

Probabilistic Analysis

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PA & RA

▶ Probabilistic Analysis

- ▶ Using probabilistic models (Appendix C-3, C-4, Ch 5.4.2)
- ▶ Using indicator random variable (Ch 5.2)
- ▶ Bucket sort (Ch 8.4)
- ▶ Hiring problem (Ch 5.1) & on-line hiring problem (Ch 5.4.4)

▶ Randomized Algorithm

- ▶ Randomized hiring problem (Ch 5.3 p122-p124)
- ▶ Randomized Select (Ch 9.2)



Probabilistic Analysis

- ▶ Probabilistic analysis is the use of probability in the analysis of problems.
- ▶ We often use probabilistic analysis to analyze the running time of an algorithm.
 - ▶ We call it **expected running time**: it is taken over the distribution of the possible inputs, i.e., averaging the running time over all possible inputs (**average-case running time**).
- ▶ Two ways of probabilistic analysis:
 - ▶ Using probabilistic models directly
 - ▶ Using indicator random variables



Probabilistic Models

- ▶ There is a tight relationship between execution of algorithms and experiments:
 - ▶ Bernoulli Trials
 - ▶ Geometric Distribution
 - ▶ Binomial Distribution
 - ▶ Balls & Bins
 - ▶ ...



Jakob Bernoulli

Random Variable

- ▶ **A random variable X** : is a function from a finite or countably infinite sample space S to the real numbers. It associates a real number with each possible outcome of an experiment.
- ▶ **Expectation**: the expected value of a discrete random variable X is $E[X] = \sum_x x \Pr\{X = x\}$.
-- the “average” of the values X takes on.

- ▶ **Example:** The possible outcomes for one coin toss can be described by the sample space $\Omega = \{heads, tails\}$. We can introduce a real-valued random variable Y as follows:

- ▶ $Y(\omega) = \begin{cases} 1, & \text{if } \omega = head, \\ 0, & \text{if } \omega = tail. \end{cases}$
- ▶ $\Pr\{Y = y\} = \begin{cases} 0.5, & \text{if } y = 1, \\ 0.5, & \text{if } y = 0. \end{cases}$
- ▶ $E[Y] = 1 \times \Pr\{Y = 1\} + 0 \times \Pr\{Y = 0\} = 0.5$



Bernoulli Trial

- ▶ **A Bernoulli Trial:** an experiment with only two possible outcomes: success (with a probability of p) and failure ($q = 1 - p$)
- ▶ If we treat the outcome of Head as a success and Tail as a failure, coin toss is a Bernoulli trial.



- ▶ **A sequence of Bernoulli trials:**
 - ▶ Each with a probability of p for success
 - ▶ Q: How many trials occur before we obtain a success?
- ▶ Let X be the number of trials needed to obtain a success.
 - ▶ $Pr\{X = k\} = q^{k-1}p$ (geometric distribution)
 - ▶ $E[X] = \sum_{k=1}^{\infty} kq^{k-1}p = \frac{p}{q} \sum_{k=0}^{\infty} kq^k = \frac{p}{q} \times \frac{q}{(1-q)^2} = \frac{1}{p}$

Balls & Bins

► Balls and Bins

- Toss identical balls into b bins.
- Tosses are independent.
- Each toss the ball is equally likely ($= \frac{1}{b}$) to end up in any bin.
- Useful model for analyzing hashing (Ch 11) and searching (problem 5.2) etc.



1



2



3



b



Balls & Bins

- ▶ How many trials we need to have one ball fall in a given bin?
- ▶ Let X be the number of trials needed to hit a given bin.



1



2



3



b



Balls & Bins

- ▶ How many balls must be tossed until every bin contains at least one ball?
 - ▶ Consider b stages:



1



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3



b



Balls & Bins

- ▶ How many balls must be tossed until every bin contains at least one ball?
 - ▶ Consider b stages
 - ▶ At the i^{th} stage, $i - 1$ bins contains balls & $b - i + 1$ bins are empty, let n_i be the number of tosses to hit a new bin.
 - ▶ The probability of a success toss at the i^{th} stage is:
 - ▶ $E[n_i] =$
 - ▶ $E[n_1 + n_2 + \cdots + n_b] = \sum_{i=1}^b E[n_i] =$



1



2



3



b



In-class exercise

- ▶ 5.2 (b) and (d): Searching for a value x in an unsorted array A consisting of n elements.
 - ▶ Strategy: pick a random index i into A . If $A[i] = x$, then we terminate; otherwise we continue the search by picking another index, until we find an index j such that $A[j] = x$ or we have checked all elements in A .
 - ▶ Q1: Suppose there is exactly one index i such that $A[i] = x$. What is the expected number of indices into A we must pick before we find x ?
 - ▶ Q2: Suppose there are no indices i such that $A[i] = x$. What is the expected number of indices into A that we must pick before we have checked all elements of A ?
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Indicator Random Variable

► *Definition*

- Given a sample space S and an event A , the indicator random variable $I\{A\}$ or X_A associated with event A is defined as: $I\{A\} = X_A = \begin{cases} 1 & \text{if } A \text{ occurs,} \\ 0 & \text{otherwise.} \end{cases}$
- $E[X_A] = 1 * Pr\{A\} = Pr\{A\}$
- The random variable Y in the toss coin example is also an indicator random variable.

► *Probabilistic Analysis*

- Expected (Average-case) running time or the number of executions of certain statements
 - $E[X + Y] = E[X] + E[Y]$ (X and Y can be dependent)
 - Simplified the calculation by using indicator random variables
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Inversions

- ▶ **Example:** What is the expected number of inversions in an input array of size n for insertion sort? Assuming all permutations of the input array are equally likely.
 - ▶ Why?
 - ▶ The running time of insertion sort is in proportion to the number of inversions.
 - ▶ Let X be the number of total inversions, $X_{ij} = I\{i, j \text{ is a pair of inversion}\}$ (i.e., $i < j$ and $A[i] > A[j]$).
 - ▶ $X = \sum_{i=1}^{n-1} \sum_{j=i+1}^n X_{ij}$
 - ▶ $E[X_{ij}] = \frac{1}{2}$
 - ▶ $E[X] =$
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Group Exercise: 5.2 (e)

- ▶ Search A for x in order, consider $A[1], A[2], \dots, A[n]$ until either find $A[i] = x$ or reach to the end of the input array.
- ▶ Q: suppose there is exactly one index i such that $A[i] = x$. What is the average-case running time of this algorithm?

