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ELEN4011: DESIGN II

DESIGN OF A ROBUST CODEC FOR A FADING CHANNEL.

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Abstract

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1 Introduction

The aim of modern and next-generation wireless communication systems is to provide communication services with high data rates and low probability of error, this also helps in catering for numerous requests from various applications and devices [1, 2, 3]. The reliability of such communication systems is often hindered by strong shadowing, intersymbol interference (ISI) and attenuation due to the destructive addition of multipaths propagation in the transmission channel [1, 4]. Such a channel is often accurately modeled using a Rayleigh model known as the Rayleigh Fading Channel (RFC) [4]. To combat the effect of fading and scattering due to RFC, several methods have been proposed. Methods such as Selection Diversity, Equal Gain Combining and Maximal Ratio Combining were proposed. However the methods proved to be inefficient and ineffective in dealing with the requirement of high data rates as per modern communication needs [5]. This lead to the development of Multiple-Input Multiple-Output (MIMO) systems. MIMO leverages off the multipath characteristic of the Rayleigh Fading Channel. It transmits data over the multiple paths, therefore increasing the amount of information the communication system carries [6]. MIMO uses multiple transmit and receive antennas to significantly increase the data throughput and link range without additional bandwidth or transmit power [7]. However MIMO cannot achieve all this without robust Forward Error Correction (FEC) and modem schemes.

This paper presents the design of a robust codec for a fading channel. The codec is to operate at a rate of at least 2 bits/Hz over a Rayleigh Fading Channel. The design of the codec focuses mainly on FEC, modulation and MIMO. The codec input data input is expected to be at 10 Mbps. The design assumes that the data stream has already been converted from analog to digital, hence source encoding is neglected. Encryption is also left out as it does not affect the Bit Error Rate (BER) and spectral efficiency. The design also assumes that the channel coefficients of the Rayleigh Fading Channel are known, hence there is no need for channel estimation. The below sections present the mathematical description of the designed codec together with testing, results and analysis. All simulations and computations are carried out in MATLAB.

2 Design Overview

A typical digital communication system constitutes of three main components, those components being the transmitter, the channel and the receiver. Figure 1 illustrates a typical communication system consisting of MIMO architecture.

From figure 1, it can be seen that the receiver mirrors the transmitter in most of the sub-components.

2.1 The Transmitter

On the transmitter side, the source encoder takes in a raw message signal and converts it into a sequence of bits. This is done to compress the raw message data and to remove redundancy. The end result of source encoding is digital information bits with lesser bandwidth and just enough information to reconstruct the original message [8]. After source encoding, the information bits are sent to the channel encoder. The channel encoder uses FEC codes to add redundancy to the information bits. This is done to facilitate error detection and correction on the receiver side after channel transmission. The next step is modulation, the modulator takes in the coded bits from the channel encoder and maps them to a signal to be transmitted over the channel. This is done to conserve power and bandwidth. The MIMO transmitter takes in the symbols



Figure 1: System block diagram of a typical communication system with MIMO architecture

and transmits them over the channel using its multiple antennas. The MIMO transmitter uses multiple antennas to boost bandwidth and to improve signal range as mentioned earlier.

2.2 The Channel

The steps mentioned in section 2.1 are taken so as to prepare the original message to prevail in the channel and so that it is sent as efficiently as possible. The channel can be any medium, wireless or physical, into which the transmitted signal propagates. The channel distorts and adds noise to the transmitted signal, and in some instances interference occurs [8].

2.3 The Receiver

The mirrored side of the transmitter, which is the receiver, performs the opposite of all the steps mentioned in section 2.1 respectively. This done in order to recover the original sent message to the highest accuracy possible.

This paper only focuses on the part of the communication system highlighted in red on figure 1.

