# Team Notebook

# $\operatorname{UITS-O}(\operatorname{Struggle})$ - University of Information Technology and Sciences

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#### 1 Data Structures

#### 1.1 1D Segment Tree

```
/** 1. segment_tree<long long> seg_tree(a);
 * 2. seg_tree.build(1, 0, n - 1);
 * 3. Note: Make sure that the segment tree type and the
      vector type must mathc. E.g If, struct segment_tree <</pre>
          long long>
   then, vector must be vector<long long>
 * 4. Note: While using index for answer. Make sure to use
       them as (0 based)
 * 5. now, you're good to go.
template <typename T>
struct segment_tree {
 int n;
 vector<T> a. tree:
 /* Used to create the tree array */
 segment_tree (vector<T> cpy) {
   a = cpv:
   n = (int) a.size();
   tree.assign(n << 2, 0);
 /* used to build the tree */
 void build (int node, int 1, int r) {
   if (1 == r) {
     tree[node] = a[1]:
     return:
   int mid = (1 + r) >> 1:
   build(node << 1. l. mid):</pre>
   build((node << 1) + 1, mid + 1, r);
   tree[node] = tree[node << 1] + tree[(node << 1) + 1]:</pre>
 /* point sum or range sum */
 T sum (int node, int start, int end, int 1, int r) {
   if (end < 1 or r < start) {</pre>
     return 0:
   if (1 <= start and end <= r) {</pre>
     return tree[node]:
   int mid = (start + end) >> 1;
   T left sum = sum(node << 1. start. mid. l. r):</pre>
   T right_sum = sum((node << 1) + 1, mid + 1, end, l, r);
   return left_sum + right_sum;
 /* point updating value / adding value */
```

```
void update (int node, int start, int end, int id, T val)
    {
    if (start == end) {
        tree[node] = val;
        return;
    }
    int mid = (start + end) >> 1;
    if (id <= mid) {
        update(node << 1, start, mid, id, val);
    } else {
        update((node << 1) + 1, mid + 1, end, id, val);
    }
    tree[node] = tree[node << 1] + tree[(node << 1) + 1];
}
};</pre>
```

### 1.2 2D Segment Tree

```
/** 1. _2D_segment_tree<long long> _2D_seg_tree(a);
 * 2. _2D_seg_tree.buildx(1, 0, n - 1);
 * 3. Note: Make sure that the segment tree type and the
     vector type must mathc. E.g If, struct
          _2D_segment_tree <long long>
   then, vector must be vector<long long>
 * 4. Note: While using index for answer. Make sure to use
       them as (0 based)
 * 5. now, you're good to go.
template <typename T>
struct _2D_segment_tree {
 int n. m:
 vector<vector<T>> a. t:
 _2D_segment_tree (vector<vector<T>> a) {
   this -> a = a:
   this -> n = (int) a.size();
   this -> m = (int) a[0].size();
   t.assign(n << 2, vector<T>(m << 2)):
 void build v (int vx, int lx, int rx, int vv, int lv, int
   if (ly == ry) {
    if (lx == rx) {
      t[vx][vy] = a[lx][ly];
      t[vx][vv] = t[(vx << 1)][vv] + t[(vx << 1) + 1][vv];
   } else {
    int mv = (lv + rv) >> 1:
     build_v(vx, lx, rx, (vy << 1), ly, my);
```

```
build v(vx, lx, rx, (vv << 1) + 1, mv + 1, rv):
   t[vx][vy] = t[vx][(vy << 1)] + t[vx][(vy << 1) + 1];
/* Prepares the _2D segment tree
* 2D seg tree.build x(1, 0, n - 1): */
void build_x (int vx, int lx, int rx) {
 if (lx != rx) {
   int mx = (lx + rx) >> 1:
   build_x((vx << 1), lx, mx);
   build x((vx << 1) + 1, mx + 1, rx):
 build_y(vx, lx, rx, 1, 0, m - 1);
T sum_y (int vx, int vy, int tly, int try_, int ly, int ry
 if (lv > rv) {
   return (T) 0;
  if (ly == tly && try_ == ry) {
   return t[vx][vv];
 int tmy = (tly + try_) >> 1;
 return sum_y(vx, (vy << 1), tly, tmy, ly, min(ry, tmy))</pre>
    + sum_y(vx, (vy << 1) + 1, tmy + 1, try_, max(ly, tmy
         + 1), ry);
/* Returns the sum of a sub-matrix from.
* [(left_x, left_y) top_left corner] to [(right_x,
     right_y) bottom_right corner]
 * _2D_seg_tree.sum_x(1, 0, n - 1, --left_x, --right_x, --
     left_v, --right_v) -> 0 based indexing */
T sum_x (int vx, int tlx, int trx, int lx, int rx, int ly,
     int ry) {
 if (lx > rx) {
   return 0:
 if (lx == tlx && trx == rx) {
   return sum_v(vx, 1, 0, m - 1, ly, ry);
  int tmx = (tlx + trx) >> 1:
 return sum_x((vx << 1), tlx, tmx, lx, min(rx, tmx), ly,</pre>
    + sum_x((vx << 1) + 1, tmx + 1, trx, max(lx, tmx + 1),
          rx, ly, ry);
void update_v (int vx, int lx, int rx, int vy, int ly, int
     ry, int x, int y, T new_val) {
 if (lv == rv) {
   if (1x == rx) {
```

```
t[vx][vv] = new val:
     } else {
       t[vx][vy] = t[(vx << 1)][vy] + t[(vx << 1) + 1][vy];
   } else {
     int mv = (lv + rv) >> 1:
     if (y <= my) {</pre>
       update_v(vx, lx, rx, (vy << 1), ly, my, x, y, new_val
     } else {
       update_v(vx, lx, rx, (vy << 1) + 1, my + 1, ry, x, y,
     t[vx][vy] = t[vx][(vy << 1)] + t[vx][(vy << 1) + 1];
 /* Updates a particular cell of the matrix - (x_axis,
   * 2D seg tree.update x(1, 0, n - 1, --left x, --left v,
       new val): */
 void update_x (int vx, int lx, int rx, int x, int y, T
      new val) {
   if (lx != rx) {
     int mx = (lx + rx) >> 1;
     if (x \le mx) 
       update_x((vx << 1), lx, mx, x, y, new_val);
       update_x((vx << 1) + 1, mx + 1, rx, x, y, new_val);
   update_y(vx, lx, rx, 1, 0, m - 1, x, y, new_val);
};
```

