

COMPILER DESIGN

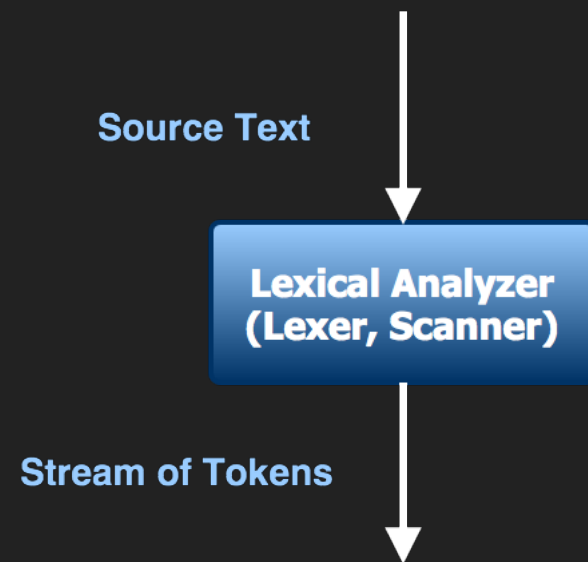
AND IMPLEMENTATION

ÁRPÁD GORETITY
BUDAPEST SWIFT MEETUP

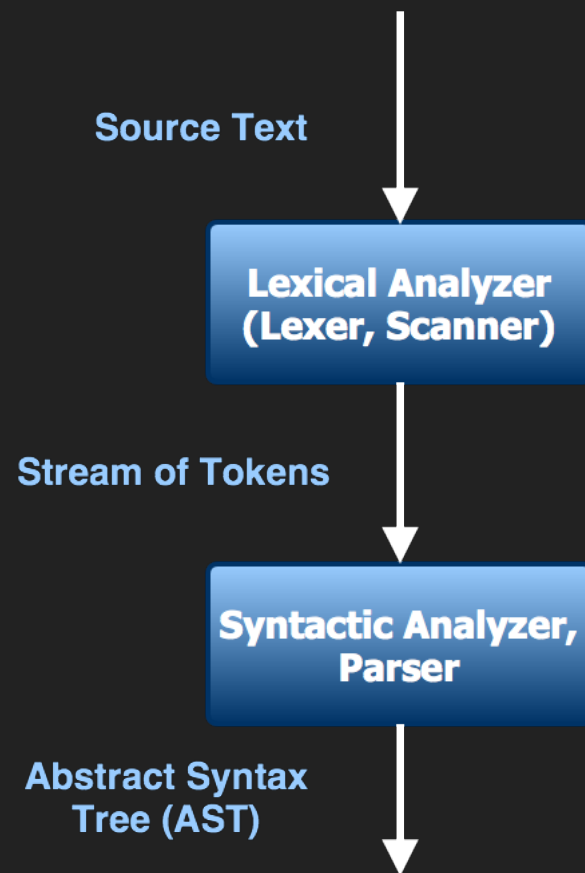
SEMANTIC ANALYSIS

PART 3

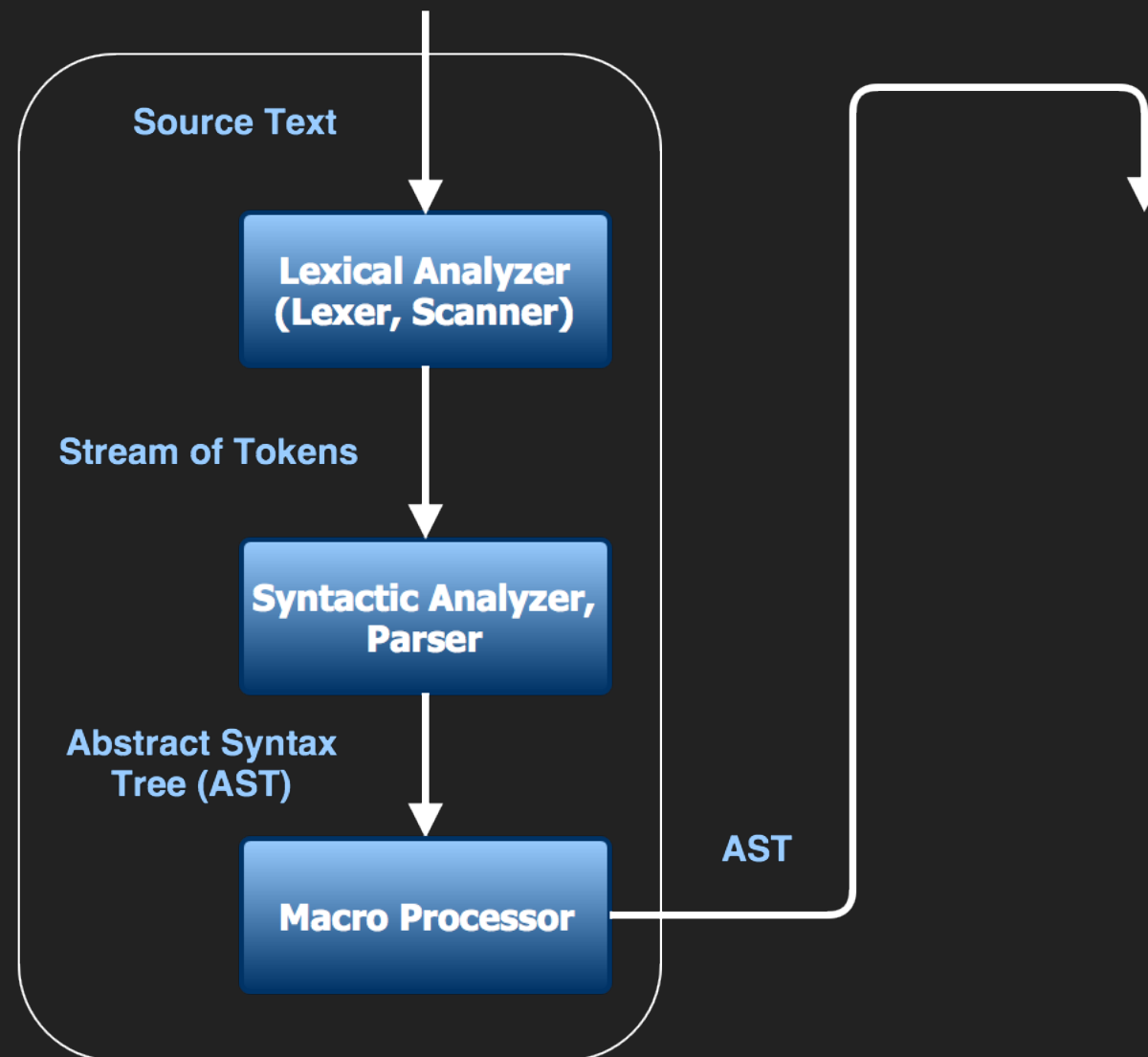
REMINDER



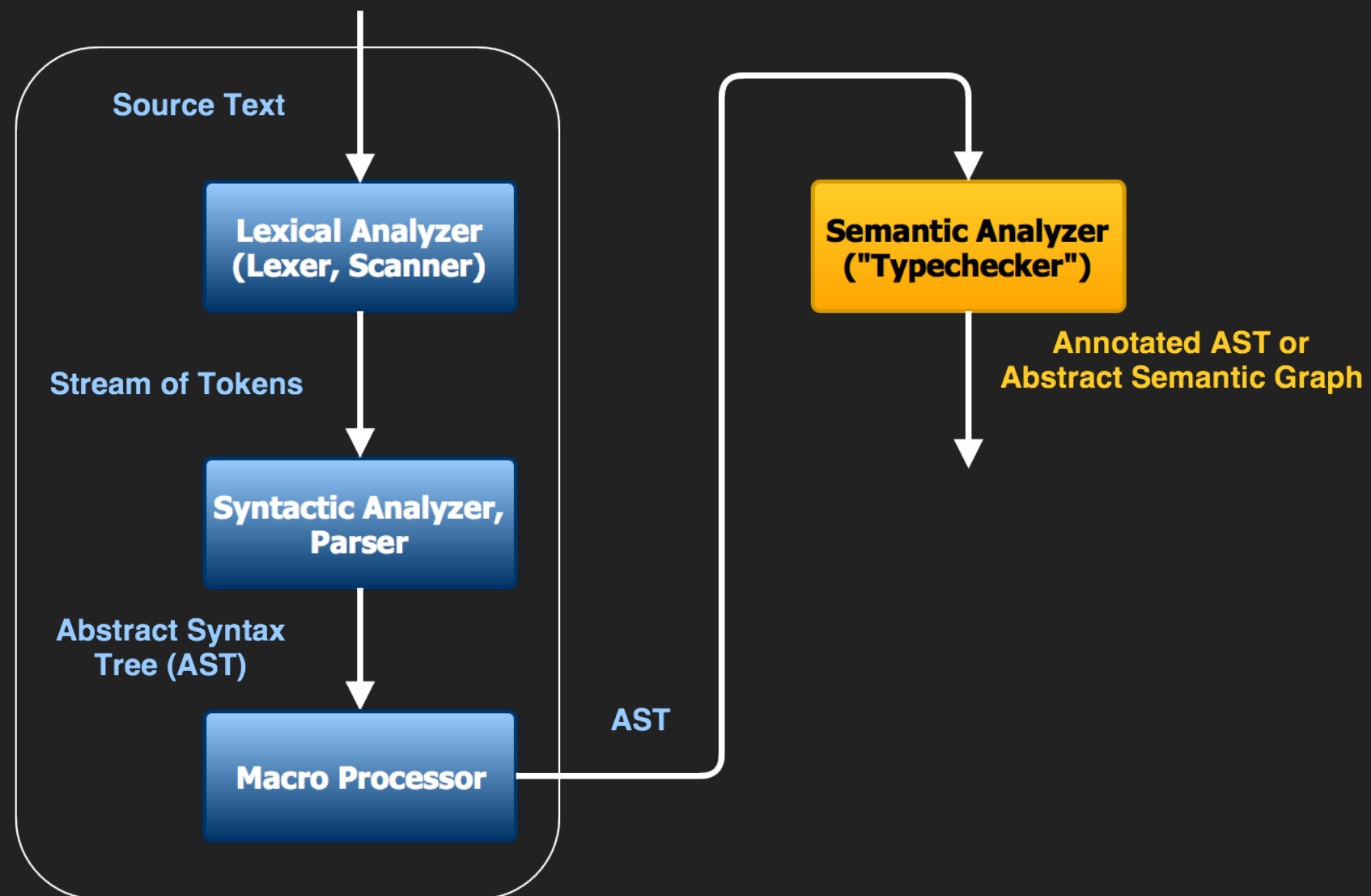
REMINDER



REMINDER



REMINDER



WHY DO WE NEED A SEMANTIC ANALYZER?

- ▶ Correct Syntax \neq Correct Program
- ▶ “Colorless green ideas sleep furiously” (Chomsky)
- ▶ Syntax: OK – verb, subject, adverbs in the right place
- ▶ Meaning: not OK – complete nonsense

WHY DO WE NEED A SEMANTIC ANALYZER?

- ▶ Reminder: **Context-Free** Grammars
- ▶ Parser does not care about the meaning of symbols
- ▶ But we do! That is **the** point of programming...
