

- a) The pivot is chosen by line 1 of RANDOMIZED-PARTITION, through randomly selecting a number from p to r . Therefore, each element is chosen with equal probability, i.e. $1/n$.

$$\mathbb{E}[X_i] = \Pr(i\text{th smallest element is chosen as the pivot}) = 1/n$$

- b) We first discuss the cost for the subcase where $X_q = 1$, i.e. the q th smallest element has been chosen as the pivot. Then, the sorting consists of 3 steps, partitioning, sorting the elements 1 to $q - 1$, and sorting the elements $q + 1$ to n . The cost for partitioning is $\Theta(n)$, the cost for sorting the elements 1 to $q - 1$ is $T(q - 1 - 1 + 1) = T(q - 1)$, and the cost for sorting the elements $q + 1$ to n is $T(n - (q + 1) + 1) = T(n - q)$. Therefore, we have, when $X_q = 1$,

$$T(n) = T(q - 1) + T(n - q) + \Theta(n).$$

We then discuss the relation between the subcases. The events $\{i\text{th smallest element is chosen as the pivot}\}$ for different i 's are disjoint, i.e. if p th smallest element has been chosen as the pivot, then q th smallest element cannot be the pivot, and the probability of their union is 1, i.e. one of them has to happen. Therefore, we have

$$T(n) = \sum_{q=1}^n X_q(T(q - 1) + T(n - q) + \Theta(n)),$$

and therefore

$$\mathbb{E}[T(n)] = \mathbb{E}\left[\sum_{q=1}^n X_q(T(q - 1) + T(n - q) + \Theta(n))\right]$$

- c) $\mathbb{E}[T(n)]$

$$= \mathbb{E}[\sum_{q=1}^n X_q(T(q - 1) + T(n - q) + \Theta(n))]$$

$$= \sum_{q=1}^n \mathbb{E}[X_q(T(q - 1) + T(n - q) + \Theta(n))], \text{ by linearity of } \mathbb{E} \text{ operator}$$

we notice that X_q only depends on choice of this pivot, and $(T(q - 1) + T(n - q) + \Theta(n))$ does not depend on choice of this pivot, and therefore, they are independent, and therefore we have:

$$= \sum_{q=1}^n (\mathbb{E}[X_q] \mathbb{E}[T(q - 1) + T(n - q) + \Theta(n)])$$

$$= \sum_{q=1}^n \left(\frac{1}{n} \mathbb{E}[T(q - 1) + T(n - q) + \Theta(n)]\right)$$

$$= \frac{1}{n} \sum_{q=1}^n (\mathbb{E}[T(q - 1)] + \mathbb{E}[T(n - q)] + \Theta(n)), \text{ by linearity of } \mathbb{E}$$

we notice that the term corresponding to $q = k$ and $q = n + 1 - k$ are identical, just in reverse order, and that if n is odd, then for $q = \frac{n+1}{2}$, $\mathbb{E}[T(q - 1)]$ and $\mathbb{E}[T(n - q)]$ are identical, therefore, by a change in the order of the summation, we have

$$= \frac{1}{n} \sum_{q=1}^n 2\mathbb{E}[T(q - 1)] + \Theta(n)$$

$$= \frac{2}{n} \sum_{q=1}^n \mathbb{E}[T(q - 1)] + \Theta(n)$$

since $T(1) = T(0) = \Theta(1)$ are ignorable, by a change of variable $q \rightarrow q - 1$, we have

$$= \frac{2}{n} \sum_{q=2}^{n-1} \mathbb{E}[T(q)] + \Theta(n)$$

$$\begin{aligned}
\text{d) } & \sum_{k=2}^{n-1} k \lg k \\
&= \sum_{k=1}^{n-1} k \lg k \\
&= \sum_{k=1}^{\lceil \frac{n}{2} \rceil - 1} k \lg k + \sum_{k=\lceil \frac{n}{2} \rceil}^{n-1} k \lg k \\
&\leq \sum_{k=1}^{\lceil \frac{n}{2} \rceil - 1} k \lg \frac{n}{2} + \sum_{k=\lceil \frac{n}{2} \rceil}^{n-1} k \lg n \\
&= \sum_{k=1}^{\lceil \frac{n}{2} \rceil - 1} k(-1 + \lg n) + \sum_{k=\lceil \frac{n}{2} \rceil}^{n-1} k \lg n \\
&= \lg n \sum_{k=1}^{n-1} k - \sum_{k=1}^{\lceil \frac{n}{2} \rceil - 1} k \\
&\leq \frac{n^2 - n}{2} \lg n - \frac{\frac{n}{2}(\frac{n}{2} - 1)}{2} \\
&\leq \frac{1}{2} n^2 \lg n - \frac{1}{8} n^2 + \frac{n}{2} \left(\frac{1}{2} - \lg n \right) \\
&\leq \frac{1}{2} n^2 \lg n - \frac{1}{8} n^2, \text{ since } n \geq 2 \rightarrow \lg n \geq 1 \rightarrow \frac{1}{2} - \lg n < 0 \rightarrow \frac{n}{2} \left(\frac{1}{2} - \lg n \right) < 0
\end{aligned}$$

e) We first notice that

$$\mathbb{E}[T(n)] = \frac{2}{n} \sum_{q=2}^{n-1} (\mathbb{E}[T(q)] + \Theta(n)) = \Theta(n) + \frac{2}{n} \sum_{q=2}^{n-1} \mathbb{E}[T(q)]$$

We assume $\Theta(n)$ in the equation above means $\exists c_1, c_2 \in \mathbb{R}^+, \text{ s.t. } \forall n \geq 2, 0 < c_1 n \leq \Theta(n) \leq c_2 n$.