#### 1.3 Disjoint Set union

```
/* Time complexity for each query: O(logN) */
/* Step 1: disjoint_set<int> dsu(n + 1); */
/* Step 2: dsu.make_set(u, v); */
vector<int> par, siz;
template<typename T>
struct disjoint_set {
  int n;
  T find_set(T v) {
   if (par[v] == v) {
     return v;
  }
  return par[v] = find_set(par[v]);
}
```

```
void init(T v) {
  par[v] = v:
  siz[v] = 1;
disjoint_set (int n) {
  this \rightarrow n = n:
  siz.assign(n + 1, 0);
  par.assign(n + 1, 0);
  for (int u = 1; u <= n; ++u) {</pre>
   init(u):
void make_set(T a, T b) {
  a = find set(a):
  b = find set(b):
  if (a != b) {
    if (siz[a] < siz[b]) {</pre>
      swap(a, b);
    par[b] = a:
    siz[a] += siz[b];
}
T find_group_size(T a) {
  a = find_set(a);
  return siz[a];
```

## 1.4 Lazy Segment Tree

```
/** 1. struct lazy_propagation <int64_t>lazy_prop(a);
 * 2. lazy_prop.build(1, 0, n - 1);
 * 3. now. vou're good to go.
template <typename T>
struct lazy_propagation {
 struct info {
  T sum = 0, prop = 0;
 };
 int n:
 vector<T> a:
 vector<info> tree;
 lazy_propagation (vector<T> cpy) {
   a = cpv:
   n = (int) a.size();
   tree.resize(4 * n);
 void build (int node, int 1, int r) {
```

```
if (1 == r) {
     tree[node].sum = a[1];
     return;
   int mid = (1 + r) >> 1;
   build(node << 1, 1, mid):</pre>
   build((node << 1) + 1, mid + 1, r);
   tree[node].sum = tree[node << 1].sum + tree[(node << 1) +</pre>
         1].sum:
 T sum (int node, int start, int end, int 1, int r, T carry
   if (end < 1 \text{ or } r < start) {
     return 0:
   if (1 <= start and end <= r) {</pre>
     return tree[node].sum + (((end - start) + 1) * carry);
   int mid = (start + end) >> 1:
   T left_sum = sum(node << 1, start, mid, 1, r, carry +</pre>
        tree[node].prop);
   T right_sum = sum((node << 1) + 1, mid + 1, end, 1, r,
        carry + tree[node].prop);
   return left_sum + right_sum;
 void update (int node, int start, int end, int 1, int r, T
   if (end < 1 or r < start) {</pre>
     return:
   if (1 <= start and end <= r) {
     tree[node].sum += ((end - start) + 1) * val;
     tree[node].prop += val;
     return;
   int mid = (start + end) >> 1:
   update(node << 1, start, mid, 1, r, val);</pre>
   update((node << 1) + 1, mid + 1, end, 1, r, val);
   tree[node].sum = tree[node << 1].sum + tree[(node << 1) +</pre>
         1].sum + (((end - start) + 1) * tree[node].prop);
};
```

