 4 will be fixed for a step and the switching of the second largest number 3 makes 4 movable again.

Steinhaus–Johnson–Trotter algorithm

Begin with $1 \xleftarrow{\quad} 2 \xleftarrow{\quad} \cdots \xleftarrow{\quad} n$.

While there exists a mobile integer, do

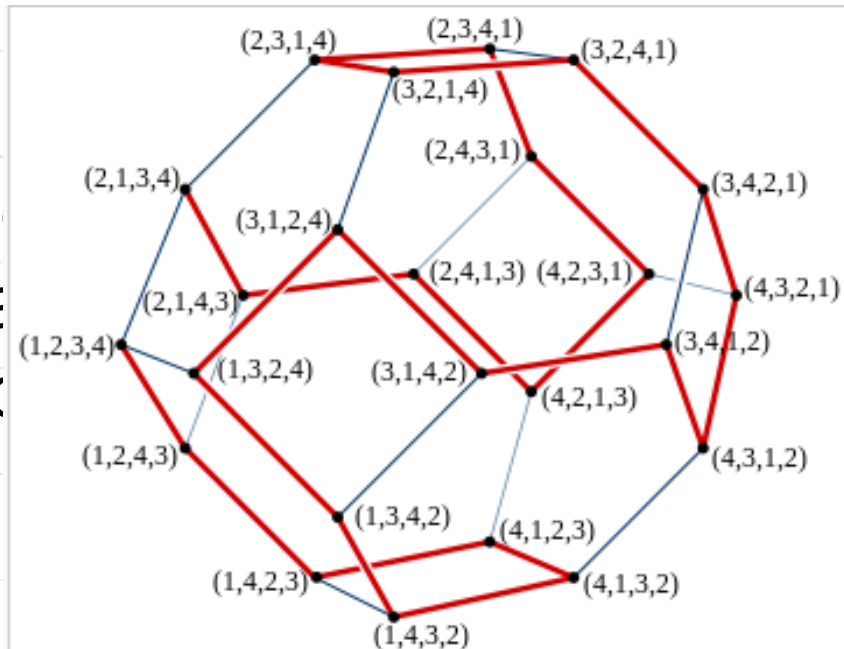
- 1) find the largest mobile integer m .
- 2) switch m and the adjacent integer its arrow points to.
- 3) switch the direction of all integers p with $p > m$.



SJT Algorithm

- **Steinhaus–Johnson–Trotter algorithm**

- Hugo Steinhaus[1958], H. F. Trotter[1962], S. M. Johnson[1963]
- Each permutation in the sequence generated differs from the previous permutation by swapping two adjacent elements of the sequence
- Equivalently, this algorithm generates a Hamiltonian path in the Cayley graph of the symmetric group



The **Hamiltonian path** in the **Cayley graph** of the symmetric group generated by the Steinhaus–Johnson–Trotter algorithm

Inversions in Permutations

- Let $i_1 i_2 \dots i_n$ be a permutation of the set $\{1, 2, \dots, n\}$. The pair (i_k, i_l) is called an **inversion** if $k < l$ and $i_k > i_l$.
- An inversion in a permutation corresponds to a pair of numbers which are out of their natural order.
- The only permutation of $\{1, 2, \dots, n\}$ with no inversions is $1 \ 2 \ \dots \ n$.
- Example: the inversions in $(3 \ 2 \ 1)$ are $(3, 2), (3, 1), (2, 1)$

The number of inversions

- Let a_j denote the **number of inversions**. It equals the number of integers which precede j in the permutation but are greater than j ; it measures how much j is out of order.
- The sequence of numbers a_1, a_2, \dots, a_n is called the inversion sequence of the permutation $i_1 i_2 \dots i_n$.
- The number $a_1 + a_2 + \dots + a_n$ measures the disorder of a permutation.

Examples of inversions

- Consider the permutation 31524. It has four inversions, namely (3, 1), (3, 2), (5, 2), (5, 4).
- Inversion records the number of integers which precede j in the permutation but are greater than j
- The inversion sequence of the permutation 31524
 - 1: 3 is larger than 1 and before 1 **1**
 - 2: 3 and 5 are larger than 2 and before 2 **2**
 - 3: no number is before 3 **0**
 - 4: 5 is larger than 4 and before 4 **1**
 - 5: no number is larger than 5 **0**

A theorem

Let b_1, b_2, \dots, b_n be a sequence of integers satisfying

$$0 \leq b_1 \leq n-1, 0 \leq b_2 \leq n-2, \dots, 0 \leq b_{n-1} \leq 1, b_n = 0.$$

Then there exists a unique permutation of $\{1, 2, \dots, n\}$ whose inversion sequence is b_1, b_2, \dots, b_n .

With the inversion sequence as 1, 2, 0, 1, 0, how to construct the permutation?

- 5: no number is larger than 5
- 4: 1 larger number before
- 3: no larger number is before 3
- 2: 2 larger numbers before
- 1: 1 larger number before

3 1 5 2 4

Algorithm I

- n : write down n .
- $n-1$: consider b_{n-1} . We are given that $0 \leq b_{n-1} \leq 1$
If $b_{n-1} = 0$, then $n-1$ must be placed **before** n . If $b_{n-1} = 1$, then $n-1$ must be placed **after** n .
- $n-2$: consider b_{n-2} . We are given that $0 \leq b_{n-2} \leq 2$
If $b_{n-2} = 0$, then $n-2$ must be placed **before** the two numbers from step $n-1$.
If $b_{n-2} = 1$, then $n-2$ must be placed **between** the two numbers from step $n-1$.
If $b_{n-2} = 2$, then $n-2$ must be placed **after** the two numbers from step $n-1$
.....
- **1**: We must place 1 after the b_{1st} number in the sequence constructed in step $n-1$.

