

CYCLES IN $kx + c$ FUNCTIONS

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ABSTRACT. This paper treats cycles in $kx + c$ functions.

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Fundamentals short and sweet

$C_{k,c}$ cycle	We consider the function $f_{k,c}(x)$ given by equation 1. A cycle $C_{k,c}$ is the sequence (v_1, v_2, \dots, v_n) of distinct positive integer, where $f_{k,c}(v_1) = v_2$ and $f_{k,c}(v_2) = v_3$ and so forth and finally $f_{k,c}(v_n) = v_1$.
Prim- itive cycle	If all members of a cycle share a same common divisor greater than one, then this cycle is referred to as a <i>non-primitive</i> cycle, otherwise it is a <i>primitive</i> cycle.
Non- reduced word	Let $C_{k,c}$ be a cycle having U odd and D even members. The non-reduced word describing this cycle is a word of length $U + D$ over the alphabet $\{u, d\}$, which has a u at those positions, where an odd member and a d where an even member is located in the cycle.
Parity vector	Analogously to the non-reduced word, the (binary) parity vector of a cycle $C_{k,c} = (v_1, v_2, \dots, v_n)$ has $n = U + D$ entries. It has a 1 at position i , if v_i is odd, and otherwise 0. We will consider a non-reduced word, for example $uuududuuuddudd$ synonymous with the parity vector $(1, 1, 1, 0, 1, 0, 1, 1, 1, 0, 0, 1, 0, 0)$ or even simpler with the binary sequence (binary word) 11101011100100.
Inter- related cycles	Two cycles are called <i>interrelated</i> if they have the same length and if they both have an equal amount of odd members, which means their non-reduced words contain an equal number of u and v . Analogously their parity vectors have the same number of zeros and ones.

1. INTRODUCTION

Starting point of our considerations is the function:

$$(1) \quad f_{k,c}(x) = \begin{cases} kx+c/2 & 2 \nmid x \\ x/2 & \text{otherwise} \end{cases}$$

Let S be a set containing two elements u and d , which are bijective functions over \mathbb{Q} :

$$(2) \quad u(x) = k \cdot x + c/2 \quad d(x) = x/2$$

Let a binary operation be the left-to-right composition of functions $u \circ d$, where $u \circ d(x) = d(u(x))$. S^* is the monoid, which is freely generated by S . The identity element is the identity function $id_{\mathbb{Q}} = e$. We call e an *empty string*. S^* consists of

all expressions (strings) that can be concatenated from the generators u and d . Every string can be written in precisely one way as product of factors u and d and natural exponents $e_i > 0$:

$$e, u^{e_1}, d^{e_1}, u^{e_1}d^{e_2}, d^{e_1}u^{e_2}, u^{e_1}d^{e_2}u^{e_3}, d^{e_1}u^{e_2}d^{e_3}, \dots$$

These uniquely written products are called *reduced words* over S . Using exponents $u_i, d_i > 0$, we construct strings $s_i = u^{u_i}d^{d_i}$ and concatenate these to a larger string:

$$s_1 s_2 \cdots s_n = u^{u_1}d^{d_1}u^{u_2}d^{d_2} \cdots u^{u_n}d^{d_n}$$

Note that each string s_i is a reduced word, since $u_i, d_i > 0$. Let us evaluate this (large) string by inputting a natural number v_1 . If the result is again v_1 then we obtain a cycle:

$$u^{u_1}d^{d_1}u^{u_2}d^{d_2} \cdots u^{u_n}d^{d_n}(v_1) = d^{d_n}(u^{u_n}(\cdots d^{d_2}(u^{u_2}(d^{d_1}(u^{u_1}(v_1)))))) = v_1$$

We write the sums briefly as $U = u_1 + \dots + u_n$ and $D = d_1 + \dots + d_n$. The cycle contains $U + D$ elements. We summarize this fact to the following definition 1.1:

Definition 1.1. *A cycle consists of $U + D$ elements, where $U = u_1 + \dots + u_n$ is the number of its odd members and $D = d_1 + \dots + d_n$ the number of its even members.*

Moreover we define $A = a_1 + \dots + a_n$ with

$$a_i = 2^{\sum_{j=1}^{i-1} u_j + d_j} \cdot (k^{u_i} - 2^{u_i}) \cdot k^{\sum_{j=i+1}^n u_j}$$

Theorem 1.2 calculates the smallest member of the cycle $C_{k,c}$, which in line with definition 1.1 consists of U odd and D even members [1]:

Theorem 1.2. *The smallest number v_1 belonging to a cycle $C_{k,c}$ having U odd and D even members is:*

$$v_1 = \frac{c \cdot A}{(k-2)(2^{U+D} - k^U)}$$

Example 1.3. *We consider a cycle $C_{3,11}$ that has $U + D = 8 + 6 = 14$ elements and choose $(u_1, u_2, u_3, u_4) = (3, 1, 3, 1)$ and $(d_1, d_2, d_3, d_4) = (1, 1, 2, 2)$. Its smallest element is $v_1 = 13$ and we obtain all elements by evaluating the strings: $v_2 = u(v_1)$, $v_3 = u(v_2)$, $v_4 = u(v_3)$ and $v_5 = d(v_4)$ and so forth. It applies:*

$$uuud \circ ud \circ uuudd \circ udd(v_1) = u^3d \circ ud \circ u^3d^2 \circ ud^2(v_1) = s_1 \circ s_2 \circ s_3 \circ s_4(v_1) = v_1$$

This cycle is $(v_1, v_2, v_3, \dots, v_{14}) = (13, 25, 43, 70, 35, 58, 29, 49, 79, 124, 62, 31, 52, 26)$. We calculate v_1 directly as follows:

$$v_1 = \frac{11 \cdot 11609}{(3-2)(2^{8+6} - 3^8)} = \frac{11 \cdot 11609}{9823} = 13$$

In this case $11609 = A = a_1 + a_2 + a_3 + a_4 = 4617 + 1296 + 3648 + 2048$:

$$\begin{array}{llll} a_1 = 2^0 & (3^3 - 2^3) & 3^{1+3+1} & = 4617 \\ a_2 = 2^{3+1} & (3^1 - 2^1) & 3^{3+1} & = 1296 \\ a_3 = 2^{3+1+1+1} & (3^3 - 2^3) & 3^1 & = 3648 \\ a_4 = 2^{3+1+1+1+3+2} & (3^1 - 2^1) & 3^0 & = 2048 \end{array}$$

2. CONDITIONS FOR CYCLES

A positive integer k is called a *Crandall number*, if there exists a cycle $C_{k,1}$ and the following very fundamental theorem 2.1 is well known, see [2], [3]:

Theorem 2.1. *Every Wieferich number is a Crandall number. In other words, if k is a Wieferich number, then a cycle $C_{k,1}$ exists.*

In conformity with definition 1.1, let us consider a cycle $C_{k,c}$ consisting of U odd integers and D even integers. The theorem 2.2 specifies the following cycle restriction:

Theorem 2.2. *A cycle only exists if the inequality $2^{U+D} - k^U > 0$ holds.*

The following theorem details the condition for the existence of a cycle [4]:

Theorem 2.3. *A cycle $C_{k,c}$ only exists if the integer c and the difference $2^{U+D} - k^U$ are not coprime: $\gcd(c, 2^{U+D} - k^U) > 1$.*

For a sequence (v_1, v_2, \dots, v_n) of numbers, let us define a (binary) parity vector consisting of n elements, which has a 1 at position i , if v_i is odd, and otherwise 0. This vector corresponds to the non-reduced word over the alphabet $\{u, d\}$ as introduced in section 1. This word has a u at each position which in the vector is occupied by a 1, and it has a v at a position at which in the vector is a 0.

Let $0 \leq x_1 < x_2 < \dots < x_U \leq U-1$ be all positions (the indexing is zero-based) in the parity vector occupied by 1 or equivalently all positions in the word at which there is a u . We can detail theorem 2.3 as follows by theorem 2.4:

Theorem 2.4. *A cycle $C_{k,c}$ only exists if the divisibility $2^{U+D} - k^U \mid c \cdot z$ holds, where $z = \sum_{i=1}^U 3^{U-i} 2^{x_i}$.*

Example 2.5. We refer to $C_{3,11} = (13, 25, 43, 70, 35, 58, 29, 49, 79, 124, 62, 31, 52, 26)$ again. The corresponding parity vector is $(1, 1, 1, 0, 1, 0, 1, 1, 1, 0, 0, 1, 0, 0)$ and the non-reduced word is $uuuduuuuddudd$. The indices are $(x_1, \dots, x_8) = (0, 1, 2, 4, 6, 7, 8, 11)$ and therefore $z = 3^7 2^0 + 3^6 2^1 + 3^5 2^2 + 3^4 2^4 + 3^3 2^6 + 3^2 2^7 + 3^1 2^8 + 3^0 2^{11} = 11609$.

Correctly it applies that $2^{8+6} - 3^8 \mid 11 \cdot 11609$, more specifically it is $9.823 \mid 127.699$ and $9.823 \cdot 13 = 127.699$.

Another restriction given by theorem 2.6:

Theorem 2.6. The number of cycles $C_{k,c}$ is alsway less than or equal to the number of cycles $C_{k,a \cdot c}$ where a is an odd number.

3. BOUNDARY FEATURES OF CYCLES

Here the text will be insertet here: Characteristics M and N (based on Halbeisen / Hungerbühler) and Darrels existing findings and our new findings

4. CONSTRUCTING ONE CYCLE FROM ANOTHER

text about interrelation of cycles comes here

11609,22325,38399,62510,31255,51794,25897,43757,70547,110732,55366,27683,46436,23218

And the reduced form (divided by $19 \cdot 47$):

13,25,43,70,35,58,29,49,79,124,62,31,52,26

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