#### 1.5 Monotonic Stack (Increasing)

```
// C++ code to implement the approach
#include <bits/stdc++.h>
using namespace std;
```

```
// Function to build Monotonic
// increasing stack
void increasingStack(int arr[], int N) {
    // Initialise stack
   stack<int> stk:
   for (int i = 0; i < N; i++) {</pre>
       // Either stack is empty or
       // all bigger nums are popped off
       while (stk.size() > 0 && stk.top() > arr[i]) {
           stk.pop():
       }
       stk.push(arr[i]);
   int N2 = stk.size():
   int ans[N2] = { 0 };
   int i = N2 - 1:
   // Empty Stack
   while (!stk.empty()) {
       ans[i] = stk.top();
       stk.pop();
       j--;
   // Displaying the original array
   cout << "The Array: ";</pre>
   for (int i = 0: i < N: i++) {</pre>
       cout << arr[i] << " ";
   cout << endl:
   // Displaying Monotonic increasing stack
   cout << "The Stack: ":</pre>
   for (int i = 0; i < N2; i++) {</pre>
       cout << ans[i] << " ";
    cout << endl:
// Driver code
int main() {
   int arr[] = { 1, 4, 5, 3, 12, 10 };
   int N = sizeof(arr) / sizeof(arr[0]);
   // Function Call
   increasingStack(arr, N);
```

```
return 0;
```

#### 1.6 PBDS with Deletion in Multiset

```
/* Necessary includes */
#include <ext/pb ds/assoc container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
/* The data structure */
typedef tree<11, null_type, less< 11 >, rb_tree_tag,
    tree order statistics node update> ordered set:
/* Functionalities */
ordered_set s; //declaring pbds
s.insert(x); //taking input
s.find(x)==s.end() // search for a present or not
s.order_of_key(x) //postion of x in sorted set
*s.find by order(r): // value presnt at index r
s1.insert({x.cnt++}): //insert in multiset
s1.erase(s1.lower_bound({x,-1})); //erase in multiset
s1.find_by_order(x)->first //value of index x in multiset
```

## 1.7 Policy Based Data Structure

```
/** 1. Firstly, place the header files and namespace and set
      the data type and comparator.
 * 2. ordered set X:
 * 3. X.insert(8):
 * 4. *X.find_by_order(1)
         Note: finds the kth largest or the kth smallest
      element (counting from zero)
              i.e. The element at the position i (powerful)
 * 5. X.order of kev(3)
        Note: finds the number of items in a set that are
     strictly smaller than our item
              i.e. The position of the current element (
     powerful)
 * Note: This will exactly work like set, multiset, map [
     also can use their functionalities.
    -> Not possible : erasing elements with their value (in
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
```

#### 1.8 Sparse Table

```
/* Used for answering queries, but can answer R(Min/Max)Q in
     0(1) */
/* Step 1: sparse_table<int> st(a, N); */
/* Step 2: cout << st.min_query(1, r) << '\n'; */
template <tvpename T>
struct sparse_table {
 int N:
 int n, k;
 vector<T> a:
 vector<T> logs;
 vector<vector<T>> st;
 void gen_logs () {
   logs[1] = 0:
   for (int i = 2; i <= N; i++) {
     logs[i] = logs[i/2] + 1:
   }
 // builds the table
 void proc () {
   st[0] = a:
   for (int i = 1; i <= k; ++i) {</pre>
     for (int j = 0; j + (1 << i) - 1 < n; ++j) {
       /* Change this line according the question */
       st[i][j] = min(st[i-1][j], st[i-1][j+(1 << (i-1)[j]))
            1))]);
 // builds the structure
 sparse_table (vector<T> a, int N) {
```

```
this -> a = a:
   this \rightarrow N = N:
   n = (int) a.size();
   logs.assign(N + 1, 0);
   gen_logs();
   k = logs[n]:
   st.assign(k + 1, vector < T > (n + 1, 0));
   proc():
  /* This function is used for only min/max query O(1); */
 T min guery (int 1, int r) {
   int p = logs[r - l + 1]:
   return min(st[p][1], st[p][r - (1 << p) + 1]);</pre>
 /* This function is used for the rest of the queries - in
       O(\log(n)) */
 T sum_query (int 1, int r) {
   T sum = 0:
   for (int i = k: i >= 0: --i) {
     if ((1 << i) <= r - 1 + 1) {
       sum += st[i][1];
       1 += 1 << i:
   return sum;
 }
};
```