- ▶ We need something to check correctness of meaning

KINDS OF SEMANTIC ANALYSIS

- ▶ Type Checking and/or Type Inference
- ▶ Scope Resolution
- ▶ Data Flow Analysis (DFA)
- ▶ Control Flow Analysis (CFA)
- ▶ ...and possibly many more...

STEPS OF SEMANTIC ANALYSIS

- ▶ These steps are strongly interrelated
- ▶ Type checking and scope resolution: Typing Context
- ▶ Scope resolution and DFA:
 - ▶ Which constants and variables are accessible where?
 - ▶ When are pointers valid? – Lifetimes (Rust's affine types)
- ▶ DFA and CFA: which computation produces which value?
Warnings: unused variable (DFA), unreachable code (CFA)

THEORETICAL BACKGROUND

- ▶ Formal logic
 - ▶ Lambda calculus
- ▶ Formal semantics: Denotational, Operational
- ▶ Type theory
 - ▶ Curry-Howard correspondence
 - ▶ Formal Type Systems (e.g. Hindley-Milner)
- ▶ DFA, CFA: Graph theory and algorithms (DFG, CFG)

THEORETICAL BACKGROUND

- ▶ Each of these subjects may span multiple semesters if you do a Computer Science degree at university...
- ▶ This is **not** an attempt to equip you with such detailed knowledge
- ▶ Focus is on some basic implementation techniques
 - ▶ Scope Resolution and Type Checking – **inevitable**
 - ▶ Type Inference – **local** inference is easy

THEORETICAL BACKGROUND

- ▶ DFA, CFA: lengthy, sometimes hard
 - ▶ So not today
 - ▶ **Except:** loop control flow statements (`break`, `continue`)
 - ▶ These are trivial to check for basic correctness