Note that we can write $\forall n \geq 2$ instead of $\forall n \geq n_0$, for some n_0 potentially larger than 2, since we know from the algorithm that for any $n \geq 2$, the partition step is approximately proportional to n .

We prove $\mathbb{E}[T(n)] \in \Theta(n \lg n)$ in two parts, $\mathbb{E}[T(n)] \in O(n \lg n)$ and $\mathbb{E}[T(n)] \in \Omega(n \lg n)$.

We first prove $\mathbb{E}[T(n)] \in O(n \lg n)$, by finding a , s.t. $\forall n \geq 2, \mathbb{E}[T(n)] \leq a n \lg n$. We prove this by induction.

We first cover the base cases, i.e. select a large enough s.t. $\mathbb{E}[T(2)] \leq a 2 \lg 2$. Note that $\mathbb{E}[T(2)]$ is a constant, which we don't have to explicitly write down, and therefore here a is well-defined.

We then assume $\mathbb{E}[T(k)] \leq a k \lg k$ for all $2 \leq k < n$.

$$\begin{aligned}
\frac{2}{n} \sum_{q=2}^{n-1} \mathbb{E}[T(q)] &\leq \frac{2}{n} \sum_{q=2}^{n-1} a q \lg q \leq \frac{2a}{n} \left(\frac{1}{2} n^2 \lg n - \frac{1}{8} n^2 \right) \leq a n \lg n - \frac{a}{4} n \\
\mathbb{E}[T(n)] &= \Theta(n) + \frac{2}{n} \sum_{q=2}^{n-1} \mathbb{E}[T(q)] \leq \Theta(n) + a n \lg n - \frac{a}{4} n \leq a n \lg n + c_2 n - \frac{a}{4} n \quad (*)
\end{aligned}$$

We let $a \geq 4c_2$, then we have

$$a \geq 4c_2 \rightarrow c_2 n - \frac{a}{4} n \leq 0 \rightarrow (*) \leq a n \lg n$$

In conclusion, we have successfully found $a = \max\left(\frac{\mathbb{E}[T(2)]}{2}, 4c_2\right)$.

We then prove $\mathbb{E}[T(n)] \in \Omega(n \lg n)$, by finding b , s.t. $\forall n \geq 2, \mathbb{E}[T(n)] \geq b n \lg n$.

We first find a lower bound for $\sum_{k=2}^{n-1} k \lg k$.

Since $k \lg k$ is monotonically increasing in k , we have

$$\sum_{k=2}^{n-1} k \lg k \geq \int_1^{n-1} x \lg x \, dx = \frac{1}{2} x^2 \lg x - \frac{1}{4 \lg 2} x^2 \Big|_{x=1}^{x=n-1} > \frac{1}{2} (n-1)^2 \lg(n-1) - \frac{1}{4 \lg 2} (n-1)^2 \quad (**)$$

We notice that $\lg(n) - \lg(n-1) = \lg\left(\frac{n}{n-1}\right) = \lg\left(1 + \frac{1}{n-1}\right) \leq 1$, so we have:

$$(**) \geq \frac{1}{2} (n-1)^2 (\lg(n) - 1) - \frac{(n-1)^2}{4 \lg 2} > \frac{1}{2} n^2 \lg n - n \lg n - \frac{1}{2} (n-1)^2 - \frac{(n-1)^2}{2} > \frac{1}{2} n^2 \lg n - 2n^2$$

We proceed our proof of $\mathbb{E}[T(n)] \geq b n \lg n$, by induction on n .

We first cover the base cases, i.e. select b large enough s.t. $\mathbb{E}[T(2)] \geq b 2 \lg 2$.

We then assume $\mathbb{E}[T(k)] \geq b k \lg k$ for all $2 \leq k < n$.

$$\frac{2}{n} \sum_{q=2}^{n-1} \mathbb{E}[T(q)] \geq \frac{2}{n} \sum_{q=2}^{n-1} b q \lg q \geq \frac{2b}{n} \left(\frac{1}{2} n^2 \lg n - 2n^2 \right) \geq b n \lg n - 4b n$$

$$\mathbb{E}[T(n)] = \Theta(n) + \frac{2}{n} \sum_{q=2}^{n-1} \mathbb{E}[T(q)] \geq \Theta(n) + b n \lg n - 4b n \geq b n \lg n + (c_1 - 4b)n \quad (***)$$

We let $b \leq \frac{c_1}{4}$, then we have

$$(c_1 - 4b)n \geq 0 \rightarrow (***) \geq b n \lg n$$

In conclusion, we have successfully found $b = \min\left(\frac{\mathbb{E}[T(2)]}{2}, \frac{c_1}{4}\right)$