## 2 Graph Theory

## 2.1 Bellman Ford

```
/* Write this in main function */
/* Time complexity: 0(V * E)
 * In case, E = V^2 then, 0(V^3) */
/* Add edges both ways for undirected graph */
/* g.push_back({u, v, w}); */
/* g.push_back({v, u, w}); */
int cyc_node = -1;
const int inf = (int) 1e9;
vector<int> dist(n + 1, inf);
vector<int> par(n + 1, -1);
auto check_neg_cycle = [&] () {
  if (cyc_node == -1) {
    return printf("No neg cycle\n"), 0;
  }
  int u = cyc_node;
  for (int i = 1; i <= n; ++i) {</pre>
```

```
u = par[u];
 vector<int> path;
 for (int cur = u; ; cur = par[cur]) {
   path.push_back(cur);
   if (1 < (int) path.size() and cur == u) {</pre>
  break:
   }
 reverse(path.begin(), path.end());
 for (auto p : path) {
   printf("%d ", p):
 printf("\n");
 return 0:
auto bellman ford = [&] (int src) {
 dist[src] = 0:
 for (int i = 1: i <= n: ++i) {
   cvc node = -1:
   for (int j = 0; j < m; ++j) {</pre>
  int u = g[j][0];
  int v = g[i][1];
  int w = g[i][2];
  if (dist[u] < inf) {</pre>
   if (dist[u] + w < dist[v]) {</pre>
   dist[v] = max(-inf, dist[u] + w):
   par[v] = u:
   cvc_node = v;
  }
 // remove this if not needed
 check_neg_cycle();
bellman_ford(1);
```

## 2.2 Dijkstras

```
/* Used to extract the shortest path from source(u) to
    destination(v) */
/* Time complexity: O(V + E log V) */
/* Step 1: dijkstra<int> dij(g, n, --src); */
/* Step 2: dij.proc_tab(); */
/* Step 3: cout << dij.get_dist(--v) << '\n'; */
/* Step 4: auto path = dij.get_path(v); */
template <typename T>
struct dijkstra {
```

```
int n:
int src:
// Change this (inf) according to the question
const T inf = (T) 1e16:
vector<int> par, seen;
vector<T> dist:
vector<vector<array<int, 2>>> g;
dijkstra (vector<vector<array<int, 2>>> g, int n, int src)
  this \rightarrow g = g;
  this \rightarrow n = n:
  this -> src = src:
 // Remove this (par) if not needed
  par.assign(n, -1);
 seen.assign(n, false);
 dist.assign(n, inf);
// Processes the distance table
void proc tab () {
 multiset<arrav<T. 2>> ms:
 dist[src] = 0;
  ms.insert({0, src});
  while (!ms.empty()) {
   auto u = *ms.begin();
   ms.erase(ms.begin());
   if (!seen[u[1]]) {
     seen[u[1]] = true;
     for (auto ch : g[u[1]]) {
       if (dist[u[1]] + ch[1] < dist[ch[0]]) {</pre>
         dist[ch[0]] = dist[u[1]] + ch[1]:
         /* Here saving the previous node as parent if
              this is giving less cost */
         par[ch[0]] = u[1];
         ms.insert({dist[ch[0]], ch[0]});
     }
// Returns the shortest distance from source to
     destination
T get_dist (int dest) {
 return dist[dest]:
// Returns the shortest path from source to destination
vector<int> get_path (int dest) {
 vector<int> path;
 for (int v = dest: v != -1: v = par[v]) {
    path.push_back(v + 1);
```

5

```
reverse(path.begin(), path.end());
return path;
};
```

#### 2.3 Floyd Warshall

```
/* Time complexity: O(n ^ 3)
* Computes the all pair shortest paths */
auto floyd_warshall = [&] () {
 /* In order to work with this algorithm, the graph needs
      to be represented in adjacency matrix form. */
 /* init the (d) array, if there doesn't exists a path b/w
      u-v then set them to infinity */
 const int inf = (int) 1e9;
 for (int k = 1; k \le n; ++k) {
   for (int i = 1; i <= n; ++i) {
     for (int j = 1; j <= n; ++j) {
      /* if there exists both of these path or not */
      if (d[i][k] < inf and d[k][i] < inf) {</pre>
        d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
    }
};
```

#### 2.4 Kruskals MST

```
/* Make sure to write this algorithm in main function */
/* Before this, make sure to write the dsu algo */
/* Time complexity: O(mlog(m)) - only for sorting, others
    done in constant time */
disjoint_set<int> dsu(n + 1);
auto kruskals = [&] () {
 long long min_cost = 0;
 /* The edges must be sorted in asc according to their
      weights */
 sort(p.begin(), p.end());
 for (int i = 0; i < m; ++i) {</pre>
   /* p[i][0] = cost, p[i][1] = u, p[i][2] = v; */
   if (dsu.find_set(p[i][1]) != dsu.find_set(p[i][2])) {
     min cost += p[i][0]:
     dsu.make_set(p[i][1], p[i][2]);
 }
 return min_cost;
```