REPRESENTATION OF SEMANTICS – AN EXAMPLE

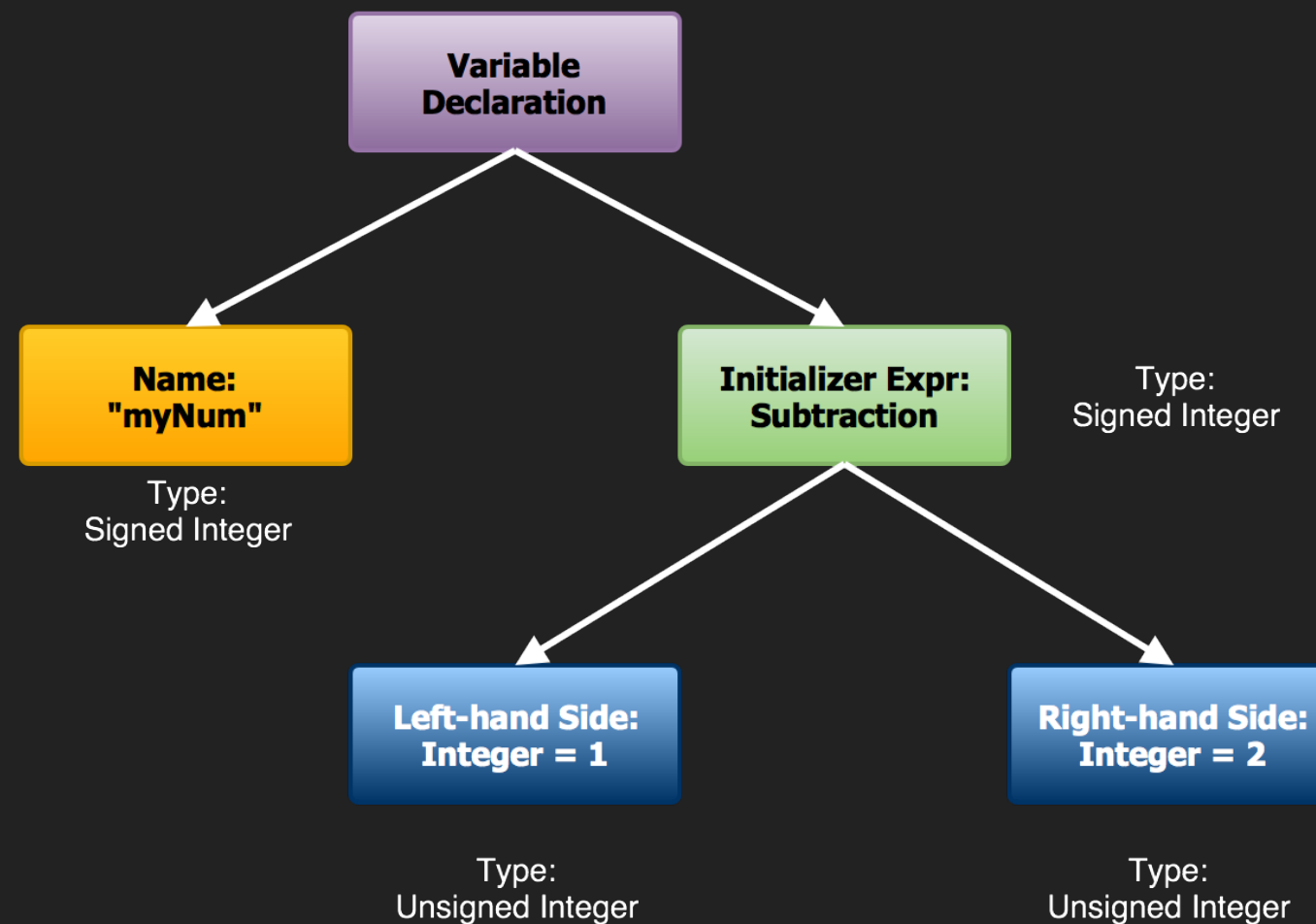
► Scopes

- Stack of sets containing variable names (or indices)
- Search starts at top of stack, going downwards

```
{                                     // []
  var a = 42                          // [ ('a') ]
  {
    var b                            // [ ('a'), ('b') ]
    var c                            // [ ('a'), ('b'), ('c') ]
  }                                  // [ ('a') ]
}                                    // []
```

REPRESENTATION OF SEMANTICS – AN EXAMPLE

- **Types:** Type annotations on AST nodes



LOCAL TYPE INFERENCE

- ▶ Relatively easy: similar to naive code generation
- ▶ Only needs e.g. recursive traversal of the AST
- ▶ Bottom-up: types propagate from leaves (children) towards the root (parent)
 - ▶ e.g. basic mathematical operators, return value of function calls
- ▶ Top-down: types propagate from parents towards children
 - ▶ e.g.: argument types of a function literal passed to a higher-order function

LET'S GET TO WORK!