

CSE316

Computer Systems Security

Lab 3. Leakage-resilient AKE

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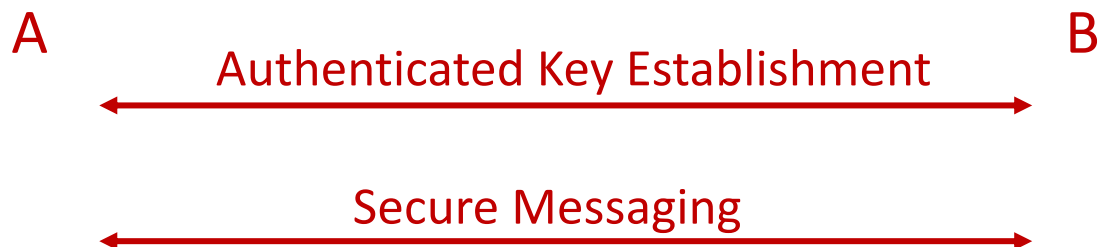
Outline

- Background
- Method

1. Background

Authenticated key exchange

- Security protocols protect communications.



- Phase 1, Authenticated key establishment
 - Establish authenticated shared keys
- Phase 2, Secure messaging
 - Use AES, HMAC, etc., to protect messages under shared keys.

Authenticated key establishment **underlies** the security of communications.

Authenticated key exchange

- Authenticated key establishment
 - A cryptographic mechanism that provides two or more parties communicating over an open network with a shared secret key.
- Category
 - Key transport protocols
 - The shared secret key is created by one party and securely transmitted to the second party.
 - Key agreement (or exchange) protocols
 - Both parties contribute information which is used to derive the shared secret key.
 - Have wide application.

AKE in international standards

- Transport Layer Security (TLS): Handshake Protocol
 - PSK key exchange mode
 - Certificate-based key exchange mode
- IEEE 802.15.6 (WBAN): Authenticated Association
 - Password Authenticated Association
 - Display Authenticated Association
- Bluetooth 5.0: Secure Simple Pairing
 - Numeric Comparison
 - Just Works
 - Out Of Band
 - Passkey Entry

Are those AKA protocols secure enough to underlie the security of communications?

Side channel attacks

- Kerckhoffs' principle
 - A cryptographic system should be secure even if everything about the system, except the **key**, is public knowledge.
- In AKE
 - Long-term keys (pre shared keys, private keys, or password) are kept secretly.
 - Provable security based on certain computational complexity assumptions.



Auguste Kerckhoffs

Side channel attacks can learn information about **long-term keys**.

Side channel attacks

- Why and how SCAs leak long-term keys?
 - Every cryptographic algorithm is ultimately implemented on a physical device that **affects the environment around it in measurable ways**
 - SCAs exploit
 - the time taken by a particular implementation of a cryptographic algorithm,
 - the amount of power consumed,
 - the electromagnetic radiation,
 - ...

2. Method

Leakage resilient methods

- Hardware level
 - Build hardware that leaks as few information as possible
 - E.g., shielding electromagnetic radiation
- Software/Algorithmic
 - Masking/blinding
 - Concealing secret with random data
 - Circuit compilers, inner product extractor, secret sharing
 - Other countermeasures
 - Introducing noise in the side channels
 - Random delays, random order execution, dummy operations

We only introduce the simplest ones here

Additive mask

- is usually used in ECC
- Parameters: n is the group order, P is a base point
- To mask the exponent d of $Q = d \cdot P$
- Generate a small random number k
- Blind d with kn : $d' = d + kn$
- Compute $Q = d' \cdot P$

Corollary of Fermat's little theorem:

Let G be a finite group of order m . Then for $g \in G$ and integer x , it holds that $g^x = g^{[x \bmod m]}$

Exponent splitting

- is a variant of the additive mask
- Represent d with a random k and $d' = d - k$
- To mask the exponent d of $Q = d \cdot P$
- Compute $R_1 = k \cdot P$
- Compute $R_2 = d' \cdot P$
- Compute $Q = R_1 + R_2$

Multiplicative mask

- is a multiplicative analogue of exponent splitting
- Parameters: n is the group order, P is a any group element.
- Represent d with a random k and $d' = k^{-1}d$
- To mask the exponent d of $Q = d \cdot P$
- Compute $R = k \cdot P$
- Compute $Q = d' \cdot R$

Tasks

- Design and realize a leakage-resilient AKE protocol
 - use the method introduced in this lecture; or
 - use your self-proposed method