# 2.5 Lowest Common Ancestor

};

```
/* Time complexity: Build O(nlog(n)), Query O(log(n)); */
/* Step 1: binary_lifting bl(n, m, g); */
/* Step 2: bl.lca(u, v); */
/* Step 3: bl.get dist(u, v): */
struct binary_lifting {
 int n. m:
 vector<int> lvl;
 vector<vector<int>> g;
 vector<vector<int>> par;
 void dfs (int v, int l, int p) {
   lvl[v] = 1:
   par[v][0] = p;
   for (auto ch : g[v]) {
    if (ch != p) {
       dfs(ch, 1 + 1, v);
 void init () {
   dfs(1, 0, -1):
   const int logn = __lg(n);
   for (int i = 1; i <= logn; ++i) {</pre>
     for (int j = 1; j <= n; ++j) {
       if (par[j][i - 1] != -1) {
         int p = par[j][i - 1];
         par[j][i] = par[p][i - 1];
 // builds the structure
 binary_lifting (int n, int m, vector<vector<int>> g) {
   this \rightarrow n = n:
   this \rightarrow m = m;
   this \rightarrow g = g;
   lvl.assign(n + 1, 0);
   const int logn = __lg(n);
   par.assign(n + 1, vector<int>(logn + 1, -1));
   init();
 // Returns the lowest common ancestor of two nodes
 int lca (int u. int v) {
   if (lvl[v] < lvl[u]) {</pre>
     swap(v, u);
```

```
int d = lvl[v] - lvl[u]:
   while (d) {
     int logd = __lg(d);
     v = par[v][logd];
     d = (1 << logd);
   if (u == v) {
     return u;
   int logn = __lg(n);
   for (int i = logn; i >= 0; --i) {
     if (par[u][i] != -1 and par[u][i] != par[v][i]) {
       u = par[u][i];
       v = par[v][i];
   return par[u][0];
 // Returns the distance between two nodes
 int get_dist (int u, int v) {
   int com_ances = lca(u, v);
   return lvl[u] + lvl[v] - (lvl[com_ances] << 1);</pre>
};
```

#### 3 Misc

#### 3.1 PBDS and Modular Arithmetic

```
//Policy based data-structure
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;

typedef tree< long long, null_type, less_equal<long long>,
    rb_tree_tag, tree_order_statistics_node_update >
    ordered_set;

//change l1 to any data type
//less_equal for multiset increasing order
//less for set increasing order
//greater_equal for multiset decreasing order
//greater for set decreasing order
//cout<<*X.find_by_order(1)<<endl; // iterator to the k-th
    largest element
//cout<<X.order_of_key(-5)<<endl; // number of items in a
    set that are strictly smaller than our item</pre>
```

```
//Number theory related
const int MOD = 1e9+7:
int gcd ( int a, int b ) { return __gcd ( a, b ); }
int lcm ( int a, int b ) { return a * ( b / gcd ( a, b ) );
inline void normal(int &a) { a %= MOD: (a < 0) && (a += MOD)
inline int modMul(int a, int b) { a %= MOD, b %= MOD; normal
    (a), normal(b); return (a*b)%MOD; }
inline int modAdd(int a, int b) { a %= MOD, b %= MOD; normal
    (a). normal(b): return (a+b)%MOD: }
inline int modSub(int a, int b) { a %= MOD, b %= MOD; normal
    (a), normal(b); a -= b; normal(a); return a; }
inline int modPow(int b, int p) { int r = 1; while(p) { if(p
    &1) r = modMul(r, b); b = modMul(b, b); p >>= 1; }
inline int modInverse(int a) { return modPow(a, MOD-2); }
inline int modDiv(int a, int b) { return modMul(a,
    modInverse(b)): }
```

## 4 Number Theory

### 4.1 Binary Exponentiation

```
/* Time complexity: O(power)
 * Step 1: binary_expo<int>(base, power, mod); */
template<typename T, typename X>
T binary_expo (T val, T power, X m) {
   T output = 1;
   while (power) {
     if (power & 1) {
       output = T((output * 1LL * val) % m);
     }
     val = (val * 1LL * val) % m;
     power >>= 1;
   }
   return output;
}
```

## 4.2 Binomial Coefficients nCr

```
/* Time complexity: log(MOD - 2) */
/* Step 1: bin_coeff<int> bcoef(MAXN, MOD); */
/* Step 2: bcoef.nCr(n, r); */
template <typename T>
struct bin_coeff {
```

```
T n. m:
 vector<T> fact:
 void gen_fact () {
   fact[0] = fact[1] = 1:
   for (int i = 2; i <= n; ++i) {</pre>
     fact[i] = (1LL * fact[i - 1] * i) % m:
 bin coeff (T n. T m) {
   this \rightarrow n = n:
   this \rightarrow m = m:
   fact.resize(n + 1):
   gen_fact();
 T inv (T val. T power) {
   T output = 1;
   while (power) {
     if (power & 1) {
       output = T((output * 1LL * val) % m):
     val = (val * 1LL * val) % m;
     power >>= 1;
   return output;
 T nCr (T N. T R) {
   return (fact[N] * 1LL * inv((fact[R] * 1LL * fact[N - R])
         % m, m - 2)) % m;
}:
```

#### 4.3 Catalan Number

```
long long catalan[n + 1];

// Initialize first two values in table
catalan[0] = catalan[1] = 1;

// Fill entries in catalan[] using recursive formula
for (int i = 2; i <= n; i++) {
   catalan[i] = 0;
   for (int j = 0; j < i; j++) {
      catalan[i] += catalan[j] * catalan[i - j - 1];
   }
}</pre>
```