3 Channel Encoding

As mentioned earlier, transmitting data over a noisy channel often leads to the information being corrupted. To minimise the occurrence of errors during transmission the information has to somehow be protected. This protection is achieved through either the use of FEC codes or Automatic Repeat Requests (ARQ). ARQ requires retransmission of bits if they are in error. This leads to delays in communication, wasteful retransmissions that require additional power and bandwidth and also interference with other users [9]. All these make ARQ not suitable for real-time applications especially in interactive media applications [10]. FEC on the other hand uses error correcting codes to correct errors that occur during transmission. It does this

by adding redundant bits on the information bits, this helps to detect and correct errors as mentioned earlier. For this design FEC is chosen due to the fact that it offers improved BER, higher throughput and minimal delays compared to its counterpart [10, 11].

FEC codes are split up into two categories: Convolutional Codes and Linear Block Codes. Error correcting capability and/or performance of an FEC coding scheme is measured by its code rate r . The code rate represented by equation 1 below.

$$r = \frac{k}{n} \quad (1)$$

The code rate presents the proportion of the data stream that is useful to the one that has redundancy, represented by k and n respectively in equation 1.

3.1 Convolutional Codes

Convolutional Codes are mostly used in applications that require real time error correction. They are generated by combining the input data stream bits in a series manner. The input bits are stored in a fixed length shift register and are combined with the use of modulo-2 adders [12]. The output of the convolutional encoder at any point is a combination of the previous bits in the shift register.

One type of a Convolutional Codes are Turbo codes. Turbo codes are known to perform extremely well compared to any other FEC code. They achieve within a fraction of a dB of the Shannon Capacity in certain channels [9]. The major disadvantage of Convolutional Codes is their complexity. The complexity curve steepens even further when trying to decode them [9]. The simplest decoder which is the Viterbi also increase in complexity and it gets much difficult to implement for large codewords [13].

3.2 Linear Block Codes

In Linear Block Codes, the resultant codeword is formed as a linear combination of two codewords. The two codewords are the original message bits having size k and the parity bits of size $n - k$, where n is the size of the resulting codeword.

There exist three main types of of Linear Block Codes used in coding theory: Bose Chaudhuri Hochquenghem (BCH), Reed Solomon (RS) and Low Density Parity Check (LDPC) codes. LDPC codes are regarded to be in the same league as Turbo codes with regards to performance. However, like Turbo codes they have major draw backs. LDPC codes require large sums of computation power, have low flexibility and have high encoding complexities [14, 15]. RS codes, derived from BCH codes, are well known for their suitability when it comes to correcting burst errors. Compared with BCH, they were found to perform poorly in RFCs [16, 17]. Due to arguments raised in section 3.1 and in this section BCH codes are chosen as the suitable candidate for this design.

BCH codes form a class of cyclic codes constructed using the principles of the Galois Finite Fields. Unlike RS codes, they are able to correct multiple random errors and unlike LDPC and Turbo codes, they are relatively easier to implement and decoding is also simple [18, 17]. Other great advantages of this coding scheme is its flexibility of choosing the number of errors correctable by the code and its great error correcting capabilities [17, 19, 20].

3.2.1 BCH Encoder and Decoder

BCH(n, k) codes are a generalisation of Hamming Code. Unlike Hamming code, BCH(n, k) posses the ability to correct multiple errors. They are described mathematically as follows:

For any positive integers $m \geq 3$ and $t < 2^{m-1}$ there exists a BCH(n, k) code with the following characteristics:

$$\text{Block length:} \quad n = 2^m - 1 \quad (2)$$

$$\text{Parity check bits:} \quad n - k \leq mt \quad (3)$$

$$\text{Minimum distance:} \quad d_{min} \geq 2t + 1 \quad (4)$$

A BCH(n, k) code described by equations 2–4 is described as a t -error correcting BCH.