Comments on Algorithm I

- The algorithm can determine the unique permutation of $\{1, 2, \dots, n\}$ whose inversion sequence is b_1, b_2, \dots, b_n .
- The **disadvantage** of this algorithm is that the location of each integer in the permutation are not known until the very end; only the relative positions of the integers remain fixed throughout the algorithm.
- The inversion sequence of the permutation 31524 is 1, 2, 0, 1, 0.
 - 1: 3 is larger than 1 and before 1 **1** 1 larger number is before “1”
 - 2: 3 and 5 are larger than 2 and before 2 **2** 2 larger number is before “2”
 - 3: no number is before 3 **0** 0 larger number is before “3”
 - 4: 5 is larger than 4 and before 4 **1**
 - 5: no number is larger than 5 **0**

Example

- Determine the permutation of $\{1, 2, 3, 4, 5, 6, 7, 8\}$ whose inversion sequence is 5, 3, 4, 0, 2, 1, 1, 0.
- $b_1=5$: put 1 in the 6th empty location
- $b_2=3$: put 2 in the 4th empty location
- $b_3=4$: put 3 in the 5th empty location
- $b_4=0$: put 4 in the 1st empty location
- $b_5=2$: put 5 in the 3rd empty location
- $b_6=1$: put 6 in the 2nd empty location
- $b_7=1$: put 7 in the 2nd empty location
- $b_8=0$: put 8 in the 1st empty location

					1		
			2		1		
			2		1	3	
4			2		1	3	
4			2	5	1	3	
4		6	2	5	1	3	
4		6	2	5	1	3	7
4	8	6	2	5	1	3	7

Algorithm II: Inversion Method

- Begin with n empty locations which label $1, 2, \dots, n$ from left to right
- 1: put 1 in location number $b_1 + 1$
- 2: put 2 in the $(b_2 + 1)$ st empty location
-
- k : (general step) counting from the left we put k in the $(b_k + 1)$ st empty location.
- n : put n in the one remaining empty location.

Comments on Algorithm II

- The algorithm can determine the unique permutation of $\{1, 2, \dots, n\}$ whose inversion sequence is b_1, b_2, \dots, b_n .
- The **advantage** of this algorithm is that the location of each integer in the permutation can be determined.
 - Algorithm I: insert the integers one by one
 - Algorithm II: find the locations of the integers one by one.

Switches and inversions

- Bring the permutation 361245 to 123456 by successive switches of adjacent numbers.
- The inversion sequence is 220110. Hence there will be 6 times of switch.

3	6	1	2	4	5
3	1	6	2	4	5
1	3	6	2	4	5
1	3	2	6	4	5
1	2	3	6	4	5
1	2	3	4	6	5
1	2	3	4	5	6

Industrial Support

—Java

```
for(List<String> list : Permutation.of(Arrays.asList("a", "b", "c")))
    System.out.println(list);
}
```

```
for(List<Integer> list : Combination.of(Arrays.asList(1, 2, 3, 4, 5), 3))
    System.out.println(list);
}
```

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Test Run

String Permutations

Dr. James McCaffrey

Code download available at: [TestRun2006_12.exe](#) (161 KB)

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Improved Permutations with the BigInteger Data Type

The major challenge when working with permutations is that the factorial function gets very, very large very, very quickly. The BigInteger data type was introduced in the .NET Framework 4.0; it enables the writing of effective permutation code.

By James McCaffrey ■ 09/04/2012



Finite Fields and Their Applications

Volume 17, Issue 1, January 2011, Pages 51–67



Ann. Comb. 14 (2010) 3–16

DOI 10.1007/s00026-010-0042-9

Published online February 16, 2010

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Annals of Combinatorics

On constructing permutations of finite fields ☆

Amir Akbary^a, Dragos Ghioca^b, Qiang Wang^c

^a Def Menemui Matematik (Discovering Mathematics)
^b Def Vol. 32, No. 2: 51–56 (2010)
^c Sch

New Recursive Circular Algorithm

Sharmila Karim¹, Zurni On
Khairil Iskandar Othman⁴, a

¹²³ College of Art and S

¹⁷ Hoang Chi Thanh et al

From Permutations to Iterative Permutations

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Permutations Generated by Stacks and Deques

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A Two-level Algorithm for Generating Multiset Permutations

Tadao Takaoka

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Christchurch, New Zealand

ISSN:2231-0711

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IJCSET [July 2012] Vol 2, Issue 7,1310-1315

rmutations in O(1) time for each
memory requirement. There a

Common Permutation generators

In C++ standard library, `next_permutation`, `prev_permutation`, could generate permutations in lexicographic order.

```
#include <algorithm>
```

```
bool next_permutation( iterator start, iterator end );
```

```
bool prev_permutation( iterator start, iterator end );
```

The `next_permutation()` function attempts to transform the given range of elements `[start,end)` into the next lexicographically greater permutation of elements. If it succeeds, it returns `true`, otherwise, it returns `false`.

http://www.slyar.com/blog/stl_next_permutation.html

Thinking Question

- Find the 2020th permutation with nine numbers of 1-9 in lexicographic order?
- We will post the GOOD HW as post session, if you do not want your HW to be posted, please inform me before Monday (Sep.28) noon.
- OJ Tasks
 - Lexicographic Order and SJT
 - Office Hour
 - Sep.28 after class TA: Kai Su

Thank you!