#### 4.4 Chinese Remainder Theorem

```
struct Congruence {
    long long a, m;
};

long long chinese_remainder_theorem(vector<Congruence> const
    & congruences) {
    long long M = 1;
    for (auto const& congruence : congruences) {
        M *= congruence.m;
}

long long solution = 0;
    for (auto const& congruence : congruences) {
        long long a_i = congruence : congruences) {
        long long M_i = M / congruence.m;
        long long N_i = mod_inv(M_i, congruence.m);
        solution = (solution + a_i * M_i % M * N_i) % M;
}
return solution;
}
```

#### 4.5 Euler Totient (precomputation)

```
/* Time complexity: 0(n*log*log(n)) */
void phi_1_to_n(int n) {
  vector<int> phi(n + 1);
  for (int i = 0; i <= n; i++) {
    phi[i] = i;
  }
  for (int i = 2; i <= n; i++) {
    if (phi[i] == i) {
      for (int j = i; j <= n; j += i) {
        phi[j] -= phi[j] / i;
      }
    }
}</pre>
```

## 4.6 Euler Totient (single)

```
/* Time complexity: O(sqrt(n))
    * Returns phi(n) */
int phi (int n) {
    int result = n;
    for (int i = 2; i * i <= n; i++) {
        if (n % i == 0) {</pre>
```

```
while (n % i == 0) {
    n /= i;
}
    result -= result / i;
}
if (n > 1) {
    result -= result / n;
}
return result;
```

#### 4.7 Extended GCD

```
int gcd(int a, int b, int& x, int& y) {
   if (b == 0) {
      x = 1;
      y = 0;
      return a;
   }
   int x1, y1;
   int d = gcd(b, a % b, x1, y1);
   x = y1;
   y = x1 - y1 * (a / b);
   return d;
}
```

## 4.8 Fermats Primality test

```
/* Time complexity: O(k*log(n))
  * Where, (k) is number of iterations and
  * (n) is the number to check for primality */
template <typename T>
bool fermat (T n, int iter=5) {
  if (n < 4) {
    return n == 2 or n == 3;
  }
  for (int i = 0; i < iter; i++) {
    T a = 2 + rand() % (n - 3);
    if (binary_expo<T>(a, n - 1, n) != 1) {
      return false;
    }
  }
  return true;
}
```

### 4.9 Matrix Exponentiation

```
// C++ program to find value of f(n) where f(n)
// is defined as,
// F(n) = F(n-1) + F(n-2) + F(n-3), n >= 3
// Base Cases :
// F(0) = 0, F(1) = 1, F(2) = 1
// Time Complexity: O(logN)
// Step 1: findNthTerm(n)
// A utility function to multiply two matrices
// a[][] and b[][]. Multiplication result is
// stored back in b[][]
void multiply(int a[3][3], int b[3][3])
   // Creating an auxiliary matrix to store elements
   // of the multiplication matrix
   int mul[3][3]:
   for (int i = 0: i < 3: i++)
       for (int j = 0; j < 3; j++)
           mul[i][j] = 0;
           for (int k = 0; k < 3; k++)
              mul[i][j] += a[i][k]*b[k][j];
   // storing the multiplication result in a[][]
   for (int i=0; i<3; i++)</pre>
       for (int j=0; j<3; j++)</pre>
           a[i][j] = mul[i][j]; // Updating our matrix
// Function to compute F raise to power n-2.
int power(int F[3][3], int n)
   int M[3][3] = \{\{1,1,1\}, \{1,0,0\}, \{0,1,0\}\}:
   // Multiply it with initial values i.e with
   // F(0) = 0, F(1) = 1, F(2) = 1
   if (n==1)
       return F[0][0] + F[0][1];
   power(F, n/2);
   multiply(F, F);
   if (n\%2 != 0)
       multiply(F, M);
```

```
// Multiply it with initial values i.e with
   // F(0) = 0, F(1) = 1, F(2) = 1
   return F[0][0] + F[0][1] ;
// Return n'th term of a series defined using below
// recurrence relation.
// f(n) is defined as
// f(n) = f(n-1) + f(n-2) + f(n-3), n>=3
// Base Cases :
// f(0) = 0, f(1) = 1, f(2) = 1
int findNthTerm(int n)
   int F[3][3] = \{\{1,1,1\}, \{1,0,0\}, \{0,1,0\}\}\};
   //Base cases
   if(n==0)
       return 0:
   if(n==1 || n==2)
      return 1:
   return power(F, n-2);
```

#### 4.10 Mobius Function

```
for (i = 0; i < MU_MAX;i++) {
    mu[i] = 1;
}

for (i = 2; i <= sqroot; i++) {
    if (mu[i] == 1) {
        // for each factor found, swap (+) and (-)
        for (j = i; j <= MU_MAX; j += i) {
            mu[j] *= (-1LL);
        }
        // square factor = 0
        for (j = i * i; j <= MU_MAX; j += i * i) {
            mu[j] = 0;
        }
    }
}</pre>
```