To encode, BCH codes uses a generator polynomial $g(x)$ defined in the Galois Field $GF(2^m)$. The polynomial is the lowest degree polynomial over $GF(2)$ which has $\alpha, \alpha^2, \dots, \alpha^{2t}$ as its roots. If $\Phi(x)_i$ is a minimal polynomial of α_i , then $g(x)$ is given by equation 5 below,

$$g(x) = LCM(\Phi(x)_1, \Phi(x)_3, \dots, \Phi(x)_{2t-1}) \quad (5)$$

A more detailed mathematical derivation of how equation 5 is obtained can be found here [21]. The codeword $c(x)$ is formed by first dividing the message polynomial $m(x)$ and obtaining the remainder $r(x)$. The remainder $r(x)$ is then added to original message polynomial $m(x)$ to make the output codeword $c(x)$ as described by equation 6–7.

$$r(x) = \text{mod}\{m(x), g(x)\} \quad (6)$$

$$c(x) = m(x) + r(x) \quad (7)$$

To decode BCH codewords, there a four main steps that are followed:

- i Syndrome computation
- ii Determine coefficients of the error locator polynomial.
- iii Find the roots of the error locator polynomial.
- iv Correct the error

The steps above are described mathematically and in more detail by [14, 22, 23]. In this design BCH(127, 64) code with parameters described in table 1 is used.

Table 1: BCH(127,64) paramaters

Parameter	m	n	k	t	r	$g(x)$
Value	6	127	64	6	0.50	$1 + x + x^3 + x^4 + x^5 + x^7 + x^{10} + x^{13} + x^{15} + x^{16} + x^{19} + x^{20} + x^{21} + x^{22} + x^{23} + x^{26} + x^{29} + x^{33} + x^{34} + x^{35} + x^{39} + x^{40} + x^{42}$

From the table it can be seen that BCH(127, 64) has a code rate of 0.50 as per equation 1 and it can correct up-to 10 errors. The encoder and decoder are both implemented in MATLAB. The choice of BCH(127,64) is further supported in the results section.

4 Digital Modulation

As mentioned in section 2, digital modulation is performed such that bandwidth and power of a communication system are conserved. This also plays a role in minimising the BER. Digital modulation involves impressing information on a sinusoidal carrier, by adjusting its physical characteristics. The physical characteristics being either the amplitude, frequency, phase or a combination [24]. There exist various modulation schemes that include but are not limited to Amplitude Shift Keying (ASK), Frequency Shift Keying (FSK), Phase Shift Keying and Quadrature Amplitude Modulation (QAM).

4.1 Amplitude Shift Keying

In ASK, also known as ON/OFF keying, modulation is achieved by varying the amplitude of the carrier signal. The amplitude varies in accordance with the information bit stream whilst the frequency and phase are kept constant. ASK is mainly used for optical fibre and point-to-point communication [25]. Binary ASK (BASK) is when only two amplitudes are used to represent either a 1 or 0 of the information stream. BASK is simple to implement, however it has very poor bandwidth efficiency and is very susceptible to noise [26].

4.2 Frequency Shift Keying

In this scheme, information bit stream is transmitted through discrete frequency changes of the carrier signal. FSK becomes Binary FSK (BFSK), when only two discrete frequency are used to modulate. It is popularly used in caller ID and remote metering. Like BASK, BFSK is simple to generate and simple to demodulate. However its major draw backs are its poor spectral efficiency and poor BER performance [24].

4.3 Phase Shift Keying

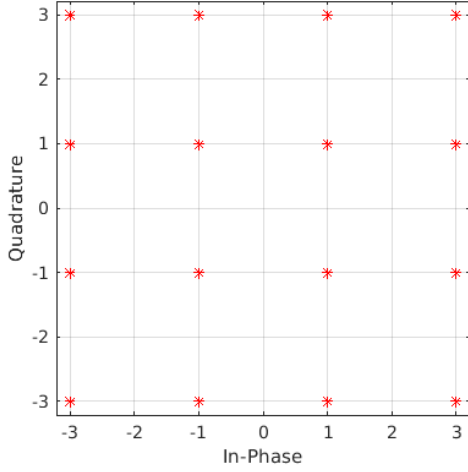
The PSK digital modulation scheme involves encoding the information bits to the carrier signal by varying the phase. There exist a number of popular PSK modulation schemes such as the Binary PSK (BPSK), Differential PSK (DPSK), Quadrature PSK (QPSK) and M-ary PSK (MPSK). Of the four, only the MPSK and QPSK are widely used. QPSK offers higher bandwidth efficiency compared to BASK, BFSK and BPSK, this because it can encode two bits per symbol. The only drawback is its complex antenna design [26]. MPSK on the other hand offers an even better spectral efficiency and greater bandwidth compared to QPSK. This due to the increase in bits per symbol offered by M-ary modulation schemes.

