

LAST TIME

operational semantics

TODAY

axiomatic semantics / Hoare logic
Floyd-Hoare logic

fact(x)

y = 1

while (x ≠ 1)

y = y * x

x = x - 1

return y

Theorem

for any $x > 0$ $x > 1000$ precondition
fact(x) returns $x!$
postcondition $> x$

Hoare triple

$\{x > 0\}$ $y = \text{fact}(x)$ $\{y = x!\}$
precondition Code/program postcondition

Hoare triple

$\{P\} S \{Q\}$

if P is true initially, and we execute S,
then if S terminates, it is in a state in Q

possible

pre / postcondition

$$P, Q : \text{State} \rightarrow \{\text{true}, \text{false}\}$$

Equivalently $P, Q \subseteq \text{State}$

a state is $\text{Var} \rightarrow \underbrace{\mathbb{Z}}_{\text{int}}$

$$P \wedge Q$$

$$\neg P$$

$$P \vee Q$$

etc.

$$P \Rightarrow Q$$

$$\boxed{\text{skip}} \quad \{P\} \text{ skip } \{P\}$$

$$\boxed{\text{asn}} \quad \{P[x \mapsto a]\} x := a \{P\}$$

Weakest precondition

e.g.

$$\{y \geq 10\} x := y \{ \overset{y}{x} \geq 10 \}$$
$$=$$
$$\{ \text{true} \} x := 5 \{ \overset{5}{x} \geq 5 \}$$
$$=$$
$$\{y \geq 0\} x := y \{x \geq 10\}_x$$
$$\{ \text{false} \} x := \underline{4} \{ \overset{4}{x} \geq 5 \}$$

Sequential
Composition

$$\frac{\{P\} S_1 \{R\} \quad \{R\} S_2 \{Q\}}{\{P\} S_1 ; S_2 \{Q\}}$$

if

$$\frac{\{P \wedge b\} S_1 \{Q\} \quad \{P \wedge \neg b\} S_2 \{Q\}}{\{P\} \text{if } b \text{ then } S_1 \text{ else } S_2 \{Q\}}$$

Weakening
(strengthening)

$$\frac{\{P\} S \{Q\} \quad P' \subseteq P \quad P' \Rightarrow P}{\{P'\} S \{Q\}}$$

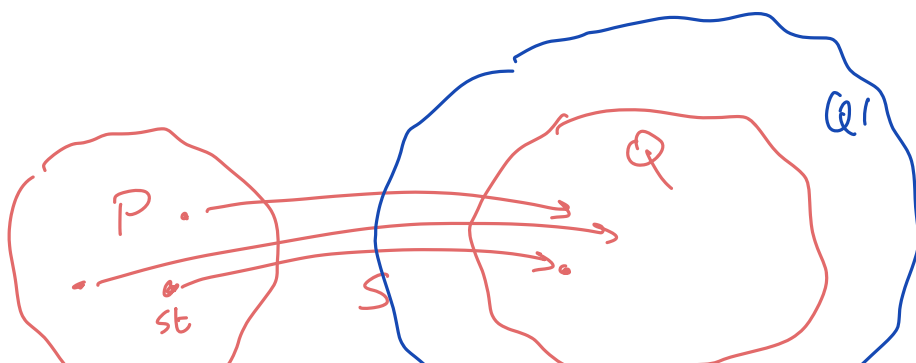
$$\{P'\} S \{Q\}$$

$$Q' \supseteq Q$$

$$\{P\} S \{Q\}$$

$$Q' \Leftarrow Q$$

$$\{P\} S \{Q'\}$$



while

$$\{P \wedge b\} S \{P\}$$

$$\{P\} \text{ while } b \text{ do } S \{P \wedge \neg b\}$$

loop
inductive invariant

$$\vdash \{P\} S \{Q\}$$

$$\models \{P\} S \{Q\}$$

vdash

$$\{m \geq 0 \wedge n > 0\}$$

$$\Gamma := 1$$

$$\{m \geq 0 \wedge n > 0 \wedge \Gamma = 1\}$$

$$i := 0$$

$$\{m \geq 0 \wedge n > 0 \wedge \Gamma = 1 \wedge i = 0\} \Rightarrow \{\Gamma = n^i \wedge i \leq m\}$$

$$\text{while } (i < m)$$

$$\Gamma := \Gamma * n$$

$$i := i + 1$$

$$\{\Gamma = n^i \wedge i \leq m \wedge i \geq m\}$$

$$\{\Gamma = n^m\}$$

$$\{\Gamma = n^i \wedge i \leq m \wedge \overset{\text{loop condition}}{i < m}\}$$

$$\Gamma := \Gamma * n$$

$$\{\Gamma = n^{i+1} \wedge i \leq m \wedge \underbrace{i+1 < m}_{\text{loop condition}} \Rightarrow i \leq m\}$$

$$\underbrace{i}_{x} := \underbrace{i+1}_a$$

$$\{\Gamma = n^i \wedge i \leq m\}$$

$$\{\Gamma = n^{i+1} \wedge \underline{i+1 \leq m}\}$$

III

$$\{\Gamma = n^{i+1} \wedge \underline{i < m}\}$$

$\{n > 0\}$

$y := 1$

$z := 0$

while $\{z \neq n\}$

$z := z + 1$

$y = y * z$

$\{y = z! \wedge z = n\}$

$\{y = n!\}$

$\{y = z!\}$

$y := 1 \quad z := 0$

$y = 0! = 1$ ✓

$y := 1 \quad z := 1$

$y = 1! = 1$

$y := 2 \quad z := 2$

$y = 2! = 2$

$y := 6 \quad z := 3$

$y = 3! = 6$

invariant

$\{n > 0\} y := 1 ; z := 0 \quad \{I\}$

$\{I \wedge z \neq n\} z := z + 1 ; y = y * z \quad \{I\}$

$\{I \wedge z = n\} \text{ skip } \quad \{y = n!\}$

loop inv
↓