#### 4.11 Trivial nCr

```
/* Used when there's no mod used
  * Time complexity: 0(k)
  * Step 1: ncr_triv<int>(n, r) */
template <typename T>
T ncr_triv (T n, T k) {
  T ncr = 1;
  if (n - k < k) {
    k = n - k;
  }
  for (T i = 0; i < k; ++i) {
    ncr *= (n - i);
    ncr /= (i + 1);
  }
  return ncr;
}</pre>
```

#### 4.12 nCr Table

```
long long ncr[maxn] [maxn] = {0};
const int mod = (int) 1e9 + 7;
void init () {
  ncr[0][0] = 1;
  for (int i = 1; i < maxn; i++) {
    ncr[i][0] = 1;
    for (int j = 1; j < i + 1; j++) {
       ncr[i][j] = (ncr[i - 1][j - 1] + ncr[i - 1][j]) % mod;
    }
}</pre>
```

## 5 Strings

#### 5.1 Aho Corasick

```
const int K = 26;
struct Vertex {
  int next[K];
  bool leaf = false;
  int p = -1;
  char pch;
  int link = -1;
  int go[K];

Vertex(int p=-1, char ch='$') : p(p), pch(ch) {
```

```
fill(begin(next), end(next), -1);
   fill(begin(go), end(go), -1);
 }
}:
vector<Vertex> t(1):
void add_string(string const& s) {
 int v = 0:
 for (char ch : s) {
   int c = ch - 'a';
   if (t[v].next[c] == -1) {
    t[v].next[c] = t.size();
     t.emplace_back(v, ch);
   v = t[v].next[c];
 t[v].leaf = true;
int go(int v, char ch);
int get_link(int v) {
 if (t[v].link == -1) {
   if (v == 0 || t[v].p == 0) {
     t[v].link = 0;
   } else {
     t[v].link = go(get_link(t[v].p), t[v].pch);
 return t[v].link;
int go(int v, char ch) {
 int c = ch - 'a':
 if (t[v].go[c] == -1) {
   if (t[v].next[c] != -1) {
    t[v].go[c] = t[v].next[c];
     t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
 return t[v].go[c];
```

### 5.2 Knuth Morris Pratt

```
/* Total complexity: O(n + m) */
/* Application (string problems) : */
```

```
/* 1. Used to extract matched positions */
/* 2. Used to know if we have a match or not */
/* Here, tab[j] denotes the length of the prefix which
    mathces with the suffix corresponding to index (j) */
/* Step 1: kmp<int> km(full_string, pattern_to_search_for);
/* Step 2: auto ids = km.pos(); */
template <typename T>
struct kmp {
 int n, m;
 string s, t;
 vector<T> tab:
 // Creating the prefix length table
 void proc () {
  int i = 0:
   for (int j = 1; j < m;) {</pre>
     if (t[i] == t[i]) {
       tab[j] = i + 1;
      i += 1, j += 1;
     } else {
      if (i) {
        i = tab[i - 1]:
      } else {
        j += 1:
 // initializing everything
 kmp (string s, string t) {
   this \rightarrow s = s;
   this \rightarrow t = t;
   n = (T) s.size():
   m = (T) t.size();
   tab.assign(m, 0);
   proc();
 // Returns all the starting positions where we have a
 // If we have a match we continue.
 // Otherwise, we look in the previous index of the table
      to save time.
 vector<T> pos () {
   int i = 0;
   int j = 0;
   vector<T> ids:
   while (i < n) {
    if (s[i] == t[j]) {
      i += 1, j += 1;
     } else {
```

```
if (j) {
    j = tab[j - 1];
} else {
    i += 1;
}
}
// If pattern found take the index
if (j == m) {
    ids.push_back(i - m);
    j = tab[j - 1];
}
return ids;
}
```

#### 5.3 Manachers

```
/* Time complexity: O(N) */
/* While solving this problem always try to solve this on
     the basis of the generated answer it is returning */
/* Which means, always try to solve on the basis of
     converted string -> #a#b#a# */
/* Return the length of a palindrome from left side.
    defining (i) as the middle of that palindrome*/
/* manachers<int> man(s): */
/* auto ans = man.ret_ans(); */
template <typename T>
struct manachers {
 int n;
 vector<int> p:
 void manac odd (string s) {
   n = (int) s.size();
   s = "(" + s + ")":
   p.assign(n + 2, 0);
   int 1 = 1, r = 1;
   for (int i = 1: i <= n: ++i) {</pre>
     p[i] = max(0, min(r - i, p[1 + (r - i)]));
     while (s[i - p[i]] == s[i + p[i]]) {
      p[i] += 1;
     if (r < i + p[i]) {
      1 = i - p[i];
      r = i + p[i];
 manachers (string t) {
   string s = "";
```

```
for (auto c : t) {
    s += string("#") + c;
}
manac_odd(s + "#");
}
vector<T> ret_ans () {
    return vector<T>(p.begin() + 1, p.end() - 1);
}
;;
```