4.4 Quadrature Amplitude Modulation

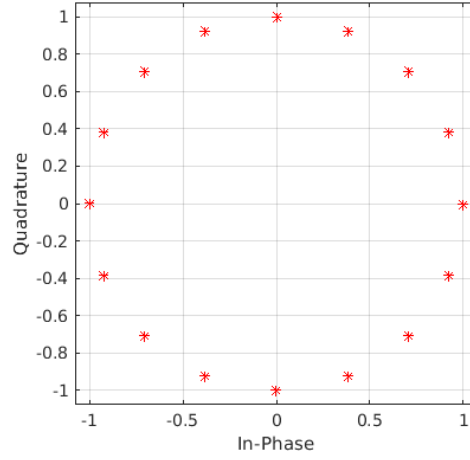
QAM is a modulation scheme which is carried out by impressing information on the amplitude of two carrier signals. Similar to MPSK, it possesses the ability to perform high order modulations. MQAM has advantage over MPSK due to its use of both the phase and amplitude to modulate, this gives it greater distances between constellations [27], this is illustrated by figure 2.

From the figures 2 it can be seen that 16-QAM has greater distances between constellations compared to 16-MPSK.

QAM exploits the fact that sine and cosine of the same frequency have a 90° phase shift. This allows it to send two symbols using the same frequency doubling the symbol rate without doubling the bandwidth [28]. It has also been proven in literature that QAM performs better than MPSK in Rayleigh Fading Channels [27, 29, 30]. Due to arguments raised in sections



(a) Constellation of 16 QAM modulation scheme



(b) Constellation of 16 PSK modulation scheme

Figure 2: 16-Ary Modulation scheme constellations

4.1–4.3 and in this section, QAM modulation is chosen as the suitable modulation scheme for this design.

4.4.1 QAM Modulation and Demodulation

As previously mentioned, QAM modulates by impressing information on both the phase and amplitude. It exploits the fact that $\cos(2\pi f_c t)$ and $\sin(2\pi f_c t)$ are orthogonal. The modulated QAM signal takes the form:

$$s(t) = A_I \sqrt{\frac{2}{T_s}} \cos(2\pi f_c t) + A_Q \sqrt{\frac{2}{T_s}} \sin(2\pi f_c t) \quad (8)$$

for $0 \leq t \leq T_s$ where T_s is time in between symbols

A_I and A_Q are the information bearing discrete amplitude of the two quadrature carriers. Alternatively the symbol waveform can be expressed as:

$$s(t) = E \sqrt{\frac{2}{T_s}} \cos(2\pi f_c t - \theta) \quad (9)$$

Where $E = \sqrt{A_I^2 + A_Q^2}$ and $\theta = \tan^{-1}(\frac{A_Q}{A_I})$ present the location of the symbol on the constellation, determined by the $\log_2(M)$ bits to symbol mapper.

To demodulate, QAM demodulators use a coherent demodulator as illustrated by figure 3.

The distorted received symbol $r(t)$ is passed on to the demodulator as illustrated by figure 3. The symbol is then passed to the I and Q channels where it is mixed with the local generated carriers $\phi_I(t) = \sqrt{\frac{2}{T_s}} \cos(2\pi f_c t)$ and $\phi_Q(t) = \sqrt{\frac{2}{T_s}} \sin(2\pi f_c t)$. The mixed outputs are the passed to the Low Pass Filter (LPF) that rejects all high frequency components. The resulting signal is a baseband one, the baseband signal is then passed to the matched filter to improve the SNR [31]. The symbol demapper is then used to decode the received symbol and match it to the nearest corresponding symbol on the MQAM constellation by calculating the relative distance of the received symbol to all the symbols on the constellation.

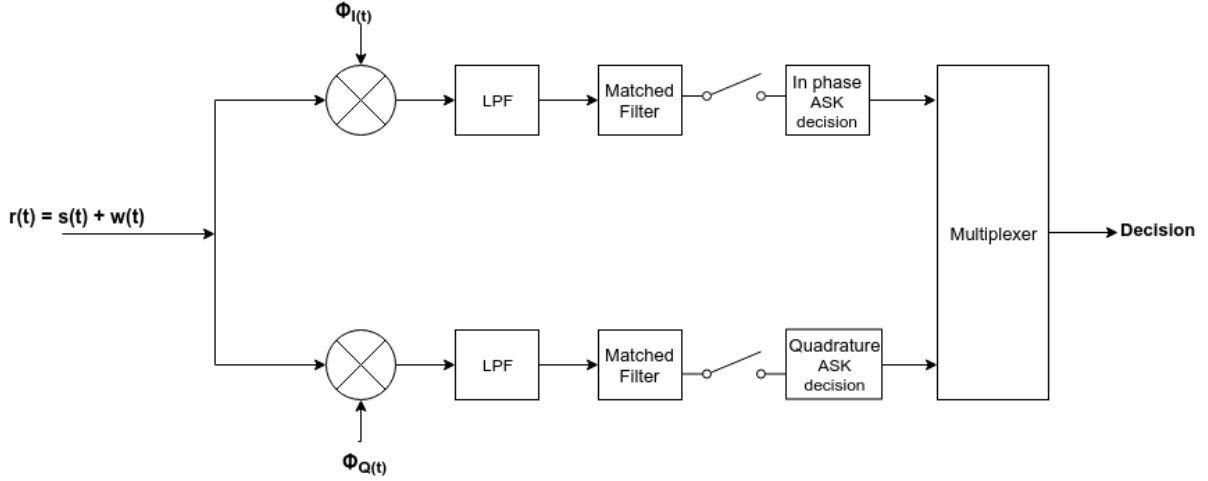


Figure 3: Coherent demodulator for QAM

For this design 16-QAM is chosen. 16-QAM is chosen because it is just high enough to give a good spectral efficiency whilst doing minimal damage on the BER [5]. The decision is also influenced by the desired spectral efficiency of the design. Spectral efficiency is given by equation 10.

$$S_E = \frac{1}{2} r \log_2(M) \alpha_M \quad (10)$$

Where: M is the modulation order, α_M is the MIMO coefficient and r is given by equation 1.

Using the above equation the S_E with the subsystems implemented in section 3 and this section is $1.008\alpha_M$ bits/Hz. With an appropriate value of α_M the desired spectral efficiency can be met. The decision of 16-QAM is further supported in the results section.

5 Channel Model

As previously mentioned a channel adds noise, distorts and attenuates information bearing signals that propagate through it, leading to corruption of data. This design focuses on the two popular channel models, Additive White Gaussian Noise (AWGN) and Rayleigh Fading Channel (RFC).

5.1 AWGN

AWGN is common to every communication channel. It is a statistically random noise characterised by the effect of many random processes that occur in nature [32]. The probability density function of AWGN follows a Gaussian distribution, as the name suggests. Given a sent signal $x(t)$, the received signal $r(t)$ that propagated through an AWGN channel is given by:

$$r(t) = x(t) + n(t) \quad (11)$$

Where $n(t)$ is AWGN noise with zero mean, $E\{n(t)\} = 0$.

The noise term $n(t)$ has power N_0 , assuming that $x(t)$ has average power E_b and the channel has bandwidth W , the SNR of $r(t)$ is give by:

$$SNR = \frac{E_b}{N_0 W} \quad (12)$$

Equation 12 shows that AWGN affects the SNR which in turn affects the BER.

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