#### 5.4 String Hashing (double)

```
#include <bits/stdc++.h>
using namespace std;
const int mod = (int) 1e9 + 7;
int add mod (int a. int b) {
 int res = (a + b) \% mod:
 res += (res < 0 ? mod : 0);
 return res:
int sub mod (int a. int b) {
 int res = (a - b) % mod:
 res += (res < 0 ? mod : 0);
 return res:
int mult mod (int a. int b) {
 int res = (a * 1LL * b) % mod;
 res += (res < 0 ? mod : 0):
 return res:
template<typename T, typename X>
T binary_expo (T val, T power, X m) {
 T output = 1:
 while (power) {
   if (power & 1) {
     output = T((output * 1LL * val) % m);
   val = (val * 1LL * val) % m:
   power >>= 1;
 return output:
int main () {
/* ios::sync_with_stdio(false); */
```

```
/* cin.tie(0): */
string s;
cin >> s:
/* This block of code completely double hashes the string
int p1 = 31, p2 = 53;
int n = (int) s.size():
vector<int> pref_hash1(n);
vector<int> pref hash2(n):
pref hash1[0] = (s[0] - 'a') + 1:
pref_hash2[0] = (s[0] - 'a') + 1;
/* The inverse array is needed to substract the substring'
    s hash */
vector<int> p_pow1(n), inv1(n);
vector<int> p_pow2(n), inv2(n);
p_pow1[0] = inv1[0] = 1;
p pow2[0] = inv2[0] = 1:
for (int i = 1: i < n: i++) {</pre>
 p_pow1[i] = (p_pow1[i - 1] * 1LL * p1) % mod;
 p_pow2[i] = (p_pow2[i - 1] * 1LL * p2) % mod;
 inv1[i] = binary_expo<int>(p_pow1[i], mod - 2, mod);
 inv2[i] = binary_expo<int>(p_pow2[i], mod - 2, mod);
 pref_hash1[i] = add_mod(pref_hash1[i - 1], mult_mod((s[i])
       - 'a' + 1), p_pow1[i]));
 pref_hash2[i] = add_mod(pref_hash2[i - 1], mult_mod((s[i]))
       - 'a' + 1), p pow2[i]));
/* This function returns the hash-1 of the substring of
* Moreover, this function also uses 0 based indesing */
auto substring_hash1 = [&] (int 1, int r) {
 int res = pref_hash1[r];
 if (0 < 1) {
   res -= pref_hash1[1 - 1];
 res = mult_mod(res, inv1[1]);
 return res:
/* This function returns the hash-1 of the substring of
 * Moreover, this function also uses 0 based indesing */
auto substring_hash2 = [&] (int 1, int r) {
 int res = pref hash2[r]:
 if (0 < 1) {
   res -= pref hash2[1 - 1]:
 res = mult_mod(res, inv2[1]);
```

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```
return res;
};

/* This block of code quering for each substring hash*/
int q;
cin >> q;
while (q--) {
  int 1, r;
  cin >> 1 >> r;
  --1, --r;
  cout << substring_hash1(1, r) << '\n';
  cout << substring_hash2(1, r) << '\n';
}
return 0;</pre>
```

#### 5.5 Trie

```
struct trie {
 struct node {
   bool endmark:
/* Change the size to (10) if working with digits */
   node* next[26];
   node () {
     endmark = false;
  /* Change the limit to 10 if working with digits */
     for (int i = 0; i < 26; ++i) {</pre>
      next[i] = NULL;
 } * root;
 /* trie tri: */
 trie () {
   root = new node();
 /* tri.insert(s); */
 /* inserts a string the the trie */
 void insert (string s) {
   node* curr = root;
   for (auto ch : s) {
     /* change 'a' according to the problem statement */
     int id = ch - 'a';
     if (curr -> next[id] == NULL) {
      curr -> next[id] = new node();
     curr = curr -> next[id];
   curr -> endmark = true:
```

```
/* return if a string is present in the list or not */
 /* cout << (tri.search(t) ? "YES\n" : "NO\n") << '\n'; */
 bool search (string s) {
   node* curr = root:
   for (auto ch : s) {
     /* change 'a' according to the problem statement */
     int id = ch - 'a';
     if (curr -> next[id] == NULL) {
      return false:
     curr = curr -> next[id]:
   return curr -> endmark;
 void del_node (node* curr) {
/* Change the limit to 10 if working with digits */
   for (int i = 0; i < 26; ++i) {
     if (curr -> next[i]) {
       del node(curr -> next[i]):
   }
   delete(curr);
 /* tri.del(); */
 /* deletes, all the nodes. Useful in reducing memory */
 void del () {
   del_node(root);